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Concurrent Erlang



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Overview: concurrent Erlang I

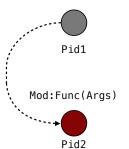
- Concurrent Erlang I
 - Creating Processes
 - Message Passing
 - Receiving Messages
 - Data in Messages
- · Concurrent Erlang II



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Creating Processes

spawn(Mod, Func, Args)



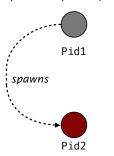
- Before
 - Code executed by Process 1
 - process identifier is Pid1
 - Pid2 = spawn(M, F, A)
- After
- A new process with Pid2 is created
- Pid2 is only known to Pid1
- Pid2 runs M:F(A)
- M:F/Arity must be exported
- Convention: we identify processes by their process ids (pids)



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Creating Processes

spawn(Mod, Func, Args)

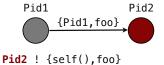


- The BIF spawn never fails
- A process terminates
 - abnormally when run-time errors occur
 - normally when there is no more code to execute



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Message Passing



- Messages are sent using the Pid! Msg expression
 - **Msg** is any valid Erlang data type
- Sending a message will never fail
- Messages sent to nonexisting processes are thrown away
- Received messages are stored in the process' mailbox

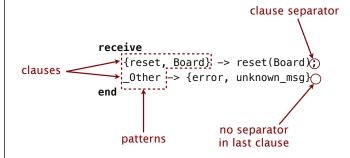
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Message Passing





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Receiving Messages

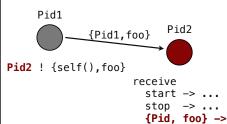
```
receive
    Pattern1 ->
         <expression 1>,
         <expression 2>,
         <expression N>;
    Pattern2 ->
         <expression 1>,
         <expression N>;
    PatternN ->
         <expression 1>,
         <expression N>
end
```

- Messages are retrieved using a receive clause
- receive suspends the process until a message is received
- · Message passing is asynchronous

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Receiving Messages



end

 Messages can be matched and selectively retrieved

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- Messages are received when a message matches a clause
- · Mailboxes are scanned sequentially.

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Receiving Messages

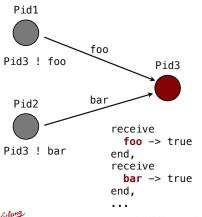
```
Pid2 ! {self(), {digit, [1,2,3,4]}}
                                                    Pid2
Pid1
              {Pid1, {digit, [1,2,3,4]}}
                                 receive
                                   {Pid, {digit, Digit}} ->
                                end
```

- If Pid is bound before receiving the message, then only data tagged with that pid can be pattern matched
- The variable **Digit** is bound when receiving the message



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Receiving Messages: selective



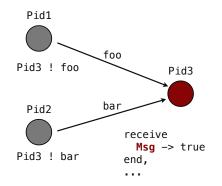
- · The message foo is received, followed by the message bar
- This is irrespective of the order in which they were sent or stored in the mailbox

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Receiving Messages: non-selective

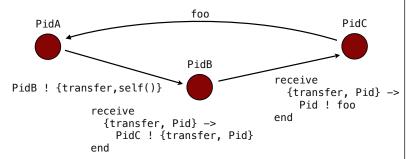


- The first message to arrive at the process Pid3 will be processed
- The variable Msg in the process Pid3 will be bound to one of the atoms foo or bar depending on which arrives first.

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Receiving Messages



- PidA sends a message to PidB containing its own Pid
- PidB binds it to variable A and sends a message to PidC
- PidC receives the message and replies directly to PidA

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Data in Messages: example

```
-module(echo).
-export([go/0, loop/0]).
go() ->
                                   loop() ->
   Pid= spawn(echo,loop,[]),
Pid ! {self(), hello},
                                   receive
                                      {From, Msg} ->
   receive
                                        From ! {self(), Msg},
                                          loop();
     {Pid, Msg} ->
     io:format("~w~n",[Msg])
                                       stop ->
   end,
Pid ! stop.
                                          true
                                     end.
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```

Summary: concurrent Erlang I

- Concurrent Erlang I
 - Creating Processes
 - Message Passing
 - Receiving Messages
 - Data in Messages
- · Concurrent Erlang II



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Overview: concurrent Erlang II

- · Concurrent Erlang I
- · Concurrent Erlang II
 - Registered Processes
 - Timeouts
 - More on Processes
 - The Process Manager



Registered Processes

register(Alias, Pid)
 Alias ! Message

- · Registers the process Pid with the name Alias
- Any process can send a message to a registered process
- The BIF registered/0 returns all registered process names
- The BIF whereis(Alias) returns the Pid of the process with the name Alias.



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Registered Processes

```
Pid1 echo
{Pid1,hello}

echo ! {self(), hello} receive {From, Msg} -> ... end

go() -> register(echo, spawn(echo, loop, [])).

loop() -> receive
{From, Msg} -> From ! {self(), Msg}, loop();
stop -> true end.
```

Message Passing

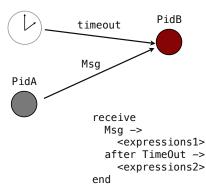


 Sending messages to nonexisting registered processes causes the calling process to terminate with a badarg

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Timeouts



- If the message Msg is received within the time TimeOut, <expressions1> will be executed
- If not, <expressions2> will be executed
- TimeOut is an integer denoting the time in milliseconds or the atom infinity



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Timeouts

```
read(Key) ->
   flush(),
   db ! {self(), {read, Key}},
   receive
      {read, R} ->
       {ok, R};
      {error, Reason} ->
       {error, timeout}
   end.
```

- If the server takes more than a second to handle the request, a timeout is generated
- Do not forget to handle messages received after a timeout

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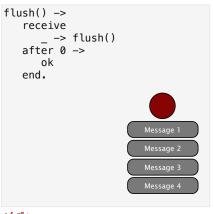
Timeouts

- send_after(T, What)
 sends the message What
 to the current process
 after T milliseconds
- The sleep(T) function will suspend the calling process for T milliseconds

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Timeouts



• flush() will clear the mailbox from all messages, stopping when it is empty.

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More on Processes: definitions

Process

A concurrent activity. The system may have many concurrent processes executing at the same time

Message

A method of communication and sharing data between processes

Timeout

A mechanism for waiting for a given period of time for an incoming message



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More on Processes: definitions

Registered Processes

Processes which have been given a name with BIFs such as register/2.

Termination

A process is said to terminate normally when it has no more code to execute.

It terminates abnormally if a run time error occurs or if someone makes it exit with a non-normal reason.



More on Processes: process skeleton start(Args) -> spawn(server, init, [Args]) Start init(Args) -> State = initialize_state(Args), loop(State). loop(State) -> Stop loop receive {handle, Msg} -> NewState = handle(Msg, State), loop(NewState); stop -> terminate(State) terminate(State) -> clean_up(State). © 1999-2011 Erlang Solutions Ltd. 25

The Process Manager

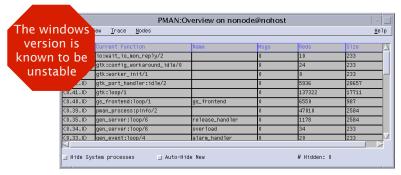
- Used to inspect the state of processes in a local or distributed Erlang system
- · Trace output for messages sent & received
- Trace output for process events such as spawn, exit and link
- · Trace output for BIF and function calls



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The Process Manager

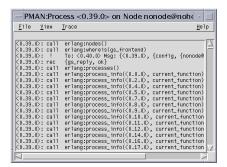


• pman:start()



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Process Manager: processes



• Prints the trace messages and process state



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Process Manager: options



- · Pick what trace messages you want to view
- Pick the inheritance level when spawning



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