# Attempts to Quantumly Solve Standard Lattice Problems: Reduction from Standard Lattice Problems to $S|LWE\rangle$ and Beyond

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## 1 Introduction

In this note, we summarize our partial results on quantumly solving standard lattice problems.

Solving standard lattice problems has been a target for designing efficient quantum algorithms for decades. Regev [Reg09] shows given a polynomial time algorithm that solves LWE<sub> $n,m,q,D_{noise}$ </sub> where  $\mathcal{D}_{noise}$  is Gaussian and m can be any polynomial, one can construct a quantum algorithm that solves standard lattice problems.

Here let us consider the following quantum variant of the LWE problem called solving LWE given LWE-like states (S|LWE).

**Definition 1** (Solving LWE given LWE-like quantum states (S|LWE))). Let n, m, q be positive integers. Let f be a function from  $\mathbb{Z}_q$  to  $\mathbb{R}$ . Let  $u \in \mathbb{Z}_q^n$  be a secret vector. The problem of solving LWE given LWE-like states  $S|LWE\rangle_{n,m,q,f}$  asks to find u given access to an oracle that outputs  $a_i$ ,  $\sum_{e_i \in \mathbb{Z}_q} f(e_i)|a_i \cdot u + e_i \pmod{q}\rangle$  on its  $i^{th}$  query, for i = 1, ..., m. Here each  $a_i$  is a uniformly random vector in  $\mathbb{Z}_q^n$ .

 $\mathsf{S}|\mathsf{LWE}\rangle_{n,m,q,\sqrt{D_{\mathsf{noise}}}}$  is easier to solve than  $\mathsf{LWE}_{n,m,q,D_{\mathsf{noise}}}$ , because we can get (classical) LWE samples by measuring  $|\mathsf{LWE}\rangle$  in computational basis. Recent work [CLZ21] shows when the noise amplitude f is of a special kind, we can solve  $\mathsf{S}|\mathsf{LWE}\rangle$  in quantum polynomial time.

**Theorem 2** ([CLZ21]). When the noise distribution f is chosen such that  $\hat{f}$  is non-negligible over  $\mathbb{Z}_q$ , then we can solve  $\mathsf{S}|\mathsf{LWE}\rangle_{n,m,q,f}$  in quantum polynomial time.

Given the 'feasibility' of solving  $S|LWE\rangle$ , one plausible roadmap towards solving standard lattice problems is first to modify Regev's reduction (from standard lattice problems to LWE) to a reduction from standard lattice problems to  $S|LWE\rangle$ , and then solve the  $S|LWE\rangle$  problem. The key point is that the noise amplitude f in  $S|LWE\rangle$  should on one hand be 'strong' enough so that the  $S|LWE\rangle$  oracle can solve standard lattice problems, but on the other hand be 'weak' enough so that the  $S|LWE\rangle$  problem is solvable by polynomial quantum algorithms.

# 2 Quantum reduction from Standard Lattice Problems to S|LWE>

In this section, we'll show how to obtain a quantum reduction from standard lattice problems to  $S|LWE\rangle$ , by modifying Regev's reduction.

## 2.1 Summary of Regev's reduction [Reg09]

Let's start by recalling the details of Regev's reduction. Many standard lattice problems can be reduced to sampling from the discrete Gaussian distribution  $(D_{L,r})$  of a nontrivial width r over the lattice L. With the help of an LWE solver, one can construct a procedure sampling from  $D_{L,r}$  given samples from  $D_{L,r\cdot c}$  with c > 1, and hence can start with samples from extremely wide  $D_{L,R}$  (which can be obtained through, say, LLL-algorithm) and end up with samples from  $D_{L,r}$  with a nontrivial (say, polynomial) width r. The precise procedure contains two subroutines:

Step 1 (Classical, uses LWE) Given an instance of  $\mathsf{CVP}_{L^*, \alpha q/(\sqrt{2}r)}$ , using  $\mathsf{poly}(n)$  samples from  $D_{L,r}$  to create LWE samples with Gaussian noise with width  $\leq \alpha q$ , and then solve it with an LWE solver which in turn solves the  $\mathsf{CVP}_{L^*, \alpha q/(\sqrt{2}r)}$  problem:

**Theorem 3** ([Reg09]). Suppose  $m \in \text{poly}(n)$ , q be an integer,  $\alpha \in (0,1)$  be a real number and  $r > \sqrt{2}q\eta_{\epsilon}(L)$  satisfying some smoothing condition with  $\epsilon \in \text{negl}(n)$ . There exists an efficient (classical) algorithm that, given an oracle that solves  $\mathsf{LWE}_{n,m,q,q\Psi_{\alpha}}$  and  $\mathsf{poly}(n,m)$  samples from  $D_{L,r}$ , solves  $\mathsf{CVP}_{L^*,\alpha q/(\sqrt{2}r)}$ , where  $\Psi_{\alpha}$  denotes the periodic Gaussian distribution and  $q\Psi_{\alpha}$  stands for scaling it by q.

Step 2 (Quantum) Using a  $\mathsf{CVP}_{L^*, \alpha q/(\sqrt{2}r)}$  solver to generate  $\mathsf{poly}(n)$  discrete Gaussian states  $|D_{L, r \cdot \sqrt{n}/(\alpha q)}\rangle = \sum_{\mathbf{v} \in L} \sqrt{\rho_{r \cdot \sqrt{n}/(\alpha q)}(\mathbf{v})} |\mathbf{v}\rangle$  and measure them to get  $\mathsf{poly}(n)$  classical samples from  $D_{L, r \cdot \sqrt{n}/\alpha q}$ :

**Theorem 4** ([Reg09]). There exists an efficient quantum algorithm that, given any n-dimensional lattice L, a number  $d < \lambda_1(L^*)/2$ , and an oracle that solves  $\mathsf{CVP}_{L^*,d}$ , outputs  $|D_{L,\sqrt{n}/(\sqrt{2}d)}\rangle$ .

These two subroutines allow us to transform the distribution  $D_{L,r}$  to a narrower distribution  $D_{L,r\cdot\sqrt{n}/(\alpha q)}$ , and hence solve the discrete Gaussian sampling problem whenever  $\alpha q/\sqrt{n} > 1$ .

## 2.2 Modifying Regev's reduction

Notice that the quantum part of the iterative algorithm actually produces discrete Gaussian states instead of just classical samples. This gives us hope to construct a procedure sampling  $|D_{L,r}\rangle$  states, given  $|D_{L,r}\rangle\langle (c>1)$  states and an S|LWE $\rangle$  solver. The procedure is as follows:

- Step 1 (Uses  $\mathsf{S}|\mathsf{LWE}\rangle$ ) Given an instance of  $\mathsf{CVP}_{L^*,\alpha q/r}$ , using  $\mathsf{poly}(n)$  discrete Gaussian states  $|D_{L,r}\rangle$  to create an  $\mathsf{S}|\mathsf{LWE}\rangle_{n,m,q,f}$  instance with certain f, and then solve it with an  $\mathsf{S}|\mathsf{LWE}\rangle_{n,m,q,f}$  solver which in turn solves the  $\mathsf{CVP}_{L^*,\alpha q/r}$  problem;
- Step 2 (Same as the quantum step in Regev's reduction) Using a  $\mathsf{CVP}_{L^*,\alpha q/(\sqrt{2}r)}$  solver to generate  $\mathsf{poly}(n)$  discrete Gaussian states  $|D_{L,r\cdot\sqrt{n}/(\alpha q)}\rangle = \sum_{\mathbf{v}\in L} \sqrt{\rho_{r\cdot\sqrt{n}/(\alpha q)}(\mathbf{v})}|\mathbf{v}\rangle;$
- Step 3 (Additional) Create arbitrarily polynomially many quantum states  $|D_{L,r'}\rangle$  from poly(n)  $|D_{L,r\cdot\sqrt{n}/(\alpha q)}\rangle$  states, where  $r\sqrt{n}/\alpha q < r' < r$ .

Step 3 appears in case the S|LWE $\rangle$  solver in step 1 needs to consume  $|D_{L,r}\rangle$  states. Step 3 can be done in multiple ways, e.g., slightly modifying the GPV discrete Gaussian sampler [GPV08] to sample  $|D_{L,r'}\rangle$  states with  $r' = r \cdot (n\omega(\sqrt{\log n}))/(\alpha q)$ . In this case we should demand  $\alpha q > n\omega(\sqrt{\log n})$ .

We are left with step 1 to close the reduction. In the sequel, we focus on doing step 1 and see the  $S|LWE\rangle$  oracle we require.

Let  $\mathbf{x}$  denote a  $\mathsf{CVP}_{L^*,\alpha q/r}$  instance. Write  $\mathbf{x} = \kappa_{L^*}(\mathbf{x}) + \mathbf{x}'$ , where  $\kappa_{L^*}(\mathbf{x})$  is the closest  $L^*$  vector to  $\mathbf{x}$ , then it is guaranteed that  $\|\mathbf{x}'\| \leq \alpha q/r$ .

According to Regev's reduction,  $\langle \mathbf{x}, \mathbf{v} \rangle + e \pmod{p} = \langle \kappa_{L^*}(\mathbf{x}), \mathbf{v} \rangle + (\langle \mathbf{x}', \mathbf{v} \rangle + e) \pmod{p}$  is an LWE instance where  $\mathbf{v}$  is a  $D_{L,r}$  sample, and e is sampled from Gaussian distribution to "smooth" the discrete Gaussian  $\langle \mathbf{x}', \mathbf{v} \rangle$ .

Here we follow the same idea to prepare |LWE⟩ state through the following steps, using the discrete Gaussian state to replace the discrete Gaussian distribution over the lattice and a pure state with Gaussian amplitudes to replace the Gaussian error. For simplicity, let's ignore the normalization factors.

1. Prepare the initial state

$$\sum_{\mathbf{v} \in L} \rho_{r\sqrt{2}}(\mathbf{v}) |\mathbf{v}\rangle \otimes \sum_{e \in \mathbb{R}} \rho_{\sqrt{2}\sigma}(e) |e \bmod q\rangle$$

 $(\sum_{e\in\mathbb{R}}$  is not well-defined, we will build a state with enough precision to replace it.)

2. Measure  $L^{-1}\mathbf{v} \mod q$  to get an outcome  $\mathbf{a}$  and a result state

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{r\sqrt{2}}(\mathbf{v}) |\mathbf{v}\rangle \otimes \sum_{e \in \mathbb{R}} \rho_{\sqrt{2}\sigma}(e) |e \bmod q\rangle$$

3. Apply a unitary to add the inner product  $\langle \mathbf{x}, \mathbf{v} \rangle$  mod q to the second register we get

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{r\sqrt{2}}(\mathbf{v}) | \mathbf{v} \rangle \otimes \sum_{e \in \mathbb{R}} \rho_{\sqrt{2}\sigma}(e) | \langle \mathbf{s}, \mathbf{a} \rangle + \langle \mathbf{x}', \mathbf{v} \rangle + e \bmod q \rangle$$
 (1)

where  $L^*\mathbf{s} = \kappa_{L^*}(\mathbf{x}) \pmod{p}$ .

4. Apply  $\mathsf{QFT}_R$  to the first register where  $R > r\sqrt{n}$  is an integer:

$$\sum_{\mathbf{v} \in \mathbb{Z}_D^n} \sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{r\sqrt{2}}(\mathbf{v}) \cdot \omega_R^{\langle \mathbf{v}, \mathbf{y} \rangle} |\mathbf{y}\rangle \otimes \sum_{e \in \mathbb{R}} \rho_{\sqrt{2}\sigma}(e) |\langle \mathbf{s}, \mathbf{a} \rangle + \langle \mathbf{x}', \mathbf{v} \rangle + e \bmod q\rangle, \tag{2}$$

5. Measure the first register to get an outcome y and a result state

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \sum_{e \in \mathbb{R}} \rho_{r\sqrt{2}}(\mathbf{v}) \rho_{\sqrt{2}\sigma}(e) \cdot \omega_R^{\langle \mathbf{v}, \mathbf{v} \rangle} | \langle \mathbf{s}, \mathbf{a} \rangle + \langle \mathbf{x}', \mathbf{v} \rangle + e \bmod q \rangle.$$
 (3)

According to Theorem 11, this state is close to:

$$|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} \rangle := \sum_{u' \in \mathbb{P}} \rho_{\sqrt{2}\sqrt{r^2 \|\mathbf{x}'\|^2 + \sigma^2}} (u') \cdot e^{2\pi i \cdot u' \cdot \theta} |\langle \mathbf{s}, \mathbf{a} \rangle + u' \bmod q \rangle, \tag{4}$$

an LWE-like state whose error distribution is Gaussian distribution with a phase, where  $\theta := \frac{r^2 \langle \mathbf{x}', \mathbf{y}'/R \rangle}{r^2 ||\mathbf{x}'||^2 + \sigma^2}, \ \mathbf{y}'/R := \mathbf{y}/R - \kappa_{(qL)^*}(\mathbf{y}/R).$ 

Hence, if one can solve **s** from  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} \rangle$ , an  $|\mathsf{LWE} \rangle$  with error distribution being Gaussian distribution with a phase, then one can solve the  $\mathsf{CVP}_{L^*, \alpha q/r}$  problem.

One caveat is this  $\mathsf{S}|\mathsf{LWE}\rangle_{n,m,q,f}$  problem has its amplitude function  $f(u) = \rho_{\sqrt{2}\sqrt{r^2\|\mathbf{x}'\|^2 + \sigma^2}}(u) + e^{2\pi i \cdot u \cdot \theta}$  which depends on  $\mathbf{x}'$  and known  $\mathbf{y}$ .

To eventually solve the CVP problem for  $\mathbf{x}$ , it suffices to extract either the center  $\langle \mathbf{s}, \mathbf{a} \rangle$ , or  $\|\mathbf{x}'\|$ , or the direction of  $\mathbf{x}'$  from the state 4. In the following sections, we will describe our attempts and partial results.

**Remark 5.** If there is no phase (i.e. y = 0), this state can be written as

$$\sum_{e' \in \mathbb{R}} \rho_{\sqrt{2}\sqrt{r^2 \|\mathbf{x}'\|^2 + \sigma^2}}(e') |\langle \mathbf{s}, \mathbf{a} \rangle + e' \bmod q \rangle, \tag{5}$$

an |LWE| with Gaussian error distribution. It is the phase that makes our |LWE| nonstandard.

## 3 Extracting secrets from |LWE| state

From now on our targets become extracting either the center  $\langle \mathbf{s}, \mathbf{a} \rangle$  or  $\|\mathbf{x}'\|$  or the direction of  $\mathbf{x}'$  from the state  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} \rangle := \sum_{u' \in \mathbb{R}} \rho_{\sqrt{2}\sqrt{r^2\|\mathbf{x}'\|^2 + \sigma^2}}(u') \cdot e^{2\pi i \cdot u' \cdot \theta} |\langle \mathbf{s}, \mathbf{a} \rangle + u' \mod q \rangle$  with measurement results  $\mathbf{a}$  and  $\mathbf{y}$ , where  $\theta := \frac{r^2 \langle \mathbf{x}', \mathbf{y}'/R \rangle}{r^2\|\mathbf{x}'\|^2 + \sigma^2}$ ,  $\mathbf{y}'/R := \mathbf{y}/R - \kappa_{(qL)^*}(\mathbf{y}/R)$ . If this is done then using the reduction in Section 2.2 we can solve standard lattice problems via quantum algorithm.

# 3.1 Measuring the overlap of $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} \rangle$ and uniform to approximate $\|\mathbf{x}'\|$

Start with the case where  $\mathbf{y} = \mathbf{0}$  and no phase is involved, then our state  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{0}} \rangle$  is displayed in Equation (5). An important observation is that when  $\|\mathbf{x}'\|$  is small, the mass of  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{0}} \rangle$  is in a small range, while when  $\|\mathbf{x}'\|$  is large,  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{0}} \rangle$  seems close to the uniform superposition  $|\nu\rangle := \sum_{z \in \mathbb{Z}_q} |z\rangle$ . Hence measuring the overlap between  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{0}} \rangle$  and  $|\nu\rangle$  reveals whether  $\|\mathbf{x}'\|$  is small or large, which allows us to estimate  $\|\mathbf{x}'\|$  within some precision.

Since the probability of getting  $\mathbf{y} = \mathbf{0}$  is negligible<sup>1</sup>, we need to take the phase into consideration. However, the distribution of  $\theta$  in the phase is "neutralizing" the above effect: the expectation of  $|\langle \psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} | \nu \rangle|^2$  is independent of ||x'||.

This is not surprising since this overlap measurement does not use the measurement result  $\mathbf{y}$ , then measuring the second register should give the same result as measuring the second register of Equation (1), which is equivalent to measuring the overlap between the uniform superposition and a mixture of  $\{\sum_{e\in\mathbb{R}} \rho_{\sqrt{2}\sigma}(e) | \langle \mathbf{x}, \mathbf{v} \rangle + e \mod q \}_{\mathbf{v}\in qL+L\mathbf{a}}$ , which is a constant depending on  $\sigma$  and q.

According to the above arguments, we need to find a way to utilize the information in the measurement result  $\mathbf{y}$  in order to extract information of  $\mathbf{x}'$ . To better utilize  $\mathbf{y}$ , let's first figure out the distribution of  $\mathbf{y}$ ,  $\mathbf{y}'/R$  and  $\theta$  in our favourite state  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}} \rangle$ .

#### 3.2 The distribution of y

Now we give a more detailed analysis of the distribution of  $\mathbf{y}$  obtained by measuring the register  $|\mathbf{y}\rangle$  in (Equation (2)):

$$\sum_{\mathbf{y} \in \mathbb{Z}_R^n} \sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \cdot \omega_R^{\langle \mathbf{v}, \mathbf{y} \rangle} |\mathbf{y}\rangle \otimes \sum_{e \in \mathbb{R}} \rho_{\sqrt{2}\sigma}(e) |\langle \mathbf{x}, \mathbf{v} \rangle + e \bmod q\rangle$$

Computing the reduced density matrix of the first register, we have the probability of measuring  $\mathbf{y} \in \mathbb{Z}_{R}^{n}$  approximately proportional to

Actually the distribution of  $\mathbf{y}$  is approximately proportional to  $\rho_{\sqrt{\Sigma^{-1}}/2}(\mathbf{y}'/R)$ . See Section 3.2 for more detail.

$$\sum_{t \in [-\frac{q}{2}, \frac{q}{2})} |\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \rho_{\sqrt{2}\sigma}(t - \langle \mathbf{x}, \mathbf{v} \rangle \bmod q) \cdot e^{2\pi i \cdot \langle \mathbf{v}, \frac{\mathbf{y}}{R} \rangle}|^{2}$$

$$\approx \sum_{t' \in [-\frac{q}{2}, \frac{q}{2})} |\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \cdot e^{2\pi i \cdot \langle \mathbf{v}, \frac{\mathbf{y}}{R} \rangle} \rho_{\sqrt{2}\sigma}(t' - \langle \mathbf{x}', \mathbf{v} \rangle)|^{2}$$
(6)

where we can drop mod q in the approximation since we set the parameters so that, with overwhelming probability over the randomness of e and  $\mathbf{v}$ , t' can be written as  $t' = \langle \mathbf{x}', \mathbf{v} \rangle + e$  without mod q.

One can compute with a little effort that in Equation (6) the term associated with a fixed t' is

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \cdot e^{2\pi i \cdot \langle \mathbf{v}, \frac{\mathbf{y}}{R} \rangle} \rho_{\sqrt{2}\sigma}(t' - \langle \mathbf{x}', \mathbf{v} \rangle)$$

$$= \sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2\Sigma}}(\mathbf{v} - \mathbf{m}_{t'}) \cdot e^{2\pi i \cdot \langle \mathbf{v}, \mathbf{y}/R \rangle}$$

$$=_{(1)} \sum_{\mathbf{w} \in (qL)^*} \rho_{\sqrt{\Sigma^{-1}/2}}(\mathbf{w} - \mathbf{y}/R) \cdot e^{2\pi i \langle \mathbf{w}, L\mathbf{a} - \mathbf{m}_{t'} \rangle} \cdot e^{2\pi i \cdot \langle \mathbf{m}_{t'}, \mathbf{y}/R \rangle}$$

$$\approx_{(2)} \rho_{\sqrt{\Sigma^{-1}/2}}(\mathbf{y}'/R) \cdot e^{2\pi i \langle \kappa_{(qL)^*}(\mathbf{y}/R), L\mathbf{a} - \mathbf{m}_{t'} \rangle} \cdot e^{2\pi i \cdot \langle \mathbf{m}_{t'}, \mathbf{y}/R \rangle}$$
(7)

where  $\mathbf{m}_{t'} := \frac{r^2 t'}{r^2 \|\mathbf{x}'\|^2 + \sigma^2} \mathbf{x}'$ ,  $\Sigma := r^2 I - \frac{r^4 \mathbf{x}' \mathbf{x}'^T}{r^2 \|\mathbf{x}'\|^2 + \sigma^2}$ ,  $\Sigma^{-1} = \frac{I}{r^2} + \frac{\mathbf{x}' \mathbf{x}'^T}{\sigma^2}$  and  $\rho_{\sqrt{\Sigma}}(\mathbf{z}) := e^{-\pi \mathbf{z}^T \Sigma^{-1} \mathbf{z}}$  (without normalization).

(1) in Equation (7) is due to the Poisson Summation Formula. (2) in Equation (7) can be proved by directly applying the generalized tail bound Corollary 9 for multi-variate Gaussian, proved in the appendix, with  $\Sigma$  having two singular values  $r^2$  and  $r^2 \cdot \frac{\sigma^2}{r^2 \|\mathbf{x}'\|^2 + \sigma^2}$ .

Hence, the distribution of  $\mathbf{y}$  is approximately proportional to  $\rho_{\sqrt{\Sigma^{-1}}/2}(\mathbf{y}'/R)$  that only depends on  $\mathbf{y}'$ . Therefore the distribution of  $\mathbf{y}/R$  can be seen as ellipsoids centered at lattice points of  $(qL)^*$  whose direction of major axes is  $\mathbf{x}'$ .

Moreover, one can prove that  $|(qL)^* + \mathbf{y}'/R \cap \mathbb{Z}_R^n|$  is the same for all  $\mathbf{y}'/R$ , since  $(qL)^*$  is a suplattice of  $\frac{1}{q}\mathbb{Z}^n$  and therefore the cube  $[-1/2,1/2)^n$  can be viewed as containing an integer number of parallelepiped  $\mathcal{P}((qL)^*)$ . Hence, the distribution of  $\mathbf{y}'/R$  is proportional to  $\rho_{\sqrt{\Sigma^{-1}}/2}(\mathbf{y}'/R)$ , i.e.,  $\mathbf{y}'/R$  follows a multivariate Gaussian distribution. So we can bound the length of  $\mathbf{y}'/R$ :

$$\|\mathbf{y}'\|/R \le \sqrt{n} \cdot \frac{\sqrt{r^2 \|\mathbf{x}'\|^2 + \sigma^2}}{2\sigma r}$$
(8)

It follows that  $\theta = \frac{r^2 \langle \mathbf{x}', \mathbf{y}'/R \rangle}{r^2 \|\mathbf{x}'\|^2 + \sigma^2}$  in the phase of the amplitude of  $|\psi_{\langle \mathbf{s}, \mathbf{a} \rangle, \mathbf{y}}\rangle$  follows the Gaussian distribution  $\rho_{\beta}$  where  $\beta := \frac{r\|\mathbf{x}'\|}{2\sigma\sqrt{r^2\|\mathbf{x}'\|^2 + \sigma^2}}$ .

#### 3.3 Where can we find the secret?

Observe that the distribution of  $\mathbf{y}/R$  contains the information we want. To be more specific, the shape of the support of  $\mathbf{y}/R$  can be seen as ellipsoids centered at lattice points of  $(qL)^*$ , and the

direction and the length of their major axes are related to  $\mathbf{x}'$ . It seems plausible that we can utilize  $\mathbf{y}$  by extracting information about the secret from the distribution of  $\mathbf{y}/R$ .

In fact, the distribution of  $\mathbf{y}/R$ ,  $\mathbf{y}'/R$  and  $\theta$  all contains information about  $\mathbf{x}'$ :

- 1. The width of  $\mathbf{y}'/R$  is inversely related to  $\|\mathbf{x}'\|$ . However  $\mathbf{y}'/R$  cannot be obtained directly. (Obtaining  $\mathbf{y}'/R$  from  $\mathbf{y}$  is an instance of  $\mathsf{CVP}_{L^*,\frac{q\sqrt{n}\sqrt{\sigma^2+\alpha^2q^2}}{2\sigma r}}$ , which is harder than the  $\mathsf{CVP}_{L^*,\alpha q/r}$  problem we're aiming to solve. )
- 2. The shape of the support of  $\mathbf{y}/R$  is related to the direction of  $\mathbf{x}'$ . However these ellipsoids are cut by the boundaries of the cube  $[-1/2, 1/2)^n$ , leading to a troublesome support of  $\mathbf{y}/R$ .
- 3. The width of the distribution of  $\theta$  is positively related to  $\|\mathbf{x}'\|$ . However  $\theta$  cannot be obtained directly either.

# $4 \quad ext{Bypassing} \ket{\mathsf{LWE}}$

The above attempt inspires us to use the distribution of our measurement results to recover useful information. Here we no longer insist on first reducing standard lattice problems to  $S|LWE\rangle$ . In fact, we only need to give an algorithm that solves CVP using polynomial discrete Gaussian states. Combining the algorithm with step 2 and step 3 of our plan, we can get an iterative algorithm for standard lattice problems.

Given an instance  $\mathbf{x}$  of CVP, we begin with discrete Gaussian state

$$\sum_{\mathbf{v}\in L} \rho_{\sqrt{2}r}(\mathbf{v}) |\mathbf{v}\rangle$$

Again we measure  $\mathbf{a} := L^{-1}\mathbf{v} \mod q$  to get our favorite state

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) |\mathbf{v}\rangle$$

We apply a unitary on the state to send  $\langle \mathbf{x}, \mathbf{v} \rangle \mod q$  to the phase and obtain

$$\sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \cdot e^{\frac{2\pi i \langle \mathbf{x}, \mathbf{v} \rangle}{q}} | \mathbf{v} \rangle$$
(9)

Apply  $\mathsf{QFT}_R$  for  $R > r\sqrt{n}$  and we can get

$$|\psi\rangle := \sum_{\mathbf{y} \in \mathbb{Z}_R^n} \sum_{\mathbf{v} \in qL + L\mathbf{a}} \rho_{\sqrt{2}r}(\mathbf{v}) \cdot e^{2\pi i \frac{\langle \mathbf{x}, \mathbf{v} \rangle}{q}} \cdot e^{2\pi i \frac{\langle \mathbf{y}, \mathbf{v} \rangle}{R}} |\mathbf{y}\rangle$$
(10)

From Poisson summation formula,

$$|\psi\rangle = \sum_{\mathbf{y} \in \mathbb{Z}_R^n} \sum_{\mathbf{w} \in (qL)^*} \rho_{\frac{1}{\sqrt{2}r}} \left( \mathbf{w} - \frac{\mathbf{x}'}{q} - \frac{\mathbf{y}}{R} \right) \cdot e^{2\pi i \left( \langle \mathbf{w}, L\mathbf{a} \rangle + \frac{\langle \mathbf{s}, \mathbf{a} \rangle}{q} \right)} |\mathbf{y}\rangle$$
 (11)

Then we measure  $|\mathbf{y}\rangle$ . The resulting vector  $\mathbf{y}/R$ , when parsed as a rational vector in  $[-1/2, 1/2)^n$ , is expected to stay with a radius of  $\sqrt{n}/2r$  around  $(qL)^* - \frac{\mathbf{x}'}{q}$ .

Here is an intuitive idea of estimating  $\mathbf{x}'$ . We collect many samples of  $\mathbf{y}/R$  and then take the average. We expect the average to be  $-\frac{\mathbf{x}'}{q}$ , which is enough for solving CVP.

Unfortunately, our intuition is not valid. To be more specific, when r is large, say exponential, then the length of the shift  $\frac{\mathbf{x}'}{q}$  is less than  $\frac{\alpha q}{r} \cdot \frac{1}{q} = \frac{\alpha}{r}$ , which is negligible and can not be detected by efficient algorithms. We can also start from some special lattices such that initially r is small, say polynomial, but then the intersection between the boundary of  $[-1/2,1/2)^n$  and the balls of radius  $\sqrt{n}/2r$  around  $(qL)^* - \frac{\mathbf{x}'}{q}$  becomes annoying and thus the average of  $\mathbf{y}/R$  is not  $-\frac{\mathbf{x}'}{q}$ .

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# A Appendix

## A.1 An extension of Banaszczyk's Gaussian tail bounds over lattices

Recall Banaszczyk's Gaussian tail bounds:

**Lemma 6** (Lemma 1.5 [Ban93]). For any n-dimensional lattice L,  $\mathbf{c} \in \mathbb{R}^n$ , and  $r \geq \frac{1}{\sqrt{2\pi}}$ ,

$$\rho((L-\mathbf{c}) \setminus r\sqrt{n}B_2^n) < 2\left(r\sqrt{2\pi e} \cdot e^{-\pi r^2}\right)^n \rho(L).$$

We extend this tail bounds' RHS to an aribitrary shift of the lattice:

**Lemma 7.** For any n-dimensional lattice L, such that  $\lambda_1(L) > 3\sqrt{n}$ , and any  $\mathbf{y} \in \mathbb{R}^n$  such that  $dist(\mathbf{y}, L) < \sqrt{n}$ , we have

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2^{-n}\rho(L - \mathbf{y}). \tag{12}$$

*Proof.* First we prove that since  $\lambda_1(L) > 3\sqrt{n}$ , we have  $\rho(L) < 1 + 2^{-n}$ . To do so, we apply Lemma 6 with  $\mathbf{c} = \mathbf{0}$  and  $r\sqrt{n} = \lambda_1(L)/2$ , which gives

$$\rho(L \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2 \left( \frac{\lambda_1(L)}{2\sqrt{n}} \sqrt{2\pi e} \cdot e^{-\pi \left(\frac{\lambda_1(L)}{2\sqrt{n}}\right)^2} \right)^n \cdot \rho(L)$$

$$= 2 \cdot e^{n \ln(\lambda_1(L)/\sqrt{n}) - \pi \lambda_1(L)^2/4 + n \ln \sqrt{\pi e/2}} \cdot \rho(L)$$
(13)

Let  $\lambda_1(L) = x \cdot \sqrt{n}$ , then consider the function

$$f(x) := \ln(x) - \pi x^2 / 4 + \ln(\sqrt{\pi e/2}) \tag{14}$$

The derivative of f is

$$f'(x) = 1/x - \pi x/2 \tag{15}$$

Therefore when  $x > \sqrt{2/\pi}$ , f is decreasing. When x > 3, f(x) < -5.24.

Hence if  $\lambda_1(L) > 3\sqrt{n}$ ,

$$\rho(L \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2 \cdot e^{-5.24n} \cdot \rho(L),$$

which means  $\rho(L) < 1 + 2^{-n}$ 

We continue proving Lemma 7 by applying Lemma 6 with  $\mathbf{c} = \mathbf{y}$  and  $r\sqrt{n} = \lambda_1(L)/2$ . This gives

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2 \left(\frac{\lambda_1(L)\sqrt{\pi e}}{\sqrt{2n}}\right)^n \cdot e^{-\pi\lambda_1(L)^2/4} \rho(L)$$

$$<_{(1)} 3 \left(\frac{\lambda_1(L)\sqrt{\pi e}}{\sqrt{2n}}\right)^n \cdot e^{-\pi\lambda_1(L)^2/4}$$
(16)

where (1) uses  $\rho(L) < 1 + 2^{-n}$ .

Let  $\mathbf{y}' = \mathbf{y} - \kappa_L(\mathbf{y})$ , then  $\|\mathbf{y}'\| = dist(\mathbf{y}, L) < \sqrt{n}$ . Then

$$\frac{\rho((L-\mathbf{y}) \setminus \frac{\lambda_{1}(L)}{2} \cdot B_{2}^{n})}{\rho(\mathbf{y}')} < 3e^{n\ln\left(\frac{\lambda_{1}(L)\sqrt{\pi e}}{\sqrt{2n}}\right) - \pi\lambda_{1}(L)^{2}/4 + \pi\|\mathbf{y}'\|^{2}} 
< 3e^{n\ln(\lambda_{1}(L)/\sqrt{n}) - \pi\lambda_{1}(L)^{2}/4 + n\pi + n\ln\sqrt{\pi e/2}} 
< (1)3e^{n\ln(3) - \frac{9}{4}n\pi + n\pi + n\ln\sqrt{\pi e/2}} 
= 3e^{n(\ln 3 - \frac{5}{4}\pi + \ln\sqrt{\pi e/2})},$$
(17)

where (1) is obtained by taking the derivative similar as before: let  $\lambda_1(L) = x \cdot \sqrt{n}$ , then consider the function

$$g(x) := \ln(x) - \pi x^2 / 4 + \ln(\sqrt{\pi e/2}) + \pi \tag{18}$$

The derivative of g is

$$g'(x) = 1/x - \pi x/2 \tag{19}$$

Therefore when  $x > \sqrt{2/\pi}$ , g is decreasing. When x > 3, g(x) < -2.1.

Hence when  $\lambda_1(L) > 3\sqrt{n}$  and  $\|\mathbf{y}'\| < \sqrt{n}$ ,

$$\frac{\rho((L-\mathbf{y})\setminus\frac{\lambda_1(L)}{2}\cdot B_2^n)}{\rho(\mathbf{y}')}<2^{-2n}.$$

Since  $\rho(L - \mathbf{y}) = \rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) + \rho(\mathbf{y}')$ , we have

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2^{-n}\rho(L - \mathbf{y}). \tag{20}$$

For technical reasons, we need a variant of Lemma 7:

**Lemma 8.** For any n-dimensional lattice L and any  $\mathbf{y} \in \mathbb{R}^n$ , such that  $\lambda_1(L) > 3dist(\mathbf{y}, L)/d$  and  $\lambda_1(L) > 3\sqrt{n}$ , we have

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2^{-n} \rho_d(L - \mathbf{y}). \tag{21}$$

Remark: we can treat d as minor axis / major axis, which is less than 1.

*Proof.* First we prove that since  $\lambda_1(L) > 3\sqrt{n}$ , we have  $\rho(L) < 1 + 2^{-n}$ . To do so, we apply Lemma 6 with  $\mathbf{c} = \mathbf{0}$  and  $r\sqrt{n} = \lambda_1(L)/2$ , which gives

$$\rho(L \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2\left(\frac{\lambda_1(L)}{2\sqrt{n}}\sqrt{2\pi e} \cdot e^{-\pi\left(\frac{\lambda_1(L)}{2\sqrt{n}}\right)^2}\right)^n \cdot \rho(L)$$

$$= 2 \cdot e^{n\ln(\lambda_1(L)/\sqrt{n}) - \pi\lambda_1(L)^2/4 + n\ln\sqrt{\pi e/2}} \cdot \rho(L)$$
(22)

Let  $\lambda_1(L) = x \cdot \sqrt{n}$ , then consider the function

$$f(x) := \ln(x) - \pi x^2 / 4 + \ln(\sqrt{\pi e/2})$$
(23)

The derivative of f is

$$f'(x) = 1/x - \pi x/2 \tag{24}$$

Therefore when  $x > \sqrt{2/\pi}$ , f is decreasing. When x > 3, f(x) < -5.24.

Hence if  $\lambda_1(L) > 3\sqrt{n}$ ,

$$\rho(L \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2 \cdot e^{-5.24n} \cdot \rho(L),$$

which means  $\rho(L) < 1 + 2^{-n}$ 

We continue proving Lemma 8 by applying Lemma 6 with  $\mathbf{c} = \mathbf{y}$  and  $r\sqrt{n} = \lambda_1(L)/2$ . This gives

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2 \left(\frac{\lambda_1(L)\sqrt{\pi e}}{\sqrt{2n}}\right)^n \cdot e^{-\pi\lambda_1(L)^2/4} \rho(L)$$

$$<_{(1)} 3 \left(\frac{\lambda_1(L)\sqrt{\pi e}}{\sqrt{2n}}\right)^n \cdot e^{-\pi\lambda_1(L)^2/4}$$
(25)

where (1) uses  $\rho(L) < 1 + 2^{-n}$ .

Let  $\mathbf{y}' = \mathbf{y} - \kappa_L(\mathbf{y})$ , then  $\|\mathbf{y}'\|/d = dist(\mathbf{y}, L)/d < \lambda_1(L)/3$ . Then

$$\frac{\rho((L - \mathbf{y}) \setminus \frac{\lambda_{1}(L)}{2} \cdot B_{2}^{n})}{\rho_{d}(\mathbf{y}')} < 3e^{n \ln\left(\frac{\lambda_{1}(L)\sqrt{\pi e}}{\sqrt{2n}}\right) - \pi\lambda_{1}(L)^{2}/4 + \pi \|\mathbf{y}'\|^{2}/d^{2}} 
< 3e^{n \ln(\lambda_{1}(L)/\sqrt{n}) - \pi\lambda_{1}(L)^{2}/4 + \pi\lambda_{1}(L)^{2}/9 + n \ln\sqrt{\pi e/2}} 
< (1)3e^{n \ln(3) - \frac{9}{4}n\pi + n\pi + n \ln\sqrt{\pi e/2}} 
= 3e^{n(\ln 3 - \frac{5}{4}\pi + \ln\sqrt{\pi e/2})},$$
(26)

where (1) is obtained by taking the derivative similar as before: let  $\lambda_1(L) = x \cdot \sqrt{n}$ , then consider the function

$$g(x) := \ln(x) - 5\pi x^2 / 36 + \ln(\sqrt{\pi e/2})$$
(27)

The derivative of g is

$$g'(x) = 1/x - 5\pi x/18 \tag{28}$$

Therefore when  $x > \sqrt{\frac{18}{5\pi}}$ , g is decreasing. When x > 3, g(x) < -2.1.

Hence when  $\lambda_1(L) > 3dist(\mathbf{y}, L)/d$  and  $\lambda_1(L) > 3\sqrt{n}$ 

$$\frac{\rho((L-\mathbf{y})\setminus \frac{\lambda_1(L)}{2}\cdot B_2^n)}{\rho_d(\mathbf{y}')}<2^{-2n}.$$

Since  $\rho_d(L - \mathbf{y}) = \rho_d((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) + \rho_d(\mathbf{y}')$ , we have

$$\rho((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2^{-n} \rho_d(L - \mathbf{y}). \tag{29}$$

Corollary 9. For any n-dimensional lattice L, any  $y \in \mathbb{R}^n$  and any symmetric and positive matrix  $\Sigma$  whose smallest singular value is  $a^2$  and whose largest singular value is  $b^2$ , such that  $\lambda_1(L) > \frac{3b}{a}dist(\mathbf{y}, L)$  and  $\lambda_1(L) > 3\sqrt{n}/a$ , we have

$$\rho_{\Sigma^{-1}}((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) \le \rho_{\frac{1}{a}}((L - \mathbf{y}) \setminus \frac{\lambda_1(L)}{2} \cdot B_2^n) < 2^{-n}\rho_{\frac{1}{b}}(L - \mathbf{y}) \le 2^{-n}\rho_{\Sigma^{-1}}(L - \mathbf{y}). \tag{30}$$

## A.2 Smoothing of Gaussian with a phase

We generalize [Reg09, Claim 3.9] to handle Gaussian function with a phase.

**Theorem 10.** Let L be a lattice,  $\mathbf{u} \in \mathbb{R}^n$  be any vector, r, s > 0 be any real numbers,  $t := \sqrt{r^2 + s^2}$ . Consider the function Y on  $\mathbf{x} \in \mathbb{R}^n$  as the convolution of

1.  $\mathbf{y}$  with support  $L + \mathbf{u}$  and amplitude  $h(\mathbf{y}) := \rho_r(\mathbf{y}) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$  for some fixed  $\mathbf{z} \in \mathbb{R}^n$  such that  $d(\mathbf{z}, L^*) < \frac{t}{rs} \sqrt{n}$ ;

2. A noise vector taken from  $\rho_s$ .

Suppose  $\frac{rs}{t}\lambda_1(L^*) > 3\sqrt{n}$ . Then  $Y(\mathbf{x}) \approx \rho_t(\mathbf{x}) \cdot e^{2\pi i \cdot (r/t)^2 \langle \mathbf{z} - \kappa_{L^*}(\mathbf{z}), \mathbf{x} \rangle}$ .

*Proof.* The function Y can be written as

$$Y(\mathbf{x}) = \sum_{\mathbf{y} \in L + \mathbf{u}} h(\mathbf{y}) \rho_{s}(\mathbf{x} - \mathbf{y})$$

$$= \sum_{\mathbf{y} \in L + \mathbf{u}} \exp\left(-\pi \left(\frac{\|\mathbf{y}\|^{2}}{r^{2}} + \frac{\|\mathbf{x} - \mathbf{y}\|^{2}}{s^{2}}\right)\right) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$$

$$= \exp\left(-\frac{\pi}{r^{2} + s^{2}} \|\mathbf{x}\|^{2}\right) \sum_{\mathbf{y} \in L + \mathbf{u}} \exp\left(-\pi \left(\frac{t}{rs}\right)^{2} \cdot \|\mathbf{y} - \frac{r^{2}}{t^{2}} \mathbf{x}\|^{2}\right) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$$

$$= \rho_{t}(\mathbf{x}) \cdot \sum_{\mathbf{y} \in L + \mathbf{u}} \exp\left(-\pi \left(\frac{t}{rs}\right)^{2} \cdot \|\mathbf{y} - \frac{r^{2}}{t^{2}} \mathbf{x}\|^{2}\right) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$$

$$(31)$$

For any  $\mathbf{y} \in \mathbb{R}^n$ , let  $g(\mathbf{y}) := \rho_{\frac{rs}{t}}(\mathbf{y}) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$ . Then

$$\hat{g}(\mathbf{w}) = \rho_{\frac{t}{rs}}(\mathbf{w} - \mathbf{z})$$

Then

$$\sum_{\mathbf{y} \in L + \mathbf{u}} \exp\left(-\pi \left(\frac{t}{rs}\right)^{2} \cdot \|\mathbf{y} - \frac{r^{2}}{t^{2}}\mathbf{x}\|^{2}\right) \cdot e^{2\pi i \cdot \langle \mathbf{y}, \mathbf{z} \rangle}$$

$$= \sum_{\mathbf{y} \in L} \rho_{\frac{rs}{t}} \left(\mathbf{y} + \mathbf{u} - \frac{r^{2}}{t^{2}}\mathbf{x}\right) \cdot e^{2\pi i \cdot \langle \mathbf{y} + \mathbf{u}, \mathbf{z} \rangle}$$

$$= \sum_{\mathbf{y} \in L} g(\mathbf{y} + \mathbf{u} - \frac{r^{2}}{t^{2}}\mathbf{x}) \cdot e^{2\pi i \cdot \left\langle \frac{r^{2}}{t^{2}}\mathbf{x}, \mathbf{z} \right\rangle}$$

$$= \sum_{\mathbf{w} \in L^{*}} \hat{g}(\mathbf{w}) \cdot e^{2\pi i \cdot \left\langle \mathbf{u} - (r/t)^{2}\mathbf{x}, \mathbf{w} \right\rangle} \cdot e^{2\pi i \cdot \left\langle \frac{r^{2}}{t^{2}}\mathbf{x}, \mathbf{z} \right\rangle}$$

$$= \sum_{\mathbf{w} \in L^{*}} \rho_{t/rs}(\mathbf{w} - \mathbf{z}) \cdot e^{2\pi i \cdot \left\langle (\mathbf{u}, \mathbf{w}) - \left\langle (r/t)^{2}\mathbf{x}, \mathbf{w} - \mathbf{z} \right\rangle}$$
(32)

where (1) uses Poisson Summation Formula (ignoring the normalization factor  $(rs/t)^n \det(L^*)$ ).

Applying Lemma 7 with the lattice L being  $\frac{rs}{t}L^*$  here, which is  $\frac{rs}{t}(qL)^*$  in the main theorem; the vector  $\mathbf{y}$  being  $\frac{rs}{t} \cdot \mathbf{z}$ , which is  $\frac{rs}{t} \cdot \frac{\mathbf{y}}{R}$  in the main theorem;  $\lambda_1(L)$  being  $\frac{rs}{t}\lambda_1(L^*)$  here, which is  $\frac{rs}{ta}\lambda_1(L^*)$  in the main theorem.

Recall that  $s\|\mathbf{x}'\| = \sigma$  and  $t = \sqrt{r^2 + s^2}$ .  $dist(\mathbf{y}, L)$  in Lemma 7 satisfies

$$dist(\mathbf{y}, L) < \sqrt{n} \cdot \frac{\sqrt{\sigma^2 + r^2 \|\mathbf{x}'\|^2}}{\sigma r} \cdot \frac{rs}{t} = \sqrt{n} \cdot \frac{\sqrt{(s\|\mathbf{x}'\|)^2 + r^2 \|\mathbf{x}'\|^2}}{s\|\mathbf{x}'\|r} \cdot \frac{rs}{\sqrt{r^2 + s^2}} = \sqrt{n}$$

Back to Eqn. (32), when  $\frac{rs}{t} > \frac{3\sqrt{n}}{\lambda_1(L^*)}$  and  $\|\mathbf{z}'\| < \frac{t\sqrt{n}}{rs}$  with  $\mathbf{z}' := \mathbf{z} - \kappa_{L^*}(\mathbf{z})$ , we have

$$\sum_{\mathbf{w} \in L^*} \rho_{t/rs}(\mathbf{w} - \mathbf{z}) \cdot e^{2\pi i \cdot \left( \langle \mathbf{u}, \mathbf{w} \rangle - \left\langle (r/t)^2 \mathbf{x}, \mathbf{w} - \mathbf{z} \right\rangle \right)} \approx \rho_{t/rs}(\mathbf{z}') \cdot e^{2\pi i \cdot \left( \langle \mathbf{u}, \kappa_{L^*}(\mathbf{z}) \rangle + \left\langle (r/t)^2 \mathbf{x}, \mathbf{z}' \right\rangle \right)}$$
(33)

Then  $Y(\mathbf{x}) \propto \rho_t(\mathbf{x}) \cdot e^{2\pi i \cdot (r/t)^2 \langle \mathbf{z} - \kappa_{L^*}(\mathbf{z}), \mathbf{x} \rangle}$ 

## A.3 Linear combination of continuous Gaussian with a phase

**Theorem 11.** For any  $\mathbf{x} \in \mathbb{R}^n$  such that  $\|\mathbf{x}\| > 0$ . Suppose the amplitude of  $\mathbf{v} \in \mathbb{R}^n$  is  $f(\mathbf{v}) = \rho_r(\mathbf{v}) \cdot e^{2\pi i(\langle \mathbf{v}, \mathbf{y} \rangle + w)}$  for some fixed  $\mathbf{y} \in \mathbb{R}^n$  and  $w \in \mathbb{R}$ , then the amplitude of  $u := \langle \mathbf{x}, \mathbf{v} \rangle$  is

$$g(u) = \lambda \cdot \rho_{\|\mathbf{x}\| \cdot r}(u) \cdot e^{2\pi i \cdot u \cdot \frac{\langle \mathbf{x}, \mathbf{y} \rangle}{\|\mathbf{x}\|^2}}.$$
 (34)

where  $\lambda$  is some fixed complex number.

*Proof.* Let  $\mathbf{v}' \in \mathbb{R}^n$  be any real vector such that  $\langle \mathbf{v}', \mathbf{y} \rangle = w$ . Then the amplitude of  $\mathbf{v} \in \mathbb{R}^n$  can be written as

$$f(\mathbf{v}) = \rho_r(\mathbf{v}) \cdot e^{2\pi i \langle \mathbf{v} + \mathbf{v}', \mathbf{y} \rangle}$$
(35)

For  $j \in [n]$ , let  $g_j$  denote the amplitude of  $u_j := x_j \cdot v_j$ . Then, when  $x_j = 0$ ,  $g_j = \delta_0 \cdot e^{2\pi i \cdot v_j' \cdot y_j}$ , where  $\delta$  denotes the indicator function; when  $x_j \neq 0$ ,

$$g_j(u_j) = \rho_{x_j \cdot r}(u_j) \cdot e^{2\pi i \cdot (u_j \cdot y_j/x_j) + v_j' \cdot y_j}$$
(36)

Then the Fourier transform of  $g_i$  is

$$\hat{g}_j(z) = \begin{cases} e^{2\pi i \cdot v_j' \cdot y_j} & \text{when } x_j = 0; \\ e^{-\pi r^2 (x_i \cdot z - y_i)^2} \cdot e^{2\pi i \cdot v_j' \cdot y_j} & \text{when } x_j \neq 0; \end{cases}$$
(37)

So the product of  $\hat{g}_1, ..., \hat{g}_n$  is

$$\hat{g}(z) := \prod_{j=1}^{n} \hat{g}_{j}(z) = e^{-\pi r^{2}(\|\mathbf{x}\|^{2} \cdot z^{2} - 2\langle \mathbf{x}, \mathbf{y} \rangle \cdot z + \delta)} \cdot e^{2\pi i \cdot w} = e^{-\pi r^{2}\|\mathbf{x}\|^{2} \cdot (z - \theta)^{2} + \delta'} \cdot e^{2\pi i \cdot w}$$
(38)

where  $\delta$  and  $\delta'$  are some real numbers that does not depend on  $\mathbf{x}$ ,  $\theta = \frac{\langle \mathbf{x}, \mathbf{y} \rangle}{\|\mathbf{x}\|^2}$  is a real number that depends on  $\mathbf{x}$ .

Then the amplitude of  $u := \langle \mathbf{x}, \mathbf{v} \rangle \in \mathbb{R}$  is the convolution of  $g_j$ , which is the Fourier transform of  $\hat{g}$ . So the amplitude of u is

$$g(u) = \hat{g}(u) = \lambda \cdot \rho_{\|\mathbf{x}\| \cdot r}(u) \cdot e^{2\pi i \cdot u \cdot \theta}.$$
 (39)

where  $\lambda$  is some fixed complex number.