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Deterministic Finite Automata (Part 2)

Lecture 07 Day 07/31

CS 154
Formal Languages and Computability
Spring 2018

Agenda of Day 07

- Summary of Lecture 06
- Quiz 2
- Lecture 07: Teaching ...
 - Deterministic Finite Automata (Part 2)
- Prepare Your Development Environment

Summary of Lecture 06: We learned ...

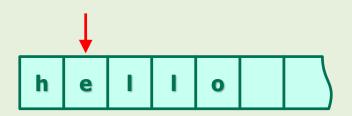
Deterministic Finite Automata

- We started constructing machines that understand formal languages.
- We'll construct several classes of machines in this course.
- Each class has different power.
- DFAs are the simplest ones.

Building blocks of DFAs:

Input tape, Control unit, Output

Input Tape



- The input tape is read-only.
- The read-head moves from left-to-right.
 - We cannot move the head back.
- Reading, consuming and scanning have the same meaning.

Any question?

Summary of Lecture 06: We learned ...

Control Unit

- Represented by transition graph.
- The number of states is finite.
- There is only one initial state.
- There can be zero or more accepting state (aka final state).

Output

- The output has two messages:
 - Accept (aka: understood, recognized, Yes)
 - Reject (aka: not understood, not recognized, No)

- DFAs configuration (aka snapshot) is the combination of the following data:
 - Timeframe of the clock
 - Input string
 - Position of the read-head
 - Current state of the transition graph
- Starting configuration is ...

Any question?

Summary of Lecture 06: We learned ...

When DFAs halt

When all input symbols are consumed.

$$h \leftrightarrow c$$

How DFAs accept a string w

- Three conditions should be satisfied:
 - The DFA halts. ≡ h
 - All symbols of w are consumed. ≡ c
 - The DFA is in an accepting state. ≡ f

$$(h \land c \land f) \leftrightarrow a$$

- For DFAs, h and c are equivalent.
- So, accepting a string condition is:

$$(c \land f) \leftrightarrow a$$

How DFAs reject a string w

 We need to negate the previous statement:

$$\sim$$
 (c \land f) \leftrightarrow \sim a
$$\equiv$$
 (\sim c \lor \sim f) \leftrightarrow \sim a

- Translation:
 - At least one symbol is NOT consumed.

OR

 The DFA is NOT in an accepting state.

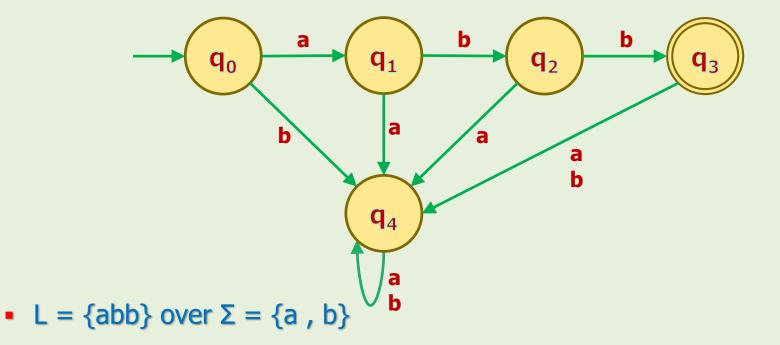
Any question?

NAME	Alan M. Turing		
SUBJECT	CS 154	TEST NO.	2
DATE	02/15/2018	PERIOD	1,2,3



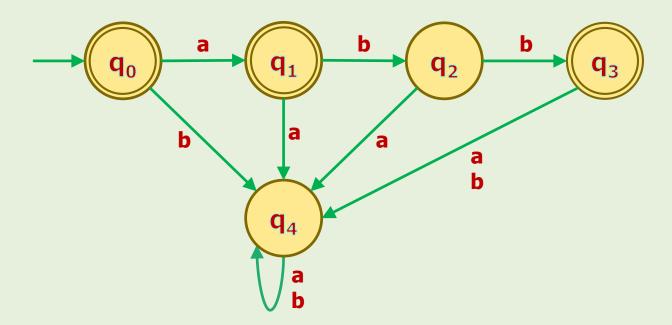
Quiz 2 Use Scantron

- What language does the following DFA accept?
- What is Σ?



Example 10

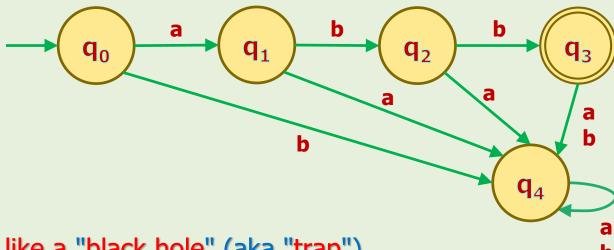
• What language does the following DFA accept?



• L = $\{\lambda, a, abb\}$ over $\Sigma = \{a, b\}$

Notes

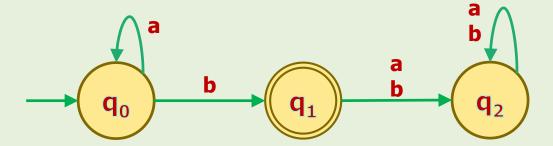
- We don't need to show the "output" and the "clock" anymore.
- 2. The role of q₄ in the previous examples:



- q₄ acts like a "black hole" (aka "trap")
 because if the machine transits to q₄,
 it gets stuck in it and would not be able to transit to other states.
- Sometimes I call it "hell"!
- Now we understand the meaning of "Go to hell!"!



- What language does the following DFA accept?
 - Represent the language by set builder method.





- Let L = $\{ad, ada, adam\}$ over $\Sigma = \{a, d, m\}$.
- Design a DFA to accept L.

1

Design Important Notes

1. The machine should be designed in such a way that:

It accepts all strings of **L**.

AND

It rejects all strings of \overline{L} .

• Have we designed the previous example correctly?



Design Important Notes

- 2. To test your machine, both groups accepted strings and rejected strings, should be picked from Σ^* .
 - We are not allowed to input strings from outside Σ*.
 - Otherwise the behavior of the machine is not guaranteed.

A million dollar question!



- Can you ever claim that your design works fine?
- Never, because Σ* is infinite!
- So, theoretically every design has potential bugs!



- Let L be the set of strings over $\Sigma = \{a, b\}$ starting with prefix ab.
 - 1. Write a set-builder for L.
 - 2. Design a DFA to accept L.



- Let L be the set of strings over $\Sigma = \{1\}$ that contains even number of 1's.
 - 1. Write a set-builder for L.
 - 2. Design a DFA to accept L.



- Let L be the set of strings over $\Sigma = \{1\}$ that contains even unary numbers.
 - 1. Write a set-builder for L.
 - 2. Design a DFA to accept L.

Design Examples: DFA over $\Sigma = \{a, b\}$



Example 16: Empty Language

Example 17: All Strings

•
$$L = \Sigma^*$$

Example 18

•
$$L = \{\lambda\}$$

Homework

•
$$L = \Sigma^+$$





- Let L = $\{a^nb^m : n \ge 0, m \ge 0\}$
- Design a DFA to accept L.



- Let L be the set of strings over $\Sigma = \{0, 1\}$ consisting substring 001.
 - 1. Write a set-builder for L.
 - 2. Design a DFA to accept L.



- Let L be the set of strings over $\Sigma = \{0, 1\}$ consisting all strings except substring 001.
 - Design a DFA to accept L.
- What is the difference between this language and the previous one?

How to Construct a DFA to Accept $\overline{\mathbf{L}}$

Algorithm

- Construct a DFA to accept L.
- Turn all accepting states to regular states.
- Turn all regular states to accepting states.

Homework: DFA Design



- For each of the following languages over Σ = {a , b}:
 - Write a set-builder.
 - 2. Design a DFA for each.
 - The set of strings that contains exactly one 'a'
 - The set of strings that contains at least one 'a'
 - The set of strings ending with prefix ab
 - All strings with no more than three 'a's
 - All strings with at least three 'a's

Homework: DFA Design



- Design a DFA over $\Sigma = \{a, b\}$ for each of the following languages:
 - 1. All strings with exactly one 'a' and exactly two 'b's
 - 2. All strings with at least one 'a' and exactly two 'b's
 - 3. All strings with exactly two 'a's and more than two 'b's

Prepare Your Development Environment



JFLAP (Java Formal Language and Automata Package)

- We'll use JFLAP tools in this course to develop and test our automata.
- Official website: http://www.jflap.org/
- Download it from Canvas: Files/Misc/JFLAP.jar
- Or Download the stable version 7.0
 (thick version with SVG May 15, 2011) from:
 http://www.jflap.org/jflaptmp/
- Tutorial: http://www.jflap.org/tutorial/

JFLAP Demo



JFLAP



Basic Features

- Creating states
- Defining a state as initial state or final state
- Creating transitions
- Deleting
- Shift-Enter to create multiple transitions
- Multiple running (testing your design)
- Debugging: step-by-state
- Saving machines (xml file)

Other Features

- Changing state's name or label
- Adding comment
- Changing edges shape
- Selecting multiple objects and moving
- Zoom-in and out

References

- Linz, Peter, "An Introduction to Formal Languages and Automata, 5th ed.," Jones & Bartlett Learning, LLC, Canada, 2012
- Kenneth H. Rosen, "Discrete Mathematics and Its Applications, 7th ed.," McGraw Hill, New York, United States, 2012
- Michael Sipser, "Introduction to the Theory of Computation, 3rd ed.," CENGAGE Learning, United States, 2013 ISBN-13: 978-1133187790