## Recorsive Sequence Highwest (Global)

As for as all this theory is, it's going to be a lor more for if he actually code up the solution. The pattern will be quite similar to the solution for the Fibonacci sequence - often all, he he got base cases and recursive cases, but the solution will be more detailed. Partly because of the data structures R makes available.

To start with, white the Fibonacci sequence, and answer actually contains two ports, hell, three: the gapped version of y, and the score; will call these xaln, yaln, and score:

x: ACTAGC => xah: ACTAGC score: 1 y: ATACC => yah: A-TACC

As rentioned, unless otherwise necessary, will encode sequeces as character rectors rather than longer strings, using the char-nec() function: "ACTAGE" "A", "C", "T", "A", "C", "C" This allows us to easily concatenate sequences, or get subsequences

We'll encode an 'ensurer' as a named list, and will also store the inputs. We could encode it by haid like so:

change <-  $l_{1}$  is t(x = c("A", "C", "T", "A", "G", "C"), y = c("A", "T", "A", "C", "C"), xaln = c("A", "C", "T", "A", "G", "C") yaln = c("A", "-1, "T", "A", "C", "C") score = 4)

Though of course well be using code to create these.

As per the homework, he created a few "helper"

Functions: one that handled any base case (a tedious little sucker) and one that scored an alignment according to the scoring rules:

# Greative charges of equal leight, returns an # integer score:

score ala 2- function (xin, yin) {

3

# given two char vecs constituting a base case # (either of length of or both of length 1)

# return on 'answer' object
base-case <- Function (xin, yin) {

3

Recorsivity

Now re get to do the fun part: the recorsive case: it also takes two char rectors, x and x, and returns an onswer. If x and y are a bare

case, it just calls the base case stuction.

global\_aln <- Function (x, y) {

if  $(\log th(x) == 0 \mid length(y) == 0 \mid (length(x) == 1))$  {

length(y) == 1)}

retum (base\_case(x, y))

IF its not a base case, he follow the recursive pattern: ne start by computing px, ex, py, and ey as defined in our proof:

 $px \leftarrow x[1: length(x) - 1]$   $ex \leftarrow x[length(x)]$   $py \leftarrow y[1: length(y) - 1]$   $ey \leftarrow y[length(y)]$ 

Next, le get to recurse. Let's revisit our l'Atle diagram for the three oprions:

(next page)

\* X[a:5] returns X from indices a to b, inclusive.

```
\left(\begin{array}{c}A\\P_x \text{ aligned } \omega/\end{array}\right) e_x \quad \text{Score}_A = S(A) + S(e_x, e_y)
Px aligned w/ e_x score_A = S(A) + S(e_x, e_y)

e_y

e_y
              ( Px aligned w/ ) ex score = 5(c) + 5(ex, "-")
          So, he need to compute the subalignments, A, B, and C,
         recurrively.
                    A <- global-aln(px, px)

B <- global-aln(c(px, ex), px)

C <- global-aln(px, c(px, ex))
       Notice the similarity between our conceptual detration and the code!
       Now we can use these subanswers to compute the three
        potential overall answer objects, accessing parts of the subanswert
        as readed at [[]] notation
                             answera <- list(x = x, y = y,
                                                                                                                               xaln = c(A[("xaln"]], ex),
                                                                                                                                yaln = c(A[["yah"]], ey),
              consider X = X, Y = Y, X = X, Y = X, X = X, X
                                                                                                                                 score = A[["score"]] + score-aln(ex, ey))
                                                                                                                           score = B[["score"]] + score-aln("-", ex))
```

```
onswere <- lis+(x=x, y=y, xaln=c(C[["xaln"]], ex), yaln=c(C[["yaln"]], "-"), score=C[["score"]]+score_aln(ex, "-"))
```

Now we need to Figure out which of those three is the dest, based on their scores, and return that one. We'll use a dead-simple way of doing that:

```
bestonswer 2- answera [["score"]]

if (answerb[["score"]] > bestscore) {

bestanwer 2- answerb

bestscore 2- answerb[["score"]]

if (answerc[["score"]] > bestscore) {

bestanswer 2- answerc

bestanswer 2- answerc

best score 2- answerc [["score"]]
```

reton (bestonsner)

And that's it! We've timed a recursive definition (proved correct via induction) into a recursive algorithm. Lets try it.

x <- char-rec ("TATCGG")

y <- char-rec ("TCTGG")

onswer <- global-aln (x, x)

pont(onswer)

# Sneecee+,

<sup>\*</sup> Yes, there was some debugging invoked - don't for a moment that code like this springs from the torehead of zeus!

Menoization
We'd Find that evan slightly bigger instances of the problem take much longer to run.

X <- char-rec ("TATCGGT")

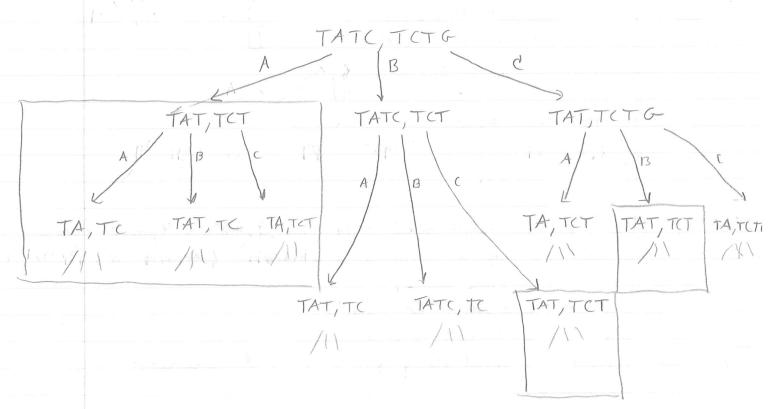
y <- char-rec ("TCTGGTCC")

onsher <- global-aln (X, y) # waiting...

proxt(asker)

This is because, much like the Folomaci solution we looked at, there is a lot of recursive work being done. While each call to the Fisc) Function resulted in two subralls, each call to our function is resulting in three. Vites!

But, also like the Fisc) Function, many of the subproblems are overlapping. Consider the call tree for X = ITATC and y = TCTG, Times.



We've highlighted three of the biggest our lapping subproblems, but there are many more to find. Obviously, memoizing our global alignment function will provide big gains. We'll use as a trey the unvector()'d varion of the inputs separated by a comma (as in "TATC, TCTG"):

```
ALN_CACHE <- hosh()

global_aln <- function (x, y) {

this call <- str_c (unvector(x), ",", unvector (y))

if (has. key (this call, ALN-CACHE)) {

return (ALN-CACHE [(this call]])

}

if (lugth(x) == Ø | length(y) == Ø | (length(x) == 1 & length(y) == 1)) {

ALN_CACHE [(this call]] <- base-cate(x, y)

return (ALN-CACHE [(this call]])

}

[rest of tund.

ALN_CACHE [(this call]] <- best as were

return (best as nor)

}
```

Now our function will run quite Fast. Just how Fast ronains to be seen.

x <- char-vec ("TATICGGTCTA")

y <- char-vec ("TCTGGTCCAC")

answer <- global-aln(x,y) # much faster

print (onswer)

Inspecting the cache

In an effort to determine how much time the alignment takes, we can take a look at the cache - since there are no loops in our recorsive call, each 'cell' of the cache represents just a few function calls once its been memoized. To do this, we're actually going to visualize the cache, for extra coolness.

Put that means we have to get the data in the cache (the values in the hash, which have the answers, inputs, and scores) into a data-trane. Unfortunately, ithe hash' package For R doesn't make this easy, so well use I stacks as an intermediary. Our strategy will be:

For each key

extract value (on 'onsur' tist)

peturn as data from the retacte)

As we do this, we'll also take any elements of each answer

17st that are characters, and unvector() them, so X = C("A", "C", "T") will become x = "ACT" and so on.

(code on next page)

dha itair

hash\_vals\_to\_df <- Function (thehash) {

tempstach <- rstach ()

For (tey in keys (the hash)) {

drsher list <- thehash [[key]]

the unvector every cleant of armedist if possible

For (i in seq (!, length (armedist)) {

if (is.character (arsner list [[i]]) {

answerlist [[i]] <- unvec\_char (answerlist [[i]))

}

tempstack <- insert\_top (tempstack, answerlist)

}

return (as. data. Frame (tempstack, strings As Factors = F))

(The beauty of R TS that It's always possible to convert

data into another format with a bit of trictiery - too bad it's never in the format we want!)

Non re can easily turn such a cache into a ricely organized data frame:

cache of <- hash rals to OF (ALN-(ACHE))
print(head (cache of))

TATCGGT... TCTGC... TATCGGT... TC-GG-... -15

We can clearly see that each row of the data. France represents a solved subproblem - Let's plot these susproblems inputs: (x and x) and just the score output: he'll do it in a grid, organized with longer x problems along the x axis, and longer y subproblems along the y (w/ longer ones toward the bottom), colored and labeled by score.

or scale to the scale

```
p 2- ggplot(cache-dt) +

geom_tile(aes(x = reorder(x, nchar(x)),

y = reorder(x, -1*nchar(x)),

fill = score)) +

geom_tex+(aes(x = reorder(x, nchar(x)),

y = reorder(x, nchar(x)),

thene_bu(16) + label = score)) +

thene(axis.text.x = element_tex+(angle = 35, hjust=1))+

(sord-equal())

Plot(p)
```

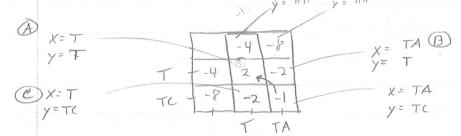
Ahah, very nice - our solved and cached sulproblems are layed out in a nice grid.

| - 4 | -8       | -12                   |                                      |
|-----|----------|-----------------------|--------------------------------------|
| 2   | -2       | -6                    | ,                                    |
| -2  | -1       | -5                    |                                      |
| -6. | -S.      | 1 -                   | 1                                    |
| -   |          |                       |                                      |
|     |          |                       |                                      |
| 1   | TAX      | RT TH                 | KC KKC                               |
|     | 2 -2 -6. | 2 -2<br>-2 -1<br>-6S. | -4-8-12<br>2-2-6<br>-2-1-5<br>-65.1. |

How big is this grid? (And thus, how may susproblems here considered?) Well, if sequice of leigth n, and I was of leigth m, the whole grid is roughly nxm, or O(nn), and this is how much work is required for the memoized solution - not great, but a lot better than the exponential function calls of the non-memoized one!

## Subproblems are organized

upper lot hand corner to be precise:



First, note that the cells along the top and bottom correspond to base cases—if we noted we cold have considered x = """ y = """ to be a base case with score of what's more interesting are the intenal set of cells: in fact, the A, B, and C subproblems for <math>x = TA, y = TC are in the three cells left and up! The solution for that cell did a recurring call for those other three.

Additionally, one of those three contributed to the bestauser' in this case it would have been subproblem A. In a sense, we can say that the solution for X=TA, Y=TC, came 'From' that cell, perhaps we'd say From X=T, From y=T.

We con visualize this too, if we encode this information in the answer lists that are cached.

For a base case, we'll simply say that the solution comes 'From' itself - it doesn't much matter, but it

does matte the eventual plot cleaner:

Sase-case c- Function (xin, xin) {

the change all answers like so:

answer (- list(x = xin, y = xin,

xaln = x aligned, yaln = yaligned,

score = score-aln(xaligned, yaligned),

Franx = xin, Frany = xin)

return(answer)

And in the recursive function, we need to determine the From information as he figure out which one is best:

global-aln <- Function (X, y) &

best owner <- answer a [["score"]]

best owner [["Fromx"]] <- A[["x"]] # new

best owner [["Fromy"]] <- A[["y"]] # new

H(answer [["score"]] > best score) {

best owner <- answer b

best owner [["Fromx"]] <- B[["x"]] # new

best owner [["Fromx"]] <- B[["y"]] # new

best owner [["Fromy"]] <- B[["y"]] # new

```
if (answer c[["score"]] > bestsore) {

bestanswer <- answerc

bestscore <- answerc [["score"]]

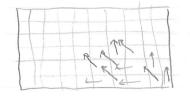
bestanswer [["fromx"]] <- C[["x"]]

bestanswer [["Fromy"]] <- C[["y"]]

o
```

This internation will end up in the plotable data Frame, so he can add a gaplot layer will some arrows representing where each 'best' come From. We also add a slight amount of jitter to the arrow and points For rendability:

Now our visualization shows, for each cell, which subproblem produced the best onswer.



etc