

3.

(1) 证  $a \vee b = b \wedge c$ ;

因为  $a \leq b$ , 所以  $a \vee b = b$ , 因为  $b \leq c$ , 所以  $b \wedge c = b$ ,  
因此  $a \vee b = b \wedge c$

(2) 因为  $a \leq b$ , 所以  $a \wedge b = a$ ,  $a \vee b = b$

又因为  $b \leq c$ , 所以  $b \wedge c = b$ ,  $b \vee c = c$ ,

于是

$$(a \wedge b) \vee (b \wedge c) = a \vee b = b;$$

$$\text{~~又~~ } (a \vee b) \wedge (b \vee c) = b \wedge c = b;$$

$$\text{所以 } (a \wedge b) \vee (b \wedge c) = (a \vee b) \wedge (b \vee c)$$

4. 因为  $a \wedge b \leq a$ ,  $c \wedge d \leq c$ , 所以由格的保序性得

$$(a \wedge b) \vee (c \wedge d) \leq a \vee c.$$

因为  $a \wedge b \leq b$ ,  $c \wedge d \leq d$ , 所以由格的保序性得

$$(a \wedge b) \vee (c \wedge d) \leq b \vee d$$

因此,  $(a \wedge b) \vee (c \wedge d)$  是  $a \vee c$  和  $b \vee d$  的下界,

$$\text{所以 } (a \wedge b) \vee (c \wedge d) \leq (a \vee c) \wedge (b \vee d)$$





12.

设  $\langle L; \leq \rangle$  是一个格,  $\#L \geq 2$ ,

用反证法: 假定存在  $l \in L$ , 使得  $l \wedge l = 0$ ,  $l \vee l = 1$ ,

则由幂等律知  $l = 0 = 1$ , 与  $\#L \geq 2$  矛盾,

所以,  $L$  中必不存在元素是它自身的补.

19.

必要性: 设存在  $i (1 \leq i \leq r)$  使  $a = b$ , 则  $b_i \leq b_1 \vee b_2 \vee \dots \vee b_r$ ,

故以  $a = b_i \leq b_1 \vee b_2 \vee \dots \vee b_r$

充分性: 设  $a \leq b_1 \vee b_2 \vee \dots \vee b_r$ ,

用反证法, 若不存在  $i (1 \leq i \leq r)$  使得  $a = b_i$ , 则由于

$a, b_1, b_2, \dots, b_r$  均是原子, 故  $a \wedge b_1 = 0$ ,  $a \wedge b_2 = 0$ ,

$\dots$ ,  $a \wedge b_r = 0$ , 所以

$a \wedge (b_1 \vee b_2 \vee \dots \vee b_r) = (a \wedge b_1) \vee (a \wedge b_2) \vee \dots \vee (a \wedge b_r) = 0$

又因为  $a \leq b_1 \vee b_2 \vee \dots \vee b_r$ , 所以  $a \wedge (b_1 \vee b_2 \vee \dots \vee b_r) = a$

而  $a$  是原子, 所以  $a \neq 0$ , 即  $a \wedge (b_1 \vee b_2 \vee \dots \vee b_r) \neq 0$

与上述结论矛盾, 所以假设错误, 故存在  $i (1 \leq i \leq r)$

使  $a = b_i$ .





25. 布尔代数有补, 故  $\alpha$  和  $\beta$  互补, 由分配律

$$f(x, y) = (x \wedge \alpha) \vee (x \wedge y) \vee (\bar{x} \wedge \bar{y})$$

$$\text{故 } f(0, 1) = 0 \vee 0 \vee 0 = 0, \quad f(0, \alpha) = 0 \vee 0 \vee \beta = \beta$$

$$f(0, \beta) = 0 \vee 0 \vee \alpha = \alpha, \quad f(0, 0) = 0 \vee 0 \vee 1 = 1$$

$$f(\alpha, 0) = \alpha \vee 0 \vee \beta = 1, \quad f(\alpha, \alpha) = \alpha \vee \alpha \vee \beta = 1$$

$$f(\alpha, \beta) = \alpha \vee 0 \vee 0 = \alpha, \quad f(\alpha, 1) = \alpha \vee \alpha \vee 0 = \alpha$$

$$f(\beta, 0) = 0 \vee 0 \vee \alpha = \alpha, \quad f(\beta, \alpha) = 0 \vee 0 \vee 0 = 0$$

$$f(\beta, \beta) = 0 \vee \beta \vee \alpha = 1, \quad f(\beta, 1) = 0 \vee \beta \vee 0 = \beta$$

$$f(1, 0) = 0 \vee 0 \vee 0 = 0, \quad f(1, \alpha) = \alpha \vee \alpha \vee 0 = \alpha$$

$$f(1, \beta) = \alpha \vee \beta \vee 0 = 1, \quad f(1, 1) = \alpha \vee 1 \vee 0 = 1$$

列出对于所有自变量  $(x, y) \in \mathcal{P}^2$  的  $f(x, y)$  之值的表:

$x$	0	0	0	0	$\alpha$	$\alpha$	$\alpha$	$\alpha$	$\beta$	$\beta$	$\beta$	$\beta$	1	1	1	1
$y$	0	$\alpha$	$\beta$	1	0	$\alpha$	$\beta$	1	0	$\alpha$	$\beta$	1	0	$\alpha$	$\beta$	1
$f(x, y)$	1	$\beta$	$\alpha$	1	1	1	$\alpha$	$\alpha$	$\alpha$	0	1	$\beta$	0	$\alpha$	1	1





20-

$f(x, y)$  的最小项标准形式

$$f(x, y) = (f(0, 0) \wedge \bar{x} \wedge \bar{y}) \vee (f(0, 1) \wedge \bar{x} \wedge y)$$

$$\vee (f(1, 0) \wedge x \wedge \bar{y}) \vee (f(1, 1) \wedge x \wedge y)$$

$$= (1 \wedge \bar{x} \wedge \bar{y}) \vee (0 \wedge x \wedge \bar{y}) \vee (1 \wedge x \wedge y)$$

$f(x, y)$  的最大项标准形式

$$f(x, y) = (f(0, 0) \vee x \vee y) \vee (f(0, 1) \vee x \vee \bar{y})$$

$$\vee (f(1, 0) \vee \bar{x} \vee y) \vee (f(1, 1) \vee \bar{x} \vee \bar{y})$$

$$= (1 \vee x \vee y) \vee (0 \vee x \vee \bar{y}) \vee (0 \vee \bar{x} \vee y) \vee (1 \vee \bar{x} \vee \bar{y})$$

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34.

因为  $\langle B; -, \vee, \wedge \rangle$  是布尔代数, 故十条基本定律在其上均成立, 所以

$$(a \vee b) \wedge (c \vee \bar{b}) = ((a \vee b) \wedge c) \vee ((a \vee b) \wedge \bar{b}) \quad (\text{分配律})$$

$$= (a \wedge c) \vee (b \wedge c) \vee (a \wedge \bar{b}) \vee (b \wedge \bar{b})$$

$$= (a \wedge c) \vee [(b \wedge c) \vee (a \wedge \bar{b})]$$

$$= (a \wedge c \wedge (b \vee \bar{b})) \vee [(b \wedge c) \vee (a \wedge \bar{b})]$$

$$= (a \wedge c \wedge b) \vee (a \wedge c \wedge \bar{b}) \vee (b \wedge c) \vee (a \wedge \bar{b})$$

$$= [(a \wedge (c \wedge b)) \vee (b \wedge c)] \vee [(a \wedge \bar{b}) \wedge c] \vee (a \wedge \bar{b})$$

$$= (b \wedge c) \vee (a \wedge \bar{b}) = (a \wedge \bar{b}) \vee (c \wedge b)$$

$$\text{即 } (a \vee b) \wedge (c \wedge \bar{b}) = (a \wedge \bar{b}) \vee (c \wedge b)$$





35.

(1)

对  $\forall a \in B_1$ , 有  $0 = a \wedge \bar{a} \in B_1$ 。由  $f$  是同态映射知  
 $f(0) = f(a \wedge \bar{a}) = f(a) \otimes f(\bar{a}) = f(a) \otimes f'(\bar{a}) = \alpha$ ,  
所以  $0 \in J$

(2)

若  $a \in J$ , 则  $f(a) = \alpha$

对  $\forall x \in B_1$ , 若  $x \leq a$  时,  $x = x \wedge a$

于是  $f(x) = f(x \wedge a) = f(x) \otimes f(a) = f(x) \otimes \alpha = \alpha$ ,  
所以  $x \in J$

(3)

对  $\forall a, b \in J$ , 有  $f(a) = \alpha$ ,  $f(b) = \alpha$

于是  $f(a \wedge b) = f(a) \otimes f(b) = \alpha \otimes \alpha = \alpha$

所以  $a \wedge b \in J$

$f(a \vee b) = f(a) \vee f(b) = \alpha \vee \alpha = \alpha$ , 所以  $a \vee b \in J$

故  $\langle J; \vee, \wedge \rangle$  是 ~~一个~~ 构成 - 代数系统

