

# SE446: Big Data Engineering

## Week 3A: MapReduce Concepts

Prof. Anis Koubaa

SE 446  
Alfaisal University

[https://github.com/aniskoubaa/big\\_data\\_course](https://github.com/aniskoubaa/big_data_course)

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جامعة الفيصل

# Today's Agenda

- 1 Introduction to MapReduce
- 2 The MapReduce Programming Model
- 3 Detailed Example: Crime Analysis
- 4 Key Concepts
- 5 When to Use MapReduce
- 6 Summary

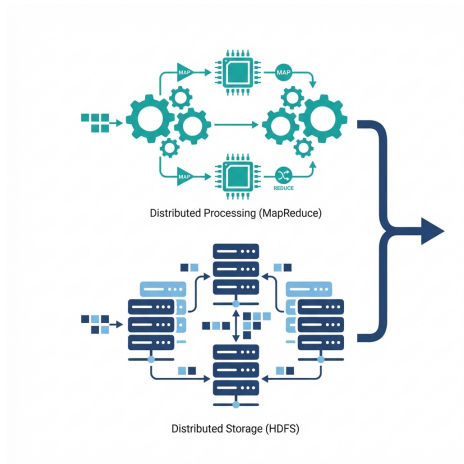
# Recap: Why HDFS Alone Is Not Enough

## HDFS provides:

- ✓ Distributed storage
- ✓ Fault tolerance
- ✓ Scalability

## But how do we process data?

- Read 10 TB from HDFS?
- Analyze on one machine?
- ✗ Bottleneck!



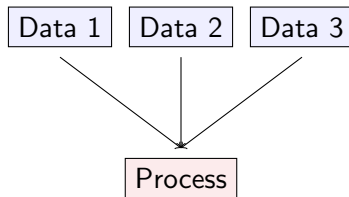
# The Big Idea: Move Computation to Data

## Traditional Approach

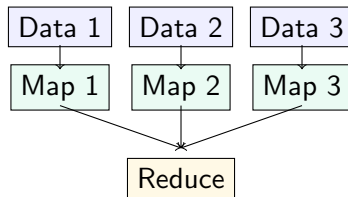
Move data to computation → Network bottleneck

## MapReduce Approach

Move computation to data → Process locally, combine results



**Traditional**



**MapReduce**

# MapReduce: Origin Story

- **2004**: Google publishes landmark paper
- **Problem**: Index the entire web (billions of pages)
- **Solution**: Simple programming model for distributed processing
- **2006**: Doug Cutting implements open-source version → Hadoop

## Key Insight

Many data processing tasks follow the same pattern:

- 1 Process each record independently (Map)
- 2 Group by some key
- 3 Aggregate groups (Reduce)

# MapReduce in One Slide

**map**(key, value)  $\rightarrow$  list(key', value')

**reduce**(key', list(value'))  $\rightarrow$  list(value'')

## Map Phase:

- Process each input record
- Emit (key, value) pairs
- Runs in parallel across nodes

## Reduce Phase:

- Receive all values for a key
- Aggregate/combine values
- Produce final output

# Crucial Concept: Mapper Input

## Common Confusion

Students often ask: "What is the key passed to the Mapper?"

It is **NOT** the same as the key you emit!

- **Input Format:** Defines how files are read.
- **TextInputFormat** (Default):
  - **Key:** Byte offset of the line (e.g., 0, 104, 256...)
  - **Value:** The actual line of text

## Typical Mapper Implementation

We usually **ignore** the input key (offset) and just process the value (text).

# The Hidden Step: Shuffle & Sort (System Phase)

## Important

Between **Map** and **Reduce**, Hadoop runs a hidden phase called **Shuffle & Sort**.

### What Shuffle does (automatically):

- 1 Collects mapper outputs (key, value)
- 2 Sends each key to the correct reducer
- 3 Sorts keys and groups all values per key

### Tiny Example:

Mapper outputs:

```
(Hello,1) (World,1) (Hello,1)  
(World,1) (Hello,1)
```

After Shuffle:

```
Hello → [1,1,1]  
World → [1,1]
```

## Key Message

You **do not write Shuffle code**.

It is done by the framework (system-level).

## Why it matters

Shuffle is often the **slowest** step (network transfer).



# Classic Example: Word Count

**Input:** 3 documents with text

**Goal:** Count occurrences of each word

**Mapper (Python):**

```
def mapper(doc_id, text):  
    for word in text.split():  
        emit(word, 1)
```

**Input:** "Hello World Hello"

**Output:**

- (Hello, 1)
- (World, 1)
- (Hello, 1)

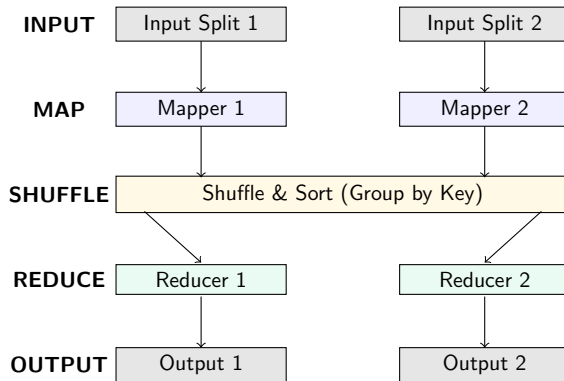
**Reducer (Python):**

```
def reducer(word, counts):  
    emit(word, sum(counts))
```

**Input:** (Hello, [1, 1])

**Output:** (Hello, 2)

# Data Flow: Step by Step



# Data Flow Example: Step-by-Step (Word Count)

## 1) Input Splits (from HDFS)

**Split 1:** "Hello World"

**Split 2:** "Hello Hadoop"

## 2) Map Output (each mapper emits pairs)

**Mapper 1 emits:**

- (Hello,1)
- (World,1)

**Mapper 2 emits:**

- (Hello,1)
- (Hadoop,1)

## 3) Shuffle & Sort (System-level, automatic)

Group all values with the same key:

- Hello → [1,1]
- World → [1]
- Hadoop → [1]

## 4) Reduce Output (final result)

- (Hello, 2)
- (World, 1)
- (Hadoop, 1)

**Map** creates pairs → **Shuffle** groups by key → **Reduce** aggregates

# Real-World Example: Chicago Crime Data

## Dataset: Chicago Crime Reports

Real-world dataset from the City of Chicago (6GB+, millions of rows).

### Key Columns:

- ID, Case Number: Unique identifiers
- Date: Timestamp of incident
- Primary Type: **Category** (THEFT, BATTERY...)
- Description: Detailed description
- Arrest: Boolean (True/False)

## Sample Data (CSV):

Date	Primary Type	Arrest
01/01/2023	THEFT	False
01/01/2023	BATTERY	True
01/02/2023	THEFT	False
01/02/2023	ASSAULT	True
...	...	...

## The Challenge

**How many crimes of each type occurred in total?**

# Stop & Think: Model the Solution

Before Looking at the Code...

Problem: **Count crimes by type**

 What should my KEY be?

*Think for 15 seconds...*

 What should my VALUE be?

*Think for 15 seconds...*

## Stop & Think: Model the Solution

Before Looking at the Code...

Problem: **Count crimes by type**



What should my KEY be?



What should my VALUE be?

→ The Answer

**Key** = Crime Type (e.g., "THEFT") — **Value** = 1 (for counting)

# Crime Count: Mapper

```
def crime_mapper(record):  
    """  
    Input: One crime record (dictionary)  
    Output: (crime_type, 1)  
    """  
    crime_type = record['Primary Type']  
    return (crime_type, 1)  
  
# Example  
record = {'Primary Type': 'THEFT', 'District': 11, ...}  
crime_mapper(record) # Output: ('THEFT', 1)
```

**Key Point:** Mapper processes ONE record at a time, emits (key, value)

# Crime Count: Reducer

```
def crime_reducer(crime_type, counts):  
    """  
    Input: crime_type (key), list of counts (values)  
    Output: (crime_type, total_count)  
    """  
    total = sum(counts)  
    return (crime_type, total)  
  
# Example  
crime_reducer('THEFT', [1, 1, 1, 1, 1])  
# Output: ('THEFT', 5)
```

**Key Point:** Reducer receives ALL values for ONE key



# Complete Data Flow (Color-Coded)

Phase	What goes in	What happens	What comes out
<b>Input</b>	CSV file	Split into individual records	One record at a time
<b>Map</b>	{"Primary Type": "THEFT"}	Extract key and emit count	("THEFT", 1)
<b>Shuffle</b>	All mapper outputs ( <i>system-controlled</i> )	<b>Group values by key</b>	"THEFT" → [1,1]
<b>Reduce</b>	"THEFT", [1,1]	Aggregate values (sum)	("THEFT", 2)

**Data shape evolution:** Record → (Key,Value) → Key→[Values] → Final Result

# Key Design Cheatsheet

## Designing Your Key is 90% of the Battle

Goal	Map Output Key	Map Output Value
<b>Counting</b> (e.g., Word Count)	Item "apple"	1 1
<b>Grouping</b> (e.g., Crimes by Type)	Group ID "THEFT"	Full Record {id: 1, loc: "Street" ...}
<b>Filtering</b> (Map-only job)	- (None)	- (None)
<b>Inverted Index</b> (Search Engine)	Keyword "Hadoop"	Document ID "doc1.txt"

## Quick Quiz: Model These Problems (1/2)

For each scenario, identify the Key and Value:

1. Total sales per store

Key: \_\_\_\_\_

Value: \_\_\_\_\_

2. Website visits per country


Key: \_\_\_\_\_

Value: \_\_\_\_\_

3. Average temperature per city

Key: \_\_\_\_\_

Value: \_\_\_\_\_

 Hint: For averages, what EXTRA info do you need?

## Quick Quiz: Model These Problems (1/2) - Answers

For each scenario, identify the Key and Value:

1. Total sales per store

Key: **Store ID**

Value: **Sale Amount**

2. Website visits per country

Key: **Country**

Value: **1**

3. Average temperature per city

Key: **City**

Value: **(temp, 1)**

✓ Need **sum** and **count** to compute average!

## Quick Quiz: Model These Problems (2/2)

### More challenging scenarios:

#### 4. Inverted Index (Search Engine)

Build an index: which documents contain each word?

**Key:** \_\_\_\_\_

**Value:** \_\_\_\_\_

#### 6. Grouping: All employees per department

Collect all records belonging to the same group.

**Key:** \_\_\_\_\_

**Value:** \_\_\_\_\_

#### 5. Filtering: Find all orders $\geq$ \$1000

Only output records matching a condition.

**Key:** \_\_\_\_\_

**Value:** \_\_\_\_\_



## Quick Quiz: Model These Problems (2/2) - Answers

### More challenging scenarios:

#### 4. Inverted Index (Search Engine)

Build an index: which documents contain each word?

**Key:** Word

**Value:** Document ID

#### 5. Filtering: Find all orders $\geq$ \$1000

Only output records matching a condition.

**Key:** (any / None)

**Value:** Full record

**i** Map-only job: emit if condition met, else skip!

#### 6. Grouping: All employees per department

Collect all records belonging to the same group.

**Key:** Department ID

**Value:** Employee record

**i** Reducer concatenates all employee records!

# Input Splits vs. HDFS Blocks

## HDFS Blocks (Storage)

- Physical division of data.
- Fixed size (e.g., 128MB).
- Replicated for safety.

## Input Splits (Processing)

- Logical division for Mapper.
- **1 Split = 1 Mapper Instance.**
- Usually, Split Size  $\approx$  Block Size.

## Exam Key Takeaway

**HDFS Blocks** are for storage reliability.

**Input Splits** control the number of mappers.

# Visual Example: 400MB File Processing

## 1. Original File (400MB)

mydata.csv (400MB)



## 2. HDFS Blocks (Storage)

Block 1  
128MB

Block 2  
128MB

Block 3  
128MB

Block 4  
16MB



## 3. Input Splits (Processing)

Split 1  
Mapper 1

Split 2  
Mapper 2

Split 3  
Mapper 3

Split 4  
Mapper 4

### Key Points:

- ✓ 4 blocks = 4 replicated chunks
- ✓ 4 splits = 4 parallel mappers
- ✓ Usually 1 block  $\leftrightarrow$  1 split



# The Shuffle Phase

## Often Overlooked but Critical!

Shuffle is the **most expensive** phase - involves network transfer

### What happens during Shuffle:

- 1 Mapper outputs are **partitioned** by key
- 2 Data is **sorted** by key
- 3 Data is **transferred** to appropriate reducer
- 4 Values for same key are **grouped** together

### Example:

- Mapper 1 emits: (A, 1), (B, 2), (A, 3)
- Mapper 2 emits: (B, 4), (A, 5)
- After shuffle:  $A \rightarrow [1, 3, 5]$ ,  $B \rightarrow [2, 4]$

After Shuffle, the reducer receives:

- "THEFT" → [1, 1, 1, 1, 1]
- "ASSAULT" → [1, 1, 1]
- "BATTERY" → [1, 1]

## ? Questions:

- 1 How many THEFT records were in the input?
- 2 What did the Mapper emit for each THEFT record?
- 3 What will the Reducer output for THEFT?

## ✓ Answers

- 1 5 THEFT records
- 2 ("THEFT", 1)
- 3 ("THEFT", 5)

## Key Insight

# Why Do We Sort Before Reducing?

The "Shuffle" phase guarantees that the Reducer receives keys in **sorted order**.

## Why is this mandatory?

- 1 **Grouping:** All values for Key  $K$  must be contiguous to be passed as a single list (iterator).
- 2 **Memory Efficiency:** The Reducer doesn't need to fit all data in RAM. It can "stream" through the sorted keys.

### Golden Rule

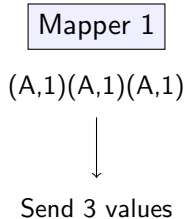
Reducers *never* see unsorted keys.

# Combiners: Local Aggregation

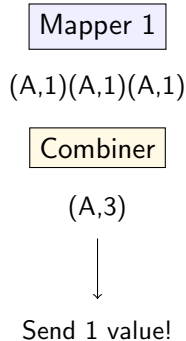
**Problem:** Shuffle transfers lots of data

**Solution:** Combine locally before shuffle!

## Without Combiner



## With Combiner



**Note:** Combiner must be associative and commutative (like sum, max, min)

# Critical Nuance: Combiners are Optional

## The "Mini-Reducer" Rule

A Combiner is an **optimization**, NOT a guarantee.

Hadoop may run the combiner:

- **Zero times**: If the spill file is small.
- **Once**: Standard case.
- **Multiple times**: During merge sorts on disk.

## Implication for Code

Your Combiner logic must be:

- ① **Associative** and **Commutative** (Order doesn't matter).
- ② Safe to run repeatedly (Idempotent-ish).
- ③ Input/Output types must match the Reducer input.

# Partitioner: Key Distribution

**Default:** Hash partitioning ( $\text{hash}(\text{key}) \% \text{num\_reducers}$ )

## Why customize?

- Ensure related keys go to same reducer
- Balance load across reducers

**Example:** Analyzing crimes by year

- Key: "2023-THEFT", "2023-ASSAULT", "2024-THEFT"
- Partitioner: Extract year, send same year to same reducer
- Result: Each reducer handles one year's data

# The Problem of Data Skew (Hot Keys)

## Scenario:

- You count words in Wikipedia.
- Word "the" appears 1 billion times.
- Word "zygote" appears 100 times.

**Result:** The reducer handling "the" runs for 5 hours. All other reducers finish in 1 minute. The job waits for the slow reducer ("Straggler").

## Impact

**1 Reducer bottleneck = Slow Job**

## Solutions:

- **Salt keys:** "the" → "the-1", "the-2" ...
- **Custom Partitioner**
- **Combine** early

# ⚡ Design Challenge: Handling Skewed Data

## Scenario: Analyze 10 Billion Tweets

- 5 billion mention “Taylor Swift” 🎵
- 5 billion mention everyone else combined

### Standard Approach:

- Key = celebrity\_name
- Value = 1

❌ **Problem:** One reducer handles 5B values while others are idle!

### 💡 Your Challenge:

How would you modify the **KEY** to distribute the load?

*Hint: Can you “spread” Taylor Swift across multiple reducers?*

### ✓ Solution: Salted Keys

Key = “Taylor Swift-1”, “Taylor Swift-2”, ..., “Taylor Swift-100”


Then run a second MapReduce job to combine the partial counts!



# Fault Tolerance: The "Secret Sauce"

**Question:** What happens if a node crashes during a 10-hour job?

## Mapper Failures

- Master detects failure.
- Reschedules the task on another node containing the data replica.
-  **Re-execution is safe** (Idempotent).
- *Note: Output is local, so it must be re-run.*

## Reducer Failures

- Master detects failure.
- Reschedules the task on another node.
- Retrieves mapper outputs again from source nodes.

*User code does NOT need to handle checkpoints or failures. The framework handles it!*

## One Output File Per Reducer

- HDFS does not support concurrent writes to a single file.
- Each Reducer writes its own part file.
- **Naming Convention:** part-r-00000, part-r-00001, ...

**Practical Tip:** If you want exactly **one** final output file, set `NumReduceTasks = 1`.  
(Warning: This kills parallelism for the Reduce phase!)

# MapReduce: Good vs Bad Use Cases

## ✓ Good for:

- Batch processing
- Large datasets (TB+)
- Embarrassingly parallel tasks
- Aggregations, counts, sums
- Log analysis
- ETL pipelines

## ✗ Bad for:

- Real-time queries
- Iterative algorithms
- Small datasets
- Interactive analysis
- Graph processing
- Machine learning

## Rule of Thumb

If you need results in  $\leq$  1 minute, MapReduce is probably wrong choice

# MapReduce Limitations

- ① **High Latency:** Minutes to hours for jobs
- ② **Disk I/O:** Intermediate results written to disk
- ③ **Only Two Phases:** Complex workflows need chaining
- ④ **No Iteration:** Each job reads from scratch

## The Evolution

MapReduce limitations led to:



- **Spark:** In-memory processing, DAG execution
- **Flink:** True streaming
- **Hive:** SQL on MapReduce


We'll cover these in later weeks!

# Key Takeaways

- 1 **MapReduce** = programming model for distributed processing
- 2 **Map**: Process records in parallel → emit (key, value)
- 3 **Shuffle**: Group by key (automatic, expensive)
- 4 **Reduce**: Aggregate values per key → final result
- 5 **Data locality**: Computation moves to data, not vice versa
- 6 **Combiner**: Local aggregation to reduce shuffle

## The MapReduce Design Formula

 <b>KEY</b>	<b>What do I want to GROUP BY?</b>
 <b>VALUE</b>	<b>What do I want to AGGREGATE?</b>

 **Write this on your exam notecard!**

# Next Session: Hands-On Practice

## Week 3B: Implementing MapReduce in Python

- Implement mappers and reducers
- Complete crime analysis exercises
- Start Milestone 2

### Pre-class preparation:

- Watch: Python MapReduce Tutorial
- Clone team repo and pull latest
- Review Chicago crime dataset