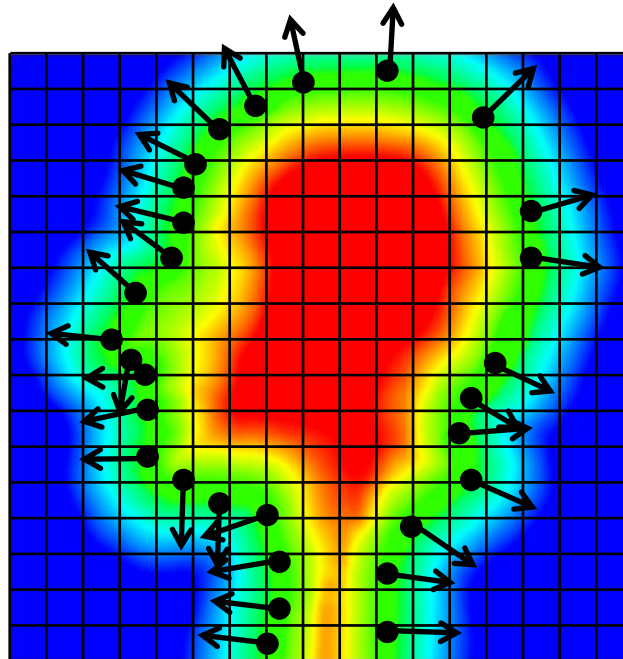


Assignment 2

Morton Codes & Hashbased Octrees

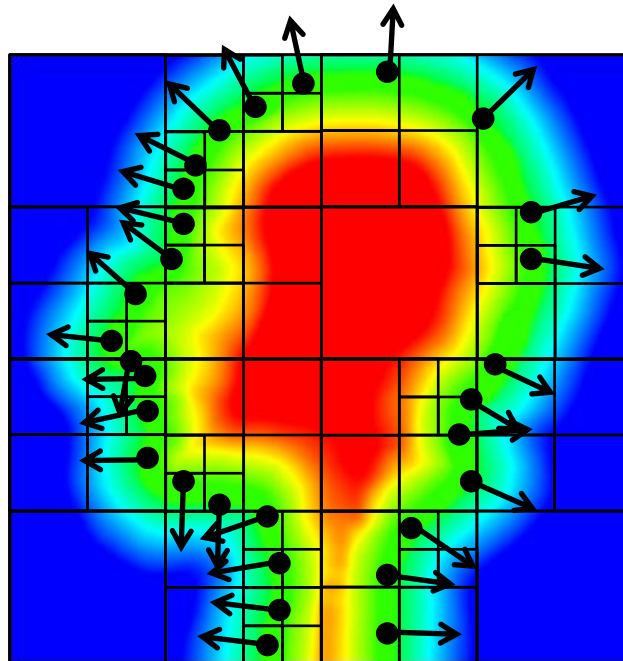
Octree

- Uniform grid inefficient
 - Cubic complexity (in 3D)
 - Many empty cells



Octree <http://en.wikipedia.org/wiki/Octree>

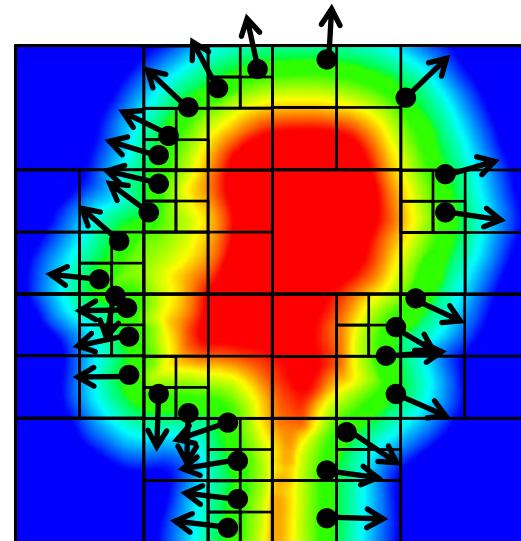
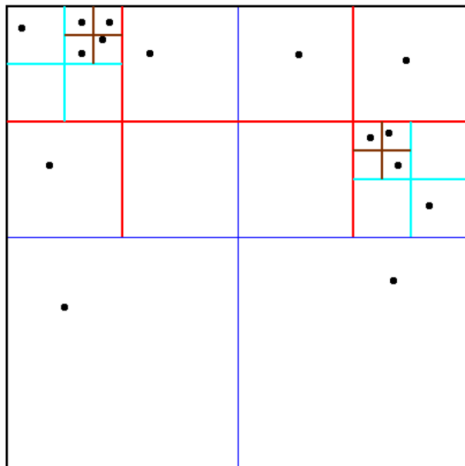
- Hierarchically split volume into 8 octants
- Visualization in 2D (quadtree)



Octree with maximum $M=5$ levels

Octree

- Applications
 - Acceleration structure
 - Next neighbor retrieval
 - Adaptive resolution



Octree <http://en.wikipedia.org/wiki/Octree>

Splitting criteria

1. Until maximum level
 - Number of maximum levels M
 - Resolution up to $2^{M-1} \times 2^{M-1} \times 2^{M-1}$
 - Typically $M=9$, larger for complex objects
2. Maintain minimum samples per node
 - Typically 1 to 5
 - For noisier inputs 5 to 10

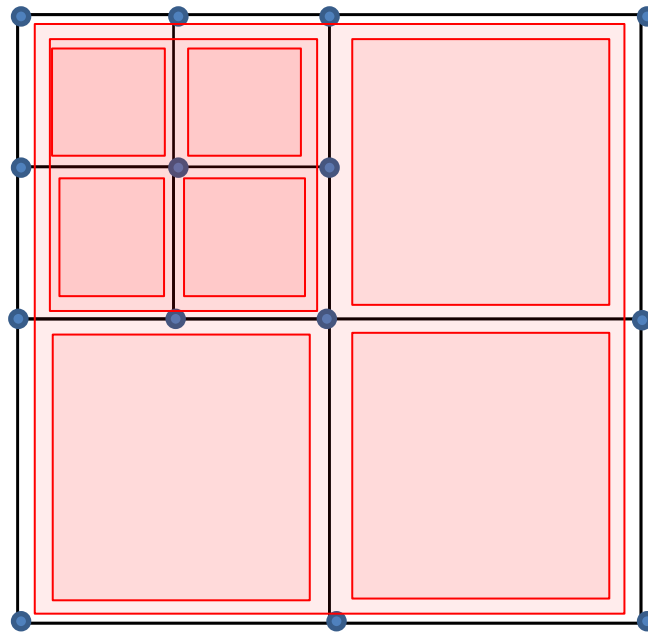
Hashtable based Octree

- Pointerless data structure with hashing
- Very fast & flexible navigation through the tree
- Following work by Lewiner et al.

http://zeus.mat.puc-rio.br/tomlew/pdfs/fastdualoctree_sgp.pdf

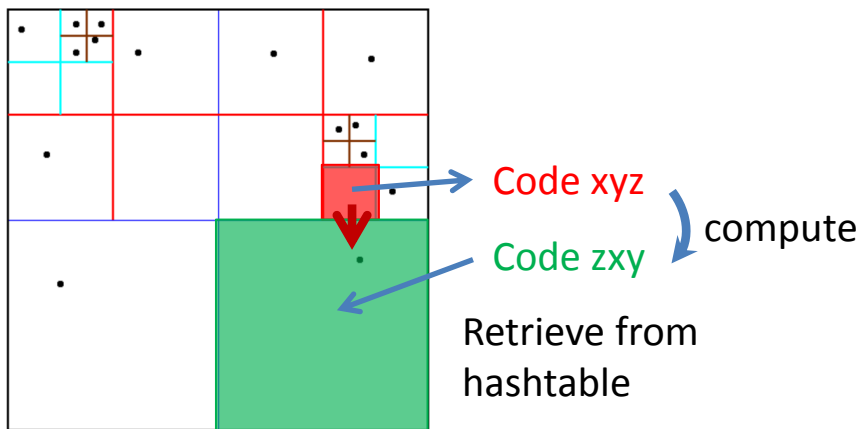
Hashbased Octree

- **Octree Cells** & **Octree Vertices** are stored in Hashtables, the hashes are given by morton codes

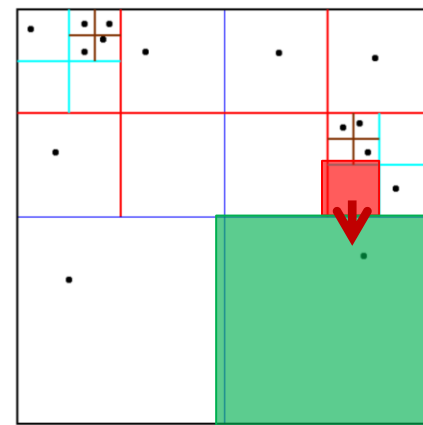


Hashbased octree

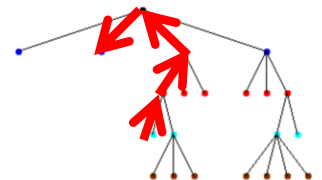
- Navigation: compute target hash and retrieve target cell instead of following links.



Hash based

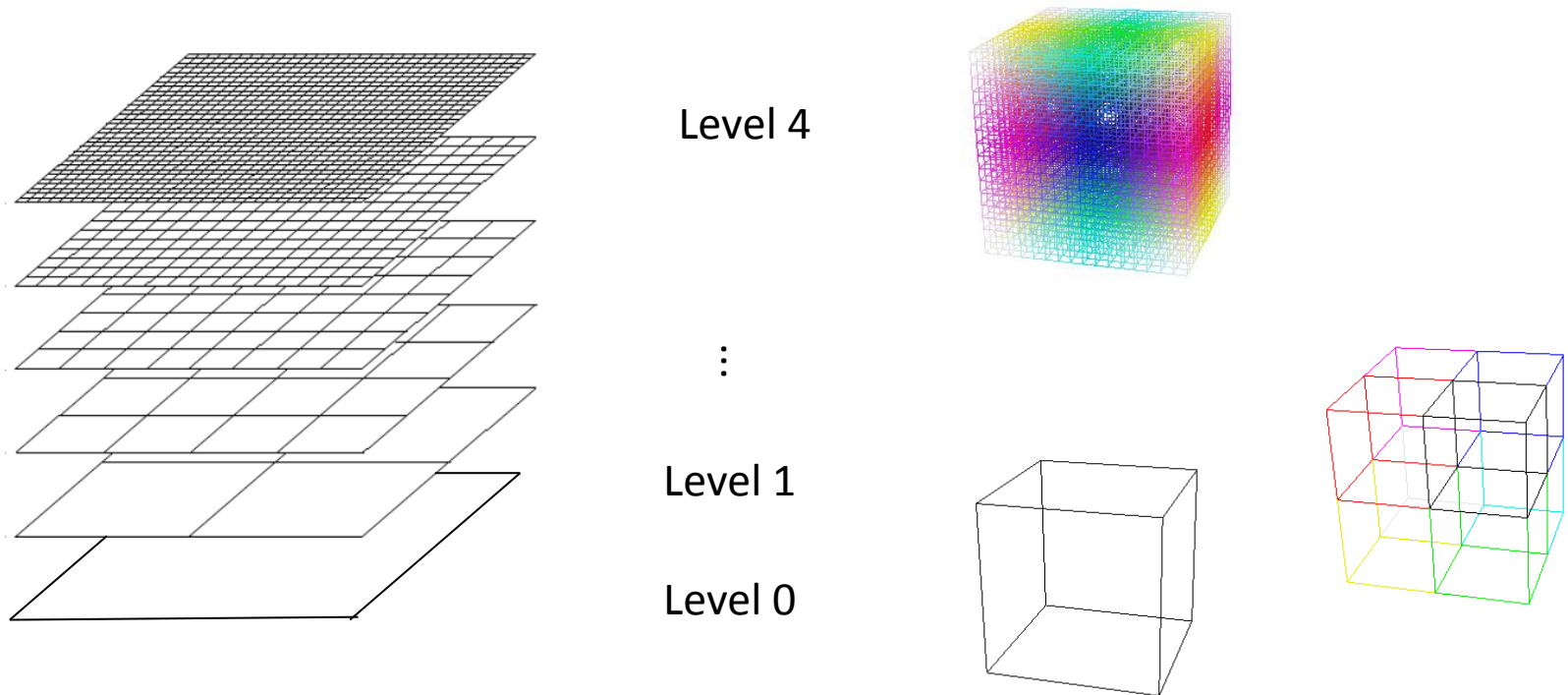


Link based



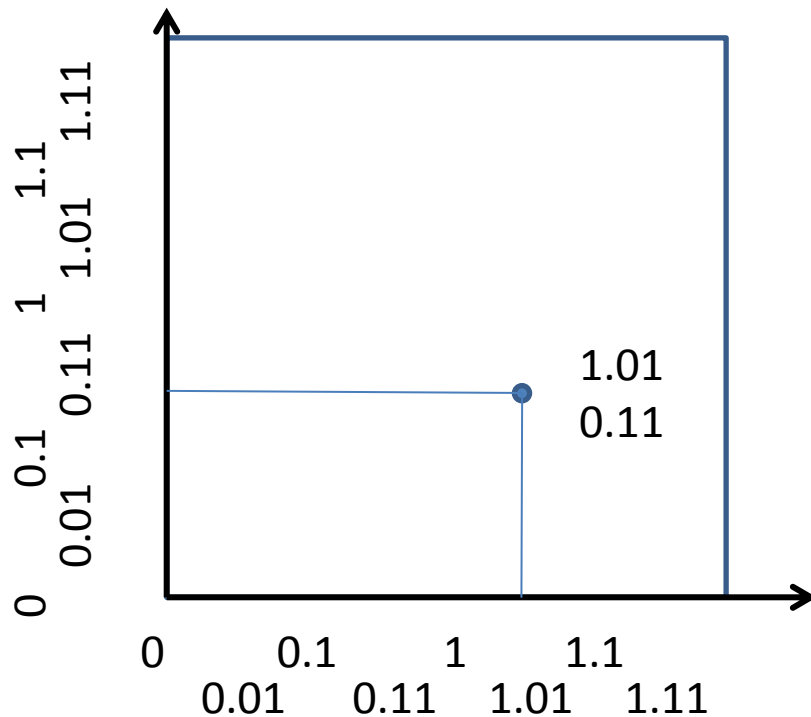
Morton Codes

- Integer that uniquely encodes level and position of a cell in a multigrid



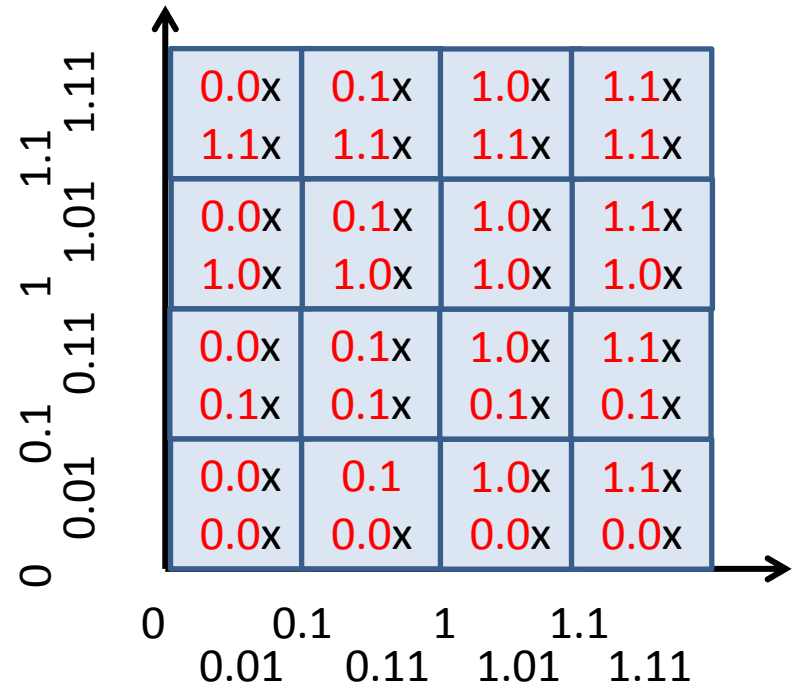
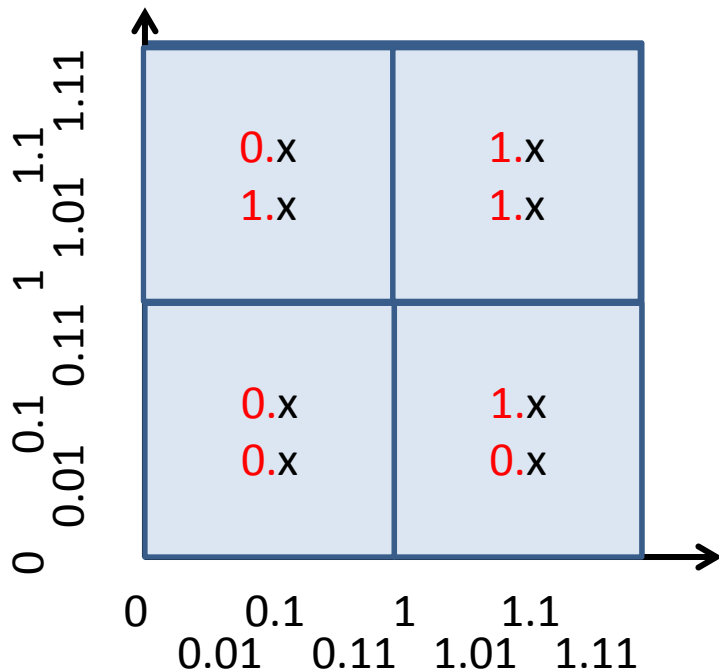
Morton codes

- Assign a coordinate system expressed with fractions in a binary system



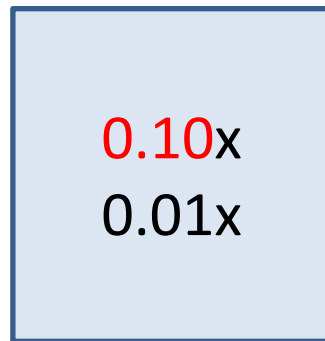
Morton codes

- Coordinates in a multigrid cell share first bits:



Morton codes

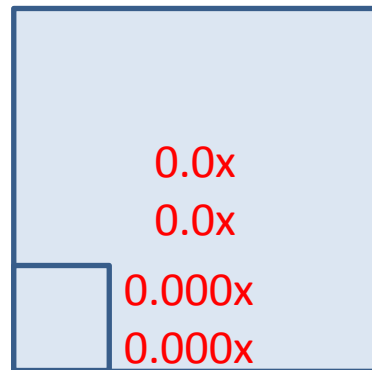
- Morton codes of multigrid cells:
 - Swizzeling coordinates + 1 delimiter
 - Our convention
 - 1 xy xy xy xy
 - 1 xyz xyz xyz xyz in 3D



1 00 10 01

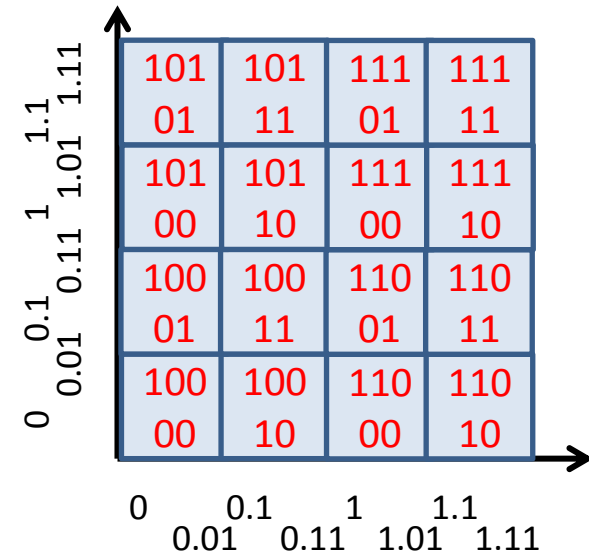
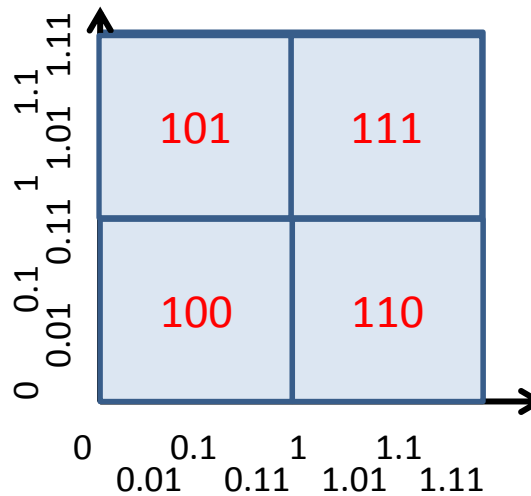
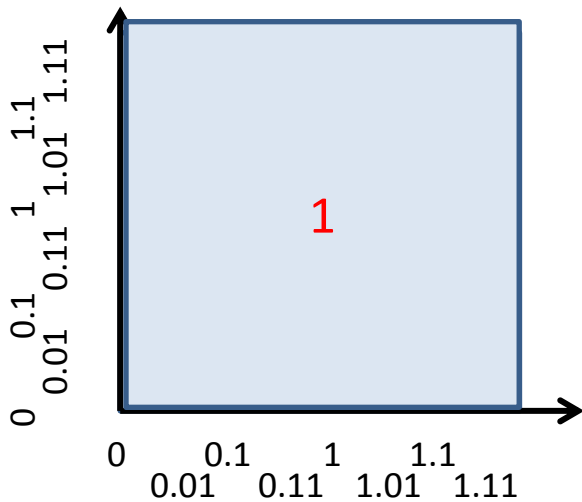
Morton codes

- 1 Delimiter:
 - 0000, 00000000 = 0
 - With delimiter: 10000 vs 100000000



Morton Codes

- Final Morton codes of the first three levels

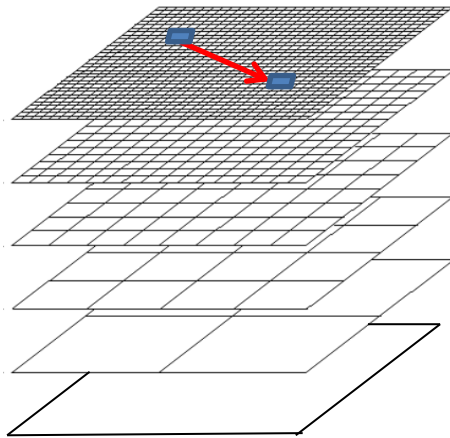


Navigation

- Via morton code manipulations
- Go to parent cell: simple bit shift
 - $C \gg 3$ in R^3 , $C \gg 2$ in R^2 .
- Go to neighbors....

Navigation („brute force“)

- From cell to cell, for a given difference vector



$$\text{red arrow} = \begin{pmatrix} 1.010 \\ -0.110 \end{pmatrix}$$

1 01 00 11 01 00

unswizzle

$$\begin{pmatrix} 0.0100 \\ 1.0110 \end{pmatrix} + \begin{pmatrix} 1.010 \\ -0.110 \end{pmatrix} = \begin{pmatrix} 1.1000 \\ 0.1010 \end{pmatrix}$$

reswizzle

1 10 11 10 01 00

- Unswizzling/reswizzling inefficient.

Navigation (good)

- Bitwise manipulations
- Dilated addition/subtraction:
 - Ignore the bits where a mask is 0.

$$x \bar{+} y = ((\underbrace{(x \mid \sim mask)}_{\text{Set masked bits 1}} + \underbrace{y \& mask}_{\text{Ignore masked bits}}) \& mask) \mid \underbrace{(x \& \sim mask)}_{\text{Reset other bits to original value}}.$$

Computes the correct value for the bits denoted by mask

$$x \bar{-} y = ((\underbrace{(x \& mask)}_{\text{Set masked bits 0}} - y \& mask) \& mask) \mid (x \& \sim mask).$$

- Intuition: masked bits are set to 1 with addition, to 0 in subtraction, such that added/subtracted bits carry to the next unmasked bit

Navigation (good)

- Example (try it out with pen and paper)

$$x = 0b\ 000\ \underline{011}\ \underline{111}$$

$$y = 0b\ \underline{111}$$

$$mask = 0b\ 100\ 100\ 100$$

$$x \oplus y = 0b\ \underline{000}\ \underline{111}\ \underline{011}$$

Legend: ignored bits

- When dropping the last term of the formula, the masked bits are set to zero

$$x \tilde{+} y = ((\underbrace{(x \mid \sim mask)}_{\text{Set masked bits 1}} + \underbrace{y \& mask}_{\text{Ignore masked bits}}) \& mask) \quad x \tilde{-} y = ((\underbrace{(x \& mask)}_{\text{Set masked bits 0}} - y \& mask) \& mask).$$

Computes the correct value for the bits denoted by mask

$$x = 0b\ 000\ 011\ 111$$

$$y = 0b\ 111$$

$$mask = 0b\ 100\ 100\ 100$$

$$x \tilde{+} y = 0b\ \underline{000}\ \underline{100}\ \underline{000}$$

Navigation (good)

- For the navigation, manipulate x bits, y bits and z bits separately:

xyz xyz xyz
 $x = 0b1\ 000\ 101\ 111$
 $y = 0b011$
xyz xyz xyz
 $mask_x = 0b\ 100\ 100\ 100$
 $mask_y = 0b\ 010\ 010\ 010$
 $mask_z = 0b\ 001\ 001\ 001$

Addition


$$\begin{aligned} Result = & (((x \sim mask_x) + (y \& mask_x)) \& mask_x) | \\ & (((x \sim mask_y) + (y \& mask_y)) \& mask_y) | \\ & (((x \sim mask_z) + (y \& mask_z)) \& mask_z) \end{aligned}$$

Subtraction

$$\begin{aligned} Result = & (((x \& mask_x) - (y \& mask_x)) \& mask_x) | \\ & (((x \& mask_y) - (y \& mask_y)) \& mask_y) | \\ & (((x \& mask_z) - (y \& mask_z)) \& mask_z) \end{aligned}$$

Navigation (good)

- To navigate through morton codes:
 - Use swizzled difference vector (no delimiter)

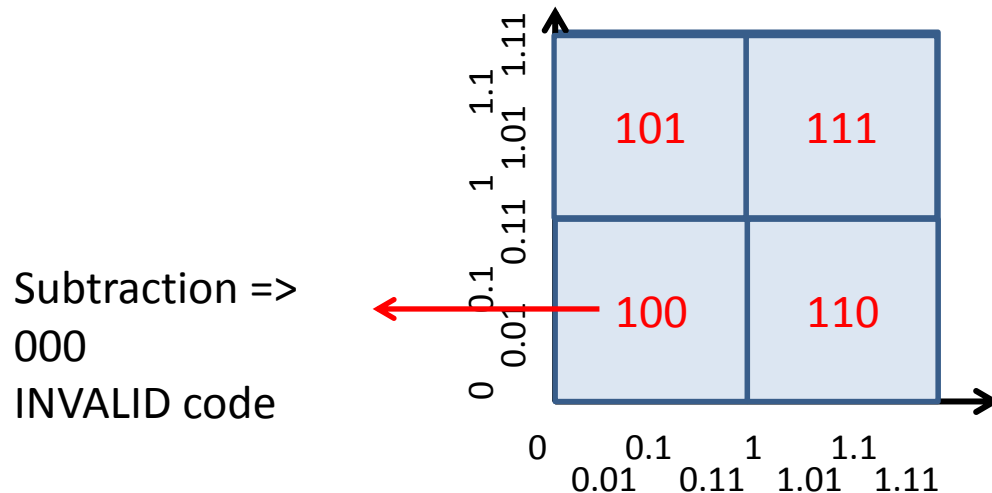

$$y = \begin{pmatrix} 1.010 \\ 0.110 \end{pmatrix} \quad 10 \ 01 \ 11 \ 00$$

- And use bit magic to modify xyz bits of a morton code separately.

$$\begin{aligned} Result = & (((x \sim mask_x) + (y \& mask_x)) \& mask_x) | \\ & (((x \sim mask_y) + (y \& mask_y)) \& mask_y) | \\ & (((x \sim mask_z) + (y \& mask_z)) \& mask_z) \end{aligned}$$

Check for Overflows!


- A level k morton code \pm a difference vector should produce a level k Morton code.



- $C = 0b1\ xyz\ xyz\ \dots\ xyz$ is a level k Morton code iff $(C \gg (3*k)) == 0b1$

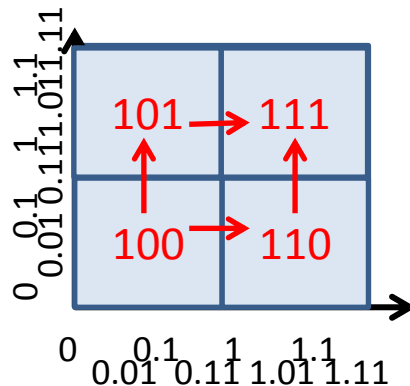
Separate Addition and Subtraction

- Addition & subtraction need to be handled separately with dilated addition / subtraction

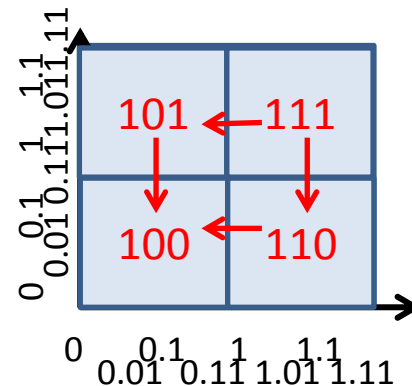


$$y = \begin{pmatrix} 1.010 \\ -0.110 \end{pmatrix}$$

- In the Java code: write separate methods to compute positive & negative neighbors



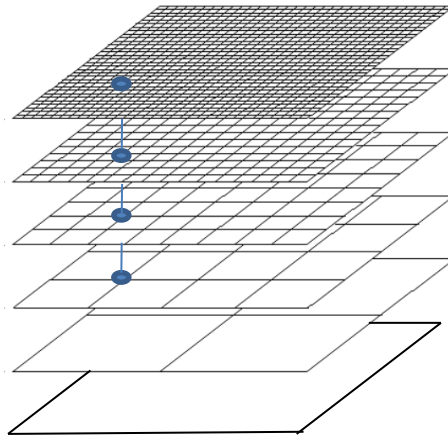
Positive neighbor



Negative neighbor

Vertex Codes

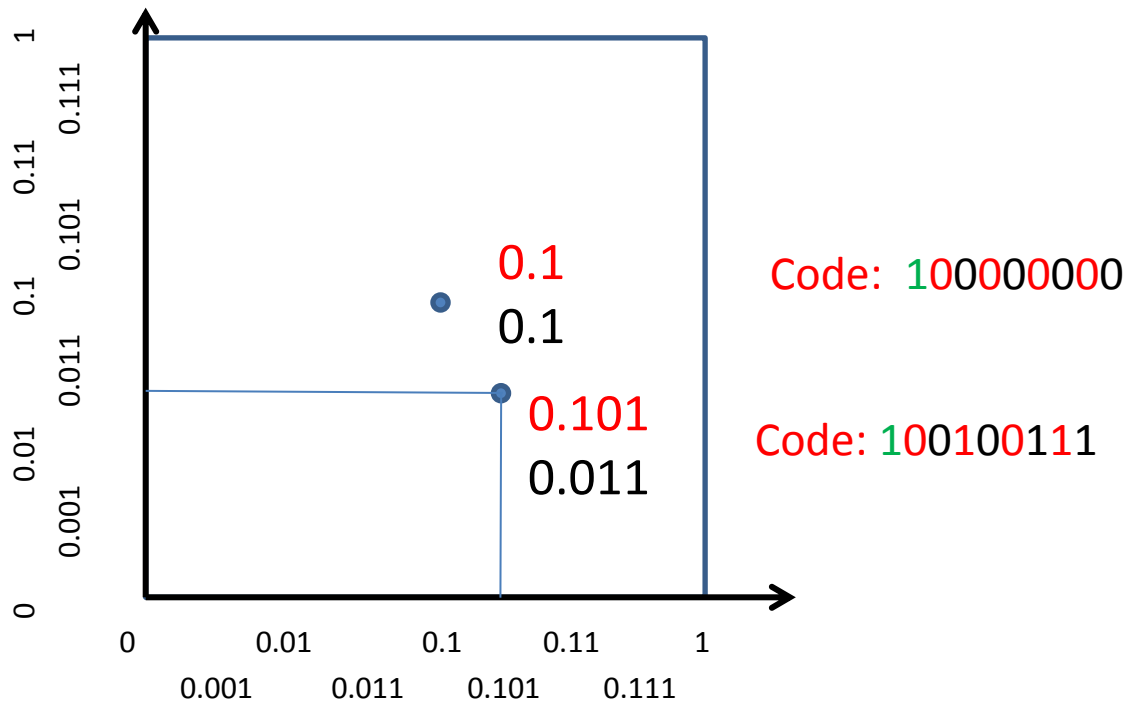
- Vertices do not belong to one fixed level, but to all levels above some level.



- To get a unique code: Restrict maximal multigrid layer and use highest level Morton codes for all vertices.
 - Equal to padding lower level code with zeros

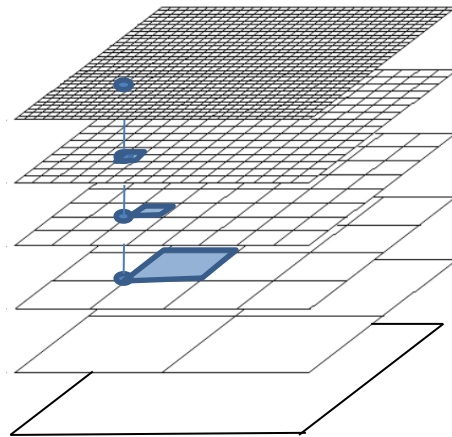
Vertex Codes

- Vertex code intuition: 1 prefix, swizzle together coordinates & pad to the maximal length.



Vertex Codes (Navigation)

- Vertex -> cell and cell -> vertex:
 - Unpadded vertex code = cell code
 - Padded cell code = vertex code



Vertex code

0b1 00 11 00 00 00

Cell code

0b1 00 11 00 00 00

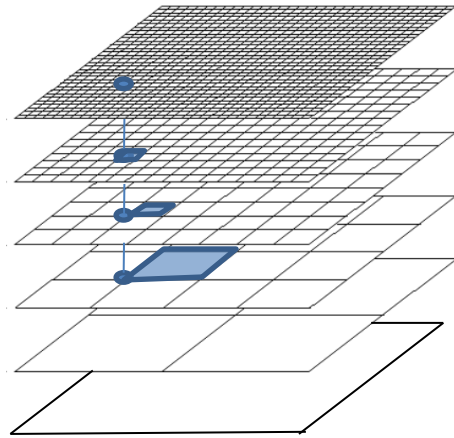
0b1 00 11 00 00

0b1 00 11 00

0b1 00 11

Vertex Codes (Navigation)

- Navigation:
 - Unpad code to appropriate cell level
 - Cell to cell navigation
 - Repad



Vertex code

0b1 00 11 00 00 00

Cell code

0b1 00 11 00 00 00

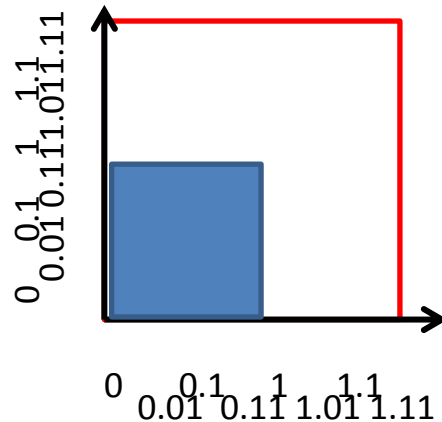
0b1 00 11 00 00

0b1 00 11 00

0b1 00 11

Hashbased Octree

- Octree: subset of multigrid
- Problem:
 - because of 1-delimiter: we cannot assign vertex codes on the red boundary



- Solution: use the codes in first quadrant only.

Codes in Hashbased Octree

	Morton Code:	level
root cell	0b1000	1
general cell	0b1 000 $xyz\ xyz \cdots xyz$	k
	$\underbrace{\hspace{10em}}_{3k\ bits}$	

	Morton Code:	maxLevel	minLevel
vertex	0b1 $xyz \cdots xyz$ 000 \cdots 000	$\max_{vertex \in cell} cell.level$	$\min_{vertex \in cell} cell.level$
	$\underbrace{\hspace{10em}}_{3 \cdot minLevel\ bits}$		
	$\underbrace{\hspace{10em}}_{tree.depth \cdot 3\ bits}$		

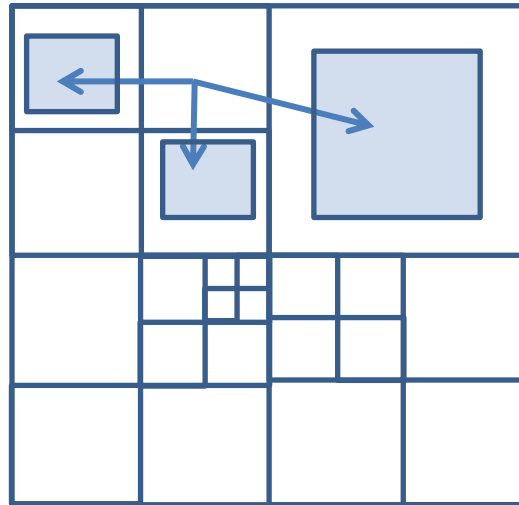
Navigation in the Hashoctree

Find neighbors

1. Compute theoretical Morton code of neighbor
2. Check existence of the associated vertex or cell, if inexistent, decide:
 - retrieve next existing vertex/cell, or
 - return null.

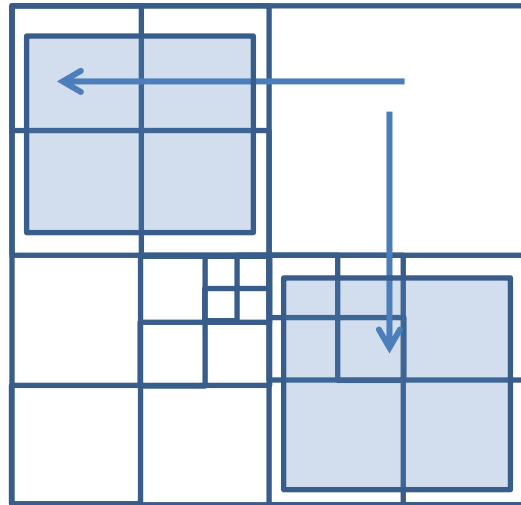
Navigation cell-> cell

- Neighbor cells: find a neighbor cell of less or equal level



Navigation cell-> cell

- Neighbor cells: find a neighbor cell of less or equal level

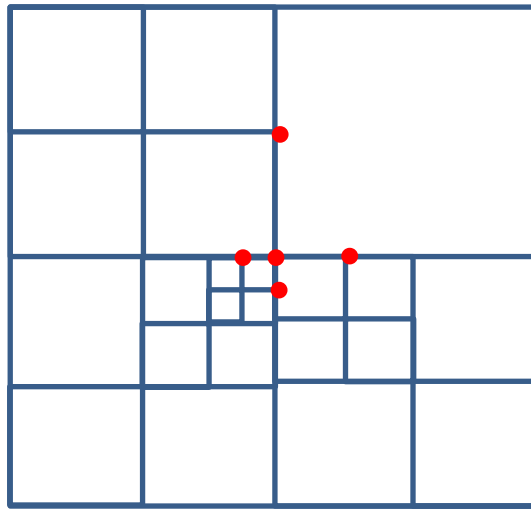


Navigation cell-> cell

- Compute the neighbor code on the same level
- Iterate until the level 0 is reached:
 - Check if the cell exists
 - If not: go to the parent.

Navigation Vertex -> Vertex

- Neighbors: closest vertex in some direction



- Unpad to finest level cell code, compute neighbor cell code, repad, check existence. Iterate up to minimum level of vertex

Bit Arithmetic

- \sim not. $|$ or. $>>$ rightshift. $<<$ leftshift
- 64 bit long:
 - First bit decides about positive, negative.
 - Negatives are coded as \sim (positive - 1).
 - $-4 = \sim(3) = \sim(0b000\dots0011) = 0b111\dots1100$
 - $-1 = 0b(111111\dots111)$
 - $\sim(-1 << k) = 0b(00\dots011111)$ can be a handy mask

Summary

- Dilated Operations:

$$x \tilde{+} y = ((\underbrace{(x \mid \sim mask)}_{\text{Set masked bits 1}} + \underbrace{y \& mask}_{\text{Ignore masked bits}}) \& mask)$$

Computes the correct value for the bits denoted by mask

$$x \tilde{-} y = ((\underbrace{(x \& mask)}_{\text{Set masked bits 0}} - y \& mask) \& mask).$$

$$\begin{aligned} Result &= (((x \mid \sim mask_x) + (y \& mask_x)) \& mask_x) \mid \\ &(((x \mid \sim mask_y) + (y \& mask_y)) \& mask_y) \mid \\ &(((x \mid \sim mask_z) + (y \& mask_z)) \& mask_z) \end{aligned}$$

- Morton Codes:

	Morton Code:	level
root cell	0b1000	1
general cell	0b1 000 $\underbrace{xyz \ xyz \ \cdots \ xyz}_{3k \text{ bits}}$	k

	Morton Code:	maxLevel	minLevel
vertex	0b1 $\underbrace{xyz \ \cdots \ xyz}_{3 \cdot \minLevel \text{ bits}}$ 000 \cdots 000	$\max_{vertex \in cell} cell.level$	$\min_{vertex \in cell} cell.level$
	$\underbrace{\hspace{10em}}_{tree.depth \cdot 3 \text{ bits}}$		

Summary

- I am paid to help
- And use the forum 😊

Questions?