## SoK: Security and Privacy of Blockchain Interoperability

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Abstract—Recent years have witnessed significant advancements in cross-chain technology. However, the field faces two pressing challenges. On the one hand, hacks on cross-chain bridges have led to monetary losses of around 3.1 billion USD, highlighting flaws in security models governing interoperability mechanisms and the ineffectiveness of incident response frameworks. On the other hand, users and bridge operators experience restricted privacy, which broadens the potential attack surface.

In this paper, we present the most comprehensive study to date on the security and privacy of blockchain interoperability. We employ a systematic literature review, yielding a corpus of 212 relevant documents, including 58 academic papers and 154 gray literature documents, out of a pool of 531 results. We systematically categorize 57 interoperability solutions based on a novel security and privacy taxonomy. Our dataset, comprising academic research, disclosures from bug bounty programs, and audit reports, exposes 45 cross-chain vulnerabilities, 4 privacy leaks, and 92 mitigation strategies. Leveraging this data, we analyze 18 notable bridge hacks accounting for over 2.9 billion USD in losses, mapping them to the identified vulnerabilities.

Our findings reveal that a substantial portion (65.8%) of stolen funds originates from projects secured by intermediary permissioned networks with unsecured cryptographic key operations. Privacy-wise, we demonstrate that achieving unlinkability in cross-chain transactions is contingent on the underlying ledgers providing some form of confidentiality. Our study offers 17 critical insights into the security and privacy of cross-chain systems. We pinpoint promising future research directions, underscoring the urgency of enhancing security and privacy efforts in cross-chain technology. The identified improvements have the potential to mitigate the financial risks associated with bridge hacks, fostering user trust in the blockchain ecosystem and, consequently, wider adoption.

Keywords: Cross-chain, Security, Privacy, Vulnerabilities, Blockchain Technology, Cryptocurrency, Financial Losses, Systematic Literature Review, Taxonomy, Mitigations.

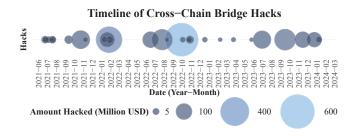


Figure 1. Timeline of cross-chain bridge hacks from May 2021 to February 2024. The dataset includes 33 bridge hacks amounting to over 3.2 billion USD

## 1. Introduction

Blockchain interoperability is paramount for realizing the full potential of blockchain technology. As the landscape evolves, interoperability is gaining momentum in use cases such as bridging liquidity fragmentation, optimizing decentralized exchange (DEX) trades, enhancing scalability through sharding [1], expanding through sidechains [2], and enabling asset exchanges and transfers across platforms [3]. Stepping back to 1996, Wegner stated that "interoperability is the ability of two or more software components to cooperate despite differences in language, interface, and execution platforms" [4]. However, achieving interoperability across blockchains - distributed systems where mutual trust is often absent – adds a dimension of complexity. Here, the challenge is not simply to sync n software components but rather to integrate n distributed systems, each with its unique challenges, encompassing safety, liveness, accountability, and centralization [5], [6]. Such orchestration is performed by interoperability mechanisms (IMs) [7]. The different transactional models, consensus mechanisms, and cryptographic primitives in the networks escalate this challenge. Despite these challenges, the domain has seen prolific contributions from scholars, providing solutions, novel architectures, and varied use cases [8]–[16]. A recurring theme in these studies underscores the pressing need for rigorous research on IM security and privacy.

The secure interoperation of different blockchains involves establishing a new security boundary that depends on the security of at least two existing networks and involves multiple design trade-offs [12], [17]–[19]. Simultaneously, the disclosure of cross-chain transactions might be sensitive and can reveal much about the entities holding the data or the data itself [20]–[22]. We underline that striking a balance

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between security and privacy is a must. While privacy is essential, accountability is equally necessary to deter and penalize misconduct and maintain protocol fairness.

Since May 2021, mounting losses due to bridge hacks have exceeded 3.1B USD and have been recurrent since then (cf. Figure 1). According to Immunefi [23], white-hat hackers have been compensated over 20M USD through bug bounty programs, preventing potential losses of a staggering 1B USD. Moreover, cross-chain bridge hacks have raised to the top of the DeFi incidents leaderboard [24]–[26], emerging as the preferred target of cybercriminals. The present scenario, as of mid-2023, paints a grim picture with rampant hacks [27]–[30]. Consequently, the total value locked (TVL) in cross-chain bridges has nose-dived from its zenith at 58B USD in early 2022 to a mere 5.5B USD in December 2023, a downfall also attributed to the decrease in asset prices in the bearish market [31], [32]. We hypothesize that intertwined cross-chain systems, in conjunction with already well-studied vulnerability-prone smart contracts, be it in bytecode or higher-level language dimensions [33], have amplified the risk exposure of these protocols. Due to the large amounts of funds involved, these protocols are attractive honeypots for attackers, making them highly soughtafter targets in the three interoperability modes we examine: asset exchanges, asset transfers, and data transfers [19], [34].

## 1.1. Research Questions and Contributions

In this paper, we systematize knowledge about the security and privacy of blockchain interoperability solutions. So far, this information has been scattered among multiple unstructured and sometimes unidentified sources. To achieve this goal, in Section 2, we first provide the relevant background knowledge on blockchain interoperability necessary to understand this work. Then, we provide three contributions answering a specific research question (RQ), pointing to a relevant section discussing it.

RQ1: What are the different security- and privacycentric goals used in blockchain interoperability, and what are the technical building blocks that guarantee them? Security-wise, we define a set of properties inspired by the distributed system literature and explore relevant security approaches used to ensure safety and liveness for each IM. In terms of privacy, we analyze the existing privacyenabling approaches for IMs and categorize them according to the guarantees of user and bridge operator anonymity, confidentiality of transactional data, and unlinkability of transactions across multiple blockchains. We present a systematization of knowledge comparing 57 IMs based on the relevant properties and approaches defined. Over half (54%) of the classified papers have been published since 2022, highlighting the timeliness of this work. Sections 3 and 4.1 answer this RQ.

RQ2: What are the cross-chain vulnerabilities, attack vectors, privacy leaks, and mitigations currently known, and how are they mapped to past incidents? Our third research inquiry requires us to investigate cross-chain attacks and privacy leaks in the literature, focusing on the vulnerabilities

that made them possible. We categorize 45 identified vulnerabilities into four distinct security layers (cf. Section 3.1) and present 4 privacy leaks. We identify possible mitigations to all vulnerabilities and leaks. We also **foster synergies between academia and industry** by examining 18 cross-chain hacks that account for more than 2.9B USD hacked and comparing them with the theoretical vulnerabilities found in academic research. We pinpoint the disparities between the existing research findings and their practical application by providing additional 6 **strategic insights**. Section 4.2 answers this RO.

RQ3: Based on the existing gaps, what are potential best practices and avenues for future research to enhance the security and privacy of cross-chain protocols? In a time when the industry is actively seeking stability, we observe that the design of cross-chain solutions remains largely ad hoc, with each solution custom-crafted for specific blockchains or applications. Through a comprehensive analysis of existing studies, we present a collection of best practices and future research avenues. We provide initial insights that protocol designers, developers, and analysts can use as a foundation for further research and development. Sections 4.3 and 4.4 answer this RQ.

Finally, Sections 5 and 6 briefly discuss the contribution to related literature and present our forward-looking remarks and insights, respectively.

## 1.2. Research Methodology

To address the research questions, we structure our methodology into two phases. Firstly, we crawled papers from 2015 using the Google Scholar keyword search. Secondly, due to the unstructured practices in the area, we included multiple gray literature resources focusing on past cross-chain hacks, audit reports, vulnerabilities, and disclosures through bug bounty programs. In total, we analyzed 242 relevant documents, 58 from Google Scholar, 111 from snowballing, forward reference search, and by setting up Google alerts, 10 audit reports from reputable entities such as *Certik, Chainsecurity, Consensys, Halborn* and *Trail of Bits*, and 63 additional gray literature documents. A more thorough description of the methodology is presented in Appendix B.

**Data and Code Availability:** The data and code for ensuring replicability are available on GitHub, accessible at the following URL: https://github.com/RafaelAPB/SoKSP BlockchainInterop.

## 2. A Primer on Blockchain Interoperability

Blockchain interoperability allows data and value to flow across different domains and is facilitated by an IM. The domains under focus in this paper are distributed ledgers  $\{l_1, l_2, ..., l_k\}$ . The IM can be designed and deployed in multiple ways, depending on the required guarantees (e.g., centralized or decentralized, distributed or non-distributed).

## 2.1. Interoperability Modes

The literature [1], [7], [34]–[37] agrees on the existence of three modes of interoperability: **asset exchanges** (**AE**), **data transfers** (**DT**), and **asset transfers** (**AT**). Different interoperation modes require distinct protocol architectures, providing different security and privacy guarantees.

Consider accounts  $\mathcal{A}_1$  and  $\mathcal{A}_2$  in ledgers  $l_1$  and  $l_2$ , respectively. Asset exchange protocols allow untrusted parties to exchange assets in different networks. For example, asset X owned by  $\mathcal{A}_1$  on  $l_1$ , can be exchanged for asset Y owned by  $\mathcal{A}_2$  on  $l_2$ . At the end,  $\mathcal{A}_2$  owns X on  $l_1$ , and  $\mathcal{A}_1$  owns Y on  $l_2$ . An asset exchange can be mediated by a trusted party or run directly between both parties through a secure off-chain communication channel.

Asset transfer protocols encompass locking, or burning, an asset in the source chain and creating (minting) a representation of that asset in the target chain – *lock-mint* or *burn-mint* pattern, respectively [11]. Once the asset is escrowed in the source chain (by either a centralized party, a multisignature, or a smart contract), the verification occurs in the target chain. It can be performed by replicating the source chain's consensus mechanism in the target chain [38], [39] or using a proof-based mechanism such as zero-knowledge proofs [40]–[42].

Data Transfers generalize interoperability. Information written in one domain is transferred (or copied) to another accompanied by proof. An example is the payload of a blockchain view [43], where DLT Gateways facilitate the process, running a gateway-to-gateway protocol [44].

## 2.2. The Source of Truth: Underlying Blockchains

Although the primary focus of this paper is not the security of blockchains, we have to recognize their role as a critical dependency for cross-chain protocols. The reasons are clear: if a transaction  $t_2$  is issued on  $l_2$  based on a rewritten transaction  $t_1$  on  $l_1$ , there is a safety violation. For instance, consider a transaction locking an asset in  $l_1$  and a representation minted in  $l_2$ . If  $t_1$  reverts, the asset in  $l_2$  becomes unbacked [45]. The probability of this is the same as of a 51% Attack on the source chain. Proof of Work-based chains subject to forks and Proof of Stake-based chains subject to long-range attacks [46] are some examples of vulnerable chains. Alternatively, chains with instant or near-instant finality, such as those using variants of PBFT [47], do not suffer from the same problem, at the cost of allowing fewer nodes.

#### 2.3. Cross-Chain Events, Transactions and Rules

The concepts of cross-chain events, transactions, and rules are important to understand this work. Transactions issued in one domain trigger internal state changes and emit events based on the operations performed. Cross-chain events are composed of native and non-native domain attributes. Native attributes are retrieved from the events emitted in the underlying domains. Non-native attributes

are additional metadata that only hold relevance in crosschain environments, such as a domain identifier, a global clock, a token price, or other off-chain information. Metadata is published on-chain by decentralized oracles, and its correctness is given by the correctness of the oracle network and according to the agreement between entities to perform cctxs.

**Definition 1** (Valid Cross-Chain Event). A cross-chain event e is valid iff its metadata is correct\*, and every local transaction  $t \in e$  is final.

The composition of multiple cross-chain events stands for state changes across several domains. We call this composition a cross-chain transaction (cctx). To evaluate the validity of a cctx, events must be verified against  $cross-chain\ rules$  that define the expected behavior. A rule for an asset transfer protocol might indicate that there must not be an event minting an asset in  $l_2$  before an event locking the corresponding asset in  $l_1$ . Given some business logic, one can create arbitrarily complex cross-chain rules.

**Definition 2** (Valid Cross-Chain Transaction). A cross-chain transaction cctxis valid iff every cross-chain event  $e \in cctx$  is valid, and all cross-chain events enforce the defined cross-chain rules.

## 3. Cross-Chain Security and Privacy Model

This section presents definitions and formalizes the relevant properties for cross-chain security and privacy. Additionally, it overviews the taxonomy of security- and privacy-enabler approaches for IMs.

#### 3.1. Security Layers

The security of a cross-chain system can be decomposed into the security of several layers. The existing literature supports similar breakdowns [48]. The *Network Layer* of forms the bedrock. It concerns the blockchains that underlie a cross-chain solution (cf. Section 2.2). Above that, the *Protocol Layer* addresses the different architectural decisions, including defining the actors, their roles, and responsibilities. Further up the stack, we encounter the *Implementation Layer*. It encompasses the entire implementation lifecycle, including off-chain (e.g., relayers, oracles) and on-chain (e.g., smart contracts, protocols) components of the IM. Finally, at the top, the *Operational Layer* specifies the procedures for deploying, upgrading, operating, and monitoring on- and off-chain components.

## 3.2. Security Properties

Based on our comprehensive literature review, we propose a set of properties that characterize a secure crosschain system, inspired by Avizienis et al. [49]. We define

\*. if metadata can be evaluated – e.g., the price of the token being transferred is within an agreed interval

three security properties for IMs: integrity, accountability, and availability.

**Definition 3** (Integrity). Consider an IM and a set of cross-chain rules  $\zeta$ . Integrity is guaranteed iff every generated cctx respects  $\zeta$ .

**Definition 4** (Accountability). An IM is accountable iff these conditions hold: 1) the metadata of any event  $e \in cctx$  is public, or at least verifiable<sup>†</sup>; 2) for every integrity violation attempt in  $\zeta$ , there is a mechanism to prove it; and 3) there is attempts (e.g., third-party, blockchain smart contract).

**Definition 5** (Availability). An IM guarantees availability iff it is always able to process (validate, issue, or relay) valid cctxs.

### 3.3. Security Approaches

The security approaches found in the literature are summarized in Table 1. Interoperability can be facilitated by **Trusted Third Parties**  $(SA_1)$  that manage the whole process end-to-end. Alternatively, in **Distributed Trust**  $(SA_2)$  approaches, trust is placed on a mechanism run within external intermediary network. In **Native State Validation**  $(SA_3)$  approaches, off-chain parties only relay information to the target chain, and the validity of cctxs is verified on-chain. Finally, **Local Verification**  $(SA_4)$  approaches rely on the end users to manually validate each other's transactions on the respective chains. These approaches can be further divided as follows (see table).

**3.3.1.**  $(SA_{11})$  **Centralization.** Trusted third parties can facilitate interoperability by holding user funds and issuing transactions in  $\{l_1, l_2, ..., l_k\}$ , or functioning solely as relay services [36], [44]. Trust is placed on the reputation of the managing party [44], [55], [61], [95]. Centralization raises concerns about safety, liveness, and fairness compromises (e.g., bankruptcy, censorship, MEV, money laundering) [59], [62]. On the other hand, it offers greater performance since no agreement between parties is required [50]. Additionally, it enforces accountability as entities must comply with KYC policies [54].

**3.3.2.**  $(SA_{12})$  **Trusted Computation.** Instead of trusting the reputation of one entity, one can leverage Trusted Execution Environments (TEE) [96] with attestable computation to manage or orchestrate cctxs [21], [60], [61]. TEE-based cross-chain solutions focus on protecting sync committee members' private keys [74], operations on key shares between operators [97], or generate proofs published on multiple chains [74], [98], [99]. Note that attestation keys are still provided by a manufacturer that might, e.g., embed malicious code into the hardware or spoof data [62], [100].

**3.3.3.**  $(SA_{21})$  **Permissionless Networks.** Compared to decentralized systems, centralized ones are simpler to implement, faster, and cheaper. However, as witnessed multiple

†. for example, private data verifiable through proving systems

times in 2023, they involve security risks [28], [101]–[103]. A solution is to rely on the distribution of power among multiple entities to enforce cross-chain rules and validate cctxs – i.e., through an intermediary distributed network that verifies and maintains proof of actions of other chains [63], [104]–[106]. The security of the IM is driven by the security of the network, both at a technical and financial level, e.g., in PoS-based networks, the overall value protected by the protocol should not exceed the stake held by the majority of validators to avoid protocol deviations [107]. For the same reason, intermediary networks using custom low-value native tokens are not recommended due to fluctuations in token price.

**3.3.4.**  $(SA_{22})$  **Permissioned Networks.** Instead of relying on a network in which anyone can join, one can opt for a more controlled environment, where validators are whitelisted and, in some cases, are controlled by reputable or trusted entities. These are based on Proof of Authority [108], Threshold Signature Schemes [71], [74] or Multi-Signatures [109]. When parties are not trusted entities, there is usually an identification service deployed, where parties register beforehand [81]. Alternatively, multi-party computation can be used to build trust in environments with mutually untrusted entities [74], [90], [92], [109]. An interesting advantage of  $\mathcal{SA}_2$  is to incorporate additional security measures in the intermediary network, such as access control and cross-chain rules enforcement, besides what is in the bridge contracts on the source and destination chains [21], [67].

**3.3.5.**  $(SA_{31})$  **Inclusion Proofs.** These approaches involve users providing verifiable evidence of actions triggered on one chain on another [38], [77]. Networks of relayers compete to relay block headers (from the source to the destination chain) used to validate user-provided proofs, e.g., Merkle Proofs [38], [69], [110]. Even if a whole network of relayers colluded, they could only disrupt the system if they possess greater mining/voting power than the rest of the source chain, equivalent to mounting a 51% attack on that network. The security of this scheme is based on the security of the light client in the target chain, which depends on the source chain's consensus mechanism. For instance, PoS light clients for the Ethereum 2.0 sync committees ([111]) [39], [40], [76] do not guarantee accountability for the bridge because no slashing mechanism is in place<sup>‡</sup>.

**3.3.6.**  $(SA_{32})$  **Validity Proofs.** Validity-proof-based bridges rely on proving systems to validate the state of the source chain's consensus mechanism within the target chain [41], [42], [65], [78], [79]. Unlike  $SA_{31}$ , there is no need to understand the consensus logic of other ledgers, as it only requires verifying a succinct zero-knowledge proof (ZKP), which is constant time on zkSNARK-based bridges [112]. Given that circuits are tailor-made to each specific program, the creation of proofs is currently the bottleneck in these

<sup>‡.</sup> https://github.com/ethereum/consensus-specs/issues/3321

TABLE 1. TWO TIER CLASSIFICATION OF SECURITY APPROACHES IN BLOCKCHAIN INTEROPERABILITY ACADEMIC STUDIES. WE PRESENT THE PRIMARY SECURITY APPROACH OF SOLUTIONS THAT EMPLOY VARIOUS.

Security Approach (Tier 1)	Security Approach (Tier 2)	IM Role	References	# (and %)
SA <sub>1</sub> Trusted Third Parties	$SA_{11}$ Centralization $SA_{12}$ Trusted Computation	Centralized Services Trusted Execution Environment	[21], [36], [50]–[59] [60]–[62]	12 (24%) 3 (06%)
SA <sub>2</sub> Distributed Trust	$SA_{21}$ Permissionless Network $SA_{22}$ Permissioned Network	Public Network Validators Whitelisted Network Validators	[63]–[65] [66]–[74]	3 (06%) 9 (18%)
SA <sub>3</sub> Native State Verification	$SA_{31}$ Inclusion Proofs $SA_{32}$ Validity Proofs $SA_{33}$ Fraud Proofs	Relayers Relayers Relayers	[38], [45], [75]–[77] [41], [42], [78], [79] None in academia	5 (10%) 4 (08%) 0 (00%)
SA <sub>4</sub> Local Verification	SA <sub>41</sub> Secret- & Time-based Locks	Off-chain Communication Channel	[80]–[94]	15 (29%)

*Note:* The table categorizes various security approaches (SAs) prevalent in blockchain interoperability research into two tiers. The first tier provides an overarching classification, while the second tier offers a finer granularity. The "IM Role" column denotes the component that takes the role of the Interoperability Mechanism (IM), and the "References" column cites specific studies or implementations that employ the particular approach. The final column quantifies the number and approximate percentage of papers adopting each method, visually represented using cell shading.

systems [40]–[42]. However, new research is improving the efficiency of proof generation, reducing memory demands, and reducing the dependence on trusted setups [105], [113]–[115].

**3.3.7.**  $(SA_{33})$  **Fraud Proofs.** Fraud proofs provide security for cross-chain protocols by optimistically accepting block headers and other proofs [116]. External watchers can submit fraud proofs to challenge this data [117]. Watchers are rewarded by presenting valid fraud proofs, and operators that sent invalid data see their stake slashed. Transactions based on this information are not final until the fraud-proof period elapses. Usually, to avoid relying on synchronous communication between parties, these periods are extended time windows ( $\sim$ 7 days) [117]. Accountability for relayers is guaranteed. To guarantee safety and liveness, there must be a correct watcher online at all times [118].

3.3.8.  $(SA_{41})$  Secret-based and time-based locks. Decentralized atomic swap protocols are commit-reveal schemes based on hash-locks and timelocks [85], or digital signatures [119], [120]. In the former, parties agree on parameters off-chain and have predefined periods in which they must act to complete the protocol. To mitigate atomicity breaches under long-standing crashes, most solutions change the assumptions of synchronous communication by inserting intermediary networks [63], [70], or focus on the usage of premiums, a collateralization technique to compensate users in case of misbehavior or crashes [80], [86], [121]. According to [122] collateralization increases the probability that the protocol is completed successfully. There are alternatives to commit-reveal schemes by avoiding the reliance on explicit time intervals. [94] reveals a secret after the counterparty has performed an agreed number of computation steps. The protocol proposed by [91] reveals a full signature once a partial signature is presented. Note that in cross-chain deal [106] and atomic swap protocols [72], [85], [86], [123], [124], safety is usually not defined in terms of atomicity, but rather if honest nodes do not end up worse off than how they started the protocol.

### 3.4. Privacy Properties

Our survey shows that privacy in cross-chain transactions is a relatively understudied area. In this section, we present the first definition and formalization of generic cross-chain privacy by defining the three most relevant properties: **unlinkability** of *cctxs*, **anonymity** of users and operators, and **confidentiality** of transactional data.

**Definition 6** (Cross-Chain Unlinkability). Consider a cctx between two related accounts  $A_1$  and  $A_2$ , where  $A_1$  might be equal to  $A_2$ . Transactions t and t' issued by  $A_1$  and  $A_2$ , on the source and destination chain, respectively, are said to be unlinked iff an external party cannot infer that t and t' are related to each other.

External parties can infer relationships between transactions using pre-trained models and heuristics, which return a similarity factor – i.e., the probability of two transactions being linked [125]. Heuristics can be related to transaction amounts, asset profiles [44], reused addresses [126], or transaction patterns [127].

**Definition 7** (Cross-Chain Anonymity). Anonymity of A holds iff 1) A cannot be linked to transactions  $t_1, ..., t_k$  it has issued in both ledgers and 2)  $t_i$  and  $t_j, \forall i, j \in [1, k]$  are cross-chain unlinkable.

**Definition 8** (Cross-Chain Confidentiality). *Cross-chain confidentiality holds iff the content of any cross-chain transaction*  $cctx_1$  *issued by an address*  $A_1$  *is indistinguishable from the content of any other cross-chain transaction*  $cctx_2$  *issued by*  $A_1$  *or any other address.* 

The notion of indistinguishability we are trying to capture is similar to IND-CPA: given two cross-chain transactions  $cctx_1$  and  $cctx_2$ , and their raw payloads  $p_1$  and  $p_2$ , respectively, an adversary cannot guess which payload p corresponds to each cctx with a probability higher than 50% (i.e., randomly).

TABLE 2. CLASSIFICATION OF PRIVACY-ENABLER APPROACHES IN BLOCKCHAIN INTEROPERABILITY STUDIES

Privacy Approach	References	# (and %)
PA <sub>1</sub> Zero Knowledge Proofs	[64]–[66], [73], [77]–[79], [81]	8 (47%)
PA <sub>2</sub> Trusted Execution Envir.	[21], [55], [60]	3 (18%)
PA <sub>3</sub> Adaptor Signatures	[87], [92]	2 (12%)
PA <sub>4</sub> Blind Signatures	[93]	1 (06%)
PA <sub>5</sub> Ring Signatures	[74]	1 (06%)
PA <sub>6</sub> Homomorphic Encryption	ı [88], [89]	2 (12%)

*Note:* The table categorizes multiple privacy-enabler approaches (PAs) in blockchain interoperability studies. The first column classifies the approach. The second column cites studies or implementations that employ the particular approach. The right-most column estimates the number and percentage of studies adopting each method.

## 3.5. Privacy-Preserving Approaches

In this section, we overview the main privacy-preserving techniques in the literature, to guarantee at least one of the identified properties. A summary is present in Table 2.

**3.5.1.**  $(\mathcal{PA}_1)$  **Zero Knowledge Proofs.** ZKPs allow proving actions and demonstrate compliance with cross-chain rules without revealing transaction details, involved parties, transaction amounts, or exchange prices [66], [77], [81], [126]. ZKPs are instrumental in validating actions and ensuring compliance with cross-chain rules without disclosing transaction specifics, involved entities, transaction amounts, or exchange rates [66], [77], [81], [126]. These proofs find their applications, e.g., within mixing services. The IM can act as a transaction mixer or leverage existing mixers within the source and destination chains [60], [66], [73].

**3.5.2.**  $(\mathcal{P}A_2)$  **Trusted Execution Environments.** TEEs facilitate computation on private data by preventing data leakage outside the secure enclaves [100]. These promote fairness, for example, in asset exchange protocols with confidential order-matching algorithms [62] (i.e., fairly matching bid and ask orders in exchanges). Due to confidentiality guarantees, the exchange rates are not public, which might help to guarantee unlinkability. TEEs also enable the enforcement of predefined checks to comply with cross-chain rules [21], [62]. A possible approach is to use TEEs as a mixing service [60], but no specific algorithm has been proposed.

**3.5.3.**  $(\mathcal{PA}_3)$  **Adaptor Signatures.** Adaptor signatures allow one party to generate a *pre-signature* on a message associated with a secret, which is guaranteed to provide the secret once the full signature is published [92]. This resembles a commit reveal scheme [87] that underlies HTLCs for asset exchanges. When Party A commits a transaction to collect Party B's assets, it reveals a secret that allows B to redeem A's assets. This protocol avoids publishing the shared secret hash, which makes HTLCs not guarantee unlinkability. It is still possible to analyze on-chain transaction amounts. However, it is unlikely that one can link transactions without knowing the cryptocurrencies exchanged and the exchange rate agreed upon off-chain by both parties.

**3.5.4.**  $(\mathcal{P}A_4)$  **Blind Signatures.** Blind signatures [128] allow users to acquire a signature on a message from a trusted third party without disclosing its specific details. As each blind signature has the same weight, they ensure the unlinkability, anonymity, and confidentiality of crosschain transactions (even with centralized IMs) [93]. They uphold fairness by shielding against centralized surveillance and eradicating the potential for custom ordering.

**3.5.5.**  $(\mathcal{P}A_5)$  **Ring Signatures.** Ring signatures provide set anonymity by obscuring users among a ring of k users, where the likelihood of exposing an individual's identity is 1/k [129]. Notably, revoking anonymity is infeasible, thereby undermining accountability. We acknowledge the existence of different ring signature protocols with trade-offs between, for example, traceability, anonymity, and linkability [130], [131]. The existing literature employs ring signatures to obfuscate sender addresses [109] and protect the identities of intermediary network members, mitigating potential threats such as Denial-of-Service (DoS) attacks [74]. However, due to the inability to identify the signer, in a cctx, the transactional data in the source chain must include the destination address, which compromises unlinkability.

**3.5.6.**  $(\mathcal{PA}_6)$  **Homomorphic Encryption.** Similarly to Adaptor Signatures, one can use homomorphic encryption to solve commit-reveal scheme linkability problems. Both [88] and [89] proposed atomic swap solutions based on HE where different secrets are deployed in each chain, attaining transaction unlinkability. Although there are tools for performing basic operations on encrypted data, more research is required to enable more intricate computations and allow general data transfers while guaranteeing confidentiality. Moreover, these protocols typically come with high on-chain computational costs, heavily influenced by the selected homomorphic functions.

## 4. Status of Cross-Chain Security and Privacy

In this section, we present the results of our work and extensively discuss the most relevant insights. Additionally, we present an extensive list of theoretical cross-chain vulnerabilities, attacks, and mitigations and map them to real-world hacks that account for more than 2.9B USD. We gather all relevant insights and propose guidelines for building secure and robust cross-chain systems.

## 4.1. Comparison Framework

We classify 51 academic papers and 6 industry solutions in light of the security and privacy models presented in the previous sections, and based on a set of performance and usability metrics relevant to both project maintainers and platform users. Industry solutions are deployed in production, accounting for more than 75% of the TVL in crosschain bridges [32]. The classification is presented in Table 3.

**4.1.1. Classification Criteria.** We present the criteria on which we base ourselves to classify the relevant IMs.

TABLE 3. CLASSIFICATION OF BLOCKCHAIN INTEROPERABILITY STUDIES IN ACADEMIA AND INDUSTRY.

50] 2019 S. 53] 2023 S. 54] 2023 S. 55] 2023 S. 56] 2023 S. 58] 2020 S. 58] 2020 S. 57] 2022 S. 51] 2021 S. 55] 2022 S. 55] 2022 S. 56] 2023 S. 59] 2020 S. 60] 2021 S. 60] 2021 S. 66] 2023 S. 67] 2022 S. 66] 2022 S. 66] 2022 S. 67] 2022 S. 68] 2020 S. 67] 2022 S. 68] 2020 S. 69] 2021 S. 69] 2022 S. 69] 2022 S. 68] 2023 S. 68] 2020 S. 68] 2021 S. 68] 2021 S. 68] 2022 S. 68] 2023 S. 68] 2020 S. 68] 2021 S. 68] 2022 S.	Security Approaches  \$SA_{11}   \$	In •	Av	Ac	Dc	Lat	Co	Deixoox: Ammeoohoo	Cf	Un	Α	D/ 1		
531         2023         S.           541         2023         S.           542         2023         S.           561         2023         S.           568         2020         S.           571         2021         S.           571         2022         S.           361         2019         S.           5515         2022         S.           361         2019         S.           661         2021         S.           661         2021         S.           661         2021         S.           661         2022         S.           661         2020         S.           671         2022         S.           681         2019         S.           772         2021         S.           773         2021         S.           771         2022         S.           761         2022         S.           762         2022         S.           771         2022         S.           772         2022         S.           771         2022         S.	$SA_{11} \\ SA_{11} \\ SA_{11} \\ SA_{11}$					Lut	CO	Privacy Approaches	CI	OII	An	IMode	PC	Imp
531         2023         S.           541         2023         S.           542         2023         S.           561         2023         S.           568         2020         S.           571         2021         S.           571         2022         S.           361         2019         S.           5515         2022         S.           361         2019         S.           661         2021         S.           661         2021         S.           661         2021         S.           661         2022         S.           661         2020         S.           671         2022         S.           681         2019         S.           772         2021         S.           773         2021         S.           771         2022         S.           761         2022         S.           762         2022         S.           771         2022         S.           772         2022         S.           771         2022         S.	$SA_{11} \\ SA_{11} \\ SA_{11} \\ SA_{11}$		0	0	0	•	0	-	-	-	-	DT	/	1
56          2023         S.           58          2020         S.           58          2020         S.           58          2020         S.           57          2022         S.           57          2022         S.           51          2021         S.           55          2022         S.           56          2019         S.           60          2021         S.           61          2021         S.           62          2019         S.           63          2020         S.           63          2022         S.           64          2023         S.           71          2021         S.           68          2019         S.           72          2021         S.           66          2023         S.           70          2022         S.           45          2019         S.           75          2022         S.           45          2019         S.           77          2022         S.           78          2023         S.	$SA_{11}$ $SA_{11}$	•	0	0	•	0	_	_1	0	0	0	AT	1	Х
58          2020         S.           21          2021         S.           57          2021         S.           57          2022         S.           51          2021         S.           55          2022         S.           36          2019         S.           59          2020         S.           60          2021         S.           62          2019         S.           63          2020         S.           63          2020         S.           64          2023         S.           67          2022         S.           68          2019         S.           71          2021         S.           68          2019         S.           69          2021         S.           69          2021         S.           45          2019         S.           45          2019         S.           45          2019         S.           45          2019         S.           48          2020         S.           80          2021         S.	$SA_{11}$	0	0	•	0	0	•	_2	0	0	0	AE	1	±
21] 2021 S. 57] 2022 S. 57] 2022 S. 56] 2021 S. 55] 2022 S. 36] 2019 S. 55] 2022 S. 36] 2019 S. 56] 2020 S. 66] 2021 S. 66] 2021 S. 66] 2021 S. 66] 2022 S. 67] 2022 S. 68] 2020 S. 67] 2022 S. 68] 2019 S. 71] 2021 S. 72] 2021 S. 73] 2023 S. 74] 2021 S. 73] 2023 S. 74] 2021 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2021 S. 79] 2021 S. 79] 2022 S. 79] 2022 S. 81] 2020 S. 81] 2020 S. 82] 2021 S. 83] 2020 S. 84] 2020 S. 84] 2021 S. 88] 2020 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2021 S.	$SA_{11}$	•	0	•	0	0	•	-	_	-	_	DT	1	±
571         2022         S.           551]         2021         S.           551]         2021         S.           551]         2022         S.           361         2019         S.           522         2023         S.           591         2020         S.           601         2021         S.           611         2021         S.           622         2019         S.           631         2020         S.           632         2020         S.           633         2020         S.           641         2022         S.           671         2022         S.           681         2019         S.           771         2021         S.           772         2022         S.           761         2022         S.           762         2022         S.           381         2020         S.           4451         2019         S.           7771         2022         S.           781         2022         S.           801         2021         S.	64 64	•	0	•	0	•	•	-	-	-	_	AE	1	<b>√</b>
511         2021         S.           5555         2022         S.           5515         2022         S.           521         2023         S.           522         2020         S.           601         2021         S.           601         2021         S.           661         2022         S.           662         2019         S.           651         2022         S.           663         2019         S.           671         2021         S.           722         2021         S.           731         2022         S.           661         2023         S.           701         2022         S.           762         2021         S.           763         2022         S.           764         2022         S.           775         2022         S.           771         2022         S.           441         2022         S.           441         2022         S.           881         2023         S.           881         2021         S.	$SA_{11}, SA_{12}$	0	0	0	0	0	0	$PA_2$	•	•	-	DT	X	±
55 s         2022         S.           36          2019         S.           52          2023         S.           59          2020         S.           60          2021         S.           60          2021         S.           61          2021         S.           62          2019         S.           63          2022         S.           63          2020         S.           63          2019         S.           72          2021         S.           73          2023         S.           74          2022         S.           66          2023         S.           76          2022         S.           45          2019         S.           45          2019         S.           45          2019         S.           45          2020         S.           79          2022         S.           80          2021         S.           81          2023         S.           82          2022         S.           83          2022         S.	$SA_{11}, SA_{21}$	•	0	0	•	-	•	-	•	0	0	DT	X	Х
36] 2019 S. 52] 2023 S. 52] 2023 S. 56] 2020 S. 66] 2021 S. 66] 2021 S. 66] 2022 S. 67] 2022 S. 68] 2019 S. 72] 2021 S. 73] 2023 S. 74] 2021 S. 76] 2022 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2022 S. 78] 2022 S. 79] 2022 S. 78] 2023 S. 79] 2022 S. 88] 2020 S. 88] 2020 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 88] 2022 S. 88] 2023 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2022 S.	$SA_{11}, SA_{21}, SA_{22}$	0	0	0	0	0	•	_	_	_	_	DT	X	/
52] 2023 S. 559] 2020 S. 569] 2020 S. 660] 2021 S. 661] 2021 S. 661] 2021 S. 662] 2019 S. 663] 2022 S. 663] 2022 S. 663] 2020 S. 67] 2022 S. 67] 2021 S. 77] 2021 S. 77] 2021 S. 77] 2021 S. 78] 2023 S. 74] 2022 S. 66] 2023 S. 76] 2022 S. 75] 2022 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2021 S. 78] 2021 S. 79] 2022 S. 88] 2020 S. 88] 2020 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 990] 2021 S. 991] 2022 S.	$SA_{11}, SA_{22}, SA_{31}$	0	•	•	0	•	•	$PA_2$	•	0	0	AT	<b>√</b>	✓
59         2020         S.           601         2021         S.           602         2019         S.           621         2019         S.           622         2019         S.           631         2020         S.           631         2020         S.           631         2020         S.           631         2022         S.           632         2019         S.           721         2021         S.           722         2021         S.           631         2023         S.           741         2022         S.           661         2023         S.           701         2022         S.           762         2022         S.           763         2022         S.           764         2022         S.           775         2022         S.           781         2020         S.           442         2020         S.           841         2022         S.           881         2022         S.           881         2022         S.     <	$SA_{11}, SA_{31}$	•	0	•	0	0	-	_1	•	•	0	DT	Х	Х
60] 2021 S. 66] 2021 S. 66] 2021 S. 66] 2022 S. 66] 2022 S. 66] 2022 S. 66] 2022 S. 67] 2022 S. 67] 2022 S. 72] 2021 S. 73] 2023 S. 74] 2021 S. 75] 2022 S. 66] 2023 S. 76] 2022 S. 77] 2022 S. 78] 2023 S. 77] 2022 S. 78] 2023 S. 79] 2022 S. 78] 2023 S. 79] 2022 S. 78] 2023 S. 79] 2022 S.	$SA_{11}, SA_{31}$	•	0	0	0	0	-	_1	•	•	0	AT	X	±
61] 2021 S. 62] 2019 S. 62] 2019 S. 63] 2022 S. 64] 2023 S. 63] 2020 S. 63] 2020 S. 63] 2020 S. 67] 2022 S. 71] 2021 S. 72] 2021 S. 73] 2023 S. 74] 2022 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2022 S. 78] 2022 S. 79] 2022 S. 78] 2022 S. 79] 2022 S. 88] 2020 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 89] 2018 S. 99] 2022 S.	$SA_{11}, SA_{41}$	0	0	0	0	0	_	- D.4	_	-	-	AE	<b>√</b>	±
62] 2019 S. 65] 2022 S. 66] 2022 S. 66] 2023 S. 66] 2020 S. 67] 2022 S. 68] 2019 S. 71] 2021 S. 72] 2021 S. 73] 2023 S. 74] 2022 S. 66] 2023 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2022 S. 77] 2022 S. 78] 2022 S. 78] 2020 S. 79] 2021 S. 81] 2020 S. 82] 2020 S. 83] 2020 S. 84] 2021 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2022 S.	$SA_{12}$	0	0	•	0	0	0	$PA_2$		$\mathbf{O}^4$	0	DT	X	±
65] 2022 S. 664] 2023 S. 664] 2023 S. 667] 2020 S. 668] 2019 S. 671] 2021 S. 672] 2021 S. 68] 2019 S. 673] 2023 S. 66] 2023 S. 66] 2023 S. 66] 2021 S. 66] 2023 S. 674] 2022 S. 66] 2021 S. 675] 2022 S. 676] 2022 S. 676] 2022 S. 677] 2022 S. 678] 2022 S. 681] 2020 S. 681] 2020 S. 681] 2020 S. 681] 2021 S. 682] 2022 S. 683] 2022 S. 684] 2021 S. 685] 2018 S. 686] 2022 S. 687] 2020 S. 688] 2021 S. 689] 2018 S.	$SA_{12}$	0	0	•	0	0	0	-,	_	-	-	AE	<b>/</b>	Х
64  2023 S. 63  2020 S. 63  2020 S. 63  2020 S. 66  2019 S. 71  2021 S. 72  2021 S. 73  2023 S. 74  2022 S. 66  2023 S. 76  2022 S. 76  2022 S. 76  2022 S. 76  2022 S. 78  2022 S. 78  2020 S. 78  2020 S. 78  2020 S. 81  2020 S. 82  2021 S. 83  2020 S. 84  2021 S. 88  2022 S. 88  2022 S. 88  2022 S. 88  2021 S. 88  2022 S. 89  2018 S. 89  2018 S. 99  2022 S. 99  2022 S. 99  2022 S.	$SA_{12}$	0	0	0	0	0	0	_2	0	0	0	AE	<b>/</b>	±
63] 2020 S. 667] 2022 S. 667] 2022 S. 77] 2021 S. 77] 2021 S. 77] 2021 S. 78] 2023 S. 79] 2021 S. 66] 2023 S. 70] 2022 S. 75] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2020 S. 77] 2022 S. 78] 2020 S. 78] 2020 S. 78] 2020 S. 78] 2021 S. 78] 2022 S. 78] 2023 S. 79] 2022 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2022 S. 99] 2022 S.	$SA_{21}$	•	•	0	•	0	0	$PA_1$	0	•	0	DT	<b>√</b>	±
667] 2022 S. 668] 2019 S. 71] 2021 S. 72] 2021 S. 73] 2023 S. 74] 2022 S. 66] 2021 S. 66] 2022 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2022 S. 78] 2022 S. 78] 2020 S. 79] 2021 S. 81] 2020 S. 82] 2020 S. 83] 2020 S. 84] 2021 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2022 S. 99] 2022 S. 99] 2022 S. 99] 2022 S.	$SA_{21}^{21}, SA_{32}$	•	•	0	•	-	0	$PA_1$		•	•	AE	X	√
68] 2019 S. 71] 2021 S. 71] 2021 S. 72] 2021 S. 73] 2023 S. 74] 2022 S. 69] 2021 S. 66] 2023 S. 75] 2022 S. 76] 2022 S. 76] 2022 S. 76] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2020 S. 79] 2020 S. 81] 2020 S. 82] 2020 S. 83] 2021 S. 84] 2021 S. 88] 2021 S. 88] 2022 S. 88] 2022 S. 88] 2022 S. 88] 2022 S. 88] 2023 S. 88] 2024 S. 88] 2025 S. 88] 2026 S. 88] 2027 S. 88] 2028 S. 88] 2028 S. 88] 2029 S. 88] 2029 S. 88] 2029 S. 88] 2029 S. 88] 2020 S. 89] 2018 S.	$SA_{21}, SA_{41}$	•	•	0	•	•	•	-	-	-	-	AE	1	Х
71] 2021 S. 72] 2021 S. 72] 2021 S. 72] 2023 S. 73] 2023 S. 74] 2022 S. 66] 2021 S. 66] 2023 S. 76] 2022 S. 76] 2022 S. 76] 2022 S. 76] 2022 S. 78] 2020 S. 78] 2020 S. 78] 2020 S. 78] 2020 S. 78] 2021 S. 88] 2020 S. 88] 2021 S. 88] 2021 S. 88] 2022 S. 88] 2022 S. 88] 2022 S. 88] 2022 S. 88] 2023 S. 88] 2024 S. 88] 2025 S. 88] 2026 S. 88] 2028 S. 88] 2028 S. 89] 2018 S.	$SA_{22}$	•	•	0	•	-	-	-	-	-	-	AT	<b>/</b>	X
72] 2021 S. 73] 2023 S. 73] 2023 S. 74] 2022 S. 66] 2021 S. 66] 2021 S. 66] 2022 S. 76] 2022 S. 77] 2022 S. 45] 2020 S. 45] 2020 S. 48] 2020 S. 78] 2022 S. 78] 2022 S. 78] 2023 S. 79] 2022 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 99] 2022 S. 99] 2022 S. 99] 2022 S.	$SA_{22}$	•	•	•	0	0	0	_	-	-	_	DT	<b>/</b>	_/
73] 2023 S. 74] 2022 S. 74] 2022 S. 76] 2021 S. 76] 2022 S. 77] 2022 S. 75] 2022 S. 75] 2022 S. 76] 2022 S. 77] 2022 S. 78] 2020 S. 78] 2020 S. 78] 2020 S. 78] 2021 S. 78] 2022 S. 78] 2023 S. 79] 2022 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 87] 2021 S. 88] 2022 S. 89] 2018 S. 99] 2022 S. 99] 2022 S. 99] 2022 S.	$SA_{22}$	•	0	0	•	0	0	-	_	-	-	DT	1	✓
74  2022 S. 69  2021 S. 69  2023 S. 70  2022 S. 75  2022 S. 75  2022 S. 75  2022 S. 76  2022 S. 77  2022 S. 38  2020 S. 42  2020 S. 44  2022 S. 80  2021 S. 81  2021 S. 82  2022 S. 83  2022 S. 83  2022 S. 84  2021 S. 85  2018 S. 86  2022 S. 87  2020 S. 87  2020 S. 88  2022 S. 89  2018 S. 90  2021 S. 91  2021 S.	$SA_{22}$	0	0	0	0	0	0	_ D.4	-	-	-	AE	<b>/</b>	±
69] 2021 S. 66] 2023 S. 66] 2023 S. 770] 2022 S. 775] 2022 S. 776] 2022 S. 777] 2022 S. 777] 2022 S. 778] 2020 S. 41] 2020 S. 41] 2020 S. 41] 2022 S. 82] 2020 S. 83] 2021 S. 84] 2021 S. 84] 2021 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 99] 2022 S.	$SA_{22}$	•	0	0	0	0	0	$PA_1$	_	•	•	AT	<b>/</b>	±
76] 2022 S. 45] 2019 S. 45] 2019 S. 38] 2020 S. 42] 2020 S. 42] 2020 S. 42] 2022 S. 78] 2023 S. 80] 2021 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S.	$SA_{22}, SA_{12}$	0	•	•	•	0	0	$PA_5$	0	0	0	AT	<b>√</b>	±
76] 2022 S. 45] 2019 S. 45] 2019 S. 38] 2020 S. 42] 2020 S. 42] 2020 S. 42] 2022 S. 78] 2023 S. 80] 2021 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S.	$SA_{22}, SA_{31}$	•	•	0	0	0	0	_i	•	0	0	DT	Х	±
76] 2022 S. 45] 2019 S. 45] 2020 S. 48] 2020 S. 42] 2020 S. 42] 2020 S. 42] 2022 S. 78] 2023 S. 78] 2022 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 99] 2018 S. 99] 2018 S.	$SA_{22}, SA_{32}$	•	•	0	0	0	0	$PA_1$	•	•	•	AT	X	Х
76] 2022 S. 45] 2019 S. 45] 2019 S. 38] 2020 S. 42] 2020 S. 42] 2020 S. 42] 2022 S. 78] 2023 S. 80] 2021 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S.	$SA_{22}, SA_{41}$	0	•	•	•	•	0	-	-	-	_	AE	<b>√</b>	±
45  2019   S.	$SA_{31}$	•	0	0	•	•	_	-	-	-	-	AE	<b>/</b>	Х
77] 2022 S. 38] 2020 S. 38] 2020 S. 442] 2020 S. 441] 2022 S. 78] 2023 S. 79] 2022 S. 80] 2021 S. 82] 2021 S. 82] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 88] 2022 S. 89] 2018 S. 99] 2018 S. 99] 2022 S. 99] 2022 S. 99] 2022 S.	$SA_{31}$	0	•	0	•	_	0	-	_	-	-	DT	<b>√</b>	<b>√</b>
38] 2020 S. 42] 2020 S. 42] 2020 S. 42] 2020 S. 78] 2022 S. 80] 2021 S. 81] 2021 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S.	$SA_{31}, SA_{21}$	<b>O</b> <sup>6</sup>	•	•	•	0	•	-	-	-	-	AT	<b>√</b>	/
42] 2020 S. 41] 2022 S. 78] 2023 S. 79] 2022 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S.	$SA_{31}, SA_{21}$	$\mathbf{O}_{0}$	0	•	•	•	0	$PA_1$	0	•		AT	X	X
41] 2022 S. (78] 2023 S. (78] 2023 S. (78] 2022 S. (80] 2021 S. (81] 2021 S. (82] 2022 S. (83] 2022 S. (83] 2022 S. (84] 2021 S. (85] 2018 S. (86] 2022 S. (87] 2020 S. (88] 2022 S. (89] 2018 S. (90] 2022 S. (91] 2021 S. (92] 2022 S. (93] 2022 S. (94] 2022 S.	$SA_{31}, SA_{33}$	•	•	0	•	•	•	-	-	-	-	DT	<b>/</b>	/
78] 2023 S. 79] 2022 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S. 93] 2022 S.	$SA_{32}, SA_{22}$	•	•	0	•	0	0	-	-	-	_	DT	<b>/</b>	✓
79] 2022 S. 80] 2021 S. 80] 2021 S. 81] 2021 S. 82] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S. 93] 2022 S.	$SA_{32}, SA_{31}$	•	0	0	0	0	0	_	_	- 7	- 7	AT	<b>/</b>	±
80]     2021     Sz       81]     2021     Sz       82]     2022     Sz       83]     2022     Sz       84]     2021     Sz       85]     2018     Sz       86]     2022     Sz       87]     2020     Sz       88]     2022     Sz       99]     2018     Sz       991     2022     Sz       991     2022     Sz       991     2022     Sz       993     2022     Sz       941     2022     Sz	$SA_{32}, SA_{31}$	•	•	0	•	_	0	$PA_1$	0	•7	•7	AT	<b>/</b>	✓
81]       2021       S.         82]       2022       S.         83]       2022       S.         84]       2021       S.         85]       2018       S.         86]       2022       S.         87]       2020       S.         88]       2022       S.         90]       2022       S.         91]       2021       S.         92]       2022       S.         93]       2022       S.         94]       2022       S.	$SA_{32}, SA_{31}$	•	•	•	-	0	0	$PA_1$	•	•	•	AT	X	±
82] 2022 S. 83] 2022 S. 83] 2022 S. 84] 2021 S. 85] 2018 S. 86] 2022 S. 87] 2020 S. 88] 2022 S. 89] 2018 S. 90] 2022 S. 91] 2021 S. 92] 2022 S. 93] 2022 S.	$SA_{41}$	•	0		•	•	•	-	-	-	-	AE	1	✓
83]     2022     Sz       84]     2021     Sz       85]     2018     Sz       86]     2022     Sz       87]     2020     Sz       88]     2022     Sz       89]     2018     Sz       90]     2022     Sz       91]     2021     Sz       92]     2022     Sz       93]     2022     Sz       94]     2022     Sz	$SA_{41}$	•	•	•	0	0	0	$PA_1$	•	•	0	AE	/	±
84      2021     S.       85      2018     S.       86      2022     S.       87      2020     S.       88      2022     S.       89      2018     S.       90      2022     S.       91      2021     S.       92      2022     S.       93      2022     S.       94      2022     S.	$SA_{41}$	•	0		•	•	•	-	-	-	_	AE	<b>/</b>	±
85]     2018     Sz       86]     2022     Sz       87]     2020     Sz       88]     2022     Sz       89]     2018     Sz       90]     2022     Sz       91]     2021     Sz       92]     2022     Sz       93]     2022     Sz       94]     2022     Sz	$SA_{41}$	•	0	•	•	•	•	-	-	-	-	AE	1	±
86     2022     Sz       87     2020     Sz       88     2022     Sz       89     2018     Sz       90     2022     Sz       91     2021     Sz       92     2022     Sz       93     2022     Sz       94     2022     Sz	$SA_{41}$	•	•	•	•	•	•	-	-	-	-	AE	<b>/</b>	±
87]     2020     \$2.       88]     2022     \$2.       88]     2018     \$3.       89]     2018     \$3.       90]     2022     \$3.       91]     2021     \$3.       92]     2022     \$3.       93]     2022     \$3.       94]     2022     \$3.	$SA_{41}$	•	0	0	•	0	•	-	-	-	-	AE	<b>/</b>	Х
88      2022     Sz       89      2018     Sz       90      2022     Sz       91      2021     Sz       92      2022     Sz       93      2022     Sz       94      2022     Sz	$SA_{41}$	•	0		•	•	•	-	-	-	-	AE	1	Х
89]     2018     Sz       90]     2022     Sz       91]     2021     Sz       92]     2022     Sz       93]     2022     Sz       94]     2022     Sz	$SA_{41}$	•	0	0	•	0	•	$PA_3$	•	•	0	AE	1	Х
90]     2022     \$2       91]     2021     \$2       92]     2022     \$2       93]     2022     \$2       94]     2022     \$2	$SA_{41}$	•	•	0	•	•	•	$PA_6$		•	0	AE	<b>/</b>	7
91] 2021 SA 92] 2022 SA 93] 2022 SA 94] 2022 SA	$SA_{41}$	•	•	0	•	-	-	$PA_6$	0	•	0	AE	<b>/</b>	)
92] 2022 SA 93] 2022 SA 94] 2022 SA	$SA_{41}$	•	•	•	•	•	•	-	-	-	-	AE	<b>/</b>	✓
93] 2022 SA 94] 2022 SA	$SA_{41}$	•	0	0	•	0	•	_ 	_	_	_	AE	1	<b>✓</b>
94] 2022 SA	$SA_{41}$	•	0	0	0	0	•	$PA_3$	0	0	0	AE	<b>√</b>	±
	$SA_{41}$	•	0	0	0	0	0	$PA_4$	0	•	•	AE	<b>/</b>	V
1321 2023 - \$	3 A <sub>41</sub>	•	•	0	•	•		_	-	_	_	AE	<b>√</b>	V
	$SA_{11}, SA_{22}$		•	•		•	•		_			AT	/	/
1331 2023 S	$SA_{11}$ , $SA_{22}$ , $SA_{33}$		•	0		0		_	_	_	_	DT	/	
1341 2023 S	$SA_{11}, SA_{22}, SA_{33}$ $SA_{11}, SA_{22}, SA_{33}$		0	0	0	0		_	_	_	_	AT		,
	$SA_{11}, SA_{22}, SA_{33}$ $SA_{11}, SA_{22}$		•	0	•	•		_	_	_	_	AT	/	/
136] 2022 SA	$SA_{22}, SA_{32}$		0	0	0	0	0	_	_	_	_	DT		±
[137] 2023 SA			_	0	•	•	ĕ	-	_	_	_	AT	/	/
		_												
tric addressed		57(100%)	57(100%)	57(100%	(b) 55(96%)	51(89%)	50(86%)	2	23(40%)	24(42%)	23(40%	)		

The classification criteria are in Section 4.1.1. The table identifies the interoperability mode used by each study (IMode), indicating whether it supports Asset Transfers (AT), Data Transfers (DT), or Asset Exchanges (AE). Additionally, it notes if the solution is independent of privacy primitives in the underlying chains (**PC**) and if an implementation is available (**Impl**). Papers marked as (–) do not focus on the specific property. The last two rows of the table summarize the classification. We present a

implementation is available (Impl). Papers marked as (-) do not focus on the specific property. The last two rows of the table summarize the classification. We present a visual representation of 1) the number of studies addressing each metric and 2) the number of studies classified using  $\bullet$  or  $\checkmark$ .

Security approaches:  $SA_{11}$  Centralization;  $SA_{12}$  Trusted Computation;  $SA_{21}$  Permissionless Network;  $SA_{22}$  Permissioned Network;  $SA_{31}$  Inclusion Proofs;  $SA_{32}$  Validity Proofs;  $SA_{33}$  Fraud Proofs;  $SA_{41}$  Secret- & Time-based Locks.

Privacy Approaches:  $PA_1$  ZKP;  $PA_2$  TEE;  $PA_3$  Adaptor Signatures;  $PA_4$  Blind Signatures;  $PA_5$  Ring Signatures;  $PA_6$  Homomorphic Encryption.

Guarantees some privacy properties even if no privacy approach is employed, due to the use of private chains and secure communication channels (e.g., TLS).

Guarantees privacy at the application layer, not the cross-chain level. Protects the order matching protocol to guarantee fairness but transactions are published normally in

- blockchains.
- Has several open-source implementations in different technological stacks, enhancing decentralization.

  With considerable liquidity in the TEE [126] we can classify it as ●.
- <sup>5</sup> One of the few solutions being standardized in reputable standardization bodies [44].
- <sup>6</sup> Strong dependency on price oracle. It can be classified as if the oracle is robust and decentralized [7].
- <sup>7</sup> Provided it has a sufficiently large anonymity set.

#### Security Properties.

- *Integrity (In)* Integrity is enforced by the underlying cryptographic primitives which are based on the hardness of well-known problems (e.g., computing the discrete logarithm) (●); integrity is enforced under strong assumptions (e.g., trusted hardware, rational participants, parties abiding by laws) (●); integrity cannot be guaranteed under misbehaving parties (○).
- Availability (Av) Availability requires a decentralized network, but there is at least one honest off-chain party (●); availability can be temporarily compromised if any party misbehaves (●); it is based on a centralized architecture, hence there are serious concerns about availability (○).
- Accountability (Ac) The misbehaving party is identifiable and automatically punished (e.g., programmatically) (●); malicious party is identifiable, but there is no punishment or needs to be enforced by a third party (●); misbehaving parties are neither identifiable nor punished (○) (the notion of accountable safety [76]).

#### Privacy Properties.

- Unlinkability (Un) It is cryptographically infeasible to link transactions or addresses (♠); it is possible to link transactions or addresses through heuristics (♠); no mechanism is in place to unlink transactions or addresses across domains (♠).
- Anonymity (An) Both users' and operators' identities are concealed (♠); users' anonymity or operator's anonymity is enhanced (e.g., through set anonymity approaches) (♠); at most pseudo-anonymity is provided for users, and operators are known (♠).
- Confidentiality (Cf) Data confidentiality is enforced through cryptographic primitives (●); conditional confidentiality i.e., can be revoked under some circumstances, or verified by auditors. Note that IMs based on private chains partially guarantee this property (●); there is no confidentiality (○).

Governance and Performance Properties. We extend our classification of solutions with governance [138] and performance properties, as they are factors that strongly influence security and privacy [8]. We collect insights from related literature to define:

- Decentralization (Dc) Fully distributed system with a consensus algorithm to settle different views on information [43] or control relies on the end-user (♠); limited decentralization of the system, being run by a small set of verifying parties (♠); control of the system resides in less than 4<sup>§</sup> parties (can be distributed or centralized) (♠).
- Latency (Lat) Latency of a cross-chain transfer is settled before finalization time (optimistic approach) (●);
   Latency of a cross-chain transfer is finalized right after the finalization time of the slowest chain (●); Latency

§, some sources suggest that 4 is a reasonable number of non-colluding parties to secure a blockchain bridge [32].

- of a cross-chain transfer is more than the finalization time of the slowest chain, due to extra processes ran before (e.g., special account setup) or thereafter (e.g., extra transactions needed) (O).
- Cost (Co) there are no protocol fees (for the user); can be run with low-tier commercially available hardware (for IM operator) (●); variable fees depending on search and demand with an upper bound lesser or equal than 1% of the bridged value (for the user); requires at most mid-tier hardware (for IM operator) (●); variable fees depending on search and demand with more than 1% of the bridged value (for the user); requires above mid-tier and/or specialized hardware for the IM operator (○).

**Miscellaneous** (**Misc.**). We provide information that complements our assessment. **IMode** indicates the main interoperability mode supported by the IM. **PC** indicates if the IM requires ( $\nearrow$ ) a privacy-enhanced chain or permissioned blockchain to operate optimally, or if it is independent of those primitives ( $\checkmark$ ). **Impl.** refers to the project having an open-source implementation and evaluation ( $\checkmark$ ), a not-open source implementation ( $\pm$ ), or no implementation ( $\nearrow$ ).

- **4.1.2. Insights.** We now present a list of insights taken from the analysis of the literature.
  - Insight 1: A conspicuous deficit exists in the literature regarding the empirical assessment of protocol performance and associated costs. This observation is consistent with findings from other studies [8], [139]. While many solutions appear to delegate computationally intensive tasks to off-chain procedures [40]–[42], further investigation in this domain remains paramount.
  - Insight 2: All studies ensure a degree of integrity, predominantly upheld by cryptographic mechanisms. A thorough examination of these mechanisms is crucial. Research that confines its scope to specific adversarial behaviors (e.g., only rational actors) exhibits diminished integrity.
  - Insight 3: Security takes precedence over privacy. Only 17 studies (29.8%) address cross-chain privacy. At the same time, projects that dominate 75% of the market neglect cross-chain privacy, suggesting a prevailing apprehension regarding bridge security, relegating privacy to a subordinate design goal.
  - Insight 4: ZKP  $(\mathcal{PA}_1)$  emerges as the prevailing approach to guarantee privacy in IMs (53% of IMs with a privacy approach), which allows shifting trust from third parties to cryptographic protocols.
  - Insight 5: The prevailing academic literature primarily emphasizes asset exchanges (44%) and transfers (29%) between blockchains, with fewer studies addressing general data transfers (26%). This trend contrasts with recent industry developments, with a surge in bridges that facilitate data transfers (arbitrary message-passing bridges, such as [104], [140], [141]).
  - Insight 6: For cross-chain anonymity, it is imperative that two distinct transactions remain indistinguishable

when originating from the same sender and targeting the same recipient. This condition can only be achieved with the generation of unique addresses for each *cctx*, exemplified by the mechanism of stealth addresses [142].

- Insight 7: The predominant trend in the literature underscores the achievement of accountability using stake-slashing mechanisms. In contrast, a smaller subset of research advocates the legal identification of the parties involved through an ancillary identity service. Such methodologies are found predominantly in centralized frameworks within permissioned network configurations where nodes possess identifiable attributes.
- Insight 8: Our analysis elucidates that cross-chain unlinkability in AT and DT-based protocols can only be guaranteed when the underlying chains ensure confidentiality. We present a rigorous formalization in Appendix C. For AE protocols, unlinkability can be achieved through cryptographic primitives such as Adaptor Signatures. Additionally, we emphasize the need to investigate heuristics to break privacy [88], [143] and the respective mitigations.
- Insight 9: Privacy is implemented with a small set of techniques that bring considerable overhead to IMs, latency- and trust-wise. Both ZKP  $(\mathcal{PA}_1)$  and TEE  $(\mathcal{PA}_2)$  require trusted setups with highly expensive computation and specific hardware for achieving confidentiality and unlinkability [21], [126].
- **Insight 10:** Asset exchange protocols, with confidential order matching algorithms [54], [62] do not offer crosschain privacy. They only guarantee fairness in the order-matching process at the application layer.
- Insight 11: Research has shown that the privacy level provided by privacy-centric platforms [144] or applications [145] can be compromised due to risky user behaviors [126], [146]–[149]. Cross-chain systems based on similar primitives face comparable challenges.

### 4.2. Vulnerabilities, Leaks, and Mitigations

In this section, we present vulnerabilities and leaks found in cross-chain interoperability. Table 4 lists and maps each identified vulnerability/leak to corresponding mitigations. Table 5 presents all mitigations identified in the literature and proposed in this paper.

**4.2.1.** Vulnerabilities and Leaks. We have identified 45 vulnerabilities across different layers: 3 in the network layer, 22 in the protocol layer, 17 in the implementation layer, and 3 in the operational layer. Many studies focus on common vulnerabilities (e.g.,  $\mathcal{V}_{13}$ ,  $\mathcal{V}_{18}$  or  $\mathcal{V}_{19}$ ), while specific bug reports highlight more specific issues (e.g.,  $\mathcal{V}_{14}$  or  $\mathcal{V}_{40}$ ). Surprisingly, we found fewer vulnerabilities in the operational layer, which plays a significant role in cross-chain hacks (cf. Section 4.2.2). Additionally, this might suggest that academia is not addressing industry-relevant issues adequately because the same vulnerabilities are continuously occurring. We have also discovered four theoretical

TABLE 4. SECURITY VULNERABILITIES AND PRIVACY LEAKS IN CROSS-CHAIN SYSTEMS. THE COLORED CIRCLE DENOTES THE LAYER WHERE IT CAN BE FOUND (CF. SECTION 3.1).

Vulnerability/Leak	Mitigations
• $\mathcal{V}_1$ Honest mining assumption [45]	$\mathcal{M}_1$ - $\mathcal{M}_5$
$\circ$ $\mathcal{V}_2$ Absence of identity verification [45], [71], [72]	$\mathcal{M}_8$ - $\mathcal{M}_{11}$
$\circ \mathcal{V}_3$ Network isolation [38], [45], [62], [77]	$\mathcal{M}_6, \mathcal{M}_7$
$\circ$ $\mathcal{V}_4$ Outdated light client state [45], [53], [150]	$\mathcal{M}_{16}$
• $V_5$ Wrong main chain identification [6], [45], [77]	$\mathcal{M}_{18}$
$\circ V_6$ Incorrect event verification [151]–[154]	$\mathcal{M}_{12}$ - $\mathcal{M}_{14}$
$\circ$ $\mathcal{V}_7$ Acceptance of invalid consensus proofs [155]	$\mathcal{M}_{15}$
$v_8$ Absence of chain identification [156]	$\mathcal{M}_4$
$\mathcal{V}_9$ Submission of repeated inclusion proofs [21], [45], [77], [157]	.,
$\mathcal{V}_{10}$ Counterfeiting assets [45], [77], [158]	$M_{19}$ - $M_{23}$
$\mathcal{V}_{11}$ Involuntary timelock expiry [63], [85]	$M_{29}$ - $M_{30}$
$\mathcal{V}_{12}$ Unset withdrawal limits [156], [159]	$\mathcal{M}_{69}$
V <sub>13</sub> Action withhold [58], [61], [80], [86], [86], [94], [160]	$\mathcal{M}_8, \mathcal{M}_{27}, \mathcal{M}_{28}$
V <sub>14</sub> Unspecified gas limit [161]	$\mathcal{M}_{65}$
V <sub>15</sub> Resource exhaustion [45], [55], [57], [60], [65], [69]	$\mathcal{M}_{48}$ - $\mathcal{M}_{50}$
<ul> <li>\$\mathcal{V}_{16}\$ Single point of failure [156], [162]</li> <li>\$\mathcal{V}_{17}\$ Publicly identifiable operators [74]</li> </ul>	$\mathcal{M}_{7}, \mathcal{M}_{32}, \mathcal{M}_{47}$
$v_{17}$ 1 dishery identifiable operators [74] $v_{18}$ Misaligned incentive mechanisms [38], [60], [65], [122]	$\mathcal{M}_{44}$ - $\mathcal{M}_{46}$ $\mathcal{M}_{23}$ , $\mathcal{M}_{31}$ - $\mathcal{M}_{34}$
$v_{18}$ This argued incentive incentarisms [56], [66], [67] $v_{19}$ Token price volatility [45], [74], [77], [80], [82], [83]	$\mathcal{M}_{35}$ - $\mathcal{M}_{39}$
• $V_{20}$ Centralized power [65], [162], [163]	$\mathcal{M}_{32}, \mathcal{M}_{43}$
$v_{20}$ Verifier's dilemma [163]	$\mathcal{M}_{24}$ - $\mathcal{M}_{26}$
$\mathcal{V}_{22}$ Manipulation of exchange rates [29], [164]–[167]	$\mathcal{M}_{40}, \mathcal{M}_{41}$
$\circ \mathcal{V}_{23}$ Unfair transaction/event ordering [65]	$\mathcal{M}_{41}, \mathcal{M}_{42}$
$\circ \mathcal{V}_{24}^{23}$ Insecure access control [168]–[173]	$\mathcal{M}_{51}^{41}, \mathcal{M}_{52}^{42}$
$\circ$ $\mathcal{V}_{25}$ Conceed approvals to third parties [152], [174], [175]	$\mathcal{M}_{53}$
<ul> <li>V<sub>26</sub> Outdated third-party library version [176]</li> </ul>	$\mathcal{M}_{78}$
$\circ$ $\mathcal{V}_{27}$ Unsafe third party modules [151], [156], [162], [177]	$\mathcal{M}_{58}, \mathcal{M}_{78}$
$\circ$ $\mathcal{V}_{28}$ Dead code [151], [159], [176]–[180]	$\mathcal{M}_{59}$
$\circ \mathcal{V}_{29}$ Usage of non-standard naming [176], [177]	$\mathcal{M}_{79}$
$\mathcal{V}_{30}$ Inconsistent smart contract engine version [156], [162], [179]	$\mathcal{M}_{80}$
$\mathcal{V}_{31}$ Unconventional code/testing architecture [176], [179]	$\mathcal{M}_{81}$
$\mathcal{V}_{32}$ Reentrancy [156]	$\mathcal{M}_{82}$
$varphi_{33}$ Failure to emit events upon state changes [151], [162], [178]	$\mathcal{M}_{83}$
$v_{34}$ Inconsistent bridge contract interfaces [180]	$\mathcal{M}_{84}$
$v_{35}$ Out of order transaction execution [151]	$\mathcal{M}_{85}$
$\mathcal{V}_{36}$ Absence of storage gaps in smart contracts [181] $\mathcal{V}_{37}$ Integer overflow and underflow [151], [159], [162], [176]	$\mathcal{M}_{86}$
$v_{37}$ integer overflow and underflow [131], [139], [102], [170] $v_{38}$ Absence of sanity checks [156], [177]	$\mathcal{M}_{87}$ $\mathcal{M}_{87}$
$v_{38}$ Abstince of samity checks [150], [177] $v_{39}$ Code and documentation mismatched [162], [176]–[179]	$\mathcal{M}_{88}$
$\mathcal{V}_{40}$ Uninitialized variables [182]	$\mathcal{M}_{66}$
$\mathcal{V}_{41}$ Compromise of ZK algorithms' private inputs [126]	$\mathcal{M}_{67}$
$\mathcal{V}_{42}$ Other smart contract vulnerabilities [151], [162], [179]	$\mathcal{M}_{51}^{67}, \mathcal{M}_{54}^{-1}, \mathcal{M}_{56}^{-1}$
• $\mathcal{V}_{43}$ Inadequate key management [152], [183]	$\mathcal{M}_{47}, \mathcal{M}_{60}$ - $\mathcal{M}_{62}$
• $\mathcal{V}_{44}$ Physical infrastructure backdoors [50]	$\mathcal{M}_{46}, \mathcal{M}_{63}$ - $\mathcal{M}_{64}$
• $\mathcal{V}_{45}^{45}$ Social engineering-related vulnerabilities [174], [184]	$\mathcal{M}_{77}^{40}$ 03 04
- $\mathcal{L}_1$ Leakage of private data in ZK ceremony input [40]	$\mathcal{M}_{89}$
- $\mathcal{L}_2$ Linking transactions through transactional data [88], [89]	$\mathcal{M}_{90}$
$-\mathcal{L}_3$ Common secret deployment [87]	$\mathcal{M}_{91}$
- L <sub>4</sub> User-generated privacy leaks [126], [146]-[149]	$\mathcal{M}_{92}$
- L <sub>5</sub> Mapping on-chain addresses to real-world identities [126]	$M_{90}, M_{92}$
25 mapping on-chain addresses to real-world identities [120]	J7190, J719 <u>2</u>

privacy leaks. From our analysis, no privacy leak has been reported in cross-chain systems. Therefore, we could not cross-reference this information with past incidents.

**4.2.2. Analysis of Real-World Hacks.** Attacks against cross-chain bridges have proliferated in the last couple of years. Table 6 presents a classification of 14 of the most impactful attacks in the industry since July 2021, that account for more than 94% of the total value stolen from cross-chain bridges. We present general attack information, incident response-related data, the components targeted by attackers, and the vulnerabilities behind each. Appendix E presents further information and mitigations for each.

Security Approach (SA) The security approach used by the bridge.

TABLE 5. LIST OF MITIGATIONS COLLECTED IN THE LITERATURE AND PROPOSED BY OUR ANALYSIS

153   Wait full confirmation time according to the source chain consensus mechanism   M <sub>49</sub>   [36]   [45]   Wait full confirmation time according to the source chain consensus mechanism   M <sub>45</sub>   [45]   Usage of block maturity periods   M <sub>45</sub>   [45]   Usage of block chain wews   M <sub>45</sub>   [45]   Wait full confirmation mechanisms   M <sub>45</sub>   [45]   Wait full carried and the multiple destination chains   M <sub>45</sub>   [45]   Wait full carried and the multiple destination chains   M <sub>45</sub>   [45]   Wait full carried and the multiple destination chains   M <sub>45</sub>   [45]   Washed the creation of identities expensive (e.g., a ligh stake per identity)   M <sub>49</sub>   [45]   Washed the creation of identities expensive (e.g., a ligh stake per identity)   M <sub>49</sub>   [46]   Washed the creation of identities expensive (e.g., a ligh stake per identity)   M <sub>49</sub>   [47]   Reward creating fewer dentities with more stake   W <sub>46</sub>   [48]   Washed the creation of identities expensive (e.g., a ligh stake per identity)   M <sub>49</sub>   [49]   W <sub>40</sub>   W <sub>40</sub>   [49]   W <sub>40</sub>   W <sub></sub>		Wait full confirmation time according to Insertion of block maturity periods Usage of blockchain views		
(45) Usage of blocke mannity periods  (45) Usage of blocked mannity periods  (46) Usage of blocked mannity periods  (47) Usage of blocked mannity news  (48) Usage of blocked mannity news  (49) Synthenity as mart contract state or multiple destination chains  (49) Synthenity as mart contract state or multiple destination chains  (40) Physical decentralization of infrastructure  (41) Usage of current centralized and mily or state  (42) Usage of a trusted centralized and mily state  (43) Enable verifiability of state updates in light clients for different consensus mechanism  (44) Deply untime monitoring modules  (45) Enable verifiability of state updates in light clients for different consensus mechanism  (46) Enable verifiability of state updates in light clients for different consensus mechanism  (47) Unique none-old generation per request  (48) Use collateralization / Over-Collateralization mechanisms  (49) Use and develop new main chain identification mechanisms  (47) Unique none-old generation per request  (48) Use collateralization / Over-Collateralization dehriques  (49) Use of Olateralization of Over-Collateralization of collateral  (41) Embedded rules in third party network consensus mechanism  (44) Use Collateralization of Over-Collateralization of Deverollateralization of partial of an analyze uses and operators  (48) Usage of Perintal incentivized watchers that antest actions/events  (48) Usage of Veritable Timed Commitments  (49) Usage of Veritable Timed Commitments  (49) Usage of Veritable Timed Commitments  (40) Usage of Veritable Timed Commitments  (41) Embedded rules in third parties and estate mining power among them  (42) Usage of Veritable Timed Commitments  (43) Usage of Veritable Timed Commitments  (44) Usage of Orentable Usage of Well of the updated exchange rates  (45) Usage of Orentable Usage of Peritable Timed Commitments  (46) Usage of Orentable Usage of Peritable Timed Commitments  (47) Usage of Orentable Usage of Mellon was to observe price fluctuations  (48) Peritable Signate Section				
Usage of hockednain views    156  Add chain identification mechanisms   M <sub>5</sub>     157  Add chain identification mechanisms   M <sub>5</sub>     158  Increase transaction settlement time   M <sub>5</sub>     159  Increase transaction settlement time   M <sub>5</sub>     150  Physical decentralization of infrastructure   M <sub>5</sub>     150  Physical decentralization of infrastructure   M <sub>5</sub>     151  Diagration with Self Sovereign Identity (SSI) mechanisms   M <sub>5</sub>     152  Listen to events only from whitelisted smart contracts   M <sub>6</sub>     152  Listen to events only from whitelisted smart contracts   M <sub>6</sub>     153  Deploy tundrine monitoring modules   M <sub>6</sub>     155  Listen to events only from whitelisted smart contracts   M <sub>6</sub>     150  Insertion of a data availability of state updates in light clients for different consensus mechanism   M <sub>6</sub>     150  Insertion of a data availability also are the read of the state updates in light clients for different consensus mechanism   M <sub>6</sub>     150  Insertion of a data availability of state updates in light clients for different consensus mechanism   M <sub>6</sub>     151  Usage of perturbation of Collateral lization the chains of collateral lization of the protocol of Sage of Permittings and Collateral lization of Collateral lization of Collateral lization of the line in time party network consensus mechanism   M <sub>7</sub>     163  Parallelising transactions are developed techniques   M <sub>7</sub>     164  Oline Sage of Verifiable Timed Commitments   Sage of Permitting the Commitments   Sage of Permitter decentralization of parties and scatter mining power among them   M <sub>7</sub>     164  Oline support for periods of asynchrony in the execution of the protocol decentralization				
156  Add chain identification mechanisms   M3   175  Add chain identification mechanisms   M4   180  Synchronize strate contract state on multiple destination chains   M4   175  Magnetonize strate contract state on multiple destination chains   M4   190				, ,
[45]         Synchronize smart contract state on multiple destination chains         Mms         175           [38]         Increase transaction settlement time         Mms         190           [36]         Physical decentralization of infrastructure         Mms         190           [37]         Physical decentralization of infrastructure         Mms         190           [38]         Increase transactions expensive (e.g., a high stake per identity)         Mms         1-           [38]         Increase transactions of infrastructure         Mms         1-           [37]         Reward creating fewer identities with more stake         Mms         1-           [152]         Listen to events only from whitelisted smart contracts         Mms         160         160           [154]         Deploy runtime monitoring strategies at the same time         Mms         160         160         160         160           [56]         Emble verifiability of state updates in light clients for different consensus mechanism         Mms         160         160         160         160         161         161         161         161         161         161         161         162         161         161         162         162         161         161         162         162         161 <t< td=""><td></td><td>•</td><td></td><td></td></t<>		•		
185   Increase transaction settlement time   187   Increase transaction settlement time   187   Increase transaction of infrastructure   187   May   187   May   187   May   187   May   187   May   188   Increase transaction of infrastructure   187   Increase transition of infrastructure   187   May   May   187   May   Ma		,		
Flysical decentralization of infrastructure   Aff.   Usage of a trusted centralized authority to mediate cetxx     IsS   Integration with Self Sovereign Hearity (SSI) mechanisms   Aff.     IsB   Usage of a trusted centralized authority to mediate cetxx     IsB   Integration with Self Sovereign Hearity (SSI) mechanisms   Aff.     IsB   Make the creation of identities expensive (e.g., a high stake per identity)   Aff.     IsB   Deploy nuttime monitoring strategies at the same time   Aff.     IsB   Deploy nuttime monitoring strategies at the same time   Aff.     IsB   Deploy nuttime monitoring strategies at the same time   Aff.     IsB   Deploy nuttime monitoring strategies at the same time   Aff.     IsB   Deploy nuttime monitoring strategies at the same time   Aff.     IsB   Insertion of a data awailability layer   Aff.     IsB   Insertion of a data awailability layer   Aff.     IsB   Insertion of a data waralability layer   Aff.     IsB   Usage of external incentivized wardress that attest actions/events   Aff.     IsB   Usage of external incentivized wardress that actions/events   Aff.     IsB   Usage of Premiums   Aff.     IsB   Parallelizing transaction verification   Aff.     IsB   Parallelizing transaction verification   Aff.     IsB   Saperate computational-beavy transactions into multiple blocks   Aff.     IsB   Usage of Verifiable Timed Commitments   Aff.     IsB   Usage of Premiums   Aff.     IsB   Increase of Premiums   Aff.     IsB   Increase of MEV to form-run misbehaving transactions   Aff.     IsB   Increase the number of parties and scatter mining power among them   Aff.     IsB   Reduce time window for users to observe price fluctuations   Aff.     Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.     Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.   Aff.				
Usage of a trusted centralized authority to mediate certxs   Ms.				
185   Integration with Self Sovereign Identity (SSI) mechanisms   M <sub>S</sub>   162     May be the creation of identities expensive (e.g., a high stake per identity)   M <sub>S</sub>   163     M <sub>S</sub>   163   Integration with castake the castake the creation of identities with more stake   M <sub>S</sub>   163     152   Listen to events only from whitelisted smart contracts   M <sub>S</sub>   163     M <sub>S</sub>   163   Enables verifiability of state updates in light clients for different consensus mechanism   M <sub>S</sub>   163     150   Insertion of a data availability layer   M <sub>S</sub>   164     151   Unique noncerid generation per request   M <sub>S</sub>   165     152   Use and develop new main chain identification mechanisms   M <sub>S</sub>   164     153   Use and develop new main chain identification nechanisms   M <sub>S</sub>   164     154   Use and develop new main chain identification nechanisms   M <sub>S</sub>   164     152   Use Collateralization / Over-Collateralization chain sections/events   M <sub>S</sub>   164     154   Use Collateralization / Over-Collateralization chain sections/events   M <sub>S</sub>   165     155   Usage of Distributed Scheme between untrusted users and operators   M <sub>S</sub>   177     186   Separate entities at the create and verify blocks   M <sub>S</sub>   178     187   Provide support for periods of asynchrony in the execution of the protocol   M <sub>S</sub>   177     188   Usage of Perniums   M <sub>S</sub>   165     189   Usage of Perniums   M <sub>S</sub>   165     181   Usage of Verifiable Timed Commitments triggered upon failures   M <sub>S</sub>   165     183   Increase the number of parties and scatter mining power among them   M <sub>S</sub>   177     188   General mingations for well-off triggered to periodical extension of the updated exchange rates   M <sub>S</sub>   164     183   Increase the number of parties and scatter mining power among them   M <sub>S</sub>   165     184   Medic and analyze are behavior through game-theory principles   M <sub>S</sub>   165     185   Parallel asset tocking   165   165   165     186   Medic and analyze are behavior through game-theory principles   M <sub>S</sub>   165     186   187   Protocol architecture decentralization of parties and		-	$\mathcal{M}_{56}$ –	Ensure there are rigorous testing guidelines being enforced
Make the creation of identities expensive (e.g., a high stake per identity)   M <sub>s</sub>   162     Make the creation of identities with more stake   M <sub>s</sub>   191     Listen to events only from whitelisted smart contacts   M <sub>s</sub>   191     M <sub>s</sub>   192     Listen to events only from whitelisted smart contacts   M <sub>s</sub>   191     M <sub>s</sub>   192     M <sub>s</sub>   193     M <sub>s</sub>   193     M <sub>s</sub>   194     M <sub>s</sub>   195     Insertion of a data availability layer   M <sub>s</sub>   194     M <sub>s</sub>   195     Insertion of a data availability layer   M <sub>s</sub>   195     M <sub>s</sub>   195     Insertion of a data availability layer   M <sub>s</sub>   195     M <sub>s</sub>   195     Insertion of a data availability layer   M <sub>s</sub>   195     M <sub>s</sub>			$\mathcal{M}_{S7}$ –	Just like on-chain smart contracts, off-chain programs and infrastructure must be audited
Reward creating fewer identities with more stake   151   Listen to events out Prom whitelisted smart contracts   152   Listen to events only from whitelisted smart contracts   153   Listen to events only from whitelisted smart contracts   154   Deploy runtime monitoring modules   155   Employ multiple different monitoring strategies at the same time   156   Employ multiple different monitoring strategies at the same time   150   Mes   150				•
152   Listen to events only from whitelisted smart contracts   153     154   Deploy runtime monitoring modules   154     155   Employ runtime monitoring modules   154     156   Employ runtime monitoring modules   155     157   Insert no or data waitability layer   156     158   Insert no or data availability layer   157     159   Insert not or data waitability layer   157     150   Insert not or data waitability layer   157     151   Insert note of a data waitability layer   157     152   Insert note of a data waitability layer   157     153   Insert note of a data waitability layer   157     154   Insert note of a data waitability layer   157     155   Insert note of a data waitability layer   157     156   Insert note of a data waitability layer   157     157   Insert note of a data waitability layer   157     158   Insert note of a data waitability layer   157     159   Insert note of a data waitability layer   157     150   Insert note of Distributed Signature Schemes between untrusted users and operators   157     158   Insert independent computational-heavy transactions into multiple blocks   157     159   Insert independent computational-heavy transactions into multiple blocks   157     158   Insert proper of Premiums   158   158     159   Insert independent computational-heavy transactions into multiple blocks   157     158   Usage of Verifiable Timed Commitments   158   158     151   Model and analyze user behaviour through game-theory principles   158     151   Protocol architecture decentralization   158   159     151   Protocol architecture decentralization   158   159     152   Model and analyze user behaviour through game-theory principles   158     151   Protocol architecture decentralization   158   158     152   Model and analyze user behaviour through game-theory principles   158   158     151   Protocol architecture decentralization   158   158   158     152   Model and analyze user behaviour for slippage   158   158   158   158   158   158   158   158   158   158   158   158   158   158			$\mathcal{M}_{59}$ –	Linting tools to raise warnings for unused code
154   Deploy runtime monitoring modules   154   Employ multiple different monitoring strategies at the same time   M <sub>16</sub>   150   160   1				
<ul> <li>[65] Employ multiple different monitoring strategies at the same time M<sub>16</sub> [50]</li> <li>[76] Insertion of a data availability of state updates in light clients for different consensus mechanism M<sub>16</sub> [50]</li> <li>[77] Unique nonce/id generation per request M<sub>16</sub> [161]</li> <li>[48] Use and develop new main chain identification mechanisms M<sub>16</sub> [161]</li> <li>[47] Trigger automatic liquidations of collateral manual manua</li></ul>		_		
Total Finable verifiability of state updates in light clients for different consensus mechanism   M <sub>c6</sub>   150]     Total Finable verifiability layer   M <sub>c6</sub>   161     Total Final Final Education per request   M <sub>c6</sub>   162     Total Final Education per request   M <sub>c6</sub>   163     Total Final Education of collateral education per request   M <sub>c6</sub>   163     Total Final Education of collateral education   M <sub>c6</sub>   163     Total Education   Over-Collateralization rechanism   M <sub>c6</sub>   163     Total Education   Over-Collateralization   Over-Collateralization   M <sub>c6</sub>   164     Total Education   Over-Collateralization   Over-Collateralization   Over-Collateralization   M <sub>c6</sub>   164     Total Education   Over-Collateralization   M <sub>c6</sub>   164     Embedded rules in third party network consensus mechanism   M <sub>c7</sub>   175     Total Embedded rules in third party network consensus mechanism   M <sub>c7</sub>   175     Total Embedded rules   M <sub>c7</sub>   175     Embedded rules in third party network consensus mechanism   M <sub>c7</sub>   175     Embedded rules in third party network consensus mechanism   M <sub>c7</sub>   175     Embedded rules in third party network consensus mechanism   M <sub>c7</sub>   175     M <sub>c7</sub>   M <sub>c7</sub>   M <sub>c7</sub>   M <sub>c7</sub>   M <sub>c7</sub>   175     Separate entities that create and verify blocks   M <sub>c7</sub>   176     Separate entities that create and verify blocks   M <sub>c7</sub>   176     Separate entities that create and verify blocks   M <sub>c7</sub>   176     Separate entities that create and verify blocks   M <sub>c7</sub>   176     Separate entities that create and verify blocks   M <sub>c7</sub>   176     Separate entities that remount recked scatter mining power among them   M <sub>c7</sub>   177     M <sub>c7</sub>   M <sub>c7</sub>   M <sub>c7</sub>   M <sub>c7</sub>   176     Separate entities that observe price fluctuations   M <sub>c8</sub>   151     Separate lasset locking   Sea   Separate mining tone of data   M <sub>c7</sub>   190     Separate the number of parties and scatter mining power among them   M <sub>c8</sub>   176     Separate the number of parties and scatter mining power among them   M <sub>c8</sub>   176     Separate the number of parties and scatter mining power am	_			
150   Insertion of a data availability layer   171   Unique none-fide generation per request   172   Unique none-fide generation per request   173   Unique none-fide generation per request   174   Unique none-fide generation per request   175   Unique none-fide generation per request   175   Use and develop new main chain identification techniques   175   176   176   176   177   178				
<ul> <li>[77] Unique noncefd generation per request</li> <li>[77] Unique noncefd generation per request</li> <li>[45] Use and develop new main chain identification mechanisms</li> <li>[47] Trigger automatic inquidations of collateral</li> <li>[47] Trigger automatic inquidations of collateral</li> <li>[48] Usage of external incentivized watchers that attest actions/events</li> <li>[47] Usage of Extract incentivized watchers that attest actions/events</li> <li>[48] Dusage of Distributed Signature Schemes between untrusted users and operators</li> <li>[49] Dusage of Distributed Signature Schemes between untrusted users and operators</li> <li>[40] Daract independent computational-heavy transactions into multiple blocks</li> <li>[40] Daract independent computational-heavy transactions into multiple blocks</li> <li>[40] Usage of Premiums</li> <li>[41] Separate entities that create and verify blocks</li> <li>[42] Usage of Verifiable Timed Commitments</li> <li>[43] Provide support for periods of asynchrony in the execution of the protocol</li> <li>[44] Usage of Verifiable Timed Commitments</li> <li>[45] Model and analyze user behavior through game-theory principles</li> <li>[46] Model and analyze user behavior through game-theory principles</li> <li>[47] Model and analyze user behavior through game-theory principles</li> <li>[48] Datacle architecture decentralization</li> <li>[49] Usage of MEV to front-run misbehaving transactions</li> <li>[40] Datacle architecture decentralization</li> <li>[41] Adjust the amount locked account for slippage</li> <li>[42] Adjust the amount locked account for slippage</li> <li>[43] Trigger automatic liquidations to avoid getting uncollateralized</li> <li>[44] Adjust the amount locked according to the updated exchange rates</li> <li>[45] Adjust the amount locked according to the updated exchange rates</li> <li>[46] Adjust the amount locked according to the updated exchange rates</li> <li>[47] Marge multiple sources of data</li> <li>[48] Enforce predefined transaction ordering rules</li></ul>				•
[45] Use and develop new main chain identification mechanisms $M_{G_0}$ <td>_</td> <td></td> <td></td> <td></td>	_			
[47] Trigger automatic liquidations of collateral         M <sub>G7</sub> [40]           [45] Use Collacteralization / Over-Collateralization techniques         M <sub>G8</sub> [153]           [47] Usage of external incentivized watchers that attest actions/events         M <sub>T1</sub> [172]           [48] Usage of Extra aline methors between untrusted users and operators         M <sub>T1</sub> [172]           [40] Usage of Distributed Signature Schemes between untrusted users and operators         M <sub>T2</sub> =	_		$\mathcal{M}_{66}$	Performe deep optimizations once the industry and the project have reached stability
[45] Use Collateralization / Over-Collateralization techniques       M <sub>68</sub> [153]         [117] Embedded rules in third party network consensus mechanism       M <sub>70</sub> [175]         [190] Usage of External incentivized watchers that attest actions/events       M <sub>71</sub> [172]         [163] Parallelizing transaction verification       M <sub>71</sub> [172]         Isseparate entities that create and verify blocks       M <sub>71</sub> [172]         [186] Osage of Distributed Signature Schemes between untrusted users and operators       M <sub>71</sub> [172]         [186] Parallelizing transaction verification       M <sub>71</sub> [172]         [186] Usage of Premiums       M <sub>72</sub> [186]         [187] Provide support for periods of asynchrony in the execution of the protocol       M <sub>72</sub> [176]         [187] Provide support for periods of asynchrony in the execution failures       M <sub>73</sub> [176]         [187] Provide support for periods of asynchrony in the execution of the protocol architecture decentralization       M <sub>73</sub> [176]         [187] Model and analyze use behaviour through game-theory principles       M <sub>73</sub> [176]         [187] Protocol architecture decentralization       A <sub>73</sub> [176]         [88] Increase the number of parties and scatter mining power among them and analyze user behaving transactions to avoid getting uncollateralized       M <sub>73</sub> [162]         [88] Reduce time window for users to observe price fluctuations       M <sub>74</sub> [162]         [45] Adjust the amount locked according to the updated exchange	_			Dispose of private inputs used to generate the CRS in zk-based solutions
<ul> <li>[117] Usage of external incentivized watchers that attest actions/events</li></ul>		-		_
<ul> <li>[141] Embedded rules in third party network consensus mechanism</li> <li>[90] Usage of Distributed Signature Schemes between untrusted users and operators</li> <li>[163] Parallelizing transaction verification</li> <li>Insert independent computational-heavy transactions into multiple blocks</li> <li>[186] Separate entities that create and verify blocks</li> <li>[180] Usage of Premiums</li> <li>[181] Usage of Premiums</li> <li>[182] Usage of Verifiable Timed Commitments</li> <li>[183] Provide support for periods of asynchrony in the execution of the protocol</li> <li>[184] Usage of Verifiable Timed Commitments riggered upon failures</li> <li>[187] Model and analyze user behaviour through game-theory principles</li> <li>[187] Protocol architecture decentralization</li> <li>[187] Protocol architecture decentralization to account for slippage</li> <li>[45] Over-collateralization to account for slippage</li> <li>[45] Adjust the amount locked according to the updated exchange rates</li> <li>[45] Adjust the amount locked according to the updated exchange rates</li> <li>[45] Argie unitiations for (MEV), such as confidential mempools</li> <li>[45] Trigger automatic liquidations for (MEV), such as confidential mempools</li> <li>[46] Employ evolving committees rather than static ones</li> <li>[47] Employ evolving committees rather than static ones</li> <li>[48] Hide one public key annong multiple parties</li> <li>[49] Other usural analy infrastructure back-foor mitigations</li> <li>[4</li></ul>				Set appropriate withdrawal limits and implement a freezing functionality
<ul> <li>1930 Usage of Distributed Signature Schemes between untrusted users and operators M<sub>71</sub> [172]</li> <li>1941 Parallelizing transaction verification</li> <li>1951 Parallelizing transaction verification</li> <li>186 Separate entities that create and verify blocks</li> <li>186 Separate entities that create and verify blocks</li> <li>187 Separate entities that create and verify blocks</li> <li>188 Usage of Premiums</li> <li>189 Losage of Verifiable Timed Commitments</li> <li>191 Use pre-deployed refund transactions/contracts triggered upon failures</li> <li>192 Model and analyze user behaviour through game-theory principles</li> <li>193 Model and analyze user behaviour through game-theory principles</li> <li>184 Disage of MEV to front-run misbehaving transactions</li> <li>185 Parallel asset locking</li> <li>186 Deduce time window for users to observe price fluctuations</li> <li>187 Adjust the amount locked according to the updated exchange rates</li> <li>188 General mitigations for (MEV), such as confidential mempools</li> <li>189 Enforce predefined transaction ordering rules</li> <li>189 Employ evolving committees rather than static ones</li> <li>189 Hide one public key among multiple keys of other users/operators</li> <li>189 Hide one public key among multiple ways of forther users/operators</li> <li>189 Order users/operators</li> <li>180 Order users/operators</li> <li>180 Order users/operators</li> <li>181 Hide one public key among multiple investores</li> <li>180 Order users/operators</li> <li>181 Hide one public key among multiple ways of forther users/operators</li> <li>180 Order users/operators</li> <li>181 Order users/operators</li> <li>182 Order users/operators</li> </ul>				
163   Parallelizing transaction verification				-
Insert independent computational-heavy transactions into multiple blocks   M <sub>73</sub> = -   186   Usage of Premiums   M <sub>74</sub> =   80   Usage of Premiums   M <sub>75</sub> =   80   Usage of Verifiable Timed Commitments   M <sub>76</sub>   186   Usage of Verifiable Timed Commitments   M <sub>77</sub>   187   M <sub>76</sub>   188   M <sub>77</sub>   188   Provide support for periods of asynchrony in the execution of the protocol   M <sub>77</sub>   176   M <sub>77</sub>   176   M <sub>77</sub>   176   M <sub>77</sub>   176   M <sub>78</sub>   176   M <sub>78</sub>   177   M <sub>78</sub>   176   M <sub>78</sub>   178   M <sub>79</sub>   144   M <sub>78</sub>   178   M <sub>78</sub>   178   M <sub>79</sub>   146   M <sub>79</sub>   146   M <sub>79</sub>   146   M <sub>79</sub>   146   M <sub>79</sub>   M <sub>79</sub>   M <sub>79</sub>   146   M <sub>79</sub>   M <sub>79</sub>   M <sub>79</sub>   146   M <sub>79</sub>   M <sub>79</sub>   M <sub>79</sub>   146			$\mathcal{M}_{72}$ –	Treat critical fixes internally before pushing them to public repositories
[186] Separate entities that create and verify blocks $M_{\gamma_1}$ [80] Usage of Perniums $M_{\gamma_5}$ [186] Usage of Verifiable Timed Commitments $M_{\gamma_5}$ [187] Provide support for periods of asynchrony in the execution of the protocol $M_{\gamma_7}$ [91] Use pre-deployed refund transactions/contracts triggered upon failures $M_{\gamma_7}$ [172] Model and analyze user behaviour through game-theory principles $M_{\gamma_7}$ [187] Protocol architecture decentralization $M_{\gamma_7}$ [188] Increase the number of parties and scatter mining power among them $M_{\gamma_7}$ [184] Usage of MEV to front-run misbehaving transactions $M_{\gamma_7}$ [185] Reduce time window for users to observe price fluctuations $M_{\gamma_7}$ [45] Adjust the amount locked according to the updated exchange rates $M_{\gamma_7}$ [45] Adjust the amount locked according to the updated exchange rates $M_{\gamma_7}$ [45] Trigger automatic liquidations to avoid getting uncollateralized $M_{\gamma_7}$ [187] Merge multiple sources of data $M_{\gamma_7}$ [188] Enforce predefined transaction ordering rules $M_{\gamma_7}$ [188] Enforce predefined transaction ordering rules $M_{\gamma_7}$ [184] Hide one public key among multiple keys of other users/operators         [44] Order useral analy inf			$\mathcal{M}_{73}$ –	Make sure critical components are updated before an audit, not afterwards
80] Usage of Premiums   M <sub>75</sub> = -   186   Usage of Premiums   M <sub>76</sub>   186   Usage of Verifiable Timed Commitments   M <sub>77</sub>   186   Provide support for periods of asynchrony in the execution of the protocol   M <sub>77</sub>   176   M <sub>78</sub>   176   M <sub>79</sub>   177   177   M <sub>79</sub>   M <sub>79</sub>   177   M <sub>7</sub>			$\mathcal{M}_{74}$ –	Do not launch projects on top of existing ones without knowing the inner details
Use   Usage of Verifiable Timed Commitments   Use   Usage of Verifiable Timed Commitments   Use   Usage of Verifiable Timed Commitments   Use pre-deployed refund transactions/contracts triggered upon failures   M <sub>77</sub>   Use pre-deployed refund transactions/contracts triggered upon failures   M <sub>78</sub>   Use   Use pre-deployed refund transactions/contracts triggered upon failures   M <sub>78</sub>   Use   U				Fix bugs as soon as they are detected, not just leaving for the future
122   Brovide support for periods of asynchrony in the execution of the protocol   M <sub>T</sub> =   176     122   Model and analyze user behaviour through game-theory principles   M <sub>T</sub>   176     182   Protocol architecture decentralization   M <sub>T</sub>   177     183   Increase the number of parties and scatter mining power among them   M <sub>T</sub>   177     184   Usage of MEV to front-run misbehaving transactions   M <sub>T</sub>   177     185   Parallel asset locking   M <sub>T</sub>   180     186   Adjust the amount locked according to the updated exchange rates   M <sub>T</sub>   180     187   Adjust the amount locked according to the updated exchange rates   M <sub>T</sub>   180     188   General mitigations for (MEV), such as confidential mempools   M <sub>T</sub>   192     188   General mitigations for (MEV), such as confidential mempools   M <sub>T</sub>     189   Enrioree predefined transaction ordering rules   180     189   Enriore problem transaction ordering rules   180     189   Enriore problem transaction ordering rules   180     180   Choral papabilities between multiple parties   M <sub>T</sub>     181   Hide one public key among multiple keys of other users/operators   M <sub>T</sub>     184   Hide one public key among multiple way of other users/operators   180     184   Hide one public key among multiple way of other users/operators   180     185   M <sub>T</sub>   M <sub>T</sub>   M <sub>T</sub>     186   M <sub>T</sub>   M <sub>T</sub>   M <sub>T</sub>     187   M <sub>T</sub>   M <sub>T</sub>     188   M <sub>T</sub>   M <sub>T</sub>     189   M <sub>T</sub>     189   M <sub>T</sub>   M <sub>T</sub>     189		_		Follow standard practices, such as RFCs.
[91] Use pre-deployed refund transactions/contracts triggered upon failures     M <sub>78</sub> [176]       [122] Model and analyze user behaviour through game-theory principles     M <sub>79</sub> [179]       [183] Protocol architecture decentralization     A <sub>80</sub> [162]       [38] Increase the number of parties and scatter mining power among them     M <sub>81</sub> [177]       [84] Usage of MEV to front-run misbehaving transactions     M <sub>81</sub> [175]       [82] Reduce time window for users to observe price fluctuations     M <sub>82</sub> [162]       [82] Reduce time window for users to observe price fluctuations     M <sub>82</sub> [180]       [45] Adjust the amount locked according to the updated exchange rates     M <sub>82</sub> [151]       [45] Trigger automatic liquidations to avoid getting uncollateralized     M <sub>88</sub> [178]       [187] Merge multiple sources of data     M <sub>88</sub> [178]       [188] General mitigations for (MEV), such as confidential mempools     M <sub>89</sub> [178]       [188] Enforce predefined transaction ordering rules     M <sub>90</sub> [144]       [163] Overlap capabilities between multiple parties     M <sub>90</sub> [144]       [184] Hide one public key among multiplie keys of other users/operators     M <sub>92</sub> [146]       [184] Hide one public key among multiplie keys of other users/operators     M <sub>92</sub> [146]				
[122] Model and analyze user behaviour through game-theory principles     M <sub>79</sub> [179]       [187] Protocol architecture decentralization     M <sub>80</sub> [162]       [18] Increase the number of parties and scatter mining power among them     M <sub>81</sub> [177]       [84] Usage of MEV to front-run misbehaving transactions     M <sub>82</sub> [157]       [82] Parallel asset locking     M <sub>83</sub> [162]       [83] Reduce time window for users to observe price fluctuations     M <sub>83</sub> [162]       [45] Over-collateralization to account for slippage     M <sub>84</sub> [180]       [45] Adjust the amount locked according to the updated exchange rates     M <sub>85</sub> [151]       [45] Trigger automatic liquidations to avoid getting uncollateralized     M <sub>85</sub> [178]       [187] Merge multiple sources of data     M <sub>85</sub> [178]       [188] Enforce predefined transaction ordering rules     M <sub>85</sub> [178]       [188] Enforce predefined transaction ordering rules     M <sub>90</sub> [174]       [163] Overlap capabilities between multiple keys of other users/operators     M <sub>91</sub> [146]       [184] Hide one public key among multipliple keys of other users/operators     M <sub>91</sub> [146]				•
[187] Protocol architecture decentralization       Ms       [162]         [38] Increase the number of parties and scatter mining power among them       Ms       [162]         [84] Usage of MEV to front-run misbehaving transactions       Ms       [156]         [82] Parallel asset locking       Ms       [162]         [82] Reduce time window for users to observe price fluctuations       Ms       [162]         [45] Adjust the amount locked according to the updated exchange rates       Ms       [151]         [45] Trigger automatic liquidations to avoid getting uncollateralized       Ms       [156]         [45] Trigger automatic liquidations to avoid getting uncollateralized       Ms       [156]         [48] General mitigations for (MEV), such as confidential mempools       Ms       [177]         [188] Enforce predefined transaction ordering rules       Ms       [174]         [163] Overlap capabilities between multiple parties       Ms       [144]         [177] Employ evolving committees rather than static ones       Ms       [146]         [184] Hide one public key among multiple keys of other users/operators       Ms       [146]         [184] Hide one public key among multiple with static ones       Ms       [146]				
138   Increase the number of parties and scatter mining power among them   M <sub>81</sub>   [177]     139   Increase the number of parties and scatter mining power among them   M <sub>82</sub>   [156]     182   Parallel asset locking   M <sub>83</sub>   [162]     183   Reduce time window for users to observe price fluctuations   M <sub>84</sub>   [181]     185   Over-collateralization to account for slippage   M <sub>85</sub>   [151]     187   Adjust the amount locked according to the updated exchange rates   M <sub>85</sub>   [151]     187   Merge multiple sources of data   M <sub>85</sub>   [156]     188   General mitigations for (MEV), such as confidential mempools   M <sub>85</sub>   [178]     189   Enforce predefined transaction ordering rules   M <sub>87</sub>   [144]     163   Overlap capabilities between multiple parties   M <sub>90</sub>   [144]     184   Hide one public key among multiple keys of other users/operators   M <sub>92</sub>   [146]     184   Hide one public key among multiple keys of other users/operators   M <sub>92</sub>   [146]				•
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184				] Educate users for privacy-preserving practices - e.g., address reuse and unique gas prices
7				
20		Other usual web2 infrastructure backdoor Decentralization at the operational level (		

*Date* The date of the first transaction exploiting a vulnerability in the protocol.

Amount The amount in USD stolen from the cross-chain bridge. We do not include any collateral losses in other protocols.

Attacker Type (AT) We classify attackers as black or white hats based on whether they returned the funds (or both if there is at least one attacker of each type). Attackers who returned the funds, excluding agreed bounty fees, are also considered white hats in our analysis.

Number of Transactions (Txs) A range of the number of transactions issued by the attackers to exploit the bridge, including both external and internal transactions, which are transactions issued directly by the user or as a consequence of another contract execution, respectively. It does not include transactions issued before or after the attack to exchange or launder funds using DEXes (e.g., Uniswap) or mixing services.

Usage of Mixers (Mix) The usage of transaction mixers (e.g., Tornado Cash) by the attacker to launder funds either before or after attacks to break the linkability of transactions.

Discovery Time (DT) The time it took maintainers to discover the attack and trigger the corresponding incident response mechanism. Given that this information is internal to each team, we contacted each of the 14 projects and asked them to provide us with data.

Communication Time (CT) The time it took maintainers to communicate the exploit to the community. This communication was performed solely as *Tweets*. This value is the difference between the timestamp of the *Tweet* and the timestamp of the first exploit transaction.

Vulnerability Location (VL) We identify the location of each vulnerability: in the Source Chain Smart Contract – the component with the bridging logic in the source chain, responsible for escrowing funds; in the Target Chain Smart Contract – the element with the bridging logic in the source chain, responsible for verifying inclusion proofs; or in the Interoperability Mechanism – the off-chain component that enables interoperability, usually composed of validators/relayers.

Exploit Location (EL) A vulnerability in one location can originate exploits in others. We classify the location of the exploit as follows: in the Source Chain Smart Contract if the attacker stole escrowed funds; in the Target Chain Smart Contract if the attacker minted unbacked funds; or in the Business Logic Smart Contract if the attacker stole funds by exploiting the business logic contract – usually because users approved a bridge-controlled contract to manage their funds (e.g., through the approve() function in the ERC20 token standard).

- **4.2.3. Insights.** We present a list of insights taken from the analysis of cross-chain bridge hacks.
  - Insight 1: 65.8% of the total value stolen originated in bridges based on intermediary permissioned networks  $(SA_{22})$ . Projects choose  $SA_{22}$  to have finer control over the bridge. However, it also eases hackers' efforts to gain control over the infrastructure. Three hacks were

- performed on both solutions that rely on centralized services  $(SA_{11})$  and intermediary permissionless networks  $(SA_{21})$  (26.8% and 0.5% of the total value, respectively); and two projects based on fraud proofs  $(SA_{33})$  (6.9%). IMs based on Local Verification  $(SA_4)$  do not enter the leaderboard.
- Insight 2: Limiting the number of internal transactions within the same contract or the amount moved per external transaction is possible and advisable. We believe setting withdrawal limits or emergency pauses would significantly reduce this value. Even though this does not avoid an attack, it serves as a cornerstone step in the first phases of incident response.
- Insight 3: Only one hack is classified as being performed by a white hat, as almost all funds were returned. Moreover, only around 35M USD were returned (1.5% of the hacked amount). There is clearly a lack of motivation for hackers to disclose vulnerabilities. There is little transparency concerning the bounties offered [205].
- Insight 4: Notably, in 14 of the hacks authored by black hats, transaction mixers were used 5 times before the attack (35.7%) and 11 times after the attack (78.6%). The pNetwork hackers did not use any mixer and still retain the funds in their accounts [206]. In the PolyBridge #1 hack, the hacker returned a noteworthy portion of 611M USD after negotiations [207]. We believe these cases highlight the difficulty of money laundering in blockchain environments due to the inherent traceability of blockchain transactions [208].
- Insight 5: We find that the *lock-mint* model for asset transfer bridges is riskier than other approaches. Attackers target escrowed funds in the source chain. Eight hacks drained funds from the escrow in the source chain, accounting for 1.8B USD (62%). Using native tokens instead of wrapped assets is a solution that allows developers to implement one-way flows *burn-mint* models. An example is Circle's USDC announcement in October 2023. USDC is now burned in Ethereum and minted natively on Polygon [209].
- Insight 6: Two teams took 5 and 13 minutes to detect the incidents. The Ronin bridge team took 6 days. This information emphasizes the need for improvement in incident detection for swift attack detection. Work has been done designing cross-chain models to identify and visualize deviations from expected behavior [153], [154], [210].

## **4.3.** Recommendations to Cross-Chain Bridge Operators

We divide the different guidelines for cross-chain systems into three different domains.

**4.3.1. Implementation Layer.** Due to the ad-hoc design and implementation of cross-chain projects [177], [211], attacks exploiting well-known vulnerabilities are still surfacing [212]. Smart contracts, just like any software com-

TABLE 6. CLASSIFICATION OF MOST PROFITABLE CROSS-CHAIN BRIDGE HACKS GROUPED BY SECURITY APPROACH. THE AMOUNTS ARE PRESENTED IN USD. THE CELLS WITH THE VULNERABILITY NUMBER ARE FILLED WITH THE COLOR ACCORDING TO THE LAYER THEY BELONG TO (CF. SECTION 3.1). WE ADD A "SUMMARY" ROW THAT AGGREGATES INFORMATION. SPECIFICALLY, WE USE CELL SHADING TO SHOW THE PERCENTAGE OF HACKS IN WHICH EACH VULNERABILITY WAS FOUND.

Project Informat	ion	Gene	ral Attack	Informa	ation		Incid	ent Resp	Wł	iere	ľ	Mapping to	Theore	etical Vu	ılnerabili	ties
Name & Ref	SA	Date	Amount	AT	Txs	Mix	DT	CT	VL	EL	$\mathcal{V}_{44}$	$\mathcal{V}_{43}$	$\mathcal{V}_{28}$	$\mathcal{V}_{27}$	$\mathcal{V}_{24}$	$\mathcal{V}_6$
[193] Ronin	$SA_{22}$	Mar 2022	624M		0	•	6d	•	IM	SC	1	/	Х	Х	Х	Х
[194] PolyBridge #1	$SA_{22}$	Aug 2021	611M		•	0	-	•	TC	SC	X	1	1	X	Х	X
[195] BNB	$SA_{11}^{22}$	Oct 2022	566M		•	•	_	•	TC	TC	X	X	X	X	✓	X
[108] Wormhole	$\mathcal{SA}_{22}$	Feb 2022	326M		0	•	-	0	TC	TC	X	X	1	X	1	X
[196] Nomad	$\mathcal{SA}_{33}$	Aug 2022	190M		•	•	_	•	SC	SC	X	X	X	X	✓	X
[197] BXH	$\mathcal{SA}_{11}$	Oct 2021	139M		0	•	_	•	_	SC	1	<b>√</b>	X	X	X	X
[198] Multichain #2	$\mathcal{SA}_{22}$	Jul 2023	126M		0	0	_	•	IM	SC	<b>√</b> †	<b>√</b> †	X	X	X	X
[199] Harmony	$\mathcal{SA}_{22}$	Jun 2022	100M		•	•	-	•	IM	SC	1	1	X	X	X	X
[200] Qubit	$\mathcal{SA}_{11}$	Jan 2022	80M		•	•	_	•	SC	TC	X	X	X	1	<b>/</b>	X
[201] pNetwork	$\mathcal{SA}_{33}$	Sep 2021	13M		•	0	13m	•	IM	SC	X	X	X	X	X	1
[202] Thorchain #3	$\mathcal{SA}_{21}$	Jul 2021	8M		0	•	_	-	IM	SC	X	X	X	X	X	1
[198] Anyswap	$\mathcal{SA}_{22}$	Jul 2021	8M		0	•	-	•	IM	TC	X	1	X	X	X	Х
[202] Thorchain #2	$\mathcal{SA}_{21}$	Jul 2021	5M		•	•	-	•	IM	TC	X	X	X	X	✓	1
[194] PolyBridge #2	$\mathcal{SA}_{22}$	Jul 2023	4.4M		•	0	7h	•	IM	TC	X	<b>√</b>	X	X	X	X
[203] Meter	$\mathcal{SA}_{22}$	Jul 2021	4.4M		0	•	_	•	SC	TC	X	X	X	X	<b>/</b>	X
[204] Chainswap	$\mathcal{SA}_{22}$	Jul 2021	4.4M				-	•	TC	TC	X	X	1	X	✓	X
[198] Multichain #1	$\mathcal{SA}_{22}$	Jan 2022	3M		_		-	•	TC	BL	X	X	X	1	✓	X
[202] Thorchain #1	$\mathcal{SA}_{21}$	Jun 2021	140K		-	•	5m	-	IM	TC	Х	Х	Х	Х	Х	<b>√</b>
Summary		07/21 - 07/23	2.9B								22%	39%	17%	11%	44%	22%
Attacker Type (AT)	Num	ber of Transactio	ons (Txs)	Usage	of Mixe	rs (Mix)		Commun	ication	Time (	(CT)	Vulnerabil	lity/Exp	loit Loc	ation (VI	L/EL)
Black hat	O 1-			<ul><li>Not</li></ul>	used			0 ]0; 2]	hours			SC Source	e Chain	SC		
	• 10	-50		Befe	ore the a	ttack		• ]2; 4]	hours			TC Target	Chain S	SC		
Black and white ha				•	er the att			<b>1</b> ]4; 6]				IM Interop			nism	
	<b>●</b> 10	0-1000 1000		• Bef	ore and a	after the a	ttack	<ul><li> ]6; 24</li><li> &gt;= 6</li></ul>	-			BL Busine	ess Logi	c SC		
<ul> <li>No information avai</li> </ul>	•		1	† Still	to be co	nfirmed			.,			Discovery	Time (l	DT)		

ponent, are vulnerable to attacks [33], [190], [213], [214]. To address these vulnerabilities, developers must implement secure coding practices, use continuous integration [215], propose design patterns, and use tools to identify and mitigate potential security issues in the codebase. These include static analysis (e.g., Slither [192], Mythril [216], Mythx [217]), formal verification [218], fuzz testing (e.g., Echidna [219], Harvey [220]), vulnerability detection at runtime (e.g., Scribble [221]), and more recently AI tools to identify vulnerable patterns and perform analysis of control/data flow graphs [33]. It is essential that these tools focus on the composability of smart contracts and not just on the analysis of isolated components.

**4.3.2. Protocol Layer.** Decentralization is paramount for cross-chain solutions, as is evident in Table 6. An infrastructure breach or key compromise can prove catastrophic for platforms that manage substantial funds. For solutions rooted in permissioned networks, it is vital to emphasize due diligence in selecting parties and ensuring expected validator performance, protocol involvement, and compliance with Service Level Agreements (SLA) [107]. Although this approach places greater responsibility on users, a viable mitigation strategy to minimize the likelihood of incidents involves increasing the reliance on user- and Dapp-specific inputs directly submitted to the target chain [65].

We highlight some practices that are of critical importance. On-chain authentication and proof verification mechanisms must be carefully designed and audited by multiple entities. Access control to contracts with critical functionality must be guaranteed by a studied cross-chain model and architecture. Additionally, our recommendations include the usage of formal frameworks to prove the correctness of protocols (e.g., UC [222] or TLA+ [223]), game-theoretical analysis of incentivization, and slashing mechanisms.

**4.3.3.** Operational Layer. At the forefront of originating cross-chain hacks is inadequate key management (1.6B USD stolen, 55%). Although solutions like Hardware Security Modules (HSM) and hardware wallets exist, their scope is limited. To mitigate risk, implementing secure protocols for rotating validator keys could limit the potential for single points of failure. Daily withdrawal limits also emerge as a practical measure to reduce losses during an attack. Table 6 indicates that hackers prefer exploiting vulnerabilities over disclosing them. Therefore, attractive bug bounty programs serve as a compelling avenue to incentivize ethical hackers to report vulnerabilities in open-source code. However, open-source software can expose internal mechanisms and potential security flaws, which has happened twice (cf. Table 10). We, therefore, highlight the need for developers to strike a delicate balance between transparency, collaboration, and system protection. We advocate for multiple accredited entities to verify smart contracts and conduct risk assessments [32].

#### 4.4. Future Research Directions

Our research has identified several areas within the blockchain community that require further exploration and development. Firstly, there is a need for proactive prevention through continuous monitoring of all components, which is not yet a structured practice. Secondly, there is a lack of comprehensive incident response frameworks for generic cross-chain systems despite some industry-specific endeavors. Additionally, the community must focus on designing new techniques to ensure unlinkability and anonymity across diverse ledgers. Furthermore, there is a need to conduct a comprehensive study of design patterns across interoperable blockchain systems. Notably, strides have been made towards a universal data model through the ISO model adopted by Overledger [224], SATP Gateways [44], and IBC [225]. However, the path to standardization is still being explored, with multiple emerging standards. Finally, there is a significant gap in empirical studies addressing the detection of theoretical attacks and associated mitigation strategies identified in our analysis, such as cross-chain Miner Extractable Value (MEV) and oracle manipulation [65], [226]. Appendix D presents some additional considerations.

#### 4.5. Summary

The importance of comprehensive security in crosschain operations cannot be understated. Despite extensive research conducted on cross-chain security, ensuring protection across the entire stack remains imperative. Given the vast attack surface, relying solely on preventive measures, such as continuous monitoring and proactive security, does not suffice. We strongly recommend that practitioners bolster their defenses by integrating reactive security measures, including robust incident response frameworks.

Concerning privacy, current research appears to be relatively underexplored. However, as interoperable central bank digital currencies gain traction [95], [227], we foresee a more substantial impetus driving advances in cross-chain privacy solutions. Full unlinkability in a permissionless cross-chain scenario is hard to achieve since at least one entity needs to perform the mapping between transactions on both the source and destination chains. We envision that protocols that fill this gap will emerge, especially with the recent evolution of zero-knowledge technology. In practice, transaction mixers do not yield a high degree of privacy to users due to their naive practices. Solutions circumventing these limitations are necessary. Furthermore, existing mixers are being used for malicious activities. Therefore, we highlight the importance of researching how to guarantee privacy in regulated and auditable environments.

#### 5. Related Work

We compare blockchain security and privacy studies with our work in Table 7. Our study stands out by 1) employing a systematic survey grounded in recognized principles, 2) fusing essential security and privacy dimensions, and 3) encompassing findings from gray literature and the industry for comprehensive scrutiny, offering developers practical insights to enhance system robustness and privacy. In recent years, notable surveys have emerged[3], [7], [11], [228].

Our examination of security and privacy intricacies revolves around cross-chain rules, explored through various methodologies such as runtime verification [229], formal definitions upheld by off-chain relayers and smart contracts [153], and other formulations to detect discrepancies in bridge mechanisms [154]. Some studies implicitly address cross-chain rules, while others explicitly define them in discussions on oracles [230], [231], bridges [116], and gateways [44].

TABLE 7. COMPARISON OF THIS PAPER WITH OTHERS ON SECURITY OR PRIVACY OF BLOCKCHAIN INTEROPERABILITY. All labels apply to systematizations.

		Sec	curity			Pri	vacy		Mi	sc.
Ref	P	V	A	M	P	V	A	M	R	IM
[17]	Х	<b>√</b> 15	√1	<b>√</b> 18	Х	✓1	Х	√1	46	2
[13]	X	<b>√</b> 29	√7	<b>√</b> 13	X	<b>√</b> 1	Х	X	_	6
[232]	X	X	X	Х	/	X	X	<b>/</b> 4	_	29
[233]	X	<b>√</b> 11	<b>√</b> 18	√6	Х	Х	X	X	_	_
this work	1	<b>√</b> 45	<b>√</b> 18	<b>√</b> 92	1	15	X	/3	212	57

- ✓ Satisfies criteria
- P Identifies relevant properties
- X Does not satisfy criteria
- M Identifies or proposes mitigations
- A Real-World attacks or leakages
   R Number of relevant references
- V Identifies cross-chain vulnerabilities
- IM Number of interoperability mechanisms systematically studied
- Not specified by the authors

Since blockchain's inception, security has been a key research focus [234]. Studies have aimed to formalize security properties [235]–[240] and consolidate this knowledge across various layers, such as network, protocol, and operations [241]–[246]. Research has extensively delved into attacks in the same layers [234], [244], [247]–[249]. Our exploration of privacy builds on seminal work on anonymity, unlinkability, and confidentiality [125], [250]. Although there are endeavors to introduce permissioned and privacy-oriented blockchains [251], [252], privacy research in interoperability primarily focuses on asset exchanges. Our conceptual model is based on [87] and extends to cover all interoperability modes.

#### 6. Conclusion

This paper systematizes the relevant security and privacy properties and approaches in blockchain interoperability research. We correlate theoretical vulnerabilities and 14 cross-chain bridge hacks, collectively responsible for 94% of the total hacked value up to date. Our data shows that security is a more pressing concern than privacy and that permissioned intermediary networks are the most vulnerable security approach. Regarding privacy, our survey reveals a prevalent reliance on zero-knowledge technology. Although this method is promising, it requires extensive additional research before being widely adopted. We collect and propose mitigations for identified vulnerabilities and outline various research pathways, such as reliable monitoring of IMs, incident response frameworks, the need for empirical studies, and further exploration of cross-chain privacy.

## Acknowledgments

We thank Dave Pasirstein and our colleagues in the IETF's SATP working group for fruitful discussions. We appreciate the DPSS group at INESC-ID for suggestions and discussion on this paper. Rafael thanks the Fulbright Association, Thomas Hardjono, and Alex Pentland for the opportunity to research at the MIT Media Lab, where this work was partially developed. We thank Guy Zyskind and colleagues from the MIT Media Lab for fruitful discussions on blockchain privacy and interoperability. Rafael was supported by national funds through Fundação para a Ciência e a Tecnologia (FCT) with reference UIDB/50021/2020 (INESC-ID) and 2020.06837.BD, and a research scholarship from the Fulbright Association, with the support of FCT. This work was developed within the scope of the project nr. 51 "BLOCKCHAIN.PT - Agenda Descentralizar Portugal com Blockchain", financed by European Funds, namely "Recovery and Resilience Plan - Component 5: Agendas Mobilizadoras para a Inovação Empresarial", included in the NextGenerationEU funding program. Luyao Zhang is supported by the National Science Foundation China on the project entitled "Trust Mechanism Design on Blockchain: An Interdisciplinary Approach of Game Theory, Reinforcement Learning, and Human-AI Interactions (Grant No. 12201266)." Luyao Zhang is also with SciEcon CIC, a not-for-profit (NPO) organization registered in the United Kingdom, aiming to cultivate interdisciplinary research and integrated talents. We thank Kevin Liao for helping us improve the visual quality of this paper.

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# **Appendix A.** Cross-Chain Concept Formalization

A transaction t is considered final in a ledger l according to a security parameter  $\lambda$  of that network (e.g., the block containing the transaction has a minimum height), and is represented as  $final^{l}(t) \rightarrow \{0,1\}$ .

A local transaction yields a state change in the form of a key-value pair. The execution of local transactions emits events. Events act as labels or wrappers for state changes caused by local transactions.

**Definition 9** (Cross-Chain Event). A ccevent gives a cross-chain meaning to a local event. It extends a local event with metadata, representing a state change in a certain ledger. We denote  $e_{type}^{l \in \mathcal{L}}(t)$  a cross-chain event that represents a state change of type type against t.target, in domain l, emitted by transaction t, such that f inal l(t) = 1.

In our model, a *ccevent* is only created when the transaction that emits the corresponding local event is considered final. However, it might be valid or not according to  $\mathcal{R}$  that defines the expected behaviour.

**Definition 10** (Valid Cross-Chain Event). A cross-chain event  $e_{type}^{l_1 \in \mathcal{L}}(t)$  is deemed valid if and only if it follows the defined cross-chain rules related to it.

Note that the validity of an event emitted by a local transaction does not imply the validity of the corresponding cross-chain event because the latter might not comply with the defined cross-chain rules.

Formally, a *cctx* is then a composition of *n* ordered crosschain events  $\mathcal{E}$  across multiple ledgers  $\mathcal{L}$  with the same cctxid, such that  $\mathcal{E} = \{e_{type_1}^{l_1 \in \mathcal{L}}(t_1), e_{type_2}^{l_2 \in \mathcal{L}}(t_2), ..., e_{type_n}^{l_k \in \mathcal{L}}(t_i)\}$ . The validity of a cross-chain transaction is given by the conjunction of the validity of every cross-chain event in  $\mathcal{E}$ , that is evaluated against a set of rules  $\mathcal{R}$ . We consider blockchain rules  $\mathcal{R}$  to be a composition of predicates  $\mathcal{E} = \{\mathcal{E}_1(\mathcal{E}), \mathcal{E}_2(\mathcal{E}), ..., \mathcal{E}_n(\mathcal{E})\}$  over a set of events  $\mathcal{E}$ .

# Appendix B. SoK Methodology

In this section, we present further details on our systematic survey methodology. Figure 2 presents the PRISMA diagram [253] for our survey.

We conducted a systematic literature review by crawling papers using Google Scholar's keyword search. The search was limited to papers since 2015 due to the limited amount of research available before that period. The following search query was used to search for papers within our research scope:

("blockchain interoperability" OR "cross-chain") AND ("attack" OR "incident" OR "hack" OR "leaks") AND (
("security" AND ("vulnerability" OR "mitigation"))
OR "privacy")

Blockchain interoperability research has been rapidly evolving in the last couple of years. Yet, academic and peerreviewed work alone falls short of delivering the most up-todate facts on interoperability, particularly in the analysis of cross-chain hacks. We pay attention to a significant amount of material available as grey literature in online databases such as Rekt and Slowmist, and online audit reports by reputed companies in the area such as *Certik*, *Chainsecurity*, Consensys, Halborn and Trail of Bits. We also find that many incident reports are divulged through unstructured and informal means of communication, namely blog or social media posts [254]. We strive to uphold the integrity of the findings presented in this work, diligently cross-referencing information whenever possible. Therefore, to the best of our knowledge, the material presented is the most reliable and up-to-date.

To mitigate the potential for unsoundness or bias due to the reliance on gray literature, we meticulously compile data from several sources, and each resource undergoes thorough examination and assessment by at least two paper authors. Additionally, we acknowledge that our analysis of industry solutions and associated vulnerabilities may not encompass the entirety of the landscape, primarily due to limitations in the available documentation. Nevertheless, we make diligent efforts to compile all accessible information concerning projects that collectively represent over 75% of the Total Value Locked (TVL) in cross-chain solutions [32].

## Appendix C. Confidentiality as a requirement for Cross-Chain Unlinkability

From our research, we show that the degree of linkability, and consequently anonymity, yielded by an interoperability solution is tied to the capacity to keep local transactions' content confidential. This idea is also supported by [255]. We study the requirements for cross-chain unlinkability to be achieved and propose Lemma C.1, which is proved below. We denote *confidential systems* as *C* and *non-confidential systems* as *S*. Confidential systems are either private blockchains (e.g., Hyperledger Fabric, Quorum,

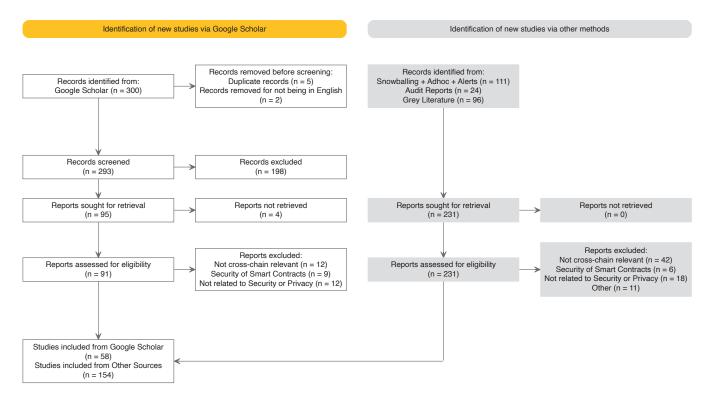


Figure 2. PRISMA Diagram depicting our methodology [253]

DAML's Canton) or public blockchains that are hardened by a privacy-preserving mechanism that hides transactional data (e.g., ZCash, Monero). Non-confidential systems are typical layer-one blockchains with no concern over privacy (e.g., Ethereum, NEAR). We denote an interoperation process (e.g., asset transfer/data transfer) by  $\rightarrow$ .

**Lemma C.1.** Cross-chain unlinkability is unlikely to be achieved without confidentiality on the underlying chains.

*Proof.* Interoperability inherently relies on linkability between transactions on a source and target blockchain. However, one must consider that this linkability is undesirable for general users, as it compromises the degree of privacy yielded by the protocol. We identify a direct relation between the confidentiality guarantees offered by one blockchain and the cross-chain anonymity and unlinkability offered by interoperability solutions built on top of these. Table 8 summarizes the comparison based on the privacy guarantees of the privacy guarantees of the ledgers, which is further explained below.

 $\mathcal{S} \to \mathcal{C}$ : On the source chain, transaction data (such as the amount, sender, and recipient) is open to everyone. However, since the destination chain is private or has privacy-preserving primitives, only authorized parties can see and link a transaction to the one issued on the source chain. To issue transactions on a permissioned network, there must be a trusted and identified IM with access to the ledger and, optionally, to private channels. Assuming this third party does not disclose information external parties cannot

TABLE 8. MAXIMUM ACHIEVABLE CROSS-CHAIN ANONYMITY BASED ON THE CROSS-CHAIN UNLINKABILITY, WHICH IS DEPENDENT ON THE CONFIDENTIALITY OF THE UNDERLYING CHAINS

Source	Target	Max. Unlinkability	Max. Anonymity
Chain	Chain	Achievable	Achievable
S	S	CC Linkability	CC Pseudonymity
C	S	CC Unlinkability	CC Anonymity
S	C	CC Unlinkability	CC Anonymity
Č	$\overset{\circ}{\mathcal{C}}$	CC Unlinkability	CC Anonymity

link transactions across chains. Cross-chain unlinkability is guaranteed.

 $\mathcal{C} \to \mathcal{S}$ : The key challenge in this setting is for an action that occurred in the source chain to be validated externally. The two options are a light client in the target chain, where the user presents a way of decrypting source blockchain data, or an interoperability mechanism that has access to the source chain and acts as a trusted party. In the former, an option might be the usage of zero-knowledge proofs, which can be verified while maintaining data confidentiality. In the latter, linkability is possible only if the trusted IM discloses information. A trusted IM can access this information, verify its validity, and issue transactions on the public chain accordingly. Since the IM must be trusted [3], there is cross-chain unlinkability.

 $\mathcal{C} \to \mathcal{C}$ : Assuming both blockchains are confidential, only authorized parties can link transactions, including a

trusted IM with access credentials on both the source and target chain. Applying the same logic above, an external observer cannot link transactions in each chain.

 $S \rightarrow S$ : The analysis of various heuristics, such as transaction amounts, addresses, or shared secrets, can enable linking transactions across multiple chains [65], [77]. Mixing services (cf. Section 3.5.1) help cover traces and break transaction linkability. In an ideal setting, these systems achieve their goals perfectly. However, in the real world, these have been studied and are shown not to be effective [126].

Interoperability requires a trusted party [3]. In centralized settings, the trusted party can hold records of transactions and corresponding mappings. Therefore, privacy concerns may arise, such as the leakage of private information or, in the worst-case scenario, sold [20]. With a trusted centralized entity that removes outdated records and does not keep track of transactions, there is unlinkability without the risk of being compromised. Cryptographic methods, such as blind signatures, might be a safeguard. We derive a direct consequence from the above Lemma C.1: since cross-chain anonymity depends on the unlinkability of *cctxs*, we extend our initial thoughts in Lemma C.2.

**Lemma C.2.** Cross-chain anonymity is unlikely to be achieved without confidentiality on the underlying chains.

*Proof.* Cross-chain anonymity is driven by cross-chain unlinkability, and cross-chain unlinkability is unlikely to be achieved without the confidentiality of the underlying chains. Therefore, cross-chain anonymity is unlikely to be achieved without confidentiality on the underlying chains due to the incapacity to provide cross-chain unlinkability under those conditions.

We derive the main consequence of the above ideas in Corollary C.2.1.

**Corollary C.2.1.** The privacy level offered by the interoperability solution is upper-bounded by the intersection of the privacy levels of the underlying chains.

# Appendix D. Future Research Directions

In this section, we delve deeper into the future research outlined in Section 4.4.

## **D.1.** Monitoring in Cross-Chain Systems

Given the inherent vulnerabilities in software systems, enhancing the robustness of cross-chain solutions becomes paramount. Initial efforts should focus on the formal verification of cross-chain protocols using an array of tools to augment the likelihood of their correctness. Concurrently, establishing rigorous security and engineering practices for both on-chain and off-chain components is essential. This includes the implementation of automated tests and the meticulous security scrutiny of software dependencies. Proactive

prevention can be achieved through the continuous monitoring of all components. Although the Hephaestus framework presents an intriguing direction [153], empirical benchmarks in real-world contexts remain an essential avenue for exploration. Studying and enhancing UX and UI practices for detecting and reporting attacks is also relevant, to increase the transparency, process, and integrity of attack/vulnerability disclosure. Although some preliminary work is done in this area [210], [256], specific solutions deployed at scale are still missing.

Note that automated frameworks for vulnerability finding often have a high degree of false positives. Additionally, the tendency for contracts to become larger and more complex is exacerbated by the newly cross-chain smart contract paradigm, which will lead to new, complex, hard-to-debug vulnerabilities. These are expressed either as native cross-chain smart contracts [68], by the orchestration of smart contracts across chains [55], or by cross-chain contract calls [257]. We foresee that the market will need cyber-security professionals with skills in several blockchains, programming languages, and attacks to protect systems from these vulnerabilities.

## **D.2.** Frameworks for Incident Response in Cross-Chain Contexts

Software platforms interfacing with the internet, especially those governing sensitive tasks like cross-chain bridges, necessitate dedicated cybersecurity oversight. The current research landscape underscores the need for enhanced operational security. Preliminary metrics for detecting bridge discrepancies exist [258], yet manual or automated responses each present their challenges. Mistaken detections, for instance, can result in bridge suspensions, impacting user experience and revenue. To date, comprehensive incident response frameworks for generic cross-chain systems remain largely uncharted, despite some industry-specific endeavors.

#### D.3. Cross-chain Privacy in Heterogeneous Systems

Our findings suggest that cross-chain privacy is predominantly maintained within underlying chains that inherently support privacy-enhancing features. The development and exploration of techniques to ensure unlinkability and anonymity across diverse ledgers remain areas of underexplored research within the scientific community.

## D.4. Blockchain interoperability design patterns

Design patterns serve as structured frameworks, enabling developers to craft secure solutions with augmented efficacy. Although each interoperability context possesses distinct characteristics, discerning common challenges and pitfalls inherent to specific interoperability solutions can yield invaluable insights. While blockchain design patterns have undergone rigorous scrutiny, a comprehensive examination of

design patterns across multifarious blockchain applications remains nascent in the current research landscape, reflecting the evolving nature of this domain.

## D.5. Data models for blockchain interoperability

Data models are fundamental to interoperability, streamlining complex mappings and varied data formats. Abstract models facilitate a semantic perspective for developers, mirroring the role of SDKs in emphasizing business logic over implementation nuances. Notable strides towards a universal data model are evident through ERC-5164, the ISO model (as adopted by Overledger [224]), SATP Gateways [44], and IBC [225]. However, the path to full standardization remains under exploration, with multiple standards emerging concurrently. The preference for open standards is evident and is crucial for achieving technical interoperability. The importance of this is underscored by initiatives such as BUNGEE [43].

## **D.6.** Empirical Investigations

The research landscape reveals a notable gap in empirical studies addressing the detection of theoretical attacks and associated mitigation strategies identified in our analysis. Additionally, there seems to be a scarcity of indepth examinations focusing on specific IMs. A couple of research trajectories stand out in terms of their pertinence and potential impact. Firstly, the identification of cross-chain Miner Extractable Value (MEV) is becoming increasingly salient due to the rapid expansion of blockchain bridges, coupled with substantial investments to enhance their usability and facilitate the onboarding of newcomers. Secondly, the empirical exploration of oracle manipulation within the cross-chain context [65], [226], [259] presents a promising direction for future investigations.

## **Summary**

The importance of comprehensive security in crosschain operations cannot be understated. Despite the extensive research conducted on cross-chain security, ensuring protection across the entire stack remains imperative. Given the vast attack surface, solely relying on preventive measures, such as continuous monitoring and proactive security, is insufficient [153]. We strongly recommend practitioners bolster their defenses by integrating reactive security measures, including robust incident response frameworks.

Regarding privacy, current research appears to be relatively underexplored. However, as interoperable central bank digital currencies gain traction – evidenced by references like [95], [227] – we foresee a more substantial impetus driving advancements in cross-chain privacy solutions. We envision that protocols filling this gap will emerge especially with the recent evolution of zero-knowledge technology, according to our findings. We also highlight the importance of researching how privacy can be guaranteed in regulated and auditable environments.

TABLE 9. Dataset of Cross-Chain Bridge Hacks ordered by  $$\operatorname{\textsc{Date}}$$ 

Bridge Name	Hack Date	Amount (Million USD)
Thorchain	June 2021	0.14
Thorchain	July 2021	5.00
Thorchain	July 2021	8.00
Thorchain	July 2021	0.08
Chainswap	July 2021	4.40
Chainswap	July 2021	0.80
Anyswap	July 2021	8.00
Poly Network	July 2021	4.40
Poly Network	August 2021	611.00
pNetwork	September 2021	13.00
BXH	October 2021	139.00
Nerve	November 2021	0.54
Multichain	January 2022	3.00
Qubit	January 2022	80.00
Wormhole	February 2022	326.00
Meter	February 2022	7.70
Ronin	March 2022	624.00
Harmony	June 2022	100.00
Nomad	August 2022	190.00
CelerNetwork	August 2022	0.24
BNB	October 2022	566.00
QANplatform	October 2022	2.00
Rubic	November 2022	1.20
pNetwork	November 2022	10.80
Rubic	December 2022	1.40
Multichain	February 2023	0.13
Allbridge	April 2023	0.57
Cellframe Network	June 2023	0.07
Multichain	July 2023	130.00
Mixin Network	September 2023	200.00
Orbit Bridge	December 2023	81.88
Socket	January 2024	3.30
	Total	3238.04

## Appendix E. Real World Cross-Chain Bridge Hacks

In Table 9 we present the entire dataset of cross-chain bridge hacks since June 2021. Additionally, Table 10 presents a more thorough analysis of each of the 18 hacks studied in Section 4.2.2. We describe the hack in detail, map to our model and present a set of mitigations that can mitigate similar attacks in the future.

TABLE 10. DESCRIPTION AND POSSIBLE MITIGATIONS FOR SOME OF THE MOST PROFITABLE CROSS-CHAIN BRIDGE HACKS.

Bridge & Refs	Mapping to our model?	Description	Mitigations
Ronin Bridge [260], [261]	The validators were compromised. The attackers compromised 5 out of 9 validators – the exact threshold.	<ul> <li>Nobody noticed for 6 days. No monitoring existed</li> <li>5 out of 9 validators were needed to approve transactions, whereas 4 were controlled by the same entity Sky Mavis.</li> <li>The Axie DAO controlled the 5th validator, however, there was a gas-free RPC node through which Sky Mavis had access to the Axie DAO validator.</li> </ul>	<ul> <li>M<sub>66</sub> – Insert monitoring procedures in the bridge.</li> <li>M<sub>60</sub> – Improve cryptographic key management (e.g., cold wallets, or multi-signatures).</li> <li>M<sub>61</sub> – Increase the number of validators, and the threshold necessary to deem a proof valid.</li> <li>M<sub>52</sub> – Audit not only snart contracts but all the infrastructure that is behind.</li> <li>M<sub>60</sub> – Set withdrewal limits.</li> <li>M<sub>70</sub> – Do not give excessive permission to individual external entities.</li> </ul>
PolyBridge [172], [262]	The contract that manages the public keys of active keepers. The user accessed this contract through the bridge one.	<ul> <li>EHICrossChainManager (the bridge contract) has a function executeCrossChainTx() that calls a destination smart contract to unlock tokens passed as an argument. Also, there is a EhICrossChainData contract that manages validators' public keys.</li> <li>The attacker accessed the function putCarEpochConPubKeyBytes(bytes) that updates the validators' public keys in the EhICrossChainData contract.</li> <li>Since EhICrossChainManager is the owner of EhICrossChainData, the all would pass the isOwnerd; check.</li> <li>Also, to trick the EhICrossChainManager contract to call that function, the attacker relied on his knowledge of the EVM. He found a hash collision where the first four bytes would match the first four bytes of the signature hash of the target function.</li> </ul>	• $M_{tg}$ – Smart contracts accessed by users should not have direct access to management smart contracts. • $M_{tj}$ – In dynamic bridges, when receiving a contract address as an argument, check that it represents a contract and, if possible, that the corresponding method being called is valid. • $M_{tg}$ – Set withdrawal limits.
PolyBridge [263]-[266]	Validators' keys were compromised	<ul> <li>3 out of 4 keys were compromised for a 3 out of 4 multi-sig.</li> <li>The attacker authorized withdrawals in multiple destination chains</li> </ul>	• $M_{60}$ – Improve cryptographic key management (e.g., cold wallets, or multi-signatures). • $M_{61}$ – Increase the number of validators, and the threshold necessary to deem a proof valid.
BNB Bridge [173], [267]	Buggy proof verification mechanism.	<ul> <li>The Cosmos' IAVL proof go implementation had a bug when checking the validity of the proof on the target chain.</li> </ul>	• $M_{41}$ – Make sure third party components/libraries have been audited by multiple entities. • $M_{68}$ – Freeze deposits and withdrawals to and from the bridge.
Wormhole [268], [269]	A bug was introduced in the proof verification component in the target chain.	<ul> <li>The bridge relies on a set of 19 guardians.</li> <li>Guardians sign evenes enrited on both chains.</li> <li>Due to a mistake, the version solama program being used didn't verify correctly the signatures.</li> <li>The attack happened a few hours after a parch was fixed in CitHub.</li> </ul>	<ul> <li>M<sub>12</sub> – Do not publish (push) critical fixes before those changes are deployed.</li> <li>M<sub>51</sub> – Perform security audits with the different versions of the libraries being used.</li> <li>M<sub>50</sub> – More than 93,000 ETH was moved back to Ethereum which could have been avoided if transfers the pusceblooked</li> </ul>
Nomad Bridge [170], [270], [271]	The verification of the lock proof would deem every message valid.	<ul> <li>The contract was initialized with 0x0 as a trusted root.</li> <li>If calling directly the process() function in the Reptica contract (the bridge smart contract in each chain), it would internally call the acceptable/Roof) function with the reasege/directorge/Hostoge state in the acceptable/Roof) function with the reasege/directorge-function from the cum 0x0, calling the acceptable/Roof) function with 0x0 as a parameter. Given that 0x0 was set as a trusted root, the validation would be successful.</li> </ul>	• $M_{2i}$ – A bug was introduced shortly after an audit by an external team. Changes in critical components of the code should not be done after an audit.
Harmony [272]–[274]	Validators' keys were compromised	<ul> <li>The bridge relied on a 2 out of 5 multi-signature and two addresses in a hot wallet were compromised.</li> <li>Humony stopped the Horizon bridge to prevent further transactions, however, the attacker was able to swap funds to ETH in the Ethereum Hotscheim.</li> </ul>	• $M_{\rm ejj}$ – Improve cryptographic key management (e.g., cold wallets, or multi-signatures). • $M_{\rm ejj}$ – Increase the number of validators, and the threshold necessary to deem a proof valid.
Qubit Finance [168], [275]	A deprecated function allowed minting tokens without a valid proof.	<ul> <li>Developers forgot to remove a deprecated function from the smart contract that allowed a zero address to call a safeTransferFrom().</li> <li>The attacker called the deposit function in the source chain (Ethereum) passing an invalid token. The transaction did not revert and the smart contract in the destination chain was instructed to mint xETH tokens to the attacker's address on BSC.</li> </ul>	• $M_{Sg}$ – Automatic tools allow identifying deprecated (or unused) functions. • $M_{Sg}$ – Code reviews should suffice to mitigate this vulnerability.
Meter [169]	Implementation bug in the deposit function in the contract deployed to the source chain.	<ul> <li>The implementation was a modified version of Chainsafe's ChainBridge. They added a new function to deal with native tokens.</li> <li>The attacker exploited the existing deposity function. The handler (called from the deposit) function) had a faulty condition to check for wrapped assets which assumed tokens were already transferred by the bridge. The attacker could pass an arbitrary amount of funds that were transferred to the attacker's address.</li> </ul>	• $M_{2i}$ – Know and understand thoroughly a codebase before forking it, especially before creating new functionality. • $M_{5i}$ – Audit code before and after changing code.
Thorchain #1 [171]	Locking a token with a name similar to ETH was interpreted as valid ETH.	<ul> <li>Logical bug in the Ethereum Bifrist smart contract caused fake tokens named "ETH" to be interpreted as real ETH.</li> </ul>	. $M_{\Omega}$ – Whitelisting valid token contracts.
Thorchain #2 [276], [277]	Implementation bug in the relayer. The attacker pretended to transfer funds, taking advantage of the deposit verification mechanism.	<ul> <li>The hacker wrapped the router with a specific smart contract implementation that tricked the bridge smart contract.</li> <li>They were able to trick the bridge smart contract into processing meg, wither tokens, even if 0 was provided.</li> <li>The vulnerability was already known (there was also a comment in the code pointing to the affected loop)</li> </ul>	<ul> <li>M<sub>Si</sub> – Fix bugs as soon as they are detected.</li> <li>M<sub>si</sub> – This attack was performed using multiple transfers, thus, a limit could have been set to limit the number of transfers – e.g., per user, per day, or both.</li> </ul>
Thorchain #3 [278]–[280]	The user forged an event deposit. Convinced the bridge that a certain action took place.	<ul> <li>The attacker noticed that if a deposit event was sent to the bridge smart contract the returnleathAssea() function would return the "deposited" ETH to the sender.</li> <li>The attacker deployed a fake contract, emitted a false event and sent it to the function, which returned real ETH to the attacker</li> </ul>	. $M_{\Omega}$ – Whitelisting valid token contracts.
Chainswap [281], [282]	Access to critical infrastructure allowed whitelisting attacker addresses.	<ul> <li>On Ethereum, the attacker was able to exploit the proxy factory contract, mining tokens directly into different addresses.</li> <li>There was a bug in the token cross-chain quota code. The signature node automatically increases the on-chain swap bridge quota. However, a logical code flaw allowed addresses that weren't whitelisted to automatically increase this amount.</li> </ul>	. $M_{\rm SI}$ – Smart contract auditing to identify vulnerabilities – in this case in the signature verification procedure.
pNetwork [283], [284]	The attacker emitted fake token burn events which were accepted by the source chain.	<ul> <li>Attackers stole BTC collateral for the pBTC-on-BSC bridge.</li> <li>A series of events was created from failty contracts, with only one being a legit peg-out request.</li> <li>The attacker produced bridge beak requests which were accepted.</li> <li>On the original side (BTC) did not validate the requests, unlooking the funds to the attacker.</li> </ul>	. $\mathcal{M}_{12}$ – Whitelisting valid token contracts. Evens from invalid contracts should not be accepted.
Anyswap [285], [286]	The attacker exploited a bug in the signature generation algorithm.	<ul> <li>The attacker reconstructed the V3 router MPC account's private key due to a bug in the implementation of the random value generation algorithm.</li> <li>Two transactions used the same (supposedly) random value to generate a signature which compromised the private key</li> </ul>	. $\mathcal{M}_{\%}$ – Follow sundard practices, namely RFC 6979
Multichain [175]	The attacker bypassed the signature verification on the target chain.	<ul> <li>The function that originated the attack was not being used anywhere, it was dead code.</li> <li>A false token contant was understood by the bridge as a Multichain incline.</li> <li>The attacker could bypass the signature verification in the ERC-20 permit/f) function, by supplying a token contract that didn't have the permit/f function implemented. This caused the fallback function to run instead of reverting.</li> <li>Multichain dapp requested from all of its users a practically infinite approval sum to save gase fees, which allowed the user to withdraw practically any funds.</li> </ul>	. $M_{S0}$ – Automatic tools allow identifying dead code. . $M_{L0}$ — Whitelisting with other control of
Multichain [27]	The attacker accessed internal off- chain infrastructure	<ul> <li>The details of this hack have not been explained yet. We refer the reader to [287] for further details.</li> <li>Either a rug pull, an internal hack, or an off-chain infrastructure compromise.</li> </ul>	We cunnot propose mitigations due to not having the details of the attack
BXH [288], [289]	The private key associated with the bridge smart contract on the destination chain was compromised	<ul> <li>Pedshield and BXH came to the conclusion that the hacker took control over the smart contract deployed at the Binance Smart Chain after getting hold of the administrator's private key</li> <li>It is suspected that this attack was an inside job using a phishing email.</li> </ul>	• $M_{\eta_0}$ – Remove a single point of failure caused by admins (or users) having excessive permissions. • $M_{\eta_0}$ – Inside job attacks can be mitigated using multi-signatures or multi-party computation to decentralize responsibilities.

## Appendix F. Meta-Review

## F.1. Summary of Paper

This paper surveys the current state of blockchain interoperability mechanisms in both academia and industry. It classifies their approaches and how they correspond to real-world attacks and vulnerabilities. Each section provides insights and principles for the design of security, privacy, and governance mechanisms of future work.

#### F.2. Scientific Contributions

• Provides a dataset for public use.

## F.3. Reasons for Acceptance

- The paper provides a well-written and clear survey of the current state of blockchain interoperability mechanisms.
- 2) The work connects attacks to causes, and shows how to categorize those attacks based on the causes.
- The authors identify gaps in current research and development trends as well as principles to avoid future pitfalls

### F.4. Noteworthy Concerns

- The insights provided in the paper are very low-level and concern either the low-level structure of existing systems, or low-level suggestions for building future ones. Coupled with the large scale of the tables and density of information, this makes it challenging to extract new high-level takeaways about how to think about the overall security of interoperability mechanisms.
- 2) Some of the definitions lack a level of formality needed to apply them to specific systems.

# Appendix G. Meta-Review Response

We emphasize that Section 4.3 provides general recommendations for operators of blockchain interoperability systems, offering high-level insights classified into the relevant layers. For example, in 4.3.1, we highlight the ad-hoc design of cross-chain protocols, which increase the attack surface for hackers, the need for monitoring, and continuous integration practices. Similarly, in 4.3.2 and 4.3.3, we highlight high-level aspects that are crucial to the area such as Service Level Agreements (SLA) for operators, attractive bug bounties, and the exposure of code through opensource. Additionally, Section 4.4 outlines Future Research Directions, presenting broader insights into areas of interest within blockchain interoperability, including considerations

on monitoring, incident response, standardization bodies, and the potential application of Miner Extractable Value (MEV) as a defense mechanism for cross-chain scenarios. We believe these sections effectively address the need for concise, high-level insights and demonstrate our commitment to providing valuable contributions to the field. Finally, we would like to emphasize that our manuscript includes 17 relevant and concise insights. Each numbered insight encapsulates key findings derived from our extensive research and interpretation of the large tables presented. We have made a big effort to make sure the relevant notations are added to each table to improve readability.