



SoftUni Team Technical Trainers







Software University

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Have a Question?





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Algorithmic Complexity



- Describes performance of the particular algorithm
 - Runtime and memory consumption based on the input size N
 - We usually care about the worst-case performance
- We measure the complexity as the Big O notation
 - Numerical function depending on the input size O(N)
 - We measure time as the number of simple steps
 - We measure memory as input data N by its type size

Algorithmic Complexity

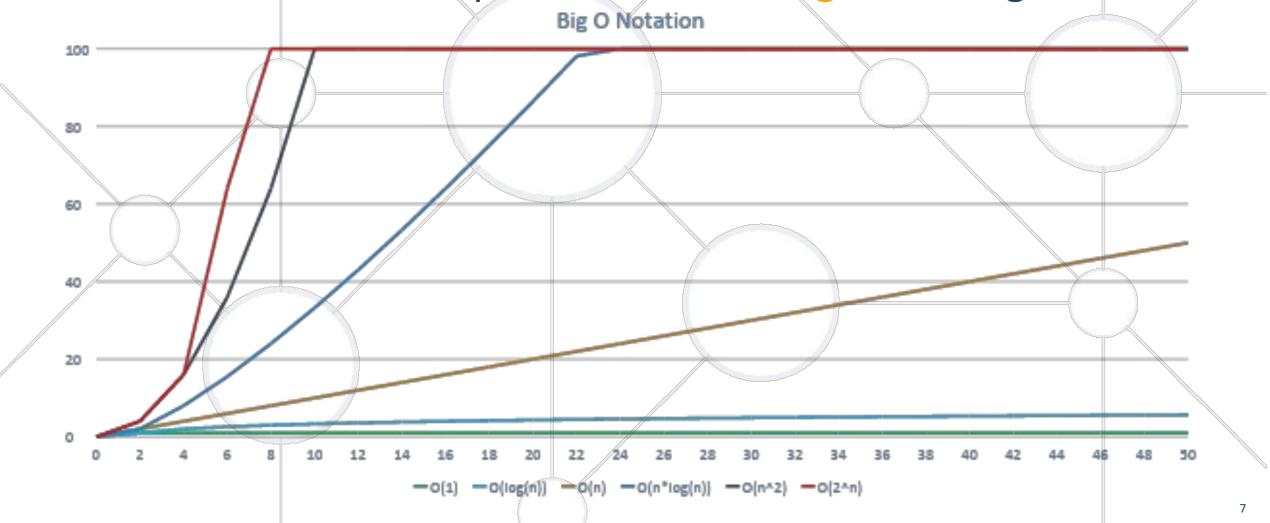


- O(1) Constant time time does not depend on N
- O(log(N)) Logarithmic time grows with rate as log(N)
- O(N) Linear time grows at the same rate as N
- O(N²), O(N³) Quadratic, Cubic grows as square or cube of N
- O(2^N) Exponential grows as N becomes the exponent worst algorithmic complexity
 - For input size of 10 1024 steps
 - For input size of 100 1267650600228229401496703205376 steps
- http://bigocheatsheet.com/

Asymptotic Functions



Below are some examples of common algorithmic growth:



Get Sum Number of Steps



Calculate maximum steps to find sum of even elements in an array

```
int getSumEven(int[] array) {
   int sum = 0;
   for (int i = 0; i < array.length; i++)
   if (array[i] % 2 == 0) sum += array[i];
   return sum;
}

Counting maximum steps is
   called worst-case analysis</pre>
```

- Assume that a single step is a single CPU instruction:
 - assignments, array lookups, comparisons, arithmetic operations

Time Complexity



Worst-case

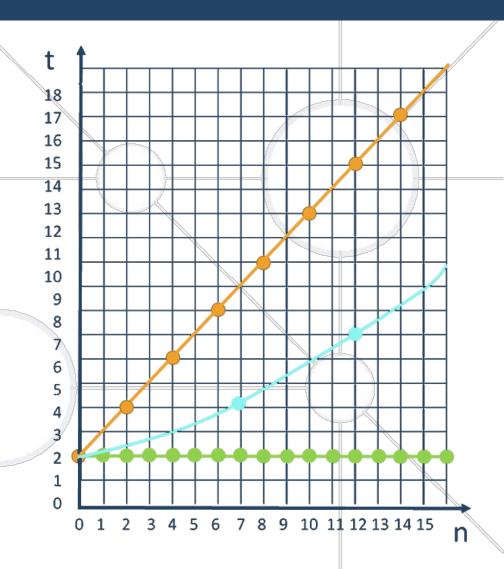
An upper bound on the running time

Average-case

Average running time

Best-case

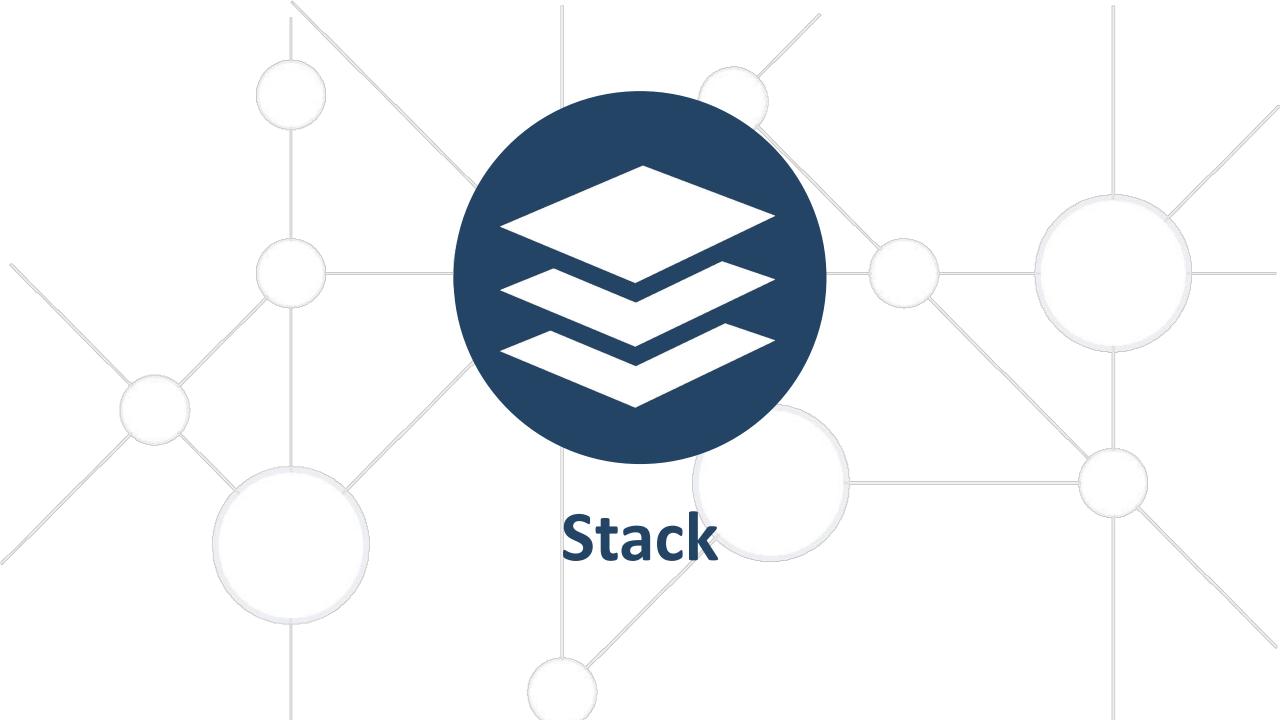
 The lower bound on the running time (the optimal case)



Stacks and Queue vs. ArrayDeque



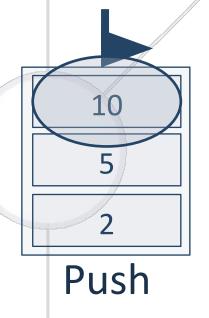
- Why don't use Stack and Queue?
 - Implementation details which make unsecure usability
 - In many cases, those structures will decrease the performance
- Why use ArrayDeque?
 - An implementation that makes the structure more secure
 - Better performance and usability
 - Methods that operate as those structures suggest

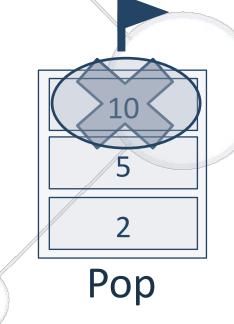


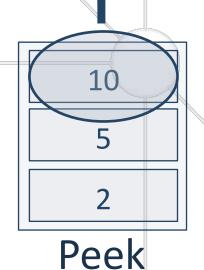
Stack Functionality



- Stacks provide the following functionality:
 - Pushing an element at the top of the stack
 - Popping element from the top of the stack
 - Getting the topmost element without removing it







ArrayDeque<E> – Java Stack Implementation Software University



Creating a Stack

ArrayDeque<Integer> stack = new ArrayDeque<>();

Adding elements at the top of the stack

```
stack.push(element);
```

Removing elements

```
Integer element = stack.pop();
```

Getting the value of the topmost element

```
Integer element = stack.peek();
```

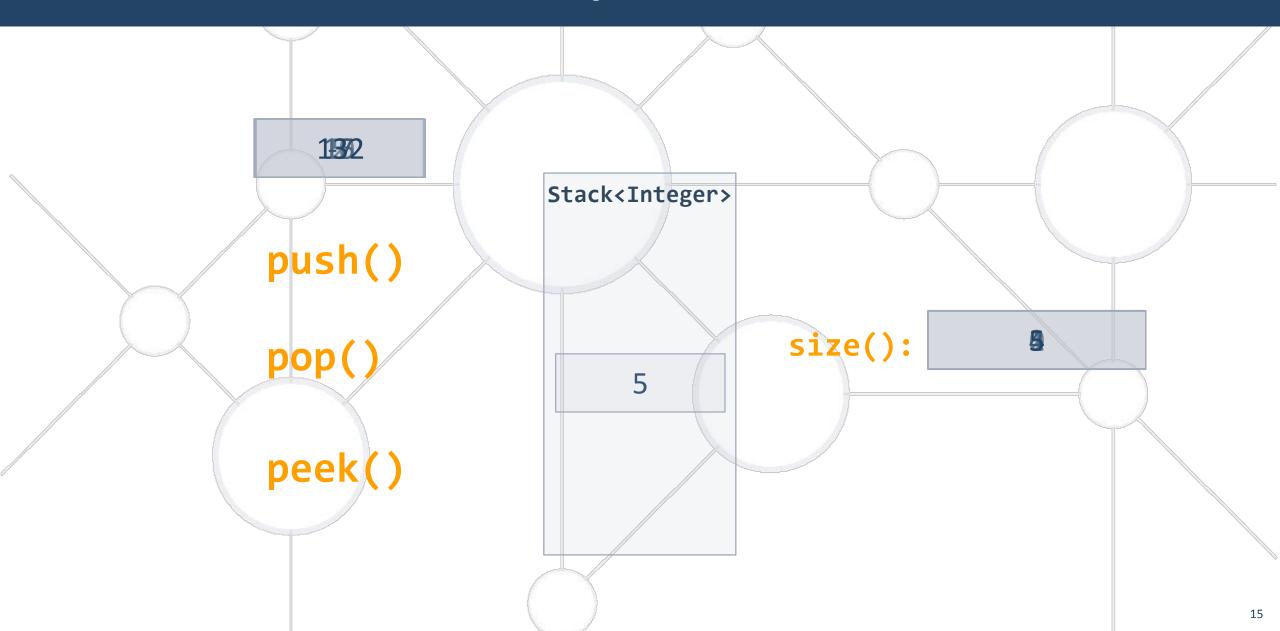
Stack – Utility Methods



```
ArrayDeque<Integer> stack = new ArrayDeque<>();
int size = stack.size();
boolean isEmpty = stack.isEmpty();
boolean exists = stack.contains(2);
```

Stack – Overview of All Operations





Problem: Browser History



- Write a program which takes 2 types of browser instructions:
 - Normal navigation: a URL is set, given by a string
 - The string "back" that sets the current URL to the last set URL

https//softuni.bg/ back https//softuni.bg/trainings/courses back https//softuni.bg/trainings/2056 back https//softuni.bg/trainings/live https//softuni.bg/trainings/live/details Home

Output



no previous URLs
https://softuni.bg/trainings/courses
https://softuni.bg/
https://softuni.bg/trainings/2056

https//softuni.bg/

https//softuni.bg/

https//softuni.bg/trainings/live

https//softuni.bg/trainings/live/details

Solution: Browser History (1)



```
Scanner scanner = new Scanner(System.in);
ArrayDeque<String> browser = new ArrayDeque<>();
String line = scanner.nextLine();
String current = "";
   continue...
```

Solution: Browser History (2)



```
while(!line.equals("Home")) {
  if(line.equals("back"))
    if(!browser.isEmpty()) { current = browser.pop();
    } else {
    System.out.println("no previous URLs");
    line = scanner.nextLine();
    continue; }
  } else {
    if(!current.equals("")) { browser.push(current); }
    current = line; }
  System.out.println(current);
  line = scanner.nextLine(); }
```

Problem: Simple Calculator



Implement a simple calculator that can evaluate simple expressions (only addition and subtraction)



Check your solution here: https://judge.softuni.org/Contests/3953/Stacks-and-Queues-Lab-RS

Solution: Simple Calculator (1)



```
Scanner scanner = new Scanner(System.in);
String[] tokens = scanner.nextLine().split("\\s+");
                                               Split by regex
Deque<String> stack = new ArrayDeque<>();
Collections.addAll(stack, tokens);
                          Adds a collection to
   continues.
                          another collection
```

Solution: Simple Calculator (2)

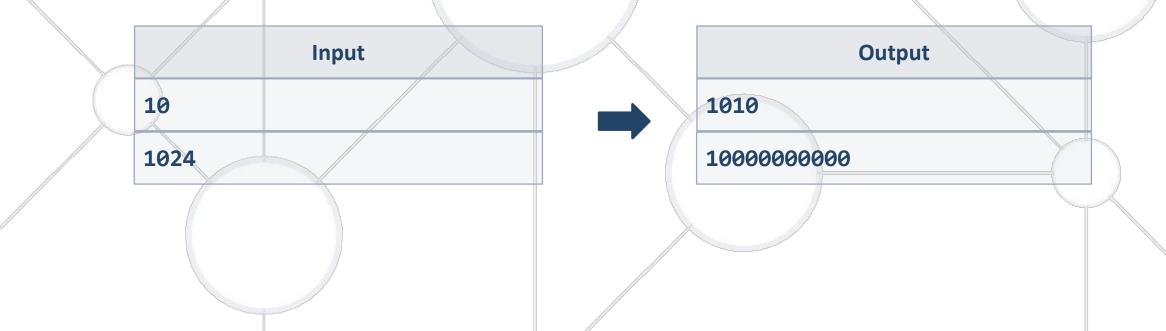


```
while (stack.size() > 1) {
  int first = Integer.valueOf(stack.pop());
  String op = stack.pop();
  int second = Integer.valueOf(stack.pop());
  switch (op)
    case "+": stack.push(String.valueOf(first + second));
      break;
    case "-": stack.push(String.valueOf(first - second));
      break;
System.out.println(stack.pop());
```

Problem: Decimal to Binary Converter



 Create a converter which takes a decimal number and converts it into a binary number



Check your solution here: https://judge.softuni.org/Contests/3953/Stacks-and-Queues-Lab-RS

Solution: Decimal to Binary Converter



```
Scanner scanner = new Scanner(System.in);
int decimal = Integer.valueOf(scanner.nextLine());
ArrayDeque<Integer> stack = new ArrayDeque<>();
// TODO: check if number is 0
while (decimal != 0)
  stack.push(decimal % 2);
  decimal /= 2;
while (!stack.isEmpty())
  System.out.print(stack.pop());
```

Problem: Matching Brackets



- We are given an arithmetical expression with brackets (with nesting)
- Goal: extract all sub-expressions in brackets



Solution: Matching Brackets (1)



```
Scanner scanner = new Scanner(System.in);
String expression = scanner.nextLine();

Deque<Integer> stack = new ArrayDeque<>();

// continue...
```

Solution: Matching Brackets (2)

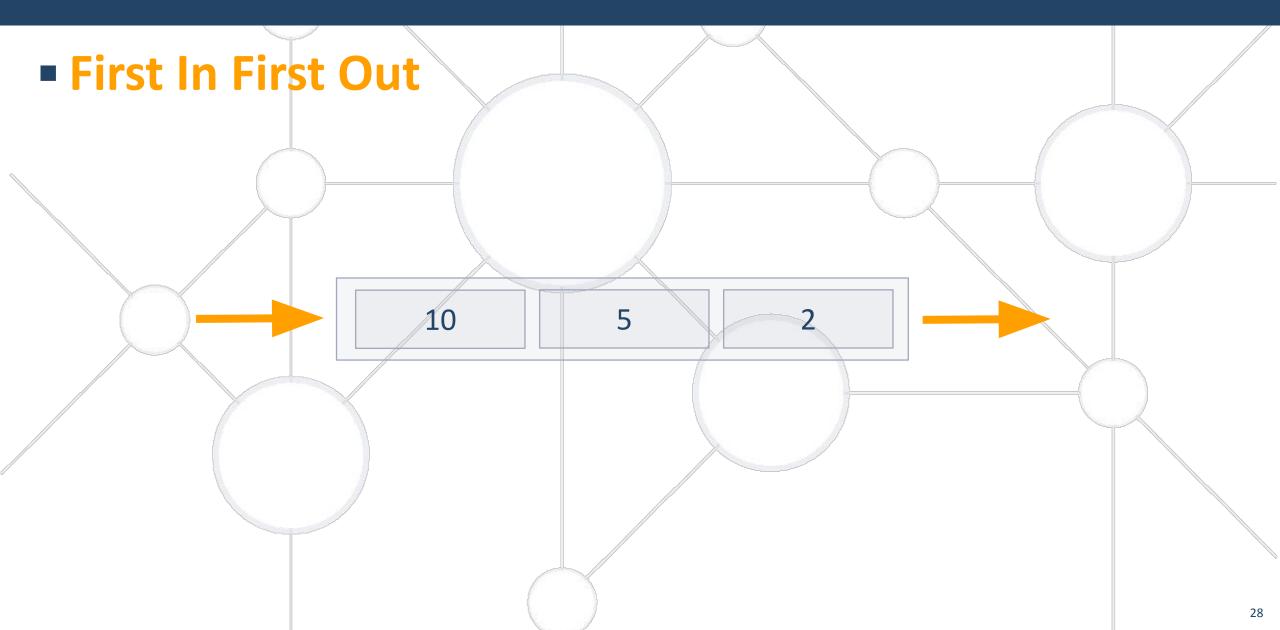


```
for (int i = 0; i < expression.length(); i++) {</pre>
  char ch = expression.charAt(i);
  if (ch == '(')
    stack.push(i);
  else if (ch == /')')
    int startIndex = stack.pop();
    String contents =
  expression.substring(startIndex, i + 1);
    System.out.println(contents);
```



Queue





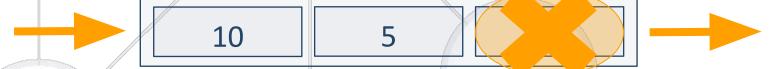
Queue – Abstract Data Type



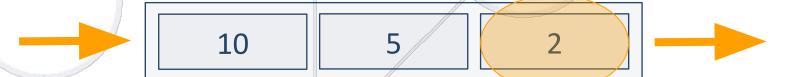
- Queues provide the following functionality:
 - Adding an element at the end of the queue



Removing the first element from the queue



Getting the first element of the queue without removing it



ArrayDeque<E> - Java Queue Implementation (1)



Creating a Queue

ArrayDeque<Integer> queue = new ArrayDeque<>();

Adding elements at the end of the queue

```
queue.add(element);
queue.offer(element);
```

- add() throws exception if queue is full
- offer() returns false if a queue is full

ArrayDeque<E> – Java Queue Implementation (2)



Removing elements

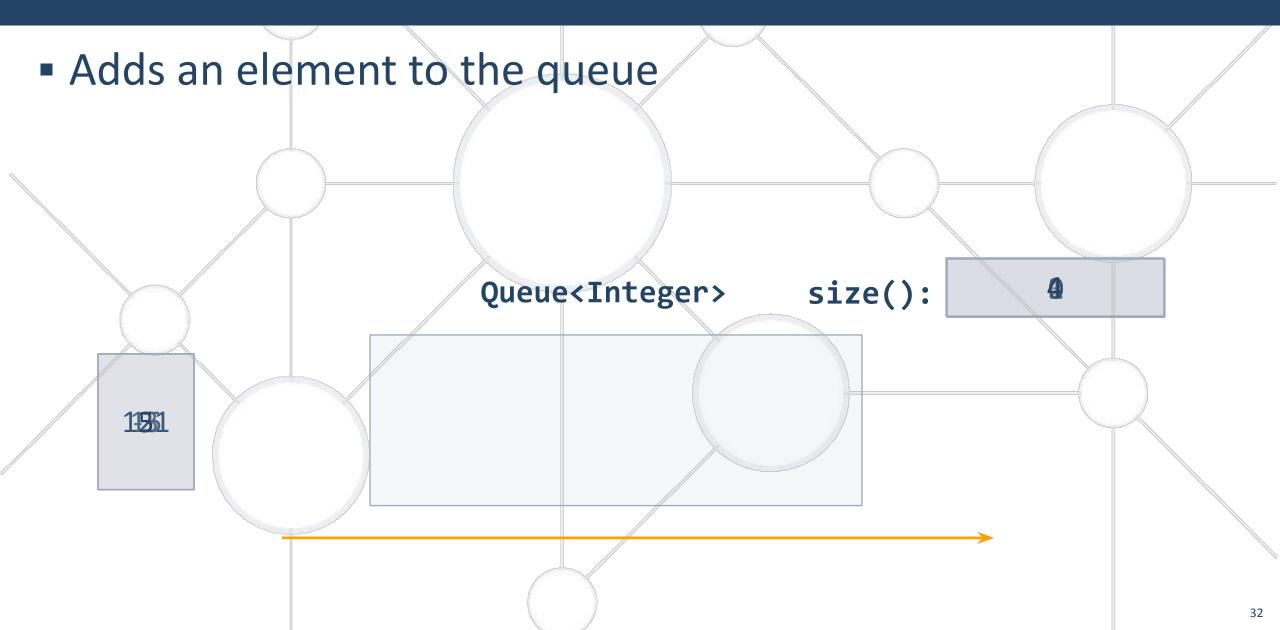
```
element = queue.remove();
element = queue.poll();
```

- remove() throws exception if queue is empty
- poll() returns null if queue is empty
- Check first element

```
element = queue.peek();
```

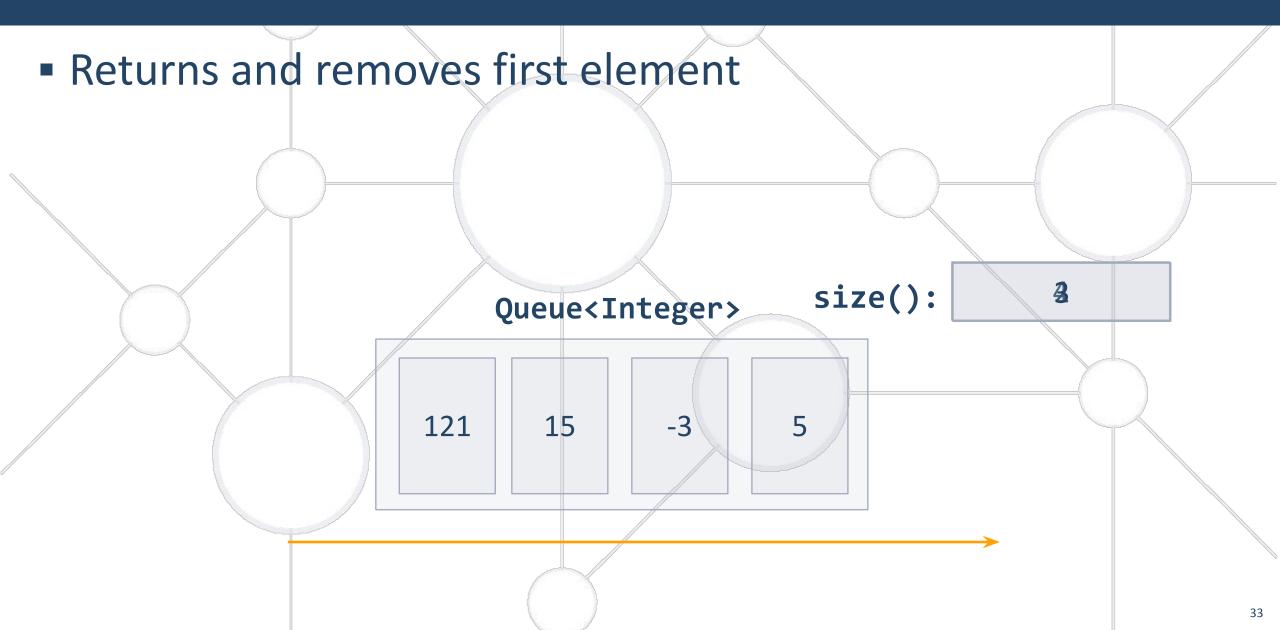
Add() / Offer()





Remove() / Poll()

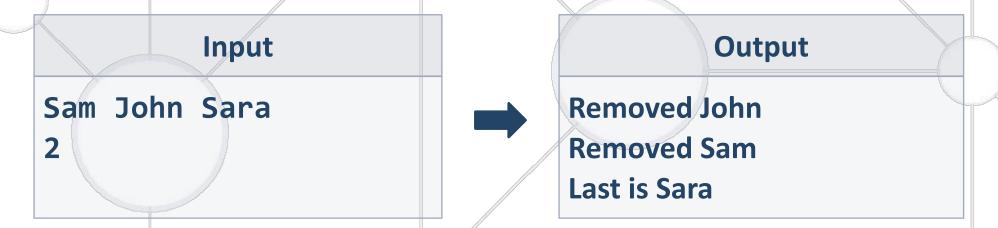




Problem: Hot Potato



- Children form a circle and pass a hot potato clockwise
- Every nth toss a child is removed until only one remains
- Upon removal the potato is passed forward
- Print the child that remains last



Solution: Hot Potato (1)



```
Scanner scanner = new Scanner(System.in);
String[] children = scanner.nextLine().split("\\s+");
int n = Integer.valueOf(scanner.nextLine());
ArrayDeque<String> queue = new ArrayDeque<>();
for (String child: children)
  queue.offer(child);
// continue...
```

Solution: Hot Potato (2)



```
while (queue.size() > 1) {
  for (int i = 1; i < n; i++)
    queue.offer(queue.poll());
  System.out.println("Removed " + queue.poll());
System.out.println("Last is " + queue.poll());
```

ArrayDeque<E> – Java Queue Implementation (3)



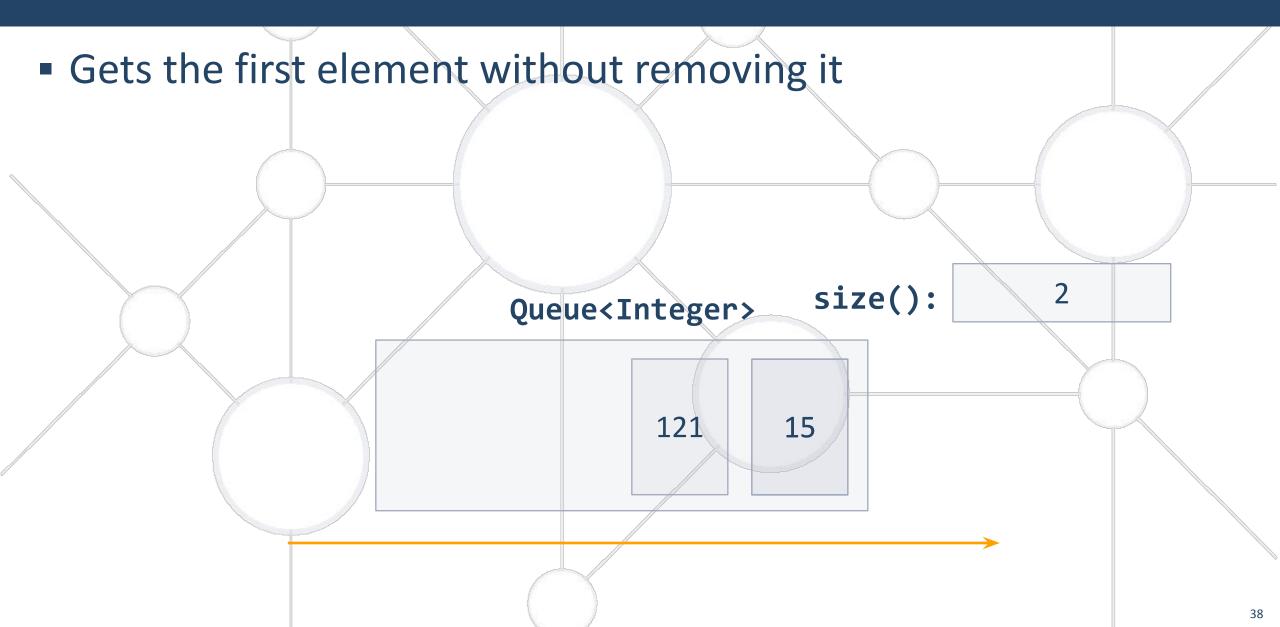
Utility Methods

```
Integer element = queue.peek();
Integer size = queue.size();
Integer[] arr = queue.toArray();
boolean exists = queue.contains(element);
```

- peek() checks the value of the first element
- size() returns queue size
- toArray() converts the queue to an array
- contains() checks if element is in the queue

Peek()

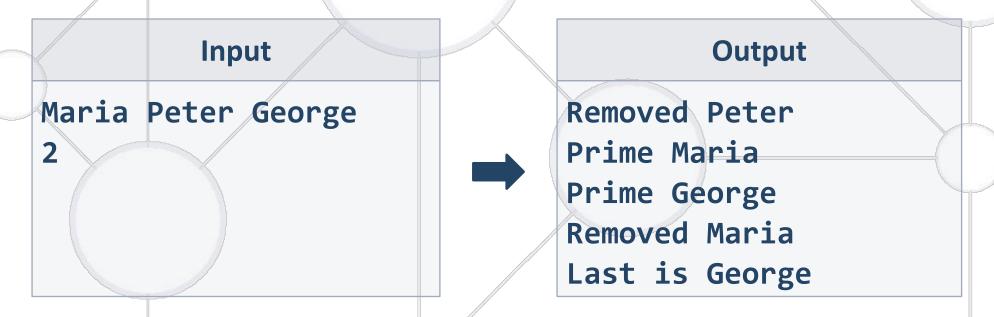




Problem: Math Potato



- Rework the previous problem so that a child is removed only on a prime cycle (cycles start from 1)
- If a cycle is not prime, just print the child's name



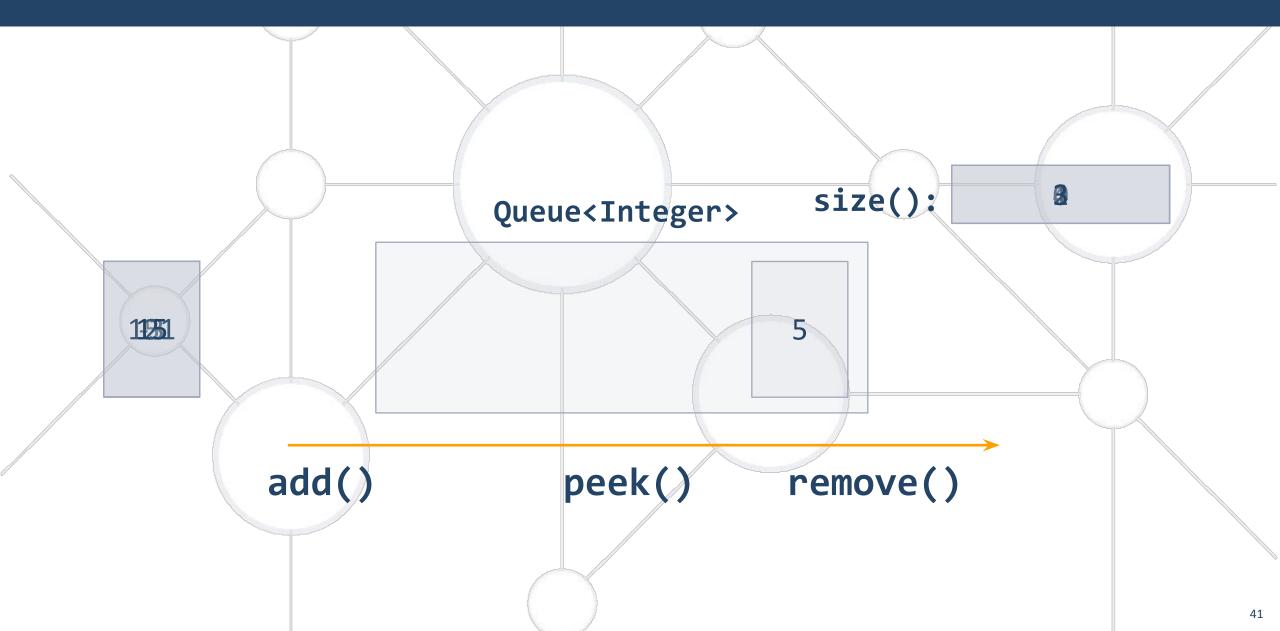
Solution: Math Potato



```
int cycle = 1;
while (queue.size() > 1) {
  for (int i = 1; i < n; i++)
    queue.offer(queue.poll());
  if (isPrime(cycle))
   System.out.println("Prime " + queue.peek());
  else
    System.out.println("Removed " + queue.poll());
  cycle++;
System.out.println("Last is " + queue.poll());
```

Queue – Overview of All Operations

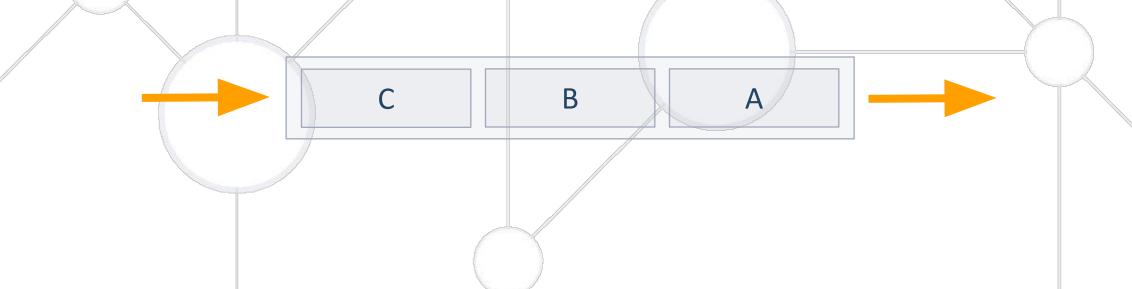




Priority Queue



- Retains a specific order to the elements
- Higher priority elements are pushed to the beginning of the queue
- Lower priority elements are pushed to the end of the queue



Summary



- Algorithmic Complexity
- Stack Last In First Out (LIFO)
 - push(), pop(), peek()
- Queue First In First Out (FIFO)
 - add(), poll(), peek()
- Priority Queue





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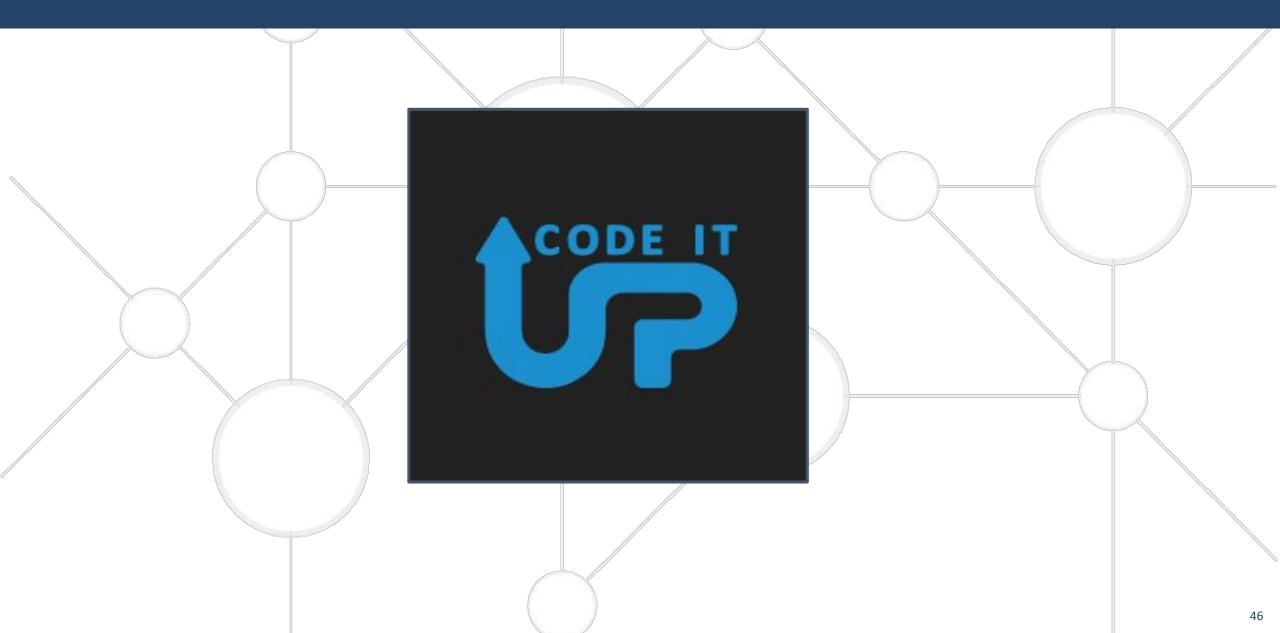






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