

第七章:

3. 证明: $\because a \leq b \quad a \wedge c \leq a \quad \therefore a \wedge c \leq b$ 同理: $a \wedge c \leq d$

$\therefore a \wedge c$ 为 b 和 d 的下界 $\therefore a \wedge c \leq b \wedge d$

$\because b \leq b \vee d, a \leq b \quad \therefore a \leq b \vee d$ 同理: $c \leq b \vee d$

$\therefore b \vee d$ 为 a, c 的上界 $\therefore a \vee c \leq b \vee d$, 得证

4. 证明: $\because a \wedge b \leq a, a \leq a \vee c \quad \therefore a \wedge b \leq a \vee c$

同理: $c \wedge d \leq a \vee c \quad \therefore a \vee c$ 为 $a \wedge b$ 和 $c \wedge d$ 的上界

$\therefore (a \wedge b) \vee (c \wedge d) \leq a \vee c$ 同理 $(a \wedge b) \vee (c \wedge d) \leq b \vee d$

$\therefore (a \wedge b) \vee (c \wedge d)$ 为 $a \vee c$ 和 $b \vee d$ 的下界

$\therefore (a \wedge b) \vee (c \wedge d) \leq (a \vee c) \wedge (b \vee d)$, 得证

12. 证明: 用反证法

① 若格中只有两个元素, 设 $a = \bar{a}$

\because 补元存在 \therefore 格有 $0, 1 \quad \therefore a = 0$ 或 1 , 无论哪种情况 $a \neq \bar{a}$

② 若格中有更多元素, 设 $a = \bar{a}$ 且 $a \neq 0, a \neq 1$

$a \vee a = 1, a \wedge a = 0$, 由幂等律, $a = 1, a = 0$, 矛盾

\therefore 不会有元素是它自身的补

19. 证明: 若 $a = b_i, (1 \leq i \leq r)$

则必有 $a \leq b_1 \vee b_2 \vee \dots \vee b_r$. 若 $a \leq b_1 \vee b_2 \vee \dots \vee b_r$

$\therefore a \wedge (b_1 \vee b_2 \vee \dots \vee b_r) = a$

由分配律, $(a \wedge b_1) \vee (a \wedge b_2) \vee \dots \vee (a \wedge b_r) = a$

$\because a \neq 0 \quad \therefore a \wedge b_i (1 \leq i \leq r)$ 不全为零

设 $a \wedge b_m \neq 0$, $\because b_m$ 为原子且 a 为原子 $\therefore a = b_m$, 得证

25. 解: $f(x, y) = (x \wedge (\bar{0} \vee y)) \vee (\bar{x} \wedge \bar{y})$

$$= [x \vee (\bar{x} \wedge \bar{y})] \wedge [\bar{0} \vee y \vee (\bar{x} \wedge \bar{y})] = [(x \vee \bar{x}) \wedge (x \vee \bar{y})] \wedge [\bar{0} \vee y \vee (\bar{x} \wedge \bar{y})]$$

$$= (x \vee \bar{y}) \wedge \{ (\theta \vee y) \vee \bar{x} \} \wedge \{ (\theta \vee y) \vee \bar{y} \} = (x \vee \bar{y}) \wedge (x \vee \bar{y} \vee \theta) \wedge 1$$

$$= (x \vee \bar{y}) \wedge (x \vee \bar{y} \vee \theta)$$

$\begin{array}{c} x \\ y \end{array}$	0	1	θ	β
0	1	1	1	1
1	0	1	θ	β
θ	β	1	1	β
β	θ	1	θ	1

26. 解: 最小项标准形式:

$$\begin{aligned} f(x, y) &= (f(0,0) \wedge \bar{x} \wedge \bar{y}) \vee (f(0,1) \wedge \bar{x} \wedge y) \vee (f(1,0) \wedge x \wedge \bar{y}) \vee (f(1,1) \wedge x \wedge y) \\ &= (1 \wedge \bar{x} \wedge \bar{y}) \vee (\theta \wedge x \wedge \bar{y}) \vee (1 \wedge x \wedge y) \end{aligned}$$

最大项标准形式

$$\begin{aligned} f(x, y) &= (f(0,0) \vee x \vee y) \wedge (f(0,1) \vee x \vee \bar{y}) \wedge (f(1,0) \vee \bar{x} \vee y) \wedge (f(1,1) \vee \bar{x} \vee \bar{y}) \\ &= (1 \vee x \vee y) \wedge (0 \vee x \vee \bar{y}) \wedge (0 \vee \bar{x} \vee y) \wedge (1 \vee \bar{x} \vee \bar{y}) \end{aligned}$$

34. 证明: 左边 = $(a \vee b) \wedge (c \vee \bar{b}) = [a \wedge (c \vee \bar{b})] \vee [b \wedge (c \vee \bar{b})]$

$$= [(a \wedge b) \vee (a \wedge c)] \vee [(b \wedge c) \vee (b \wedge \bar{b})] = [(a \wedge b) \vee (a \wedge c)] \vee [b \wedge c]$$

$$= [a \wedge (c \vee \bar{b})] \vee (b \wedge c)$$

右边 = $(a \wedge b) \vee (c \wedge b) = [(c \wedge b) \vee a] \wedge [(c \wedge b) \vee \bar{b}]$

$$= [(c \wedge b) \vee a] \wedge (c \vee \bar{b}) = [(c \wedge b) \wedge (c \vee \bar{b})] \vee [a \wedge (c \vee \bar{b})]$$

$$= [a \wedge (c \vee \bar{b})] \vee [(c \wedge b \wedge c) \vee (c \wedge b \wedge \bar{b})] = [a \wedge (c \vee \bar{b})] \vee (b \wedge c) = \text{左边}, \text{得证}$$

35. 由题意, 对 $\forall x_1, x_2 \in V_1$, 有 $f(x_1 \vee x_2) = f(x_1) \oplus f(x_2)$; $f(x_1 \wedge x_2) = f(x_1) \odot f(x_2)$

证明: $\because f(0 \vee x) = f(0) \oplus f(x) = f(x)$, $f(0 \wedge x) = f(0) \odot f(x) = f(0)$

$$\therefore f(0) \text{ 为 } V_2 \text{ 的最小元, } f(0) = \theta \quad \therefore 0 \in J$$

1) 证明: 由题意, $a \in B$, $f(a) = \theta \quad \therefore f(x \vee a) = f(x) \oplus f(a) = f(x)$

$$f(x \wedge a) = f(x) \odot f(a) = f(a) \quad \because x \leq a \quad \therefore f(x \vee a) = f(a) = f(x)$$

$$\therefore x \in J$$

13, 证明: 由 (1), (2), J 不为空

$$\text{若 } x_1, x_2 \in J, \text{ 则 } f(x_1 \vee x_2) = f(x_1) \oplus f(x_2) = 0 \vee 0 = 0$$

由 V 的封闭性, $x_1, x_2 \in B_1$, 则 $x_1 \vee x_2 \in B_1$

$$\therefore x_1 \vee x_2 \in J$$

$$f(x_1 \wedge x_2) = f(x_1) \otimes f(x_2) = 0 \otimes 0 = 0 \quad \text{同理, 有 } x_1 \wedge x_2 \in J$$

综上: $\langle J, \vee, \wedge \rangle$ 构成 - 代数系统