

Zachary Neeley

2/20/19

Lab 4

Task 3:

(Base Cases)

This case is used when $n(n+1)/2$ which will return the value without doing any recursion.

The base case is 0. For this case we know that the formula will be $0(0+1)/2$ which will not work returning a zero.

(General Cases)

With the general case we consider that $n(n+1)/2$ will always return the same value as the recursive call of $n + \text{recursiveCall}(n-1)/2$. The smallest version of n to consider in the general cases are 1. Which this case we know that $1(1-1)/2$ which would be $0/2$. Knowing this we can see that the with the next value of 2 we know that every lower value has already been tested. Because the two value will always return the same answer we know that the proof is correct.

Task 4:

(Base Cases)

The base case is when n is equal to 0 than the base case will return 0. Once this case is reached the recursion will stop.

(General Cases)

When assuming the base case is 0 we can see that $E1, n \mid (2i - 1)$ is equivalent to $(2n-1) + E_i = 1, n-1 \mid (2i-1)$. Knowing this we can know that $n = 1$ so we will return n . Knowing this we can move from 1 to any other number to use induction hypotheses. The program given which is the recursion of $n-1$, with n

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getting smaller by 1 every go around. This can be written as $n * n$ or n^2 . Knowing this we can prove that n^2 will always return the same value as the recursive call. The proof is complete.

Task 5)

Base Case)

The base cases are -1 and 0 in both the equations given. Only if negative numbers are introduced into the equation. If not then the base case will be 0!=.

General Cases)

We can assume when using induction hypothesis, that if 0 and -1 are passed to the equation then the base cases will be used. We will be looking for any values greater than 0 and less than -1. If this is the case then we can use $(E_i = 1, n \mid i(i+1))$ and the equivalent equation of $(E_i = 1, n \mid i^2)$ which having the lower and upper bound of infinity, but not including -1 and 0.

Since we are given the need equations then we can prove the case by plugging into our equations. If $n = 1$. Our first equation will give us $1+1$ and the other equation which equals 2. We can prove $n = 2$ assuming that $n = 1$ is also true using induction hypothesis. Plugging 2 into the first equation you will get 6 out and which is also the same for the second equation. Now this we can check -2 which will give us the proper output for both the equations. The proof is complete.

Task 6)

Images on next page.

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$$\text{Task 6)} \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}$$

Task 5)

$$\text{Base Case } \sum_{i=1}^n i = \frac{n(n+1)}{2}$$

$$\left(\sum_{i=1}^n i^2 \right) + \left(\sum_{i=1}^n i \right) = \frac{n(n+1)}{2} + \frac{n(n+1)(2n+1)}{6}$$

$$= \frac{3(n+1)}{6} + \frac{n(n+1)(2n+1)}{6}$$

$$= \frac{3n+3n+2n^3+3n^2}{6}$$

$$= \frac{2n^3+3n^2+4n}{6} = \frac{n^3+3n^2+2n}{3}$$

$$= \frac{n(n+1)(n+2)}{3}$$