

# COMP261 Lecture 19

Parsing 4 of 4

**Victoria**  
UNIVERSITY OF WELLINGTON  
*Te Whare Wānanga  
o te Upoko o te Ika a Māui*  
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## Admin: Test results

- See your test results online.
- Collect your test paper from School office.

## A Better parser: using patterns

- Give names to patterns to make program easier to understand and to modify
- Precompile the patterns for efficiency:
 

```
private Pattern numPat = Pattern.compile(
    "[ -+]?((\\d+([.]\\d*)?)|([.]\\d+))");
private Pattern addPat = Pattern.compile("add");
private Pattern subPat = Pattern.compile("sub");
private Pattern mulPat = Pattern.compile("mul");
private Pattern divPat = Pattern.compile("div");
private Pattern opPat =
    Pattern.compile("add|sub|mul|div");
private Pattern openPat = Pattern.compile("\\(");
private Pattern commaPat = Pattern.compile(",");
private Pattern closePat = Pattern.compile("\\)");
```

## A Better parser: using patterns

```
public Node parseExpr(Scanner s) {
    Node n;
    if (!s.hasNext()) { fail("Empty expr",s); }
    if (s.hasNext(numPat)) { return parseNumber(s); }
    if (s.hasNext(addPat)) { return parseAdd(s); }
    if (s.hasNext(subPat)) { return parseSub(s); }
    if (s.hasNext(mulPat)) { return parseMul(s); }
    if (s.hasNext(divPat)) { return parseDiv(s); }
    fail("Unknown expr",s);
    return null;
}
```

## A Better parser: multiple arguments

Allow `add(1,2,3)`, etc.

```
public Node parseAdd(Scanner s) {
    List<Node> args = new ArrayList<Node>();
    require(addPat, "Expecting add", s);
    require(openPat, "Missing '(', s);
    args.add(parseExpr(s));
    do {
        require (commaPat, "Missing ',', s);
        args.add(parseExpr(s));
    } while (!s.hasNext(closePat));
    require(closePat, "Missing ')'", s);
    return new AddNode(args);
}
```

(need new version of require, taking a Pattern instead of a String)

## Multiple arguments: Printing AST

```
NumberNode: public String toString(){
    return String.format("%.5f", value);
}

AddNode: public String toString(){
    String ans = "(" + first;
    for (Node nd : rest){ ans += " + " + nd; }
    return ans + ")";
}

SubNode: public String toString(){
    String ans = "(" + first;
    for (Node nd : rest){ ans += " - " + nd; }
    return ans + ")";
}
```

## Examples

Expr: `add(10.5,-8)`

Print → `(10.5 + -8.0)`

Value → 2.500

Expr: `add(sub(10.5,-8), mul(div(45, 5), 6.8))`

Print → `((10.5 - -8.0) + ((45.0 / 5.0) * 6.8))`

Value → 79.700

Expr: `add(14.0, sub(mul(div (1.0, 28), 17), mul(3, div(5, sub(7, 5)))))`

Print → `(14.0 + (((1.0 / 28.0) * 17.0) - (3.0 * (5.0 / (7.0 - 5.0)))))`

Value → 7.107

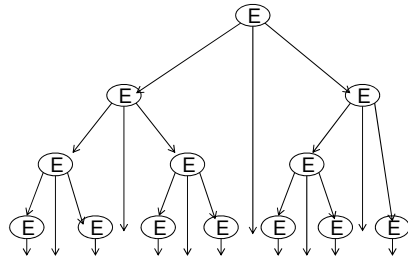
## Less Restricted Grammars

- This is enough for most of the assignment:
  - method for each Node type
  - peek at next token to determine which branch to follow
  - build and return node
  - throw error (including helpful message) when parsing breaks
  - use `require(...)` to wrap up "check then consume/return or fail"
  - adjust grammar to make it cleaner
- What happens when our grammar is not quite so helpful?
- For example:
 
$$E ::= \text{number} \mid E "+" E \mid E "-" E \mid E "*" E \mid E "/" E$$
- What are the problems, and how can you fix them?

## Ambiguous Grammars

Grammar:

$E ::= \text{number} \mid E "+" E \mid E "-" E \mid E "*" E \mid E "/" E$

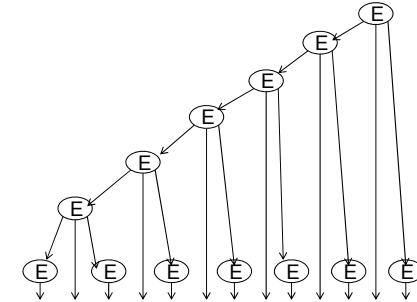


Text:  $65 * 74 - 68 + 25 * 5 / 3 + 16$

## Possible Parses

Grammar:

$E ::= \text{number} \mid E "+" E \mid E "-" E \mid E "*" E \mid E "/" E$

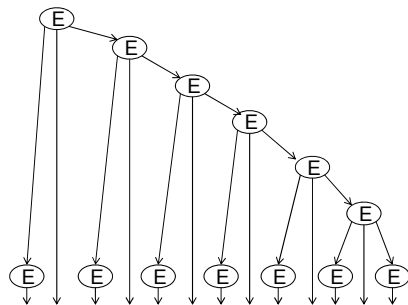


$65 * 74 - 68 + 25 * 5 / 3 + 16$

## Possible Parses

Grammar:

$E ::= \text{number} \mid E "+" E \mid E "-" E \mid E "*" E \mid E "/" E$

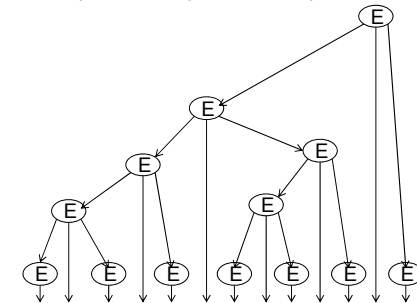


$65 * 74 - 68 + 25 * 5 / 3 + 16$

## Possible Parses

Grammar:

$E ::= \text{number} \mid E "+" E \mid E "-" E \mid E "*" E \mid E "/" E$



$65 * 74 - 68 + 25 * 5 / 3 + 16$

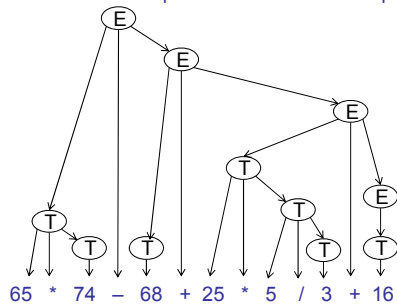
## Ambiguous Grammars

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- If a grammar allows multiple parses then we need to specify which we want (if it makes a difference)
- For example:  

$$\text{EXPR} ::= \text{TERM} \mid \text{TERM} "+" \text{EXPR} \mid \text{TERM} "-" \text{EXPR}$$

$$\text{TERM} ::= \text{number} \mid \text{number} "*" \text{TERM} \mid \text{number} "/" \text{TERM}$$



## Telling which option to follow

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$$\text{EXPR} ::= \text{TERM} \mid \text{TERM} "+" \text{EXPR} \mid \text{TERM} "-" \text{EXPR}$$

$$\text{TERM} ::= \text{number} \mid \text{number} "*" \text{TERM} \mid \text{number} "/" \text{TERM}$$

- Break into subrules, collecting the shared elements:

$$\text{EXPR} ::= \text{TERM} \text{ RESTOFEXPR}$$

$$\text{RESTOFEXPR} ::= "+" \text{EXPR} \mid "-" \text{EXPR} \mid \epsilon$$

$$\text{TERM} ::= \text{number} \text{ RESTOFTERM}$$

$$\text{RESTOFTERM} ::= "*" \text{TERM} \mid "/" \text{TERM} \mid \epsilon$$

( $\epsilon$  means "empty string")

- Transformations such as these can often turn a problematic grammar into a tractable grammar

## A more practical approach

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- Instead of  

$$\text{E} ::= \text{number} \mid \text{E} "+" \text{E} \mid \text{E} "-" \text{E} \mid \text{E} "*" \text{E} \mid \text{E} "/" \text{E}$$
- Write:  

$$\text{E} ::= \text{number} [\text{Op number}]^*$$

$$\text{Op} ::= "+" \mid "-" \mid "*" \mid "/"$$
- And the parser as:  

```

parseNum(s);
while (!s.hasNext(opPat)) {
    s.next();
    parseNum(s);
}

```

## A more practical approach

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- What about operator precedence: \* before +, etc "
- Grammar:  

$$\text{E} ::= \text{T} [ "+" | "-" ]^* \quad \text{Expression}$$

$$\text{T} ::= \text{F} [ "*" | "/" ]^* \quad \text{Term}$$

$$\text{F} ::= \text{number} \mid "(" \text{E} ")" \quad \text{Factor}$$
- Parser:  

```

public parseE(s) {
    parseT;
    while (!s.hasNext(addOpPat)) { // + or -
        s.next();
        parseT(s);
    }
}

```

## Next week: String searching

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- How can you find an occurrence (all occurrences) of a string  $s$  in a text  $t$ ?
- What is the cost?
- Can you make it faster??