Name:

ID number:

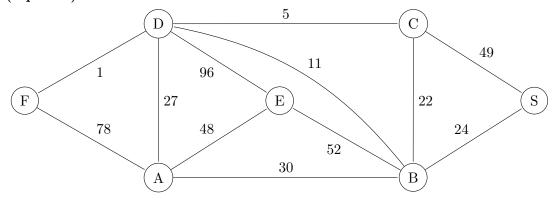
# 1. (6 points) True or False

Note: You should write down your answers in the box below.

(a)	(b)	(c)	(d)	(e)	(f)
F	Т	Т	F	F	Т

- (a) (1') A directed graph with n vertices has n-1 edges could be strongly connected. (n>1)
- (b) (1') A union tree generated by a union-by-height DSU of height h has  $\Omega(2^h)$  nodes.
- (c) (1') If we use BFS(breadth-first search) to find a path from u to v on an unweighted graph, this path will be the shortest path between u and v.
- (d) (1') If we consider the disjoint set time complexity, the time complexity of Kruskal's algorithm is  $O(|E|\alpha(|V|)\log|V|)$ .
- (e) (1') The maximum weight edge in a graph will not be in any MST.
- (f) (1') Let G be a weighted undirected graph with positive weights where edge e has weight  $w_e$  for all  $e \in E$ . And G' is a copy of G except that every edge e has weight  $w_e^2$ . Then any MST in G is also a MST in G'.

### 2. (7 points) Prim



(a) (3') Use Prim's algorithm to find the Minimal Spanning Tree of the graph. You should select S as the root node. Write the visit order of all nodes.

Solution:  $S \rightarrow B \rightarrow D \rightarrow F \rightarrow C \rightarrow A \rightarrow E$ 

- (b) (4') Let's modify the weights of some edges. Please give out the maximum and minimum weight of the edges given that won't change the MST. (You can write  $+\infty$  or  $-\infty$  if there is no maximum or minimum weight. We choose nodes by lexicographical order if we come across the same weight.)
  - AD: Maximum: \_\_\_\_\_\_ 30 \_\_\_\_ Minimum: \_\_\_\_\_ −∞
  - DE: Maximum: \_\_\_\_\_\_\_ Minimum: \_\_\_\_\_\_\_ 48

# 3. (2 points) Fill in the blanks

(a) (2') Assume we have an undirected graph G = (V, E) with 2 connected components. If |V| = 2n, the maximum of |E| is (2n-1)(n-1) and the minimum of |E| is (2n-2).

## 4. (7 points) kTrees

Given a complete undirected graph  $K_n$  on n vertices, every pair of distinct vertices i, j is connected with a **positive** weight  $w_{ij}$ .

The first k vertices  $v_1, \dots, v_k$  are chosen to serve as roots of k trees. You need to put other n-k vertices  $v_{k+1}, \dots, v_n$  to form k trees  $\{T_1, \dots, T_k\}$ , where each  $T_i$  is rooted at  $v_i$ .

The cost is defined as  $\sum_{i=1}^{k} W(T_i)$ , where  $W(T_i)$  is the sum of the weights of the edges in  $T_i$ . You need to design an efficient algorithm to find a solution of minimum cost.

(a) (4') Describe an algorithm to find a solution of minimum cost, you don't need to prove its correctness in this question.

### **Solution:**

- 1. Define a new vertex s.
- 2. Connect s and  $v_i$  for  $1 \le i \le k$ , where the weight is 0.
- 3. Run Kruskal's algorithm on the new graph G' = (V + s, E') (or run Prim starting at s), which is an alternative solution of minimum cost.
- (b) (3') Analyze the time complexity of your algorithm. You **need** to use k = O(n) to simplify your answer to O(f(n)).

### **Solution:**

- New vertex:  $\Theta(1)$ .
- Connect k edges: O(n).
- MST: |E| = n(n-1)/2 + k,  $O(n^2 \log n)$

Therefore, the time complexity is  $O(n^2 \log n)$ .