

## Programming Assignment 1 Report

COSC 581

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### Implementation:

My quicksort is a basic implementation using 3 arrays to construct the sorted list recursively [quicksort(less) + pivot + quicksort(greater)] using my implementation of median of medians pivot selection method. After being divided into subfiles of size  $r$ , I select my pivot by looping through the subfile and removing the largest and smallest element, until one is remaining (the median). I call this function the `gulag_median()`.

### R Values:

I chose  $r$  values of 3, 4, 5, 6, 7, 9, 11, 25, and 26. As discussed in class, the odd values in  $r > 3$  provide median of medians a linear time, while all even values provide a polynomial runtime. I also included values of 25 and 26 to test the impact of a larger subfile size (which should result in a slightly slower runtime with a poor pivot selection method like the `gulag_median()`).

### Methodology:

To test my quicksort, I created a program to generate pseudo-random values of a user-specified size in the same format as specified in the writeup. I compared the quicksort against the built in sorted command to ensure correctness. I measured runtime with the python time module for all  $r$  values for all file sizes to provide the following data.

### Results:

Running on my local machine, I was able to only reasonably perform tests on up to file sizes of 100 million integers. The results mostly line up with the theoretical expectations, with a few outliers probably due to system variance. Most importantly,  $r=5$  appears near the bottom and  $r=26$  appears at the top as expected. In order to see polynomial runtime begin to really take over, the test cases will need to be much larger.

