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# **Grammars**

## **(Part 1)**

**Lecture 20**  
**Day 24/31**

**CS 154**  
**Formal Languages and Computability**  
**Spring 2018**

# Agenda of Day 24

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- Solution and Feedback of Quiz 7
- Solution and Feedback of Quiz ++
- Summary of Lecture 19
- Couple of slides from previous lecture
- Lecture 20: Teaching ...
  - Grammars (Part 1)

# Solution and Feedback of Quiz 7 (Out of 20)



Metrics	Section 1	Section 2	Section 3
Average	17	17	18
High Score	20	19	20
Low Score	6	11	12

# Solution and Feedback of Quiz ++ (Out of 60)



Metrics	Section 1	Section 2	Section 3
Average	50	49	50
High Score	59	54	59
Low Score	39	39	28

# Summary of Lecture 19: We learned ...

## REGEXs

- We defined REGEXs formally as:

1.  $\emptyset$ ,  $\lambda$ , and  $a \in \Sigma$  are all REGEXs.
2. If  $r_1$  and  $r_2$  are REGEXs, then the following expressions are REGEXs too:

$$r_1 + r_2$$

$$r_1 \cdot r_2$$

$$r_1^*$$

$$(r_1)$$

3. A string is REGEX iff it can be derived from the primitive REGEXs by a finite number of applications of the rule #2.

- Between REGEXs and languages, there are the following correspondence:

1.  $L(\emptyset) = \{ \}$
2.  $L(\lambda) = \{ \lambda \}$
3.  $L(a) = \{ a \}$  for all  $a \in \Sigma$
4.  $L(r_1 + r_2) = L(r_1) \cup L(r_2)$
5.  $L(r_1 \cdot r_2) = L(r_1) L(r_2)$
6.  $L((r_1)) = L(r_1)$
7.  $L(r_1^*) = (L(r_1))^*$

- We learned how to find the language represented by a REGEX by using the above correspondences.

**Any Question?**

# Summary of Lecture 19: We learned ...

## REGEXs

- Two regular expressions  $r_1$  and  $r_2$  are **equivalent** iff ...
- ... both **represent the same language**.

$$r_1 = r_2 \leftrightarrow L(r_1) = L(r_2)$$

## Identities

- If  $r$ ,  $s$ , and  $t$  are REGEXs, and  $a, b \in \Sigma$ , then:
  - $r(s + t) = rs + rt$
  - $(s + t)r = sr + tr$
  - $(a^*)^* = a^*$
  - $(a \dots a)^* a = a (a \dots a)^*$
  - $a^* (a + b)^* = (a + b)^* a^* = (a + b)^*$

## Theorem

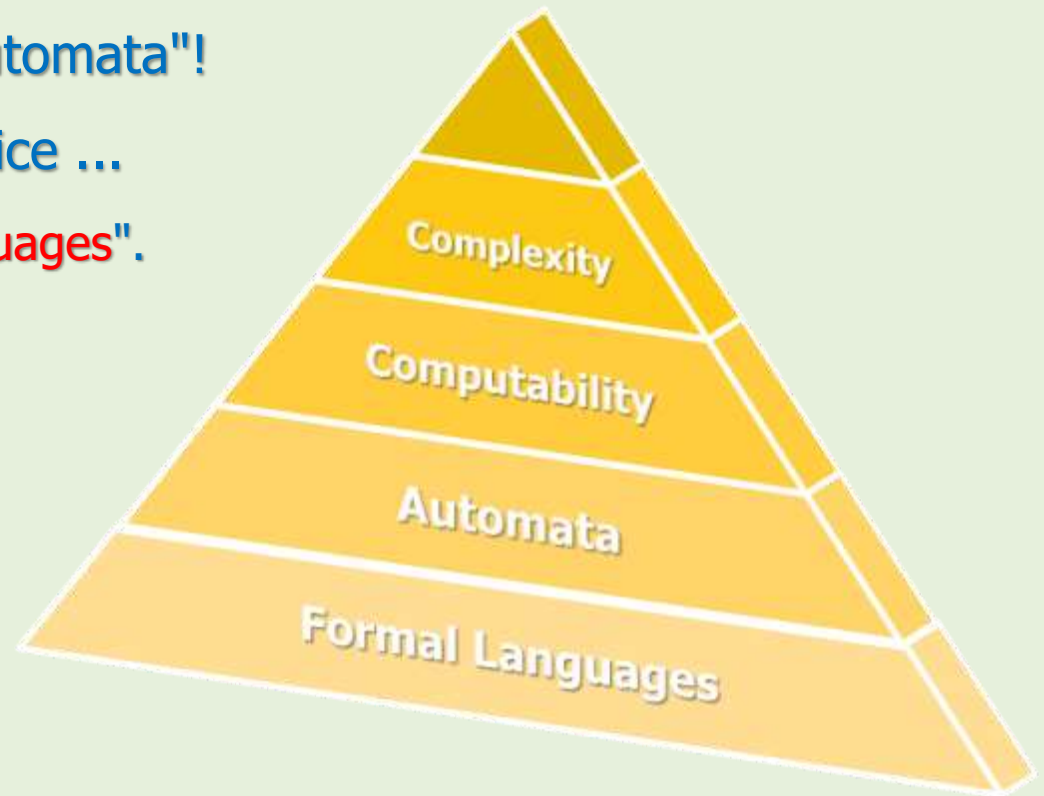
- A language is regular iff at least one REGEX represents it.
- The **limitation** of REGEXs ...
  - They just represent **regular languages**.
- While more interesting languages are **non-regular**!

**Any Question?**

# The Big Picture of the Course

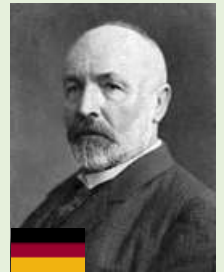
## Computer Science Foundation

- We started the semester with "Formal Languages" but we said:
  - we'd get back to it during the semester.
- We are done with the "Automata"!
- So far, we've got back twice ...
  - Introduced "Regular Languages".
  - Introduced "REGEXs".



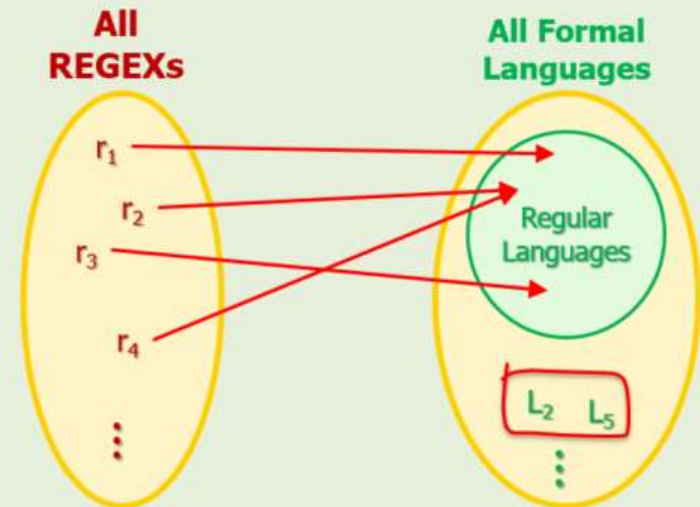
# Motivation

- So far, we've represented formal languages by two mathematical objects:
  - Sets (set builder, roster method, Venn diagrams)
  - Regular Expressions (REGEXs)



## What is the problem with them?

- Sets was introduced by a mathematician.
  - German mathematician, **George Cantor** (1845-1918)
- Sets are **NOT** practical in computer science!
- REGEXs are **limited** to regular languages.





# Objective of This Lecture

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- We need a more powerful and practical tool to represent all formal languages!

## That is Grammar!

- So, our target is to represent all formal languages!
- But we'll see the limitations of grammars later.

# Grammars

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# Introduction

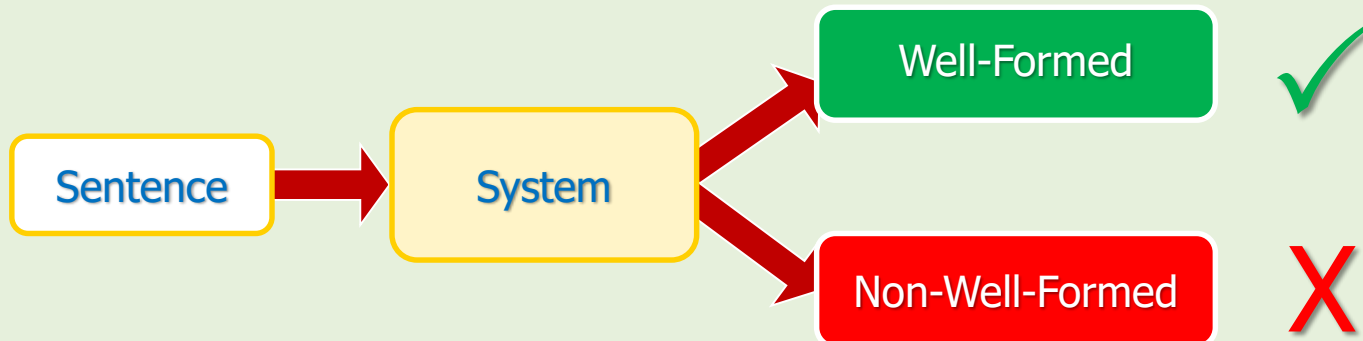
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- What would be your **reaction** if you encounter the following **English sentences**:
  - dog the runs.
  - dog runs the.
  - runs dog the.
- Even though **all words are correct English** words but the combinations do not make sense.
- The words' **positions** are not as they should be!
- In **computer science and linguistic terminology**, we say:  
These sentences are not "**well-formed**".

# What Are We Looking For?

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- We need a system to distinguish between:  
    "well-formed" and "non-well-formed"



- Let's take some examples from a "natural language" like English.
- Then we'll generalize the idea to formal languages.

# A Simple English Grammar

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- A simple rule for constructing a "sentence" in English grammar is:

$\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$

- Read " $\rightarrow$ " symbol as: "is defined as"



- This expression is called a "production rule".

- We can produce an English "sentence" by this production rule.

- The problem is:

We defined a sentence but we introduced two new "variables":

- $\langle \text{noun-phrase} \rangle$  and  $\langle \text{predicate} \rangle$ !

- So, we need to define  $\langle \text{noun-phrase} \rangle$  and  $\langle \text{predicate} \rangle$ .

# A Simple English Grammar

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$\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$

- We can define these new variables by the following **production rules**:

$\langle \text{noun-phrase} \rangle \rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle$

$\langle \text{predicate} \rangle \rightarrow \langle \text{verb} \rangle$

- Again, we introduced **new variables**  $\langle \text{article} \rangle$ ,  $\langle \text{noun} \rangle$ , and  $\langle \text{verb} \rangle$  that should be defined.
- So, we **keep going** ...

# A Simple English Grammar

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<sentence> → <noun-phrase> <predicate>

<noun-phrase> → <article> <noun>

<predicate> → <verb>

- We can define these new variables by the following **production rules**:

<article> → a

<article> → the

<noun> → dog

<noun> → boy

<verb> → runs

<verb> → walks

# A Simple English Grammar: Notes & Definitions

- Variables are defined either by other variables,
- Or they are assigned "values" (aka terminals),
- Or mix of the two (we'll see in the next slides)

## Draft Definition of Grammar

- The set of production rules is called "grammar".
  - Later, we define it formally.
- Every grammar has a "starting variable".
  - In this example, <sentence> is the starting variable.

## Simple English Grammar

1. <sentence> → <noun-phrase> <predicate>
2. <noun-phrase> → <article> <noun>
3. <predicate> → <verb>
4. <article> → a
5. <article> → the
6. <noun> → dog
7. <noun> → boy
8. <verb> → runs
9. <verb> → walks

Repeated



# A Simple English Grammar: Examples

## Example 1

- Is the following sentence "well-formed"?  
"the dog runs"

## Solution

- A sentence is well-formed if we can "derive" it from the production rules.
- We start from the "starting variable":

$\langle \text{sentence} \rangle \Rightarrow \langle \text{noun-phrase} \rangle$   
 $\quad \quad \quad \langle \text{predicate} \rangle$   
 $\Rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle \langle \text{verb} \rangle$   
 $\Rightarrow \text{the} \langle \text{noun} \rangle \langle \text{verb} \rangle$   
 $\Rightarrow \text{the dog} \langle \text{verb} \rangle$   
 $\Rightarrow \text{the dog runs}$

## Simple English Grammar

- $\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$
- $\langle \text{noun-phrase} \rangle \rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle$
- $\langle \text{predicate} \rangle \rightarrow \langle \text{verb} \rangle$
- $\langle \text{article} \rangle \rightarrow \text{a}$
- $\langle \text{article} \rangle \rightarrow \text{the}$
- $\langle \text{noun} \rangle \rightarrow \text{dog}$
- $\langle \text{noun} \rangle \rightarrow \text{boy}$
- $\langle \text{verb} \rangle \rightarrow \text{runs}$
- $\langle \text{verb} \rangle \rightarrow \text{walks}$

Repeated

# A Simple English Grammar: Examples

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## Example 1 (cont'd)

- We could derive the sentence "the dog runs" from the set of production rules (grammar), so, it is well-formed.
- We used " $\Rightarrow$ " notation for "derivation".
  - Read " $\Rightarrow$ " as: "derives"
- We can also represent the whole process as:
$$\langle \text{sentence} \rangle \overset{*}{\Rightarrow} \text{the dog runs}$$
to represent "multiple derivations" when we don't care about the middle ones.
- Now, let's take a failure example.

# A Simple English Grammar: Examples

## Example 2

- Is the following sentence "well-formed"?  
"the runs dog"

### Solution

- $\langle \text{sentence} \rangle \Rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$   
 $\Rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle \langle \text{verb} \rangle$   
 $\Rightarrow \text{the} \langle \text{noun} \rangle \langle \text{verb} \rangle$
- "runs" is NOT a "noun"!
- It fails, so, the sentence is not well-formed.
- What else can we derive from this grammar?

## Simple English Grammar

- $\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$
- $\langle \text{noun-phrase} \rangle \rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle$
- $\langle \text{predicate} \rangle \rightarrow \langle \text{verb} \rangle$
- $\langle \text{article} \rangle \rightarrow \text{a}$
- $\langle \text{article} \rangle \rightarrow \text{the}$
- $\langle \text{noun} \rangle \rightarrow \text{dog}$
- $\langle \text{noun} \rangle \rightarrow \text{boy}$
- $\langle \text{verb} \rangle \rightarrow \text{runs}$
- $\langle \text{verb} \rangle \rightarrow \text{walks}$

Repeated

# A Simple English Grammar: Examples

## Example 3

- The **set of all sentences derivable** from the grammar:

{ "a dog runs",  
"a dog walks",  
"a boy runs",  
"a boy walks",  
"the dog runs",  
"the dog walks",  
"the boy runs",  
"the boy walks" }

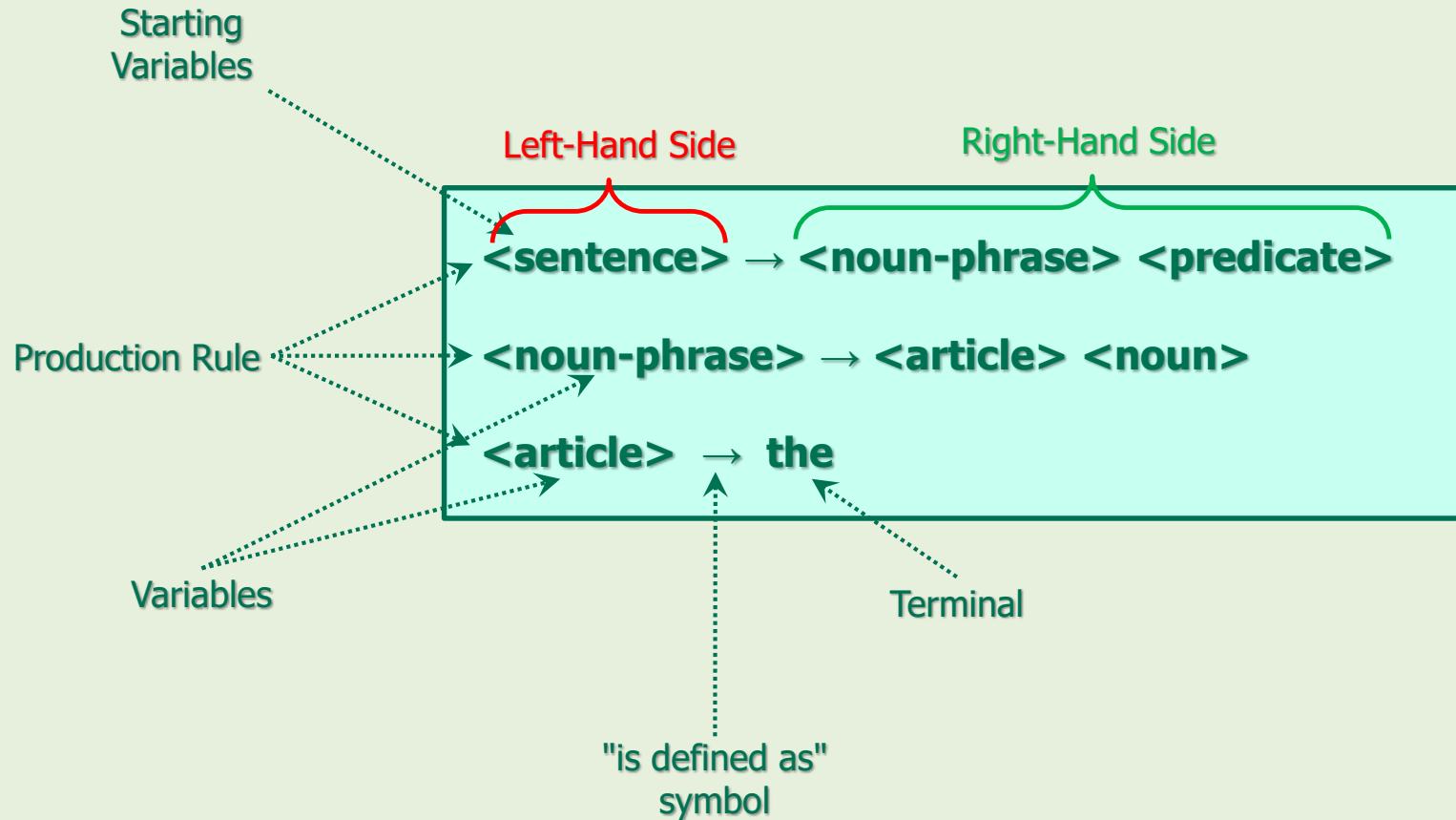
- What can we call this **set of strings**?  
**The language generated by the grammar**

## Simple English Grammar

1.  $\langle \text{sentence} \rangle \rightarrow \langle \text{noun-phrase} \rangle \langle \text{predicate} \rangle$
2.  $\langle \text{noun-phrase} \rangle \rightarrow \langle \text{article} \rangle \langle \text{noun} \rangle$
3.  $\langle \text{predicate} \rangle \rightarrow \langle \text{verb} \rangle$
4.  $\langle \text{article} \rangle \rightarrow a$
5.  $\langle \text{article} \rangle \rightarrow the$
6.  $\langle \text{noun} \rangle \rightarrow dog$
7.  $\langle \text{noun} \rangle \rightarrow boy$
8.  $\langle \text{verb} \rangle \rightarrow runs$
9.  $\langle \text{verb} \rangle \rightarrow walks$

Repeated

# Grammar Terminologies



## A Side Note About Natural Languages

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- Unfortunately, natural languages were **not developed** by mathematicians.
- That's why, there are **thousands of exceptions** in their definition!
- So, learning them perfectly as second languages is **almost impossible!**
- The only language that was **developed by a scientist** is ...  
**Esperanto!**
- For more info, please look at the appendix A at the end of this lecture note.
- Note that it is **just for your information** and is **NOT** part of this course.

# Formal Grammars

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# Formal Grammars

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- We can generalize the natural languages grammars to formal grammars.

## Example 4

- Consider the following set of "production rules":

$$\left\{ \begin{array}{l} S \rightarrow aB \\ B \rightarrow baB \\ B \rightarrow \lambda \end{array} \right.$$

- These rules are an example of "formal grammar".
- Let's explain in detail what they mean?



# Ingredients of the Production Rules

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- S and B are "variables".
  - Variables are represented by capital letters.
- The "starting variable" by default is 'S' unless we mention something else.
- 'a', 'b', and  $\lambda$  are called "terminal symbols".
  - Terminals are represented by lower-case letters.
  - $\lambda$  is our familiar "empty string".
  - Terminals can be any sequence of terminal symbols or  $\lambda$ .
- "aB" and "baB" contain both variable and terminal and is called "sentential form".

$$\left\{ \begin{array}{l} S \rightarrow aB \\ B \rightarrow baB \\ B \rightarrow \lambda \end{array} \right.$$

# How a string can be derived from a grammar?

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## Example 5

- Consider the following grammar:
  - $S \rightarrow a S b$
  - $S \rightarrow \lambda$
- Derive string "ab"

## Solution

$$S \xRightarrow{1} a S b \xRightarrow{2} a \lambda b \Rightarrow ab$$



- Could we derive the string "ab" if we had started with rule #2?

# Derivation of Strings

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## Example 6

- Consider the following grammar:
  - $S \rightarrow a S b$
  - $S \rightarrow \lambda$
- Derive string "aabb".

## Solution

$$S \xRightarrow{1} a S b \xRightarrow{1} aa S bb \xRightarrow{2} aabb$$

- We can summarize the above derivation like this:

$$S \xRightarrow{*} aabb$$

- As we said before, this notation is used when we don't care about the detail of the derivations.

# A Convention

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- When the **left-hand sides** of two or more production rules **are the same**, we can combine the right-hand sides by separating them with a vertical bar "|".

## Example 7

- Consider the following grammar:

$$S \rightarrow a S b$$

$$S \rightarrow \lambda$$

- We can represent it as:

$$S \rightarrow a S b \mid \lambda$$

# Associated Language to Grammars

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- We can apply the production rules "recursively" in any arbitrary orders.
- Therefore, a grammar can generate many strings.

## Definition



- The set of all strings produced (aka generated) by the grammar  $G$  is called the "associated language to  $G$ " and is denoted by  $L(G)$ .

# ! Associated Language to Grammars



## Example 8

- Consider the following grammar:

$$S \rightarrow a S \mid \lambda$$

- What language does it produce?
  - Show it by set-builder.

## Solution



- How about this one:

$$S \rightarrow S a \mid \lambda$$

- Is there any difference?



# Associated Language to Grammars

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## Example 9

- Consider the following grammar:

$$S \rightarrow a S b \mid \lambda$$

- What language does it produce?

## Solution



## Conclusion

- After this example, we know that grammars can represent more languages than regular languages!



# Associated Language to Grammars

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## Example 10

- Consider the following grammar:
  1.  $S \rightarrow AB$
  2.  $A \rightarrow aA \mid \lambda$
  3.  $B \rightarrow bB \mid \lambda$
- What language does it produce?

## Solution



# References

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1. Linz, Peter, "An Introduction to Formal Languages and Automata, 5<sup>th</sup> ed.," Jones & Bartlett Learning, LLC, Canada, 2012
2. Michael Sipser, "Introduction to the Theory of Computation, 3<sup>rd</sup> ed.," CENGAGE Learning, United States, 2013  
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3. The ELLCC Embedded Compiler Collection, available at: <http://ellcc.org/>

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# **Esperanto**

## **A World Without War!**

# **Appendix A**

**Spring 2018**

# Esperanto Creator's Motivation

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- Constructed in 1873 by **Polish** medical doctor, inventor, and writer, **Ludwik L. Zamenhof** (1859-1917).



- He had the dream of "**a world without war**".
- He believed this language can help international people to **communicate** easily.
- And the communications could prevent wars!

# Alphabet (5 Vowels + 23 Consonants)

Letter	English Example
a	fa <b>th</b> er
b	
c	ca <b>t</b> s
ĉ	ch <b>i</b> p
d	
e	be <b>t</b>
f	
g	g <b>o</b>
ĝ	g <b>e</b> m

Letter	English Example
h	
ĥ	kh in Persian
i	se <b>e</b>
j	y <b>e</b> s
ĵ	mea <b>s</b> ure
k	
l	
m	
n	

Letter	English Example
o	s <b>o</b>
p	
r	rolled "R"
s	
ŝ	sh <b>a</b> re
t	
u	so <b>o</b> n
ŭ	co <b>w</b>
v	
z	

# Vowels & Consonants

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- 5 Vowels: A, E, I, O, U
- 23 Consonants: the rest of alphabets
- Name the vowels by their pronunciation
- Name the consonants by: letter + o
- For example, we call "b" as "bo".
- The number of syllables of a word is the number of vowels.
- For example, "domo" (means house) has two syllables.
- The accented syllable is the second to last syllable.

# Rules

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- All "**nouns**" are ended with '**o**'.
- For example, "domo" means house.
- All "**adjectives**" are ended with '**a**'.
- For example, "doma" means domestic.
- **Adjectives** can be **placed** before or after nouns.
- For example, "dolĉa pomo" or "pomo dolĉa" means "sweet apple".
- To make a **negative** adjective, just prefix it with "**mal**".
- For example, "bona" means "good" and "malbona" means "bad".

# Rules

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- All "adverbs" are ended with 'e'.
- For example, "rapide" means "quickly".
- Adverbs can be placed before or after the verb they modify.
- For example, "Ĝi flugas rapide" means "it flies quickly".
- Or we can say, "Ĝi rapide flugas".
- To make "plural", add 'j' at the end of the nouns and adjectives.
- For example, "seĝoj" means chairs.
- Example for plural, "rapidaj aŭtoj" means fast cars.
- Note that both adjective and noun should be plural.

# Rules

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- There is only one "definite article" "la" in Esperanto.
- For example, "la domo" means "the house".
- There is no indefinite article in Esperanto.
- So, if you don't use "la", it means the article is indefinite.



# Rules

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- To make "**possessive personal pronoun**", just add 'a' to the end of personal pronoun.
- For example, "mi" means "I", "mia" means "my".
- "vi" means "you", "via" means "your".
  
- To make a **verb negative**, just add "ne" before it.
- For example, "Ŝi estas alta." means "She is tall".
- "Ŝi ne estas alta." means "She is not tall".

# Tenses

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- **Infinitives** are ended with **'i'**.
- For example, "flugi" means "to fly".
  
- **Present tense** is ended with **"as"**.
- For example, "Mi flugas" means "I fly".
  
- **Past tense** is ended with **"is"**.
- For example, "Mi flugis" means "I flied".
  
- **Future tense** is ended with **"os"**.
- For example, "Mi flugos" means "I will fly".

# Tenses

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- **Progressive Present** is ended with "**anta**" but needs "**estas**" as well.
- For example, "Mi estas fluganta" means "I am flying".


# Miscellaneous

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- To **make a question**, add "**Ĉu**" at the beginning of the sentences.
  - For example, "Mi flugas" means "I fly".
  - "**Ĉu** mi flugas" means "Do I fly".
  - It doesn't matter what tense the sentence has.
- 
- Over all, Esperanto has **16 constants rules**.
  - These were some of them, I selected for your information.

# General Info About Esperanto

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- "Esperanto" means "one who hopes".
- The flag of Esperanto ...
- It is the most successful constructed language in the world.
- Almost two million people speak in Esperanto.
- The first World Congress of Esperanto was organized in 1905 in France.
- Esperanto was recognized as a language by UNESCO in 1954.
  - It was recommended as international non-governmental organizations language in 1985.
- In 2012, Google Translate added Esperanto into its list.
- Language codes in computer: ISO 639-1, ISO 639-2, ISO 639-3

# Did You Know That ...

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- Learning Esperanto as the second language will **speed up** learning 3<sup>rd</sup> and more languages.
- **8 Nobel laureates** have been Esperantists
- Esperantists make up the **largest non-political grouping** in the British parliament.
- **Leo Tolstoy** helped found the Esperantist Vegetarian Association in 1908.
- In 1993, more than 4900 people (mainly non-Esperantists) visited the International **Esperanto Museum in Vienna**.
- Esperantists have continually **suffered oppression** from totalitarian governments

**Reference:** <http://esperanto.org/us/USEJ/world/index.html>

# References

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1. Wikipedia, Esperanto available at: <https://en.wikipedia.org/wiki/Esperanto>
2. Font for Windows: Tajpi <http://www.zz9pza.net/tajpi/en/installation/>
3. Esperanto lessons: <https://www.youtube.com/watch?v=bLx5hLag6WQ>