

CSE 2017 Data Structures and Lab

Lecture #3: Data Abstraction and List

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Different Views of Data

1. Data abstraction and encapsulation
2. Data structure
3. Abstract data type operator categories

1. Data Abstraction

- Separation of a data type's logical properties from its implementation.

LOGICAL PROPERTIES

What are the possible values?

What operations will be needed?

IMPLEMENTATION

How can this be done in C++?

How can data types be used?

Data Encapsulation

- is the separation of the representation of data from the applications that use the data at a logical level; a programming language feature that **enforces information hiding**.

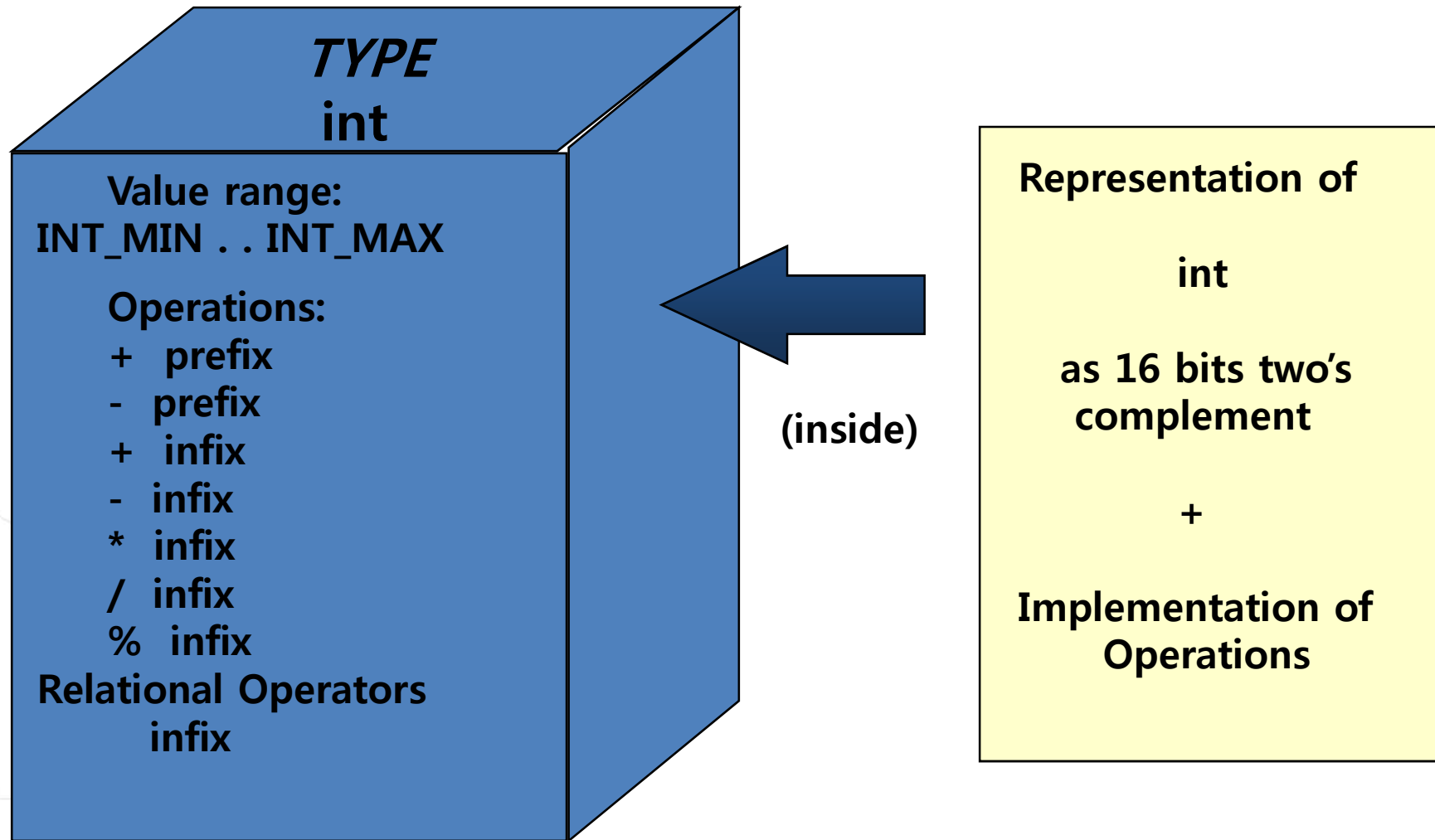
APPLICATION

```
int y;  
  
y = 25;
```

REPRESENTATION

```
0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1
```

Encapsulated C++ Data Type int



Abstract Data Type (ADT)

- A data type whose properties (domain and operations) are specified independently of any particular implementation.

List as abstract data type

List

- empty list has size 0
- insert
- remove
- Count
- Read/modify element at a position
- Specify data-type

A

A[0]	A[1]
2	4	6	7	9					...

int A[MAXSIZE];
int end = -1;
insert(2)

2. Data Structures

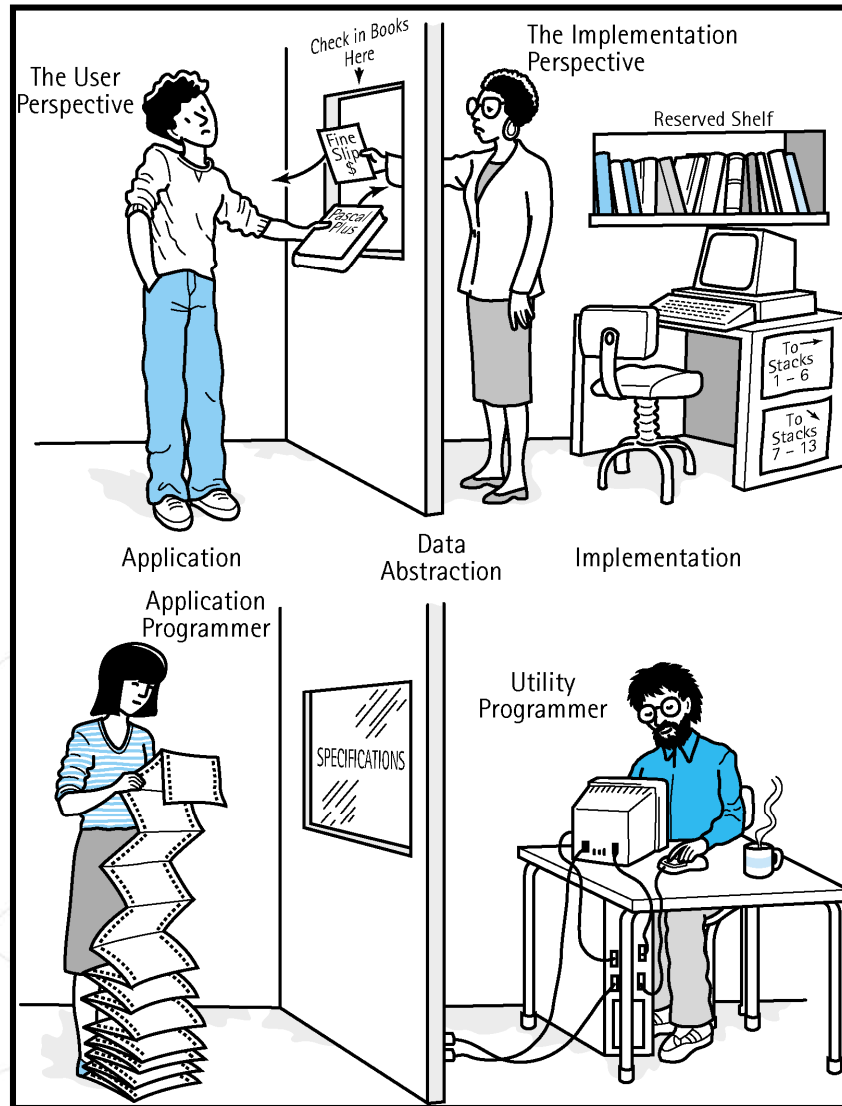
Defined by

- the logical arrangement of **data elements**
- the **set of operations** we need to access the elements

Data from 3 different levels

- *Application (or user) level*: modeling real-life data in a specific context.
- *Logical (or ADT) level*: abstract view of the domain and operations. WHAT
- *Implementation level*: specific representation of the structure to hold the data items, and the coding for operations. HOW

Logical Level



Viewing a library from 3 different levels

- *Application (or user) level:* Library of Congress, or Baltimore County Public Library.
- *Logical (or ADT) level:* domain is a collection of books; operations include: check book out, check book in, pay fine, reserve a book.
- *Implementation level:* representation of the structure to hold the “books”, and the coding for operations.

3. ADT Operator Categories

- **Constructor** -- creates a new instance (object) of an ADT.
- **Transformer** -- changes the state of one or more of the data values of an instance.
- **Observer** -- allows us to observe the state of one or more of the data values without changing them.
- **Iterator** -- allows us to process all the components in a data structure sequentially.

Lists

Lists in every-day life:

- Grocery list
- Laundry list
- To-do list
- Invitation list
- ...

CSE2017 Class List

Name	SS	E-mail	Quiz1	Lab1	Lab2	...	Midterm ...
XXX	999-99-9999	xx@xxx	20	10	10		
YYY	000-00-0000	yy@yyy	25	9	10		
ZZZ	111-11-1111	zz@zzz	24	10	8		
...							
...							
...							
UUU	222-22-222	uu@uuu	23	9	10		

What is a List?

- A list is a homogeneous collection of elements, with a **linear relationship** between elements.
- That is, each list element (except the first) has a **unique predecessor**, and each element (except the last) has a **unique successor**.

Sorted and Unsorted Lists

UNSORTED LIST

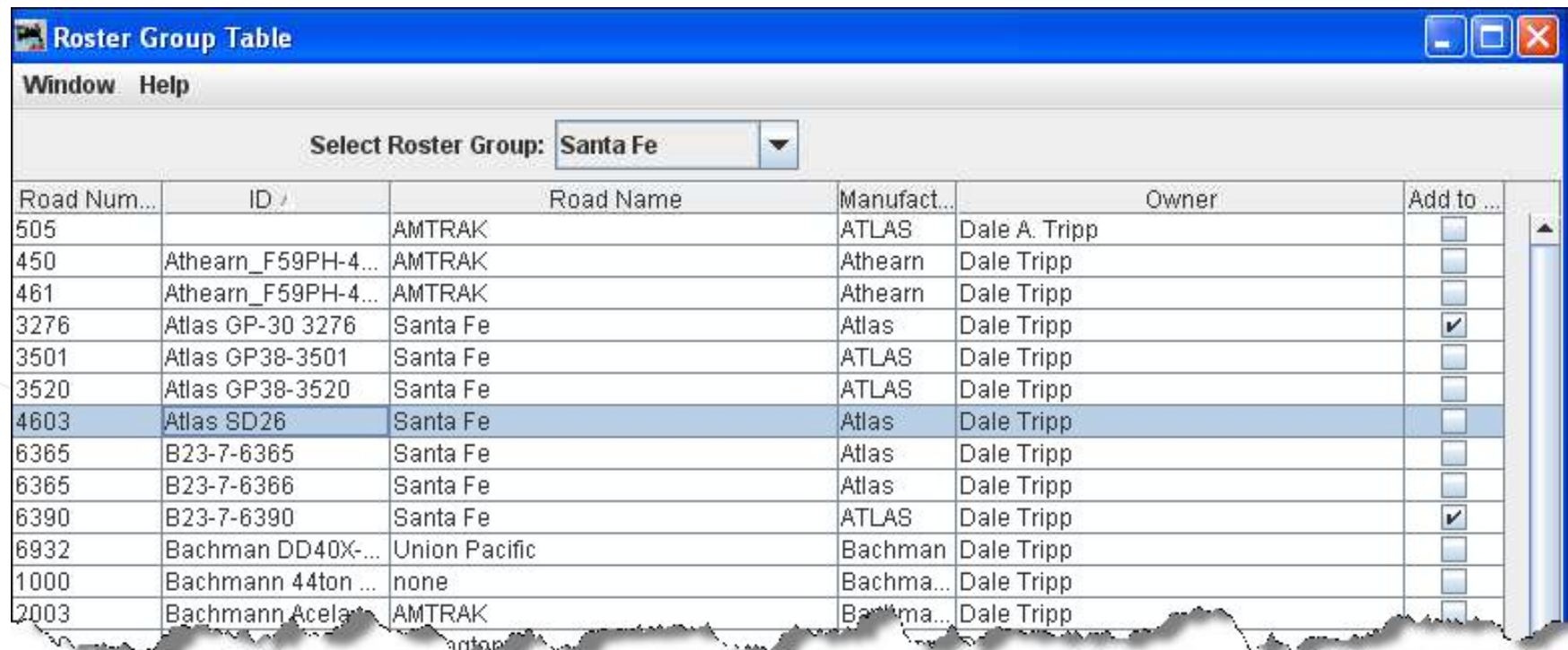
Elements are placed into the list in no particular order.

SORTED LIST

List elements are in an order that is sorted in some way -- either numerically or alphabetically by the elements themselves, or by a component of the element (called a **KEY** member) .

Key:

- A member of a record (struct or class) whose value is used to determine the logical and/or physical order of the items in a list



Roster Group Table

Window Help

Select Roster Group: Santa Fe

Road Num...	ID	Road Name	Manufact...	Owner	Add to ...
505		AMTRAK	ATLAS	Dale A. Tripp	<input type="checkbox"/>
450	Athearn_F59PH-4...	AMTRAK	Athearn	Dale Tripp	<input type="checkbox"/>
461	Athearn_F59PH-4...	AMTRAK	Athearn	Dale Tripp	<input type="checkbox"/>
3276	Atlas GP-30 3276	Santa Fe	Atlas	Dale Tripp	<input checked="" type="checkbox"/>
3501	Atlas GP38-3501	Santa Fe	ATLAS	Dale Tripp	<input type="checkbox"/>
3520	Atlas GP38-3520	Santa Fe	ATLAS	Dale Tripp	<input type="checkbox"/>
4603	Atlas SD26	Santa Fe	Atlas	Dale Tripp	<input type="checkbox"/>
6365	B23-7-6365	Santa Fe	Atlas	Dale Tripp	<input type="checkbox"/>
6365	B23-7-6366	Santa Fe	Atlas	Dale Tripp	<input type="checkbox"/>
6390	B23-7-6390	Santa Fe	ATLAS	Dale Tripp	<input checked="" type="checkbox"/>
6932	Bachman DD40X-...	Union Pacific	Bachman	Dale Tripp	<input type="checkbox"/>
1000	Bachmann 44ton ...	none	Bachma...	Dale Tripp	<input type="checkbox"/>
2003	Bachmann Acela...	AMTRAK	Bachma...	Dale Tripp	<input type="checkbox"/>

ADT Unsorted List

- Application Level



- Logical Level

Add a name on the list
Delete a name on the list
Print out all names
Retrieve a name

Sort names

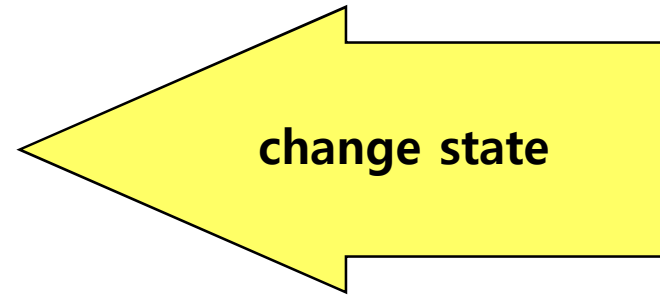
- Implementation Level

0	b	size = 5 cursor = 2
1	r	
2	y	
3	c	
4	u	
5		
6		
7		

ADT Unsorted List Operations

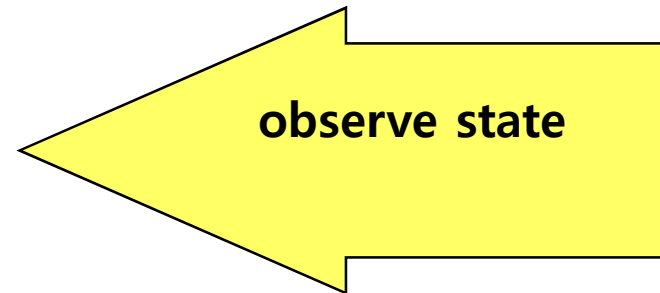
Transformers

- MakeEmpty
- InsertItem
- DeleteItem



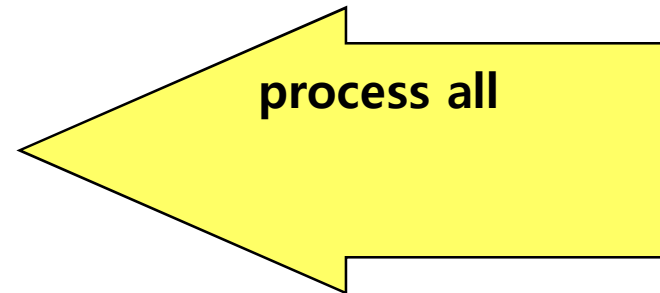
Observers

- IsFull
- LengthIs
- RetrieveItem



Iterators

- ResetList
- GetNextItem



What is a Generic Data Type?

- A **generic data type** is a type for which the operations are defined but the types of the items being manipulated are not defined.
- One way to simulate such a type for our UnsortedList ADT is via a user-defined class `ItemType` with member function `CompareTo` having enumerated type value `LESS`, `GREATER`, or `EQUAL`.

// SPECIFICATION FILE

(unsorted.h)

#include "ItemType.h"

class UnsortedType {
public :

// declares a class data type

// 8 public member functions

void MakeEmpty () ;

bool IsFull () **const ;**

int LengthIs () **const ;** *// returns length of list*

void RetrieveItem (ItemType& item, bool& found) ;

void InsertItem (ItemType item) ;

void DeleteItem (ItemType item) ;

void ResetList () ;

void GetNextItem (ItemType& item) ;

private :

// 3 private data members

int length ;

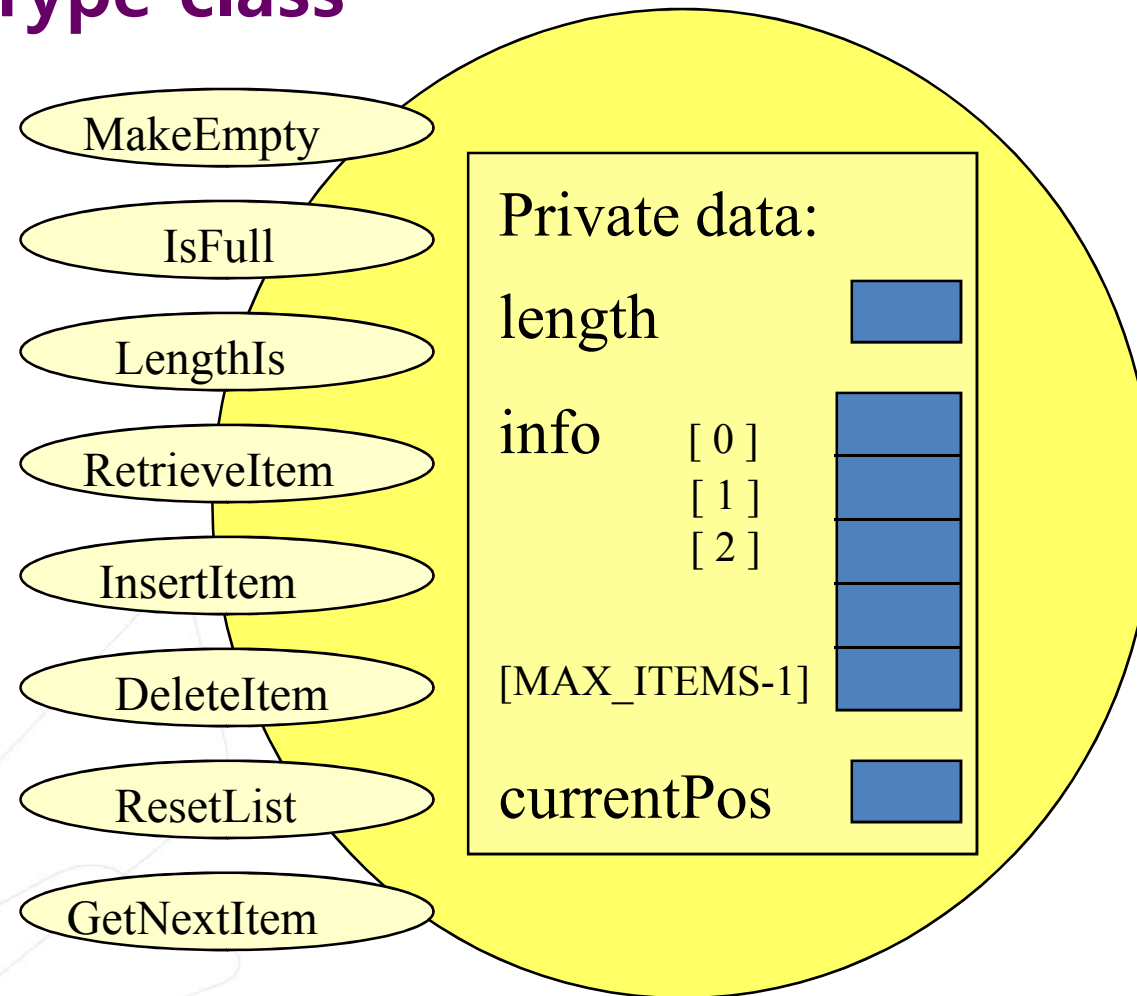
ItemType info[MAX_ITEMS] ;

int currentPos ;

};

Class Interface Diagram

UnsortedType class



// IMPLEMENTATION FILE ARRAY-BASED LIST (unsorted.cpp)

#include "itemtype.h"

void UnsortedType::MakeEmpty () {

// Pre: None.

// Post:List is empty.

length = 0 ;

}

void UnsortedType::InsertItem (ItemType item) {

// Pre: List has been initialized. List is not full. item is not in list.

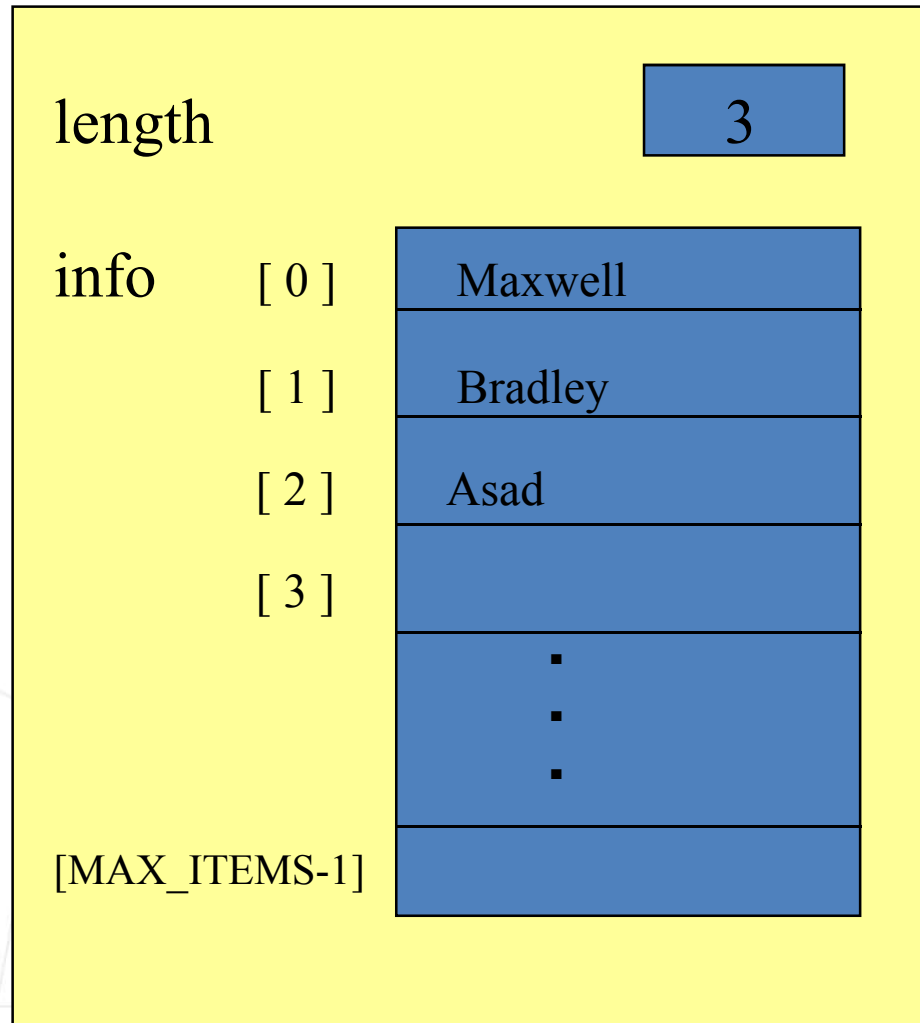
// Post: item is in the list.

info[length] = item ;

length++ ;

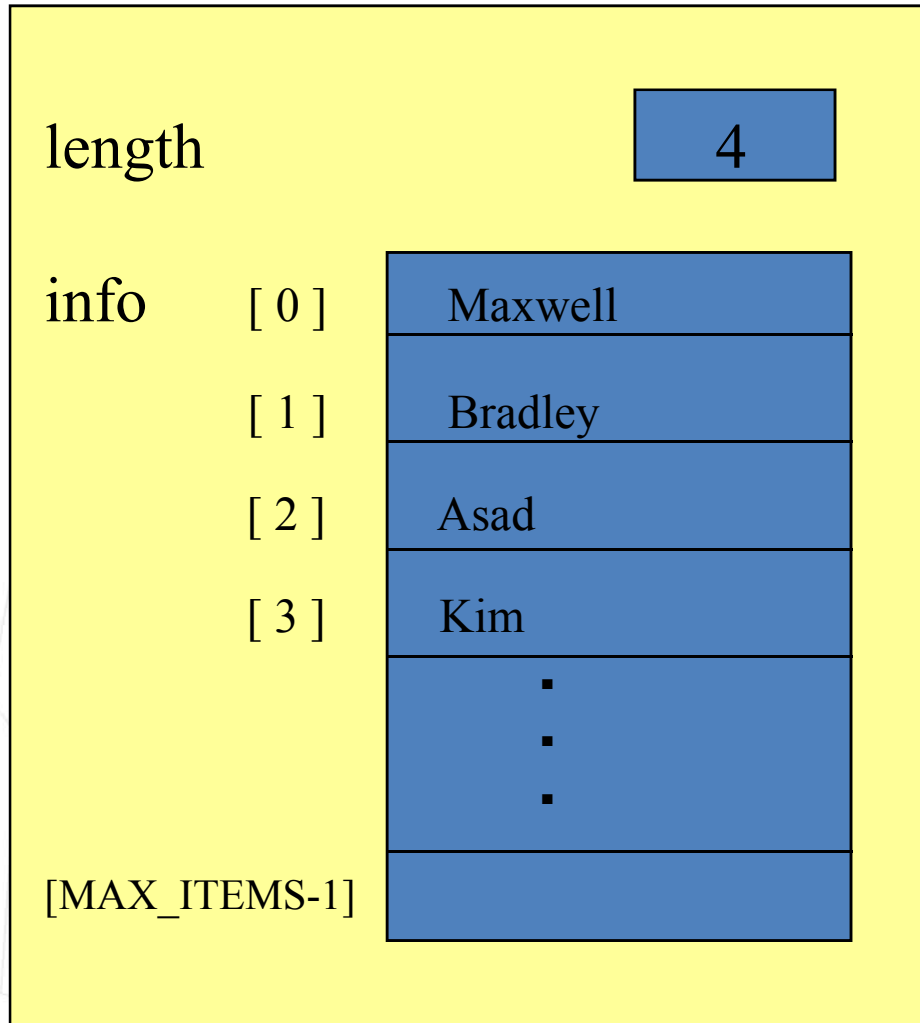
}

Before Inserting Kim into an Unsorted List



The item will be placed into the length location, and length will be incremented.

After Inserting **Kim** into an Unsorted List




```
void UnsortedType::LengthIs ( ) const {  
    // Pre: List has been inititalized.  
    // Post:Function value == ( number of elements in list ).  
    return length ;  
}
```

```
bool UnsortedType::IsFull ( ) const {  
    // Pre: List has been initialized.  
    // Post:Function value == ( list is full ).  
    return ( length == MAX_ITEMS ) ;  
}
```

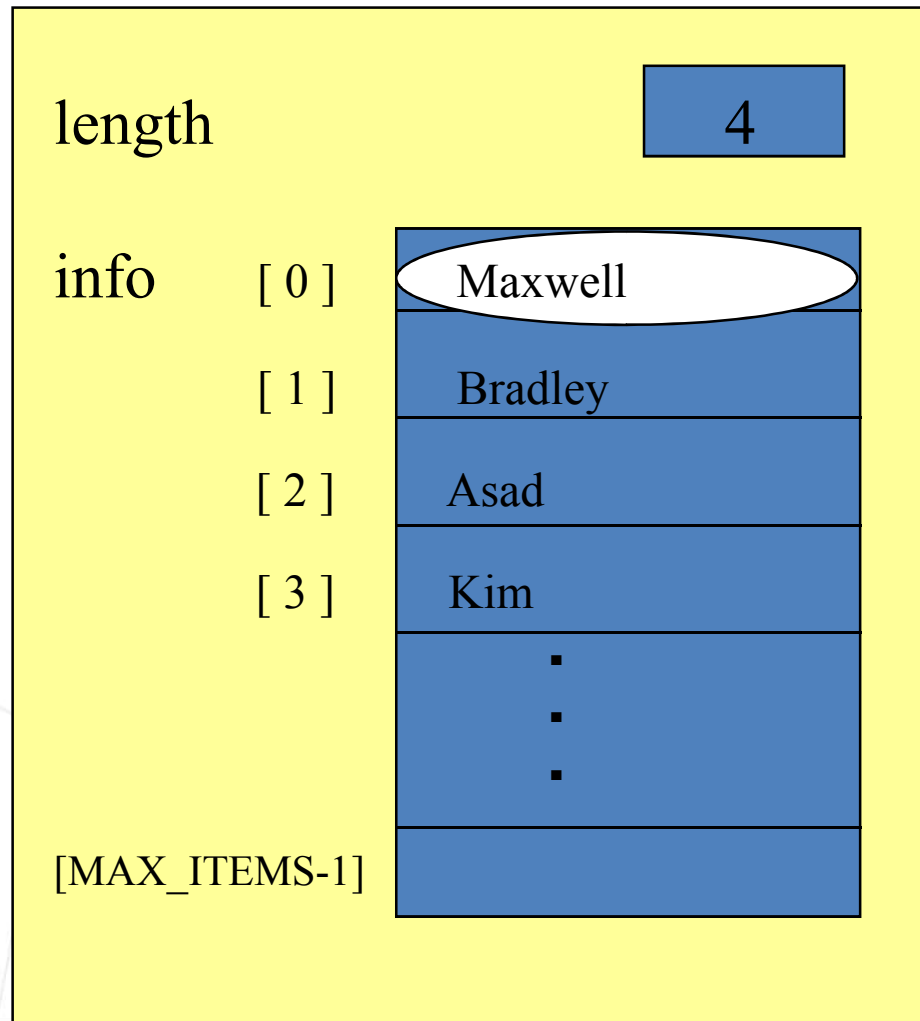
```

void UnsortedType::RetrieveItem ( ItemType& item,  bool& found ) {
    // Pre: Key member of item is initialized.
    // Post: If found, item's key matches an element's key in the list and a copy
    // of that element has been stored in item; otherwise, item is unchanged.
    bool moreToSearch ;
    int   location = 0 ;

    found = false ;
    moreToSearch = ( location < length ) ;
    while ( moreToSearch && !found ) {
        switch ( item.ComparedTo( info[location] ) ) {
            case LESS      :
            case GREATER   : location++ ;
                           moreToSearch = ( location < length ) ;
            case EQUAL     : found = true ;
                           item = info[ location ] ;
                           break ;
        }
    }
}

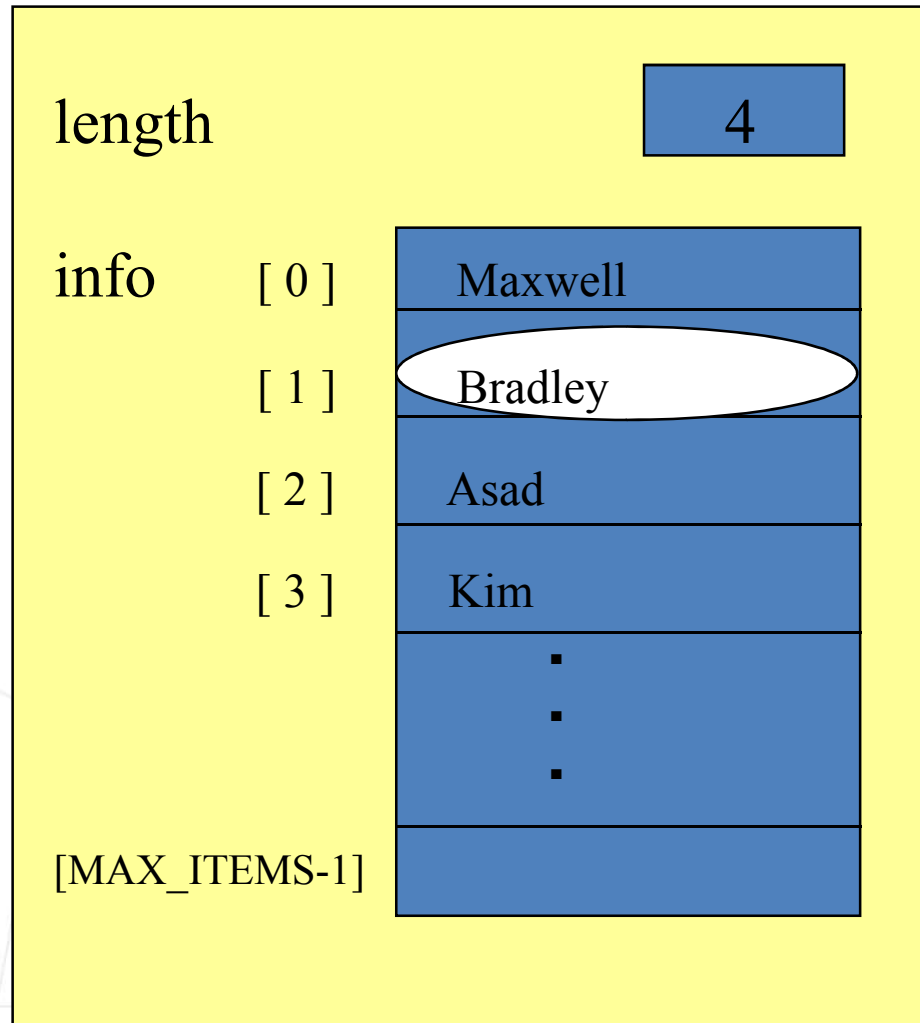
```

Retrieving Park from an Unsorted List



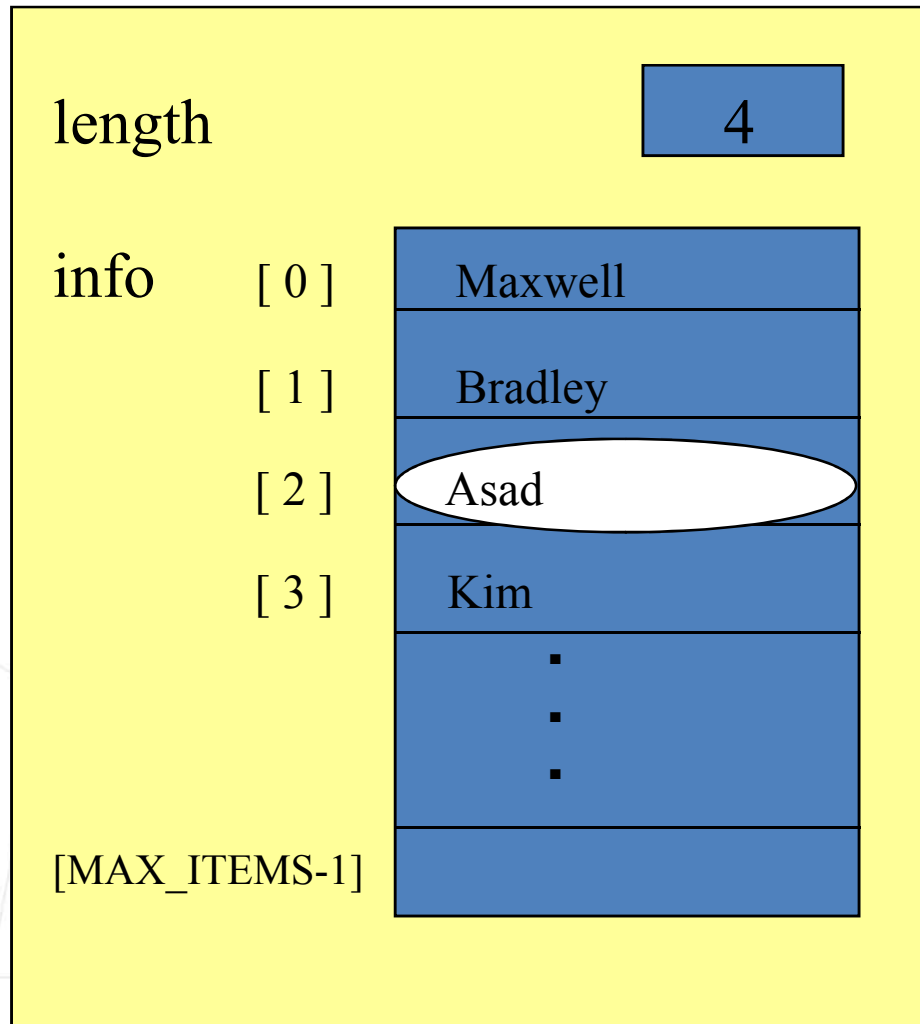
moreToSearch: true
found: false
location: 0

Retrieving **Park** from an Unsorted List



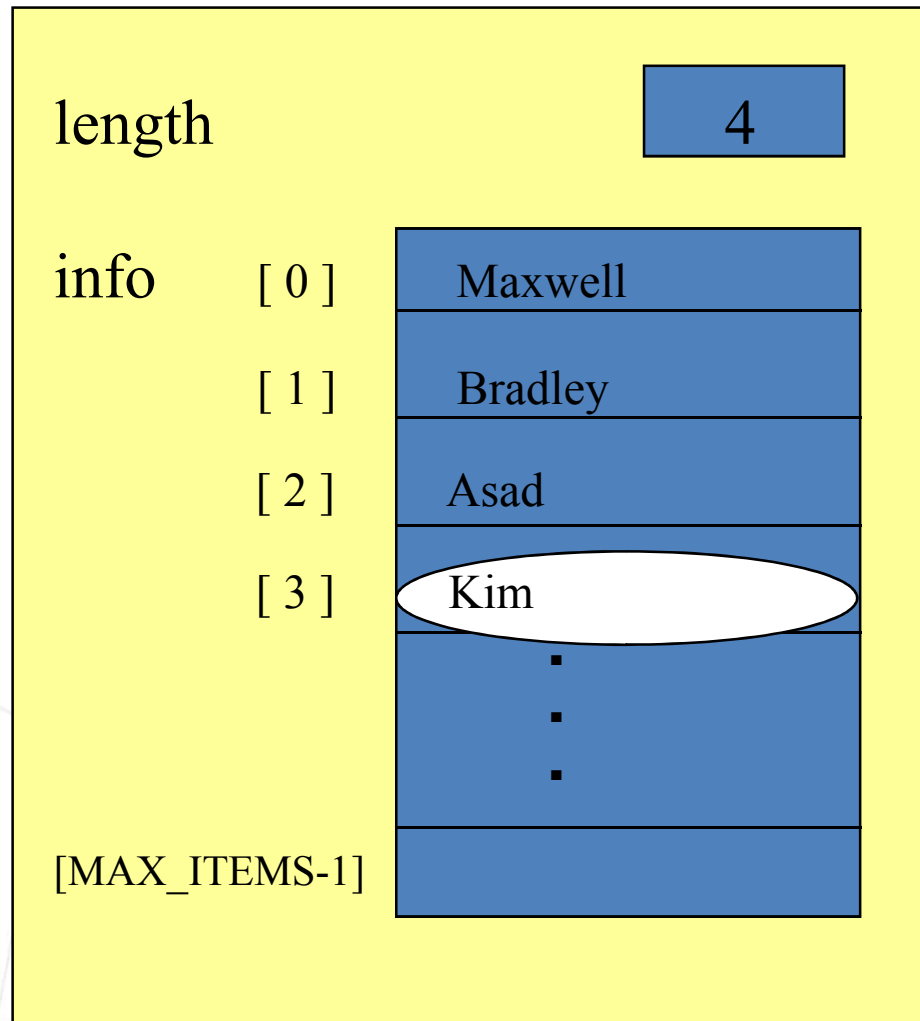
moreToSearch: true
found: false
location: 1

Retrieving Park from an Unsorted List



moreToSearch: true
found: false
location: 2

Retrieving Park from an Unsorted List



moreToSearch: true
found: false
location: 3

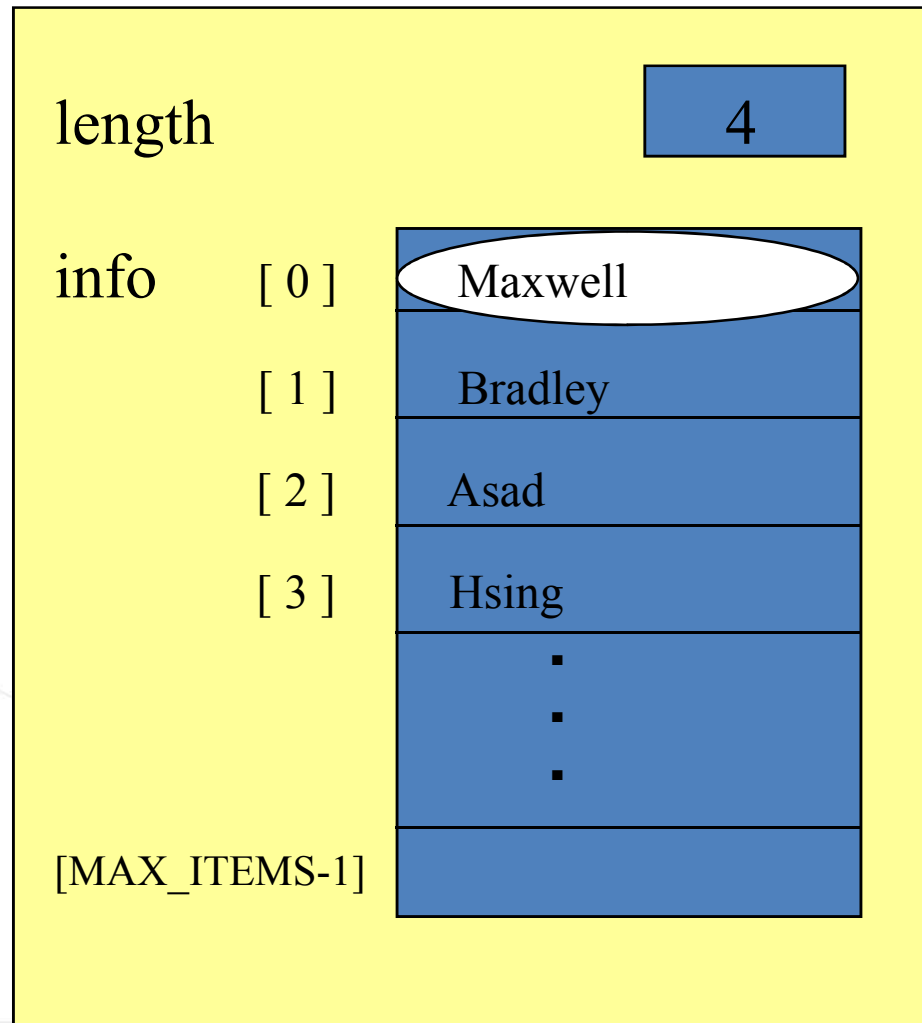
Retrieving **Park** from an Unsorted List

length		4
info	[0]	Maxwell
	[1]	Bradley
	[2]	Asad
	[3]	Kim
		▪
		▪
		▪
	[MAX_ITEMS-1]	

moreToSearch: false
found: false
location: 4

```
void UnsortedType::DeleteItem ( ItemType item ) {  
    // Pre: item's key has been initialized.  
    //      An element in the list has a key that matches item's.  
    // Post: No element in the list has a key that matches item's.  
    int location = 0 ;  
  
    while (item.ComparedTo (info [location] ) != EQUAL )  
        location++ ;  
  
    // move last element into position where item was located  
  
    info [location] = info [length - 1 ] ;  
    length-- ;  
}
```

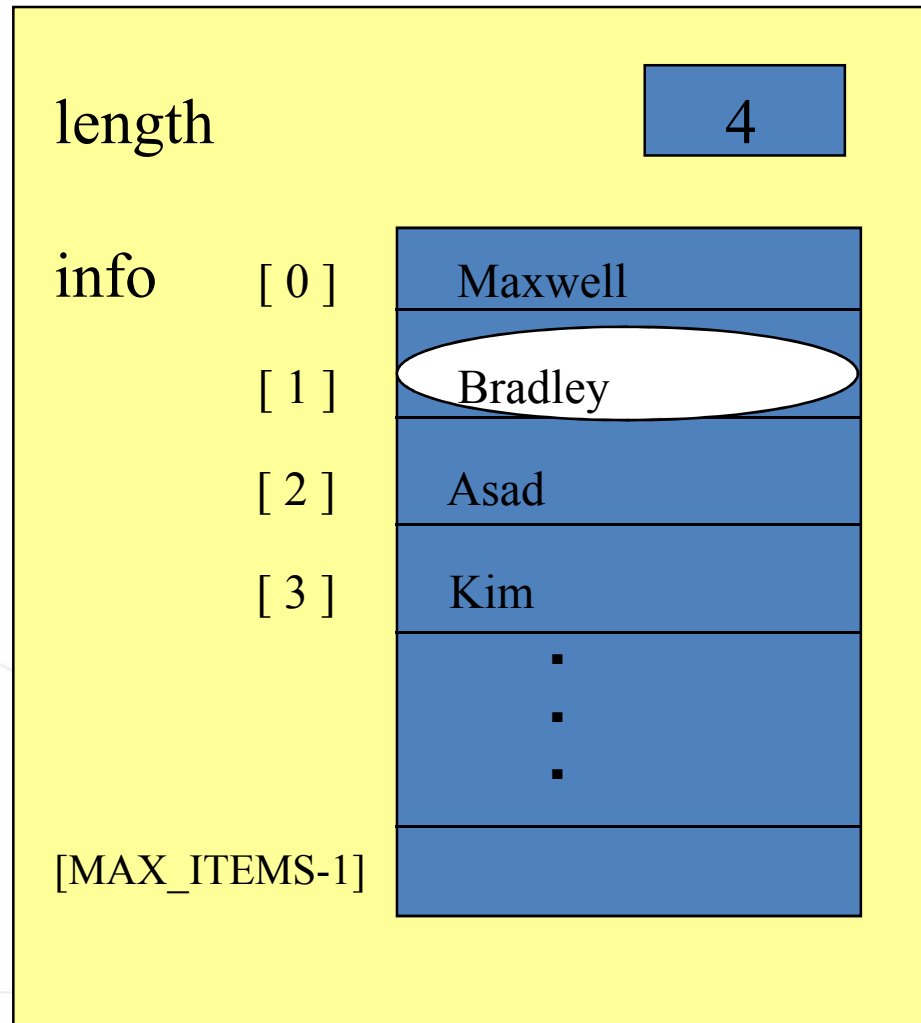

Deleting Bradley from an Unsorted List



location: 0

**Key Bradley has
not been matched.**

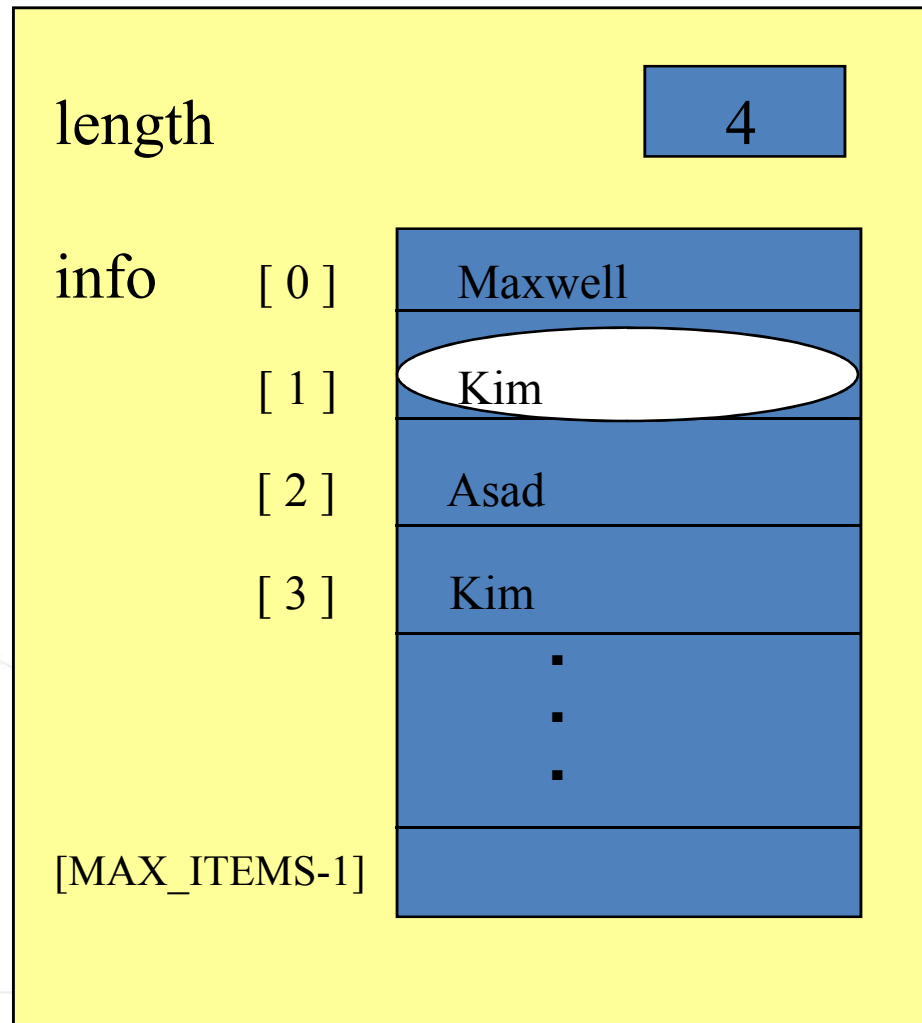
Deleting Bradley from an Unsorted List



location: 1

**Key Bradley has
been matched.**

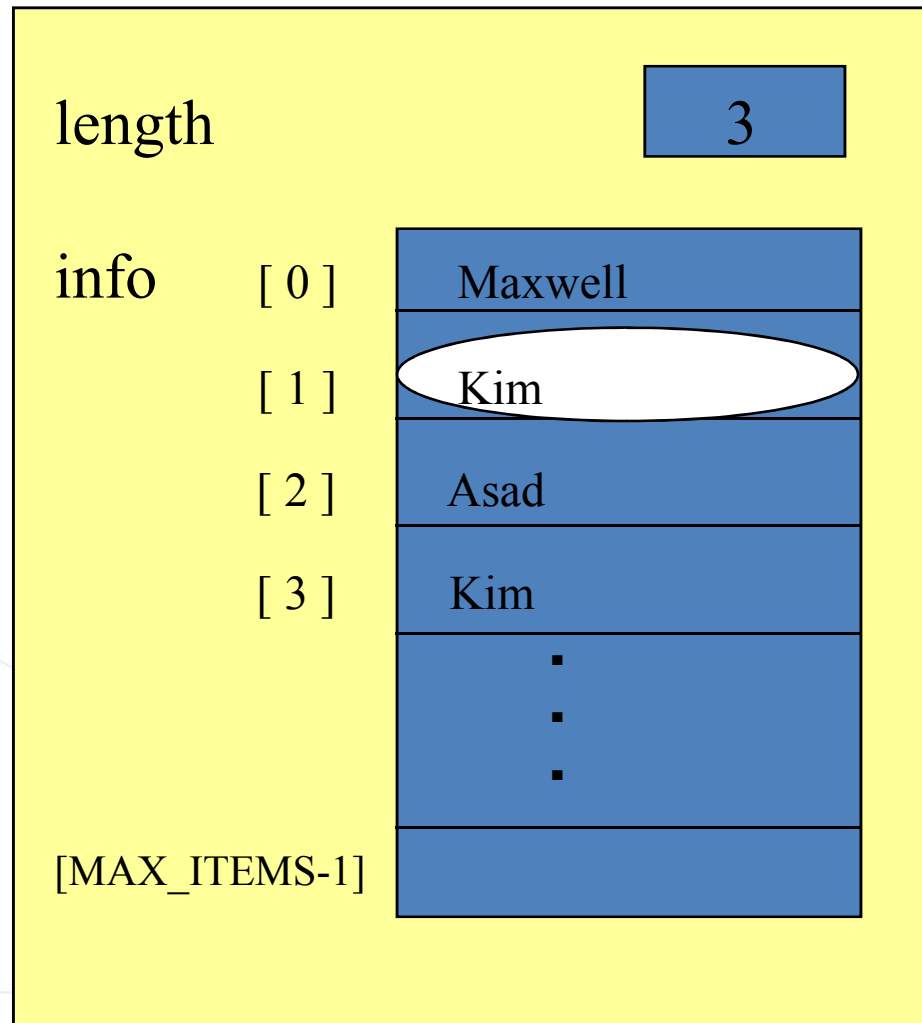
Deleting Bradley from an Unsorted List



location: 1

**Placed copy of
last list element
into the position
where the key Bradley
was before.**

Deleting Bradley from an Unsorted List



location: 1

Decrement length.

```
void UnsortedType::ResetList ( ) {  
    // Pre: List has been initialized.  
    // Post: Current position is prior to first element in list.  
    currentPos = -1 ;  
}
```

```
void UnsortedType::GetNextItem ( ItemType& item ) {  
    // Pre: List has been initialized. Current position is defined.  
    // Element at current position is not last in list.  
    // Post: Current position is updated to next position.  
    // item is a copy of element at current position.  
    currentPos++ ;  
    item = info [currentPos] ;  
}
```

Specifying class ItemType

```
// SPECIFICATION FILE ( itemtype.h )

const int MAX_ITEM = 5 ;
enum RelationType { LESS, EQUAL, GREATER } ;

class ItemType {                                // declares class data type
public :                                         // 3 public member
    functions
    RelationType ComparedTo ( ItemType ) const ;
    void Print ( ) const ;
    void Initialize ( int number ) ;

private :                                     // 1 private data member
    int value ;                               // could be any different
    type
};
```

// IMPLEMENTATION FILE (itemtype.cpp)

// Implementation depends on the data type of value.

#include "itemtype.h"

#include <iostream.h>

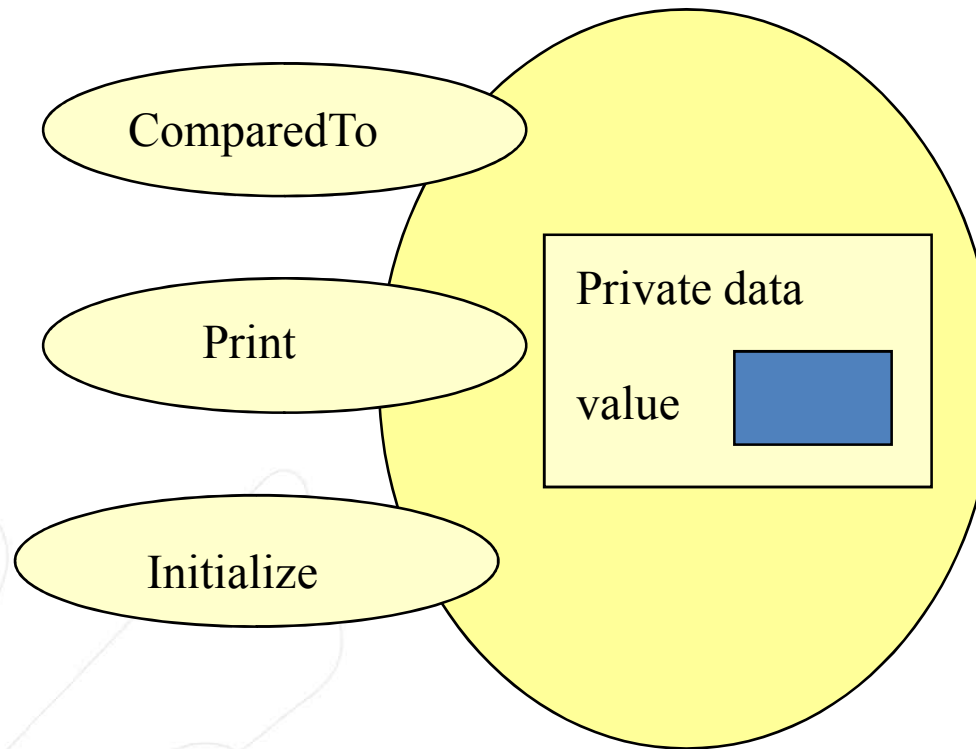
```
RelationType ComparedTo ( ItemType otherItem ) const {  
    if ( value < otherItem.value )  
        return LESS ;  
    else if ( value > otherItem.value )  
        return GREATER ;  
    else return EQUAL ;  
}
```

```
void Print ( ) const {  
    cout << value << endl ;  
}
```

```
void Initialize ( int number ) {  
    value = number ;  
}
```

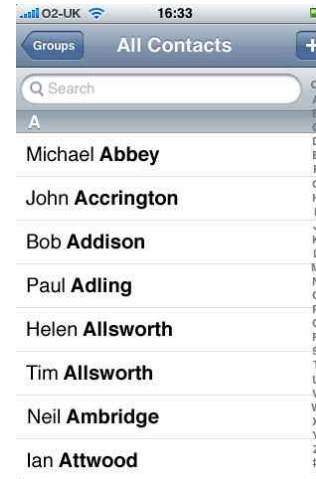
ItemType Class Interface Diagram

- class ItemType



ADT Sorted List

- Application Level



- Logical Level

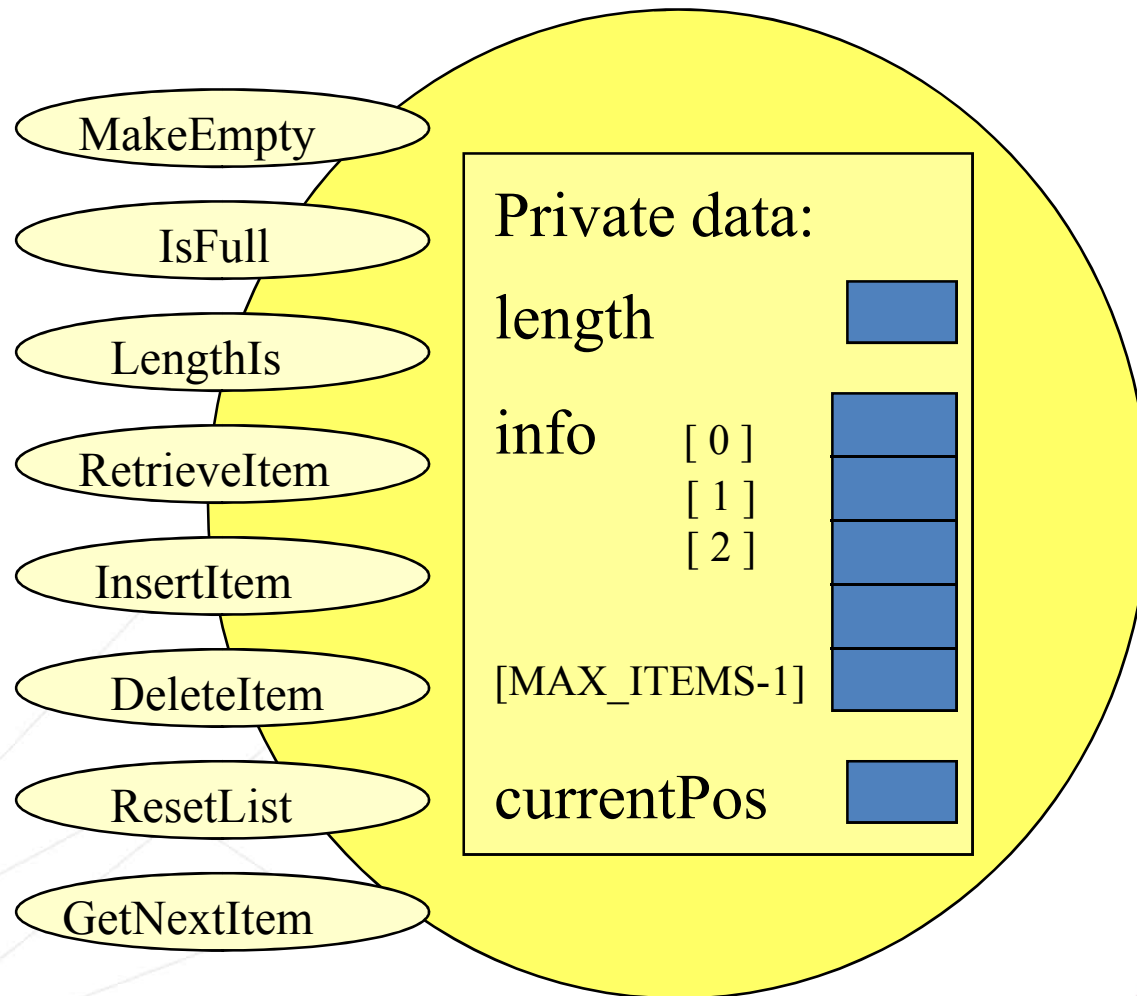
Add a name on the list
Delete a name on the list
Print out all names
Retrieve a name

- Implementation Level

0	a	size = 4 cursor = 2
1	b	
2	p	
3	r	
4		
5		
6		
7		

SortedType Class Interface Diagram

- SortedType class



Member functions

- Which member function specifications and implementations must change to ensure that any instance of the Sorted List ADT remains sorted at all times?

- **InsertItem**

- **DeleteItem**

InsertItem algorithm for SortedList ADT

- Find proper location for the new element in the sorted list.
- Create space for the new element by **moving down** all the list elements that will follow it.
- Put the new element in the list.
- Increment length.

Implementing SortedType member function InsertItem

// IMPLEMENTATION FILE

(sorted.cpp)

#include "itemtype.h" *// also must appear in client code*

void SortedType :: InsertItem (ItemType item) {

// Pre: List has been initialized. List is not full. item is not in list.

// List is sorted by key member using function ComparedTo.

// Post: item is in the list. List is still sorted.

.
. .
.

}

```

void SortedType :: InsertItem ( ItemType item ) {
    bool moreToSearch ;
    int    location = 0 ;
        // find proper location for new element
    moreToSearch = ( location < length ) ;
    while ( moreToSearch ) {
        switch ( item.ComparedTo( info[location] ) ) {
            case LESS      : moreToSearch = false ;
                           break ;
            case GREATER   : location++ ;
                           moreToSearch = ( location < length ) ;
                           break ;
        }
    } // make room for new element in sorted list
    for ( int index = length ; index > location ; index-- )
        info [ index ] = info [ index - 1 ] ;
    info [ location ] = item ;
    length++ ;
}

```

DeleteItem algorithm for SortedList ADT

- Find the location of the element to be deleted from the sorted list.
- Eliminate space occupied by the item being deleted by **moving up** all the list elements that follow it.
- Decrement length.

Implementing SortedType member function DeleteItem

```
// IMPLEMENTATION FILE    continued                (sorted.cpp)

void SortedType :: DeleteItem ( ItemType item )
// Pre: List has been initialized. Key member of item is initialized.
//      Exactly one element in list has a key matching item's key.
//      List is sorted by key member using function ComparedTo.
// Post: No item in list has key matching item's key.
//      List is still sorted.
{
    .
    .
    .

}
```

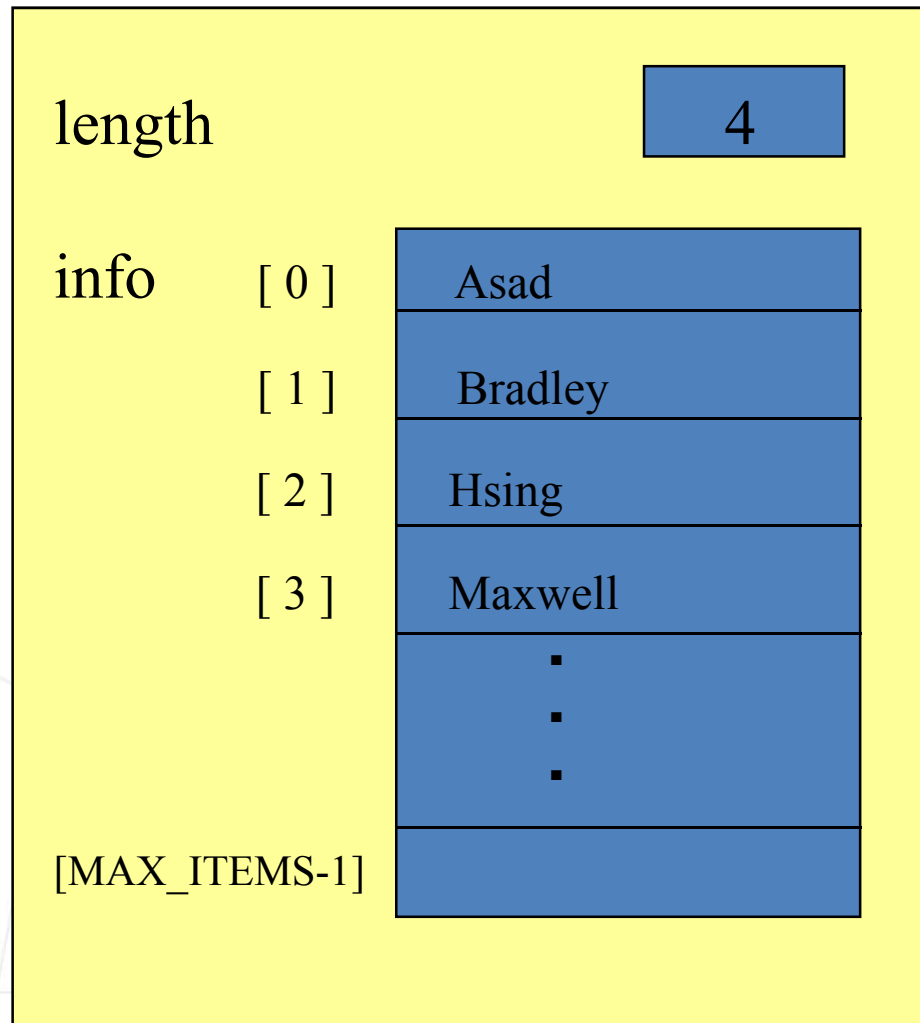


```
void SortedType :: DeleteItem ( ItemType item ) {  
    int    location = 0 ;  
    // find location of element to be deleted  
  
    while ( item.ComparedTo ( info[location] )  !=  EQUAL )  
        location++ ;  
  
    // move up elements that follow deleted item in sorted list  
  
    for ( int index = location + 1 ; index < location ; index++ )  
        info [ index - 1 ] = info [ index ] ;  
  
    length-- ;  
}
```

Improving member function RetrieveItem

- Recall that with the Unsorted List ADT we examined each list element beginning with info[0], until we either found a matching key, or we had examined all the elements in the Unsorted List.
- How can the searching algorithm be improved for Sorted List ADT?

Retrieving **Eric** from a Sorted List



The sequential search for Eric can stop when Hsing has been examined.

Binary Search in a Sorted List

- **Examines the element in the middle of the array.** Is it the sought item? If so, stop searching. Is the middle element too small? Then start looking in second half of array. Is the middle element too large? Then begin looking in first half of the array.
- **Repeat the process in the half of the list** that should be examined next.
- **Stop when item is found, or when there is nowhere else to look and item has not been found.**

```

void SortedType::RetrieveItem ( ItemType& item,  bool& found ) {
    // Pre: Key member of item is initialized.
    // Post: If found, item's key matches an element's key in the list and a
    //        copy
    //        of that element has been stored in item; otherwise, item is
    //        unchanged.
    int    midPoint ;
    int    first = 0;
           int      last = length - 1 ;
    bool moreToSearch = ( first <= last ) ;

    found = false ;
    while ( moreToSearch && !found )    {
        midPoint = ( first + last ) / 2 ; // INDEX OF MIDDLE ELEMENT
        switch ( item.ComparedTo( info [ midPoint ] ) )    {
            case LESS      : ... // LOOK IN FIRST HALF NEXT
            case GREATER   : ... // LOOK IN SECOND HALF NEXT
            case EQUAL      : ... // ITEM HAS BEEN FOUND
        }
    }
}
53
}

```

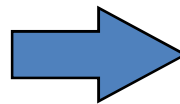
Trace of Binary Search

item = 45

15	26	38	57	62	78	84	91	108	119
----	----	----	----	----	----	----	----	-----	-----

info[0] [1] [2] [3] [4] [5] [6] [7] [8] [9]
first midPoint last

LESS

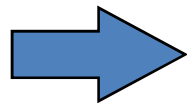


$\text{last} = \text{midPoint} - 1$

15	26	38	57	62	78	84	91	108	119
----	----	----	----	----	----	----	----	-----	-----

info[0] [1] [2] [3] [4] [5] [6] [7] [8] [9]
first midPoint last

GREATER



$\text{first} = \text{midPoint} + 1$

Trace continued

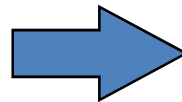
item = 45

15	26	38	57	62	78	84	91	108	119
---------------	---------------	----	----	---------------	---------------	---------------	---------------	----------------	----------------

info[0] [1] [2] [3] [4] [5] [6] [7] [8] [9]

first, last
midPoint

GREATER



first = midPoint + 1

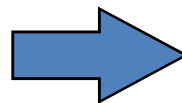
15	26	38	57	62	78	84	91	108	119
---------------	---------------	---------------	----	---------------	---------------	---------------	---------------	----------------	----------------

info[0] [1] [2] [3] [4] [5] [6] [7] [8] [9]

first,
midPoint,
last

55

LESS



last = midPoint - 1

Trace concludes

item = 45

15	26	38	57	62	78	84	91	108	119
info[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]
		last	first						

first > last → found = false


```
void SortedType::RetrieveItem ( ItemType& item, bool& found ) {  
    // ASSUMES info ARRAY SORTED IN ASCENDING ORDER
```

```
    int    midPoint ;
```

```
    int    first = 0;
```

```
        int    last = length - 1 ;
```

```
    bool moreToSearch = ( first <= last ) ;
```

```
    found = false ;
```

```
    while ( moreToSearch && !found ) {
```

```
        midPoint = ( first + last ) / 2 ;
```

```
        switch ( item.ComparedTo( info [ midPoint ] ) ) {
```

```
            case LESS      :      last = midPoint - 1 ;  
                                moreToSearch = ( first <= last ) ;  
                                break ;
```

```
            case GREATER   :      first = midPoint + 1 ;  
                                moreToSearch = ( first <= last ) ;  
                                break ;
```

```
            case EQUAL     :      found = true ;  
                                item = info[ midPoint ] ;  
                                break ;
```

```
        }
```

```
    }  
57 }
```

```
}
```

Comparison of Algorithms

- Order of Magnitude of a Function
- The **order of magnitude**, or **Big-O notation**, of a function expresses the computing time of a problem as the term in a function that increases most rapidly relative to the size of a problem.

Names of Orders of Magnitude

N: size of the problem

$O(1)$ bounded (by a constant) time

$O(\log_2 N)$ logarithmic time

$O(N)$ linear time

$O(N \cdot \log_2 N)$ $N \cdot \log_2 N$ time

$O(N^2)$ quadratic time

$O(2^N)$ exponential time

N **$\log_2 N$** **$N * \log_2 N$** **N^2** **2^N**

1	0	0	1	2
2	1	2	4	4
4	2	8	16	16
8	3	24	64	256
16	4	64	256	65,536
32	5	160	1024	4,294,967,296
64	6	384	4096	
128	7	896	16,384	

Big-O Comparison of List Operations

OPERATION	UnsortedList	SortedList
RetrieveItem	$O(N)$	$O(N)$ linear search $O(\log_2 N)$ binary search
InsertItem		
Find	$O(1)$	$O(N)$ search
Put	$O(1)$	$O(N)$ moving down
Combined	$O(1)$	$O(N)$
DeleteItem		
Find	$O(N)$	$O(N)$ search
Put	$O(1)$ swap	$O(N)$ moving up
Combined	$O(N)$	$O(N)$