



# Bitwise operators

- ① AND  $\&$
- ② OR  $|$
- ③ XOR  $\wedge$
- ④ NOT  $\sim$  (unary op)
- ⑤ Left Shift  $\ll$
- ⑥ Right Shift  $\gg$

$$\begin{array}{r} 101 \\ 111 \\ \hline 101 \end{array}$$

$$\underbrace{5 \ \& \ 7} \Rightarrow 5$$

↑  
operate on bits

28 logical and  
11 logical OR  
! logical Not

boolean expression

$$\underbrace{5 > 3}_T \quad \boxed{28} \quad \underbrace{7 > 4}_T$$

└──────────┘  
T

Q1. Check/Get the  $i^{\text{th}}$  bit of number.

$n = 13$

get last bit  $\Rightarrow n \& 1$

Get  $i^{\text{th}}$  Bit ( $n, i$ ) {

return  $(n \gg i) \& 1$

}

$n = 13$

$= \overset{3}{1} \overset{2}{1} \overset{1}{0} \overset{0}{1}$

$i = 2$

0 0 0 1 1 0 0

↑  
last-bit

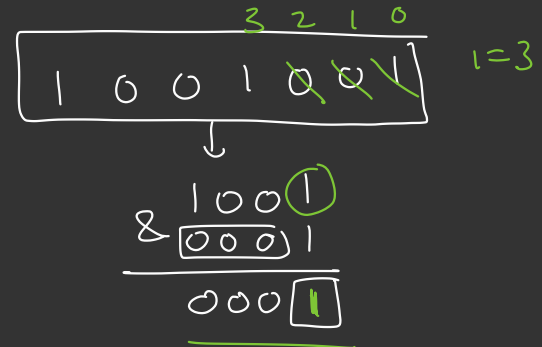
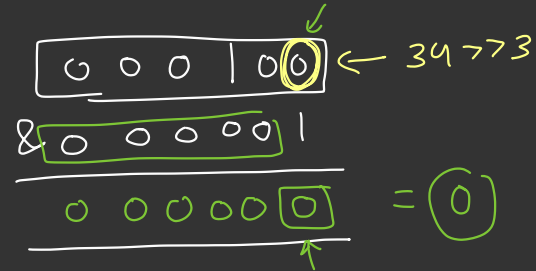
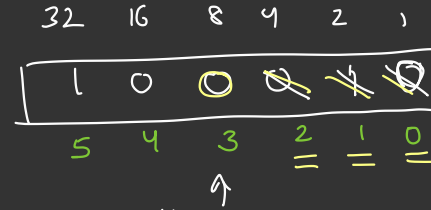
$$n = 34$$

$$i = \underline{\underline{3}}$$

$$(34 \gg \underline{\underline{3}}) \& 1$$

$$i^{\text{th}} \text{ bit} = (n \gg i) \& 1$$

$$\underline{\underline{n \gg 3}}$$

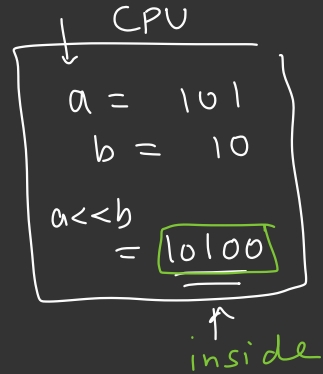


int a = 5

int b = 2

print( a << b )

↑  
(20) outside  
(display)  
==



Get  $i^{\text{th}}$  Bit ( $n, i$ ) {

return  $(n \gg i) \& 1$

↑  
always produces 0 or 1

}

$(n \& (1 \ll i)) == 0$   
 $> 0$

0 bit  
1 bit

$n = 13$   
 $i = 2$

2 1 0  
1 1 0 1

1 1 0 1  
& 0 1 0 0  
-----  
0 1 0 0

↑  
i<sup>th</sup> bit  
↓  
set bit

$> 0$

$n =$  1 0 0 1  
0 1 0 0  
-----  
0 0 0 0

$= 0$   
i<sup>th</sup> bit  
↓  
zero bit

$O(1)$

get I<sup>th</sup> Bit ( $n, i$ ) {

if ( $(n \& (1 \ll i)) > 0$ )

return 1

else

return 0

}

return

~~$n \& (1 \ll i)$~~   
No

$n =$

$i = 5$

$n =$ 

	5	4	3	2	1	0
	1	0	0	0	1	0

  
↑

$1 \ll i =$ 

0	1	0	0	0	0	0
---	---	---	---	---	---	---

  
⇒ 

0	1	0	0	0	0	0
---	---	---	---	---	---	---

1	1	0	1
0	1	0	0
0	1	0	0

  
→ 4

Q Unset  $i^{\text{th}}$  bit of a number.

→ if bit is 0 → No change  
 → " is 1 → 0

$N = 45$

5	4	3	2	1	0
1	0	1	1	0	1
(32)		8	4	1	
Retain		↓	Retain		
<hr/>					
1	0	1	0	0	1
<hr/>					

$i = 2$

$h \rightarrow$

1	0	1	1	0	1
^			1	0 0	
<hr/>			0	0 1	

$1 \ll i$



set Bit

if (checkbit(n, i) == 1) {

n = n ^ (1 << i)

}

only when  $i^{\text{th}}$  bit is 1

Q) Count total no of set bits in number N  
 [one] <sup>↑ bit == 1</sup>

N = 45    000 101101 → 1+0+1+1+0+1  
ans = 0    (4)

ans = 0+1+0+1  
 +1+0+1  
 = 4

while (N > 0) {

last\_bit = (N & 1)

ans = ans + last\_bit

N = N >> 1

↓ integer 0 or 1

}  
 print(ans)

↘ equivalent N = N/2

X 1 X X X X X

Time Complexity

O(Log N)

≤ 32

log N

$n = \underline{45}$   
 $n = \underline{22}$   
 $n = \underline{11}$   
 $n = \underline{5}$   
 $n = \underline{2}$   
 $n = \underline{1}$   
 $n = \underline{0}$

$\underline{10110}$  ~~1~~  
 $0 \quad \underline{1011}$  ~~1~~  
 $00 \quad 101$  \*  
 $000 \quad 10$  \*  
 $0000 \quad 1$  \*  
 $0000$  \*  
 $\underline{00000}$

ans = 0

$1$   
 $+$   
 $0$   
 $+$   
 $1$   
 $+$   
 $1$   
 $+$   
 $0$   
 $+$   
 $1$

STOP

$\rightarrow \boxed{n = n/2}$  Div  
code  $\rightarrow \boxed{n = n > 1}$  left shift

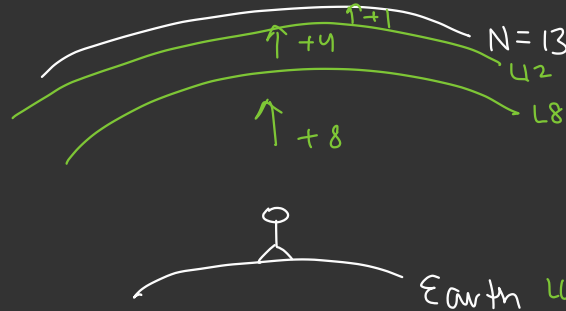
$n = \frac{100}{1} \rightarrow n = \frac{50}{2} \rightarrow n = \frac{25}{3} \rightarrow n = \frac{12}{4} \rightarrow n = \frac{6}{5} \rightarrow n = \frac{3}{6} \rightarrow n = \frac{1}{7} \rightarrow 0$

$\log_2 100$   
 $= 6.xx$   
 $= 7$

Q. There is a battle going a N levels above the surface earth. The hero which in at earth can take jumps in power of 2. Minimum jumps he will need to reach level N.

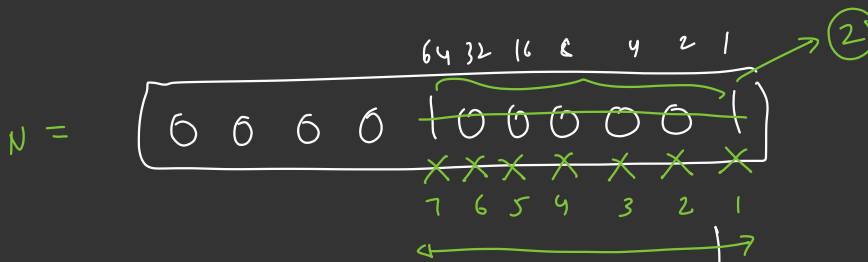
$$\text{jump sizes} = [1, 2, 4, 8, 16, \dots]$$

3 jumps



$$13 = \underbrace{1}_8 + \underbrace{1}_4 + \underbrace{1}_1$$

Set Bits  $\rightarrow$  <sup>Set</sup> count bit  
 $= 3$



How many its itr

7 iterations

$\log 65 = 6.0 \times \times$

$= (7)$

① ①

$N \rightsquigarrow N/2 \rightsquigarrow N/4 \dots 0$

↓

$\log N$

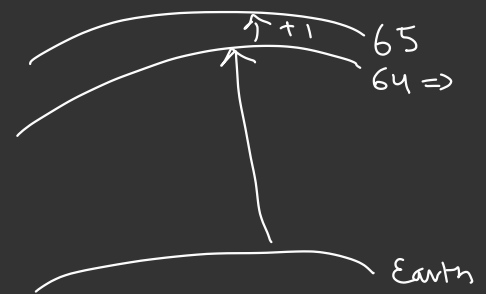
Counting set bit takes  $\log N$  iterations

No of sets bits  $\neq \log N$

count

Set Bits = (2)

2 jumps

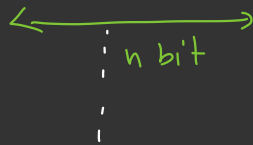


## Negative Numbers

4 bit bucket



$\Rightarrow 0$



$\Rightarrow 15$

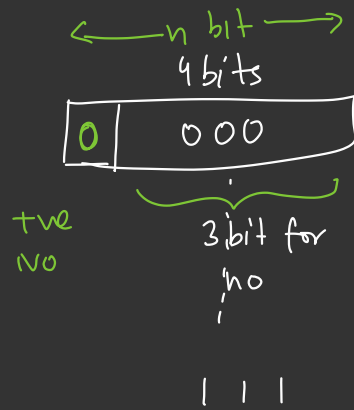
$$2 \times 2 \times 2 \times 2 = 2^4 = 16 \text{ combinations}$$

Range 0-15  
16 Numbers

Concept of sign bit was introduced



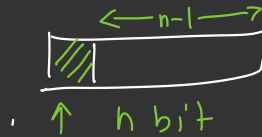
$$0 \longleftrightarrow 2^n - 1$$



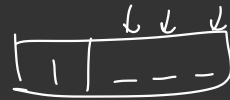
(0 - 7)

8 combinations

$$2^{4-1} - 1 = 7$$



0 to  $2^{n-1} - 1$



negative  
no

8 combinations

Range

(-8, -7, ..., -1)



4 bits

0 — 15  
└────────┘  
16 nbs



↑  
MSB  
as  
sign bit

-8 — 0 — 7  
└────────┘  
16 Numbers

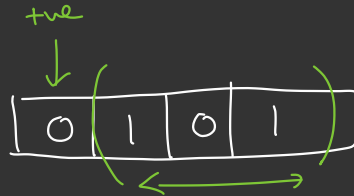


How -ve no's are actually stored?

↳ 2's complement form

↓ if we add negative no's  
in this  
form  
result of  
subtraction.

(5)

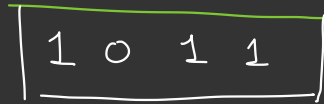


-(5)

① Flip all bits of 5

② Add 1

1 0 1 0  
+ 1

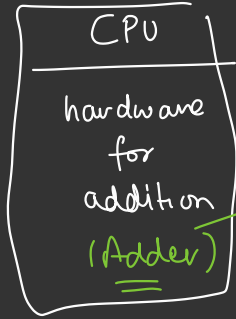


~~5 - 5~~ subtrahend

5 + (-5)

= (0)





dropped off ← 4 bit box →

(5)

5

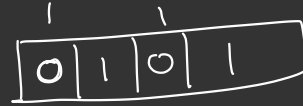
- 3

=

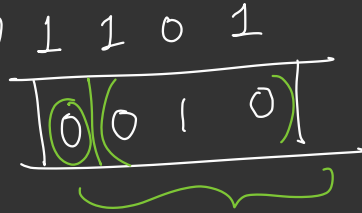
(2)

5 + (-3)

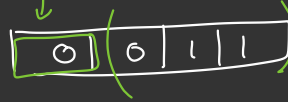
(5)



1



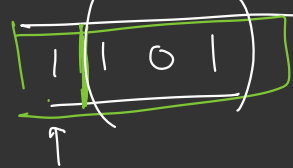
+ve



3

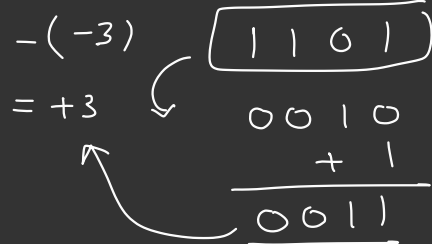
1 1 0 0

+

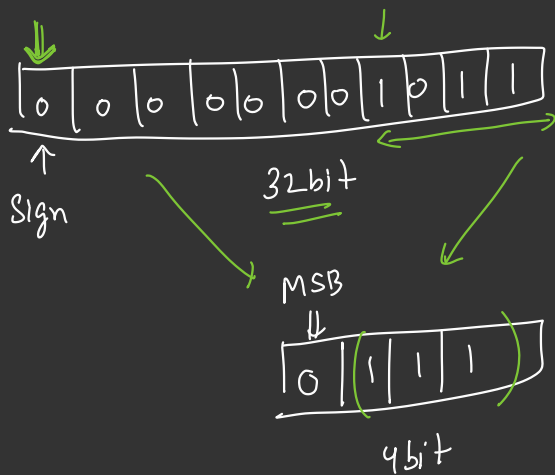


-3

(2)

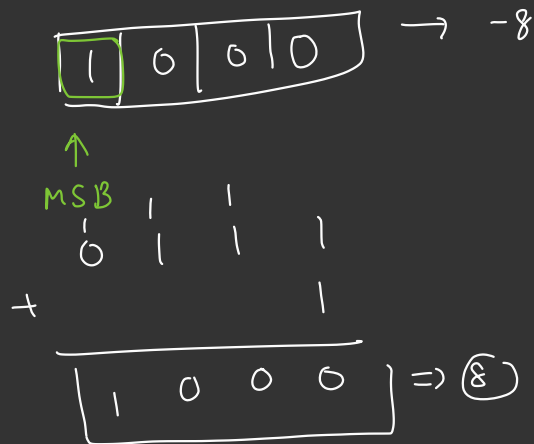


Computer



Read  $\Rightarrow$   
 this  
 NO

$-(-x)$   
 $\Rightarrow +x$

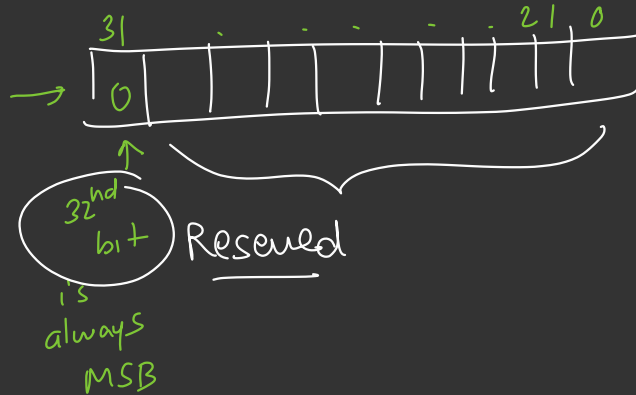


$$\begin{array}{r} 5 \\ - 5 \\ \hline 1010 \\ + 1 \\ \hline 1011 \\ \hline \end{array}$$

⑧

Java

int x = 44  
y = -76



(Standard)

getI<sup>m</sup>Bit(N, i)  
↓  

yes

~~no~~

=> 1 

-ve

  
0 

+ve

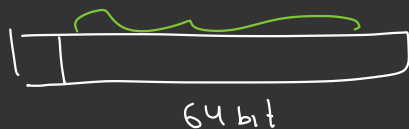
## Range of No's

int



$$\begin{array}{c} \overset{31}{-2} \longleftarrow 0 \longrightarrow \overset{31}{2-1} \\ -2 \times 10^9 \qquad \qquad \qquad \approx 2 \times 10^9 \end{array}$$

long



$$\begin{array}{c} \overset{63}{-2} \longleftarrow 0 \longrightarrow \overset{63}{2-1} \\ -2 \times 10^{18} \qquad \qquad \qquad \approx 2 \times 10^{18} \end{array}$$

Constraints  
Analyze

Q

Calc sum of all elements of an array of size N

$$\underline{A[i] \leq 10^6}, \quad \underline{N = 10^5}$$

long

```
int ans = 0  
for (int x : arr) {  
    ans = ans + x  
}
```

$$\underbrace{[10^6 + 10^6 + \dots + 10^6]}_{10^5} \approx 10^{11}$$

}



(b) given two integers, find  $a * b$

$$\begin{aligned} a &\leq 10^9 \\ b &\leq 10^9 \end{aligned}$$

X (1)  $\text{print}(a * b)$   $\xrightarrow{\text{int}}$  overflow

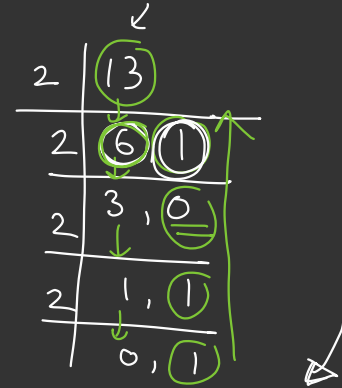
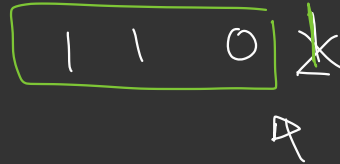
system  
assumes  
output is  
an int

X (2)  $\text{long ans} = \underline{a} * \underline{b}$   $\xrightarrow{\text{int}}$   $\text{print}(\text{ans})$   $\sim 10^{18}$

✓ (3)  $\text{long ans} = \frac{\text{long}(a) * b}{\text{long}}$   $\text{print}(\text{ans})$   $\xrightarrow{\text{long}}$

# Counting Set Bits

13 =



ans = 0

while (N > 0) {

ans = ans +  $\frac{(n \& 1)}{1}$

n = n >> 1

↑  
last bit

3

faster

while (N > 0) {

rem =  $N \% 2$

ans = ans + rem

N = N / 2

3

5      5 bits

0	0	1	0	1
---	---	---	---	---

- 5

1	1	0	1	1
---	---	---	---	---

~~$(-5) + (-3)$~~

3

0	0	0	1	1
---	---	---	---	---

-3

1	1	1	0	1
---	---	---	---	---

$(-5)$   
 $+(-3)$   


---

 $-8$

~~⊗~~

1	1	0	1	1
1	1	1	0	1

1	1	0	0	0
---	---	---	---	---

$-x \Rightarrow$ 
 $-8$

↑

MSB  
-ve

$-(-x) = +x$

↓

 $8$

$1+1+1$   
 $= 11$

$$3 + (-5)$$

$$\begin{array}{r}
 \begin{array}{cccc|cc}
 & & & & 1 & & \\
 0 & 0 & 0 & & 1 & 1 & \\
 1 & 1 & 0 & 1 & & 1 & 
 \end{array} \\
 \hline
 \boxed{1 \ 1 \ 1 \ 1 \ 0} \Rightarrow (-2) \\
 \begin{array}{r}
 00001 \\
 + \quad 1 \\
 \hline
 00010 \Rightarrow (2)
 \end{array}
 \end{array}$$

$$\begin{array}{r}
 \begin{array}{ccc|ccc}
 & & & 1 & 1 & 1 & \\
 0 & 0 & & 1 & 1 & 1 & \\
 + & & & & & 1 & 
 \end{array} \\
 \hline
 \boxed{0 \ 1 \ 0 \ 0 \ 0} \\
 \begin{array}{c}
 \uparrow \\
 \text{MSB} \\
 \text{+ve}
 \end{array}
 \quad
 \textcircled{+8}
 \end{array}$$

# Decimal to Binary

N = 13

2	13	
	↓	
2	6, 1	(1)
		✓
2	3, 0	0
		✓
2	1, 1	(1)
		✓
	0, 1	(1)
		✓

output =

1101

1000  
100  
00  
1

p = 1  
1101

$1 \times 1 = 1$   
 $0 \times 10 = 00$   
 $1 \times 100 = 100$   
 $1 \times 1000 = 1000$

while (N > 0) {

rem = N % 2

ans = ans \* 10 + rem

p = p \* 10

N = N / 2

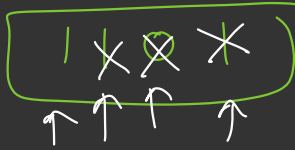
}

ans = 0, p = 1

N & 1

N = N >> 1

13



$$\begin{array}{r} 1 \times 1 \\ + 0 \times 10 \\ + 1 \times 100 \\ + 1 \times 1000 \\ \hline \underline{\underline{1101}} \end{array}$$

while(N>0){

bit = N&1

ans = ans + b\*bit

b = b\*10

N = N>>1

}