Homework 3 for Chapter 5-7

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Problem 5.9

For each process, we first calculate all the primes up to \sqrt{n} . And we allocate the each prime with the process in turn. For a process got a prime, it marks all mutiples of this prime as "not prime". In the end we use process 0 to calculate the number of all the primes up to n.

```
for (i = 2; i <= (int)(sqrt(limit)); ++i){
    if (is_prime[i])
        for(j = 2*i; j \le limit; j+=i)
            is\_prime[j] = 0;
}
count = 0;
for (i = 2; i \le limit; ++i)
    if(is_prime[i]) {
        count++;
        if(count \% p == id) {
             printf("process_\%d\catches_\%d\n", id, i);
             for (j = 2*i; j \le n; j+=i)
                 marked[j] = 1;
        }
    }
MPI_Reduce(marked, global_marked, n+1, MPI_CHAR, MPI_LOR, 0, MPI_COMM_WORLD);
// MPI_Bcast(&count, 1, MPI_INT, id, MPI_COMM_WORLD);
elapsed_time += MPI_Wtime();
if (!id) {
    count = 0;
    for (i = 2; i \le n; i++) {
        if (!global_marked[i]) count++;
    printf ("%d\primes\are\less\tan\or\equal\to\%d\n", count, n);
    printf ("Total_elapsed_time:_\%10.6f\n", elapsed_time);
}
  And the result is shown below
chenzixuan@LAPTOP-947PNMTJ:/mnt/c/Users/ChenZixuan/Desktop/Parallel-Programming-Project/hw3$
mpirun -np 3 ./eratosthenes
1000
168 primes are less than or equal to 1000
Total elapsed time:
                    0.000818
```

Problem 6.10

We define our own MPI_Bcast in the following rules. First if the current process is root process, it sends the buffer to all the other process using MPI_Send. If it is not, call MPI_Recv to receive new buffer information

from root process.

```
int my_MPI_Bcast(void *buffer, int count, MPI_Datatype datatype,
                  int root, MPI_Comm comm){
    int p, id;
    MPI_Status status;
    MPI Comm size (comm, &p);
    MPI Comm rank (comm, &id);
    if (id = root)
         for (int i = 0; i < p; ++i)
             if (i != id)
                  MPI_Send(buffer, count, datatype, i, 0, comm);
    }
    else {
         MPI_Recv(buffer, count, datatype, root, 0, comm, &status);
    return 0;
}
we test our algorithm by synchronizing certain amout of arrays. And the result is shown below
chenzixuan@LAPTOP-947PNMTJ:/mnt/c/Users/ChenZixuan/Desktop/Parallel-Programming-Project/hw3$
mpirun -np 3 ./my_MPI_Bcast 100000
Using my_MPI_Bcast: process 0 is initialized with value = 0
Using my_MPI_Bcast: process 1 is initialized with value = 1
Using my_MPI_Bcast: process 2 is initialized with value = 2
Using my_MPI_Bcast: process 1 updated with value = 0, time passed: 0.000708
Using MPI_Bcast: process 1 is re-initialized with value = 1
Using my_MPI_Bcast: process 0 updated with value = 0, time passed: 0.001026
Using my_MPI_Bcast: process 2 updated with value = 0, time passed: 0.001027
Using MPI_Bcast: process 2 is re-initialized with value = 2
Using MPI_Bcast: process 0 is re-initialized with value = 0
Using MPI_Bcast: process 2 updated with value = 0, time passed: 0.000343
Using MPI_Bcast: process 0 updated with value = 0, time passed: 0.000161
```

Problem 6.13

We divide the cell state A into p pieces, with each piece A_i maintained by a process. In the actual algorithm, each process maintain a full copy of A and the counting matrix C. Each process is responsible for calculating the living cell count around each cell of A_i , and update new cell state of A_i .

Using MPI_Bcast: process 1 updated with value = 0, time passed: 0.000724

```
//printf("pos = %d, x = %d, y = %d, count = %d \ n", i, x, y, count);
    return count;
}
int main (int argc, char *argv[]) {
    double elapsed_time; /* Parallel execution time */
    \mathbf{int} \quad i \ , \quad j \ , \quad m, \quad n \ , \quad p \ , \quad k \ ;
    int id; /* Process ID number */
    int* cell_state;
    int* count;
    MPI_Status status;
    int x, y;
    int low, high, size;
    int c;
    FILE* fp = fopen(FILE_NAME, "r");
    if (fp = NULL \mid | fscanf(fp, "%d_\%d", \&m, \&n) = EOF)
        printf("ERROR: ureadufrom file %s.\n", FILE_NAME);
    if (m > 0 \&\& n > 0)
        cell_state = (int*) malloc(m*n*sizeof(int));
        count = (int*) malloc (m*n*sizeof(int));
    else
        for(x = 0; x < m; ++x)
        for(y = 0; y < n; ++y)
             fscanf(fp, "%d", \&cell\_state[x*n+y]);
    fclose (fp);
    MPI_Init (&argc, &argv);
    MPI_Barrier (MPI_COMM_WORLD);
    elapsed time = -MPI Wtime();
    MPI_Comm_rank (MPI_COMM_WORLD, &id);
    MPI_Comm_size (MPI_COMM_WORLD, &p);
    if (p > m*n/2)
             printf("too_{\square}many_{\square}processes \n");
        MPI_Finalize();
        exit (1);
    if (argc != 3)
        if (!id)
        printf ("Command_line: | \%s | < m > n", argv [0]);
        MPI_Finalize();
        exit(1);
    }
    j = atoi(argv[1]);
    k = atoi(argv[2]);
    low = BLOCK\_LOW(id, p, m*n);
    high = BLOCK\_HIGH(id, p, m*n);
    size = BLOCK\_SIZE(id, p, m*n);
    // printf("low = %d, high = %d, size = %d \ ", low, high, size);
```

```
for(i = 0; i < j; ++i)
         MPI_Bcast(cell_state, m*n, MPI_INT, 0, MPI_COMM_WORLD);
         if (! id && (i+1) \% k == 0){
              printf("cell_{\,\square}\,state_{\,\square}\,at_{\,\square}\,iteration_{\,\square}\%d:\ \ \ \ n"\ ,\ \ i+1);
             for(x = 0; x < m; ++x){
                  for(y = 0; y < n; ++y)
                      printf("%d", cell_state[x*n+y]);
                  printf("\n");
         for (c = low; c < high; ++c)
             count[c] = count_grid(cell_state, c, m, n);
         for (c = low; c < high; ++c){
             if(!cell\_state[c] \&\& count[c] == 3)
                  cell\_state[c] = 1;
             else if (cell_state[c] && (count[c] < 2 || count[c] > 3))
                  cell\_state[c] = 0;
         if (id)
             MPI_Send(cell_state+low, size, MPI_INT, 0, 0, MPI_COMM_WORLD);
         else
             for (int pid = 1; pid < p; ++pid)
                  MPI_Recv(cell_state+BLOCK_LOW(pid, p, m*n),
                           BLOCK\_SIZE(pid, p, m*n),
                           MPI_INT,
                           pid,
                           0,
                          MPI COMM WORLD,
                          &status);
             }
    }
    MPI_Finalize ();
    return 0;
The result is shown below:
\begin { verbatim }
chenzixuan@LAPTOP-947PNMTJ:/mnt/c/Users/ChenZixuan/Desktop/Parallel-Programming-Project/hy
mpirun -np 3 ./cellular_aut
omaton 4 1
cell state at iteration 0:
1 0 0 1 1
0 \ 0 \ 0 \ 1 \ 1
0 0 0 0 1
1 1 1 0 1
cell state at iteration 1:
0 \ 0 \ 0 \ 1 \ 0
0 \ 0 \ 0 \ 0 \ 1
1 0 0 1 0
1 0 1 0 0
cell state at iteration 2:
0 0 0 0 0
```

```
0 \ 0 \ 1 \ 0 \ 1
0 \ 1 \ 0 \ 1 \ 0
0 \ 0 \ 0 \ 0 \ 0
cell state at iteration 3:
0 \ 0 \ 0 \ 0
0 \ 1 \ 1 \ 0 \ 0
1 \ 1 \ 0 \ 0 \ 0
0 \ 0 \ 0 \ 0 \ 0
\end{verbatim}
```

Problem 7.11

For Guatafson-Barsis's Law

$$\psi(n,p) \le s + (1-s)p$$

Since $s = \frac{\sigma(n)}{\sigma(n) + \phi(n)/p}$ is an implicit function of p. It increases as p increases. Actually the overall speed up still converges.