

1. Purpose

Implement two types of algorithms for multiplying two $n \times n$ matrices. Assume n is a power of 2:

1. The straight-forward $O(n^3)$ matrix multiplication algorithm.
2. Strassen's matrix multiplication algorithm.

Evaluate your different algorithms, and write a short report. Create test matrices for different values of n (4, 16, 512 and 1024).

2. Algorithm

1. Algorithm for straight-forward $O(n^3)$ matrix multiplication.

Here is the pseudocode of the straight-forward matrix multiplication.

for $i = 1$ to n

 for $j = 1$ to n

 for $k = 1$ to n

$C(i,j) = C(i,j) + A(i,k) * B(k,j)$

Therefore, the running time is $O(n^3)$.

2. Strassen's matrix multiplication algorithm.

First, divide A , B , and C into same size matrix.

$$A = \begin{bmatrix} A_{1,1} & A_{1,2} \\ A_{2,1} & A_{2,2} \end{bmatrix}, B = \begin{bmatrix} B_{1,1} & B_{1,2} \\ B_{2,1} & B_{2,2} \end{bmatrix}, C = \begin{bmatrix} C_{1,1} & C_{1,2} \\ C_{2,1} & C_{2,2} \end{bmatrix};$$

So, we can find matrix C can be solved by the below method.

$$\begin{aligned}
C_{1,1} &= A_{1,1}B_{1,1} + A_{1,2}B_{2,1} \\
C_{1,2} &= A_{1,1}B_{1,2} + A_{1,2}B_{2,2} \\
C_{2,1} &= A_{2,1}B_{1,1} + A_{2,2}B_{2,1} \\
C_{2,2} &= A_{2,1}B_{1,2} + A_{2,2}B_{2,2}
\end{aligned}$$

Second, from Strassen's method, we can decrease the calculate times from 8 to 7.

$$\begin{aligned}
M_1 &:= (A_{1,1} + A_{2,2})(B_{1,1} + B_{2,2}) \\
M_2 &:= (A_{2,1} + A_{2,2})B_{1,1} \\
M_3 &:= A_{1,1}(B_{1,2} - B_{2,2}) \\
M_4 &:= A_{2,2}(B_{2,1} - B_{1,1}) \\
M_5 &:= (A_{1,1} + A_{1,2})B_{2,2} \\
M_6 &:= (A_{2,1} - A_{1,1})(B_{1,1} + B_{1,2}) \\
M_7 &:= (A_{1,2} - A_{2,2})(B_{2,1} + B_{2,2})
\end{aligned}$$

Third, we use these 7 new sub-matrix to get the sub-matrix of C.

$$\begin{aligned}
C_{1,1} &= M_1 + M_4 - M_5 + M_7 \\
C_{1,2} &= M_3 + M_5 \\
C_{2,1} &= M_2 + M_4 \\
C_{2,2} &= M_1 - M_2 + M_3 + M_6
\end{aligned}$$

In the end, we use these four sub-array of C to get the complete answer of C.

Therefore, we get the conclusion that the running time of Strassen's matrix multiplication

$$\text{is } O(n^{\log_2 7}) = O(n^{2.807}).$$

3. Conclusion

When I use the test file from the professor, I get the correct answer which show on the

below.

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This is Matrix 1:
9.0 7.0 2.0 2.0
5.0 3.0 4.0 7.0
6.0 1.0 4.0 8.0
8.0 1.0 7.0 2.0

This is Matrix 2:
8.0 6.0 1.0 2.0
7.0 4.0 6.0 0.0
9.0 6.0 8.0 3.0
1.0 5.0 3.0 6.0

141.0 104.0 73.0 36.0
104.0 101.0 76.0 64.0
99.0 104.0 68.0 72.0
136.0 104.0 76.0 49.0

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Are matrices the same? true

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When I run the JUnit file, I got the result on the below.

<pre> When size of the matrix is 4 Strassen method using time: 1 Counter of the Strassen method: 57 When size of the matrix is 16 Strassen method using time: 4 Counter of the Strassen method: 2801 When size of the matrix is 512 Strassen method using time: 21445 Counter of the Strassen method: 47079208 When size of the matrix is 1024 Strassen method using time: 147842 Counter of the Strassen method: 329554457 </pre>	<pre> When size of the matrix is 4 Strassen method using time: 1 Counter of the Strassen method: 57 When size of the matrix is 16 Strassen method using time: 4 Counter of the Strassen method: 2801 When size of the matrix is 512 Strassen method using time: 21445 Counter of the Strassen method: 47079208 When size of the matrix is 1024 Strassen method using time: 147842 Counter of the Strassen method: 329554457 </pre>
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From the final answer of comparing these two method, we can see that the counter of the Strassen method is always smaller than straight forward method. That means the running time of Strassen method is always smaller than straight forward method. Therefore, Strassen's matrix multiplication algorithm is more efficiency than straight-forward matrix multiplication algorithm.