

Weights Vector  $[1, \dots, \dots]$   
Length  $l$

Series:  $[ \dots ]$   
Length  $n$

Transformation:

$$T: \mathbb{R}^l \Rightarrow \mathbb{R}^{n-l}$$

↑  
weights

$$\overline{X}_t = \sum_{k=1}^l \underbrace{w_k}_{l} \underbrace{X_{t-k}}_l$$

Python Indices approximately:  
 $i-l$  to  $i$  (index factor -1) w/o

$\tau, k$  inequality:

$$\tau > |w_k| = \left| \prod_{i=0}^{k-1} d-i \left( \frac{1}{k!} \right) \right|$$

By inspection of recursive formula.

Thus,  $k = k(d, \tau)$ .

Fix  $\tau \Rightarrow$  Find  $k$ .

Speed:  $O(n) \Rightarrow \prod_{i=0}^{k-1} \text{loop}$

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Recall: Indices iterated over for  $\bar{X}_t$ :

$i$  - width  $\rightarrow i$

$i \mapsto [\text{length}(\text{weights}), \text{length}(\text{series})]$

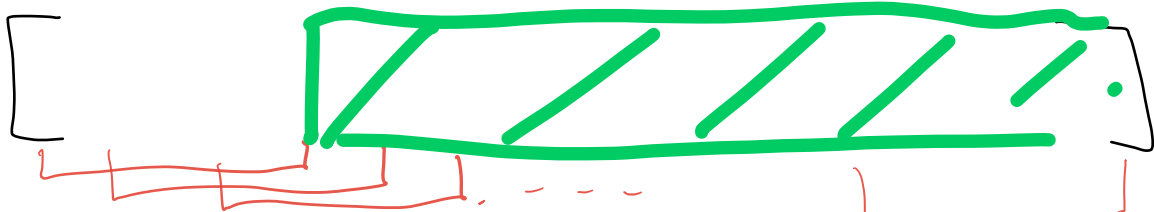
i.e. Index Bounds for Dot Prod  $\bar{X}_t$ :

$l = \text{length}(\text{weights})$ ,  $n = \text{length}(\text{series})$

$[l-l, l]$

$[l+3-l, l+3]$

$[n-l, n]$

Series: 

window size  $l$  (how many  $l$ -sized windows in  $n$ -sized Series)

 =  $\tilde{x}_t$  Series (transformed)

Forget about  $\pm 1$  python index,

In general: To drop  $k$  data pts,  
keep  $n-k$  data pts.  
( $k = l = \text{length}(\text{weights})$ ).

Recall  $\tau = \tau(k, d)$ :

```
def calc_tau(k, d):
```

```
    val = 1
```

```
    for i in range(0, k):
```

```
        val *= (d-i)
```

```
    val /= k!
```

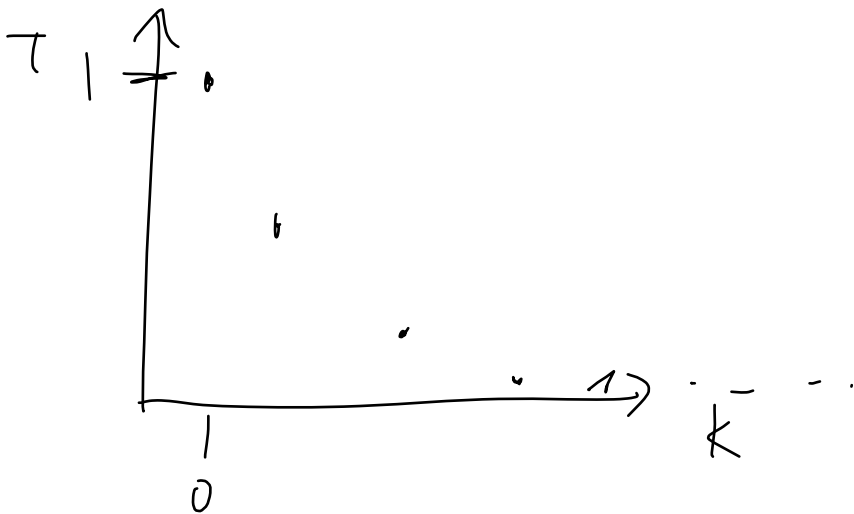
```
    return val.
```

☆ - See  
code  
file/  
Test

$\Rightarrow$  Returns  $\tau$  acceptable for max  $k$  bars.

Note: If  $k$  too big, for  $\tau$  is too large for python integer!

Effects of  $k$  vs  $\tau$ :



Variable Computation:

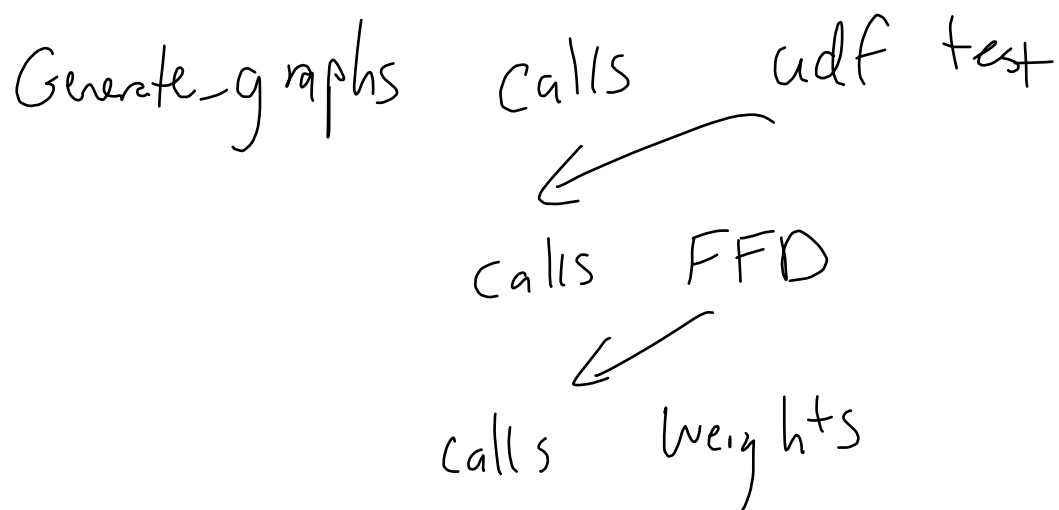
Fix  $k \rightarrow$  thrsh  $\rightarrow$  changes adfuller test?  
informs  $d \leftarrow$  ??

From previous trials: **Trend:**

**$d \uparrow \rightarrow k \downarrow$**

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Functional Process:



In our inputs:

Fix range of  $d = [0, 0.1, 1]$

(+) Fix a  $k$ ,

→ range of thresh ( $d \xrightarrow{\text{isomorphism}} \text{unique } \tau$ )

For adfuller fn:  $\text{FractDiff} \rightarrow \text{correlation} \rightarrow \text{pick best}$   
pVal

For a fixed  $d, k$ :

$\mapsto \tau$

$\rightarrow ad_{fuller} \rightarrow \text{check } pVal.$

Simply, it is  $\tau$  vs  $d$  to  
find optimal  $d$ .

Previously,  $d$  vs  $\tau$ :

$\tau(k)$	0.01	0.001	0.0001
optimal	0.5	0.3	0.2
k-drops	12	73	503

$ADF < 0.05$  : Statram

Note: do not use  $\tau$  too small:

Zero vector breaks code!

i.e.  $k$  small  $\rightarrow \tau$  big

$\rightarrow$  less drops  $\rightarrow d_{\text{optimal}}$  bigger



























