

# Hanzo DSO: Decentralized Semantic Optimization with Byzantine-Robust Prior Aggregation

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October 2025

## Abstract

We present **Hanzo DSO** (Decentralized Semantic Optimization), a protocol for sharing and aggregating experiential priors across distributed language model agents without parameter updates. DSO enables zero-training model adaptation by (i) extracting semantic advantages via local TF-GRPO, (ii) compressing priors with 1-bit quantization ( $29.5\times$  savings), (iii) aggregating contributions via byzantine-robust median voting, and (iv) synchronizing via IPFS/Arweave with on-chain registry verification. Key contributions: (a) Byzantine-robust aggregation scheme with stake-weighted quality scores, (b) Content-addressed experience registry with Merkle proofs, (c) Gossip-based P2P sync optimized for high-quality shard propagation, and (d) Integration with Hanzo Network’s PoAI consensus for attestation-based rewards. We demonstrate **15.2% improvement** in multi-agent code generation tasks when nodes share priors vs. isolated operation.

## 1 Introduction

Traditional federated learning aggregates gradients or parameters, requiring compute-intensive operations and risking catastrophic interference. Recent work on parameter-efficient adaptation (LoRA, adapters) reduces overhead but still requires training. In-context learning enables zero-shot adaptation but lacks systematic mechanisms for sharing knowledge across agents.

**Our Contribution.** Hanzo DSO introduces a decentralized protocol for sharing *experiential priors*—compressed semantic advantages extracted from agent rollouts—enabling collective intelligence without parameter updates.

## 2 System Architecture

### 2.1 Roles and Components

- **Agents:** Run local ASO (TF-GRPO), extract and compress priors
- **Registry:** On-chain smart contract storing CIDs, Merkle roots, quality scores
- **Storage:** Off-chain IPFS/Arweave for prior data
- **Aggregators:** Compute byzantine-robust median under fixed schema
- **Validators:** Verify attestations via PoAI, slash malicious submissions

## 3 Decentralized Semantic Optimization (DSO)

### 3.1 Experience Priors

Each agent/node maintains an *experience prior*  $E$ : token/embedding-level memory distilled from rollouts. Locally, nodes run **Active Semantic Optimization (ASO)** to extract *semantic advantages* from groups of rollouts (TF-GRPO). Priors are compressed (§5) and written to the on-chain *ExperienceRegistry* with Merkle proofs.

### 3.2 Training-Free GRPO as Bayesian PoE

For a base model with conditional  $p_\theta(y \mid x)$  and a set of experiences  $\{e_k\}$  mapping to token-level factors  $\phi_k(y \mid x)$ , decoding uses a *product-of-experts*:

$$p(y \mid x, E) \propto p_\theta(y \mid x) \prod_k \phi_k(y \mid x)^{\alpha_k}, \quad \alpha_k \geq 0. \quad (1)$$

Here  $\phi_k$  are distilled from group-relative semantic advantage; weights  $\alpha_k$  are learned by introspective calibration without gradient updates to  $\theta$ .

### 3.3 Distributed Aggregation

Hanzo Network aggregates *priors, not gradients*. Let node priors be  $\{E_i\}$ . We publish hashes and quality scores; the chain computes a byzantine-robust aggregate  $\hat{E} = \text{median}_q\{E_i\}$  under a fixed schema (token bins / embedding centroids). Conflicting contributions resolve by stake-weighted quorum plus quality caps.

## 4 ExperienceRegistry and P2P Sync

**Registry.** On-chain contract stores: content-addressed CID, Merkle root, schema version, quality vector, submitter, slashing bond. **Storage.** Off-chain IPFS/Arweave; local SQLite+LanceDB with Merkle verification. **Sync.** Gossip protocol with CRDT merge; priority given to high-quality shards (fee rebates bias peers to propagate them).

## 5 1-Bit Semantic Compression

Inspired by BitDelta, we store only the *signs* of per-bucket deltas plus per-matrix scales. For an experience matrix  $\Delta \in \mathbb{R}^{n \times m}$ ,

$$\hat{\Delta} = \alpha \text{Sign}(\Delta), \quad \alpha = \frac{1}{nm} \sum_{ij} |\Delta_{ij}|. \quad (2)$$

Scales are distilled by matching logits to a teacher rollout. We observe  $\approx 29.5\times$  storage savings with negligible loss in downstream utility, enabling multi-tenant caching and rapid hot-swaps of personalizations.

## 5.1 Compression Algorithm

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**Algorithm 1** BitDelta-Inspired Compression

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- 1: **input:** full-precision experience matrix  $\Delta \in \mathbb{R}^{n \times m}$
  - 2: Compute average absolute value:  $\alpha = \frac{1}{nm} \sum_{ij} |\Delta_{ij}|$
  - 3: Quantize:  $\hat{\Delta}_{ij} = \alpha \cdot \text{sign}(\Delta_{ij})$
  - 4: Store: binary signs  $\{\text{sign}(\Delta_{ij})\}$  and scalar  $\alpha$
  - 5: **return** compressed representation:  $(\{\text{sign}(\Delta_{ij})\}, \alpha)$
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## 5.2 Decompression and Application

At inference time:

$$\Delta_{ij}^{\text{approx}} = \alpha \cdot \text{sign}(\Delta_{ij}), \quad (3)$$

yielding a 1-bit per element representation plus one scalar per matrix block. This enables efficient storage and rapid loading of experience priors.

## 6 Byzantine-Robust Aggregation

### 6.1 Threat Model

Adversaries may submit:

- **Random noise:** Low-quality priors to pollute aggregate
- **Targeted attacks:** Priors designed to degrade specific tasks
- **Sybil attacks:** Multiple identities voting for malicious priors

### 6.2 Median Voting Under Schema

Fix a schema  $\mathcal{S}$ : token bins or embedding centroids. Each submission  $E_i$  is decomposed:

$$E_i = \{\Delta_i^{(s)}\}_{s \in \mathcal{S}}, \quad (4)$$

where  $\Delta_i^{(s)}$  is the advantage for schema element  $s$ . The aggregate computes:

$$\tilde{E}^{(s)} = \text{weighted-median}_i(\Delta_i^{(s)}; w_i), \quad (5)$$

with weights  $w_i$  proportional to stake times quality score  $q_i$ .

### 6.3 Quality Scoring

Quality scores  $q_i \in [0, 1]$  are estimated via:

1. **Holdout validation:** Test priors on reserved benchmarks
2. **Cross-validation:** Peer nodes sample and verify
3. **Reputation:** Historical performance on-chain

Quality below threshold  $q_{\min}$  (default 0.3) triggers bond slashing.

## 7 ExperienceRegistry Contract

### 7.1 Interface

```
interface IExperienceRegistry {
    struct Entry {
        bytes32 merkleRoot;
        string cid;           // IPFS/Arweave
        uint64  schema;
        uint64  quality;      // quantized [0, 1000]
        address submitter;
        uint256 bond;
    }
    function submit(Entry calldata e) external payable
        returns (uint256 id);
    function voteQuality(uint256 id, uint64 score) external;
    function slash(uint256 id, address challenger,
        bytes calldata proof) external;
    function aggregate(uint256[] calldata ids) external
        returns (bytes32 aggregateMerkleRoot);
}
```

### 7.2 Bond and Slashing

Submission requires bond  $D$  (default 25 \$AI). If  $q_i < q_{\min}$ , challenger submits proof (counter-examples); successful challenge burns  $\sigma D$  (default  $\sigma = 0.5$ ), refunds  $(1 - \sigma)D$  to challenger.

## 8 P2P Synchronization

### 8.1 Gossip Protocol

Nodes maintain local replica of high-quality priors:

1. Subscribe to registry events (new submissions, quality updates)
2. Fetch CID from IPFS/Arweave if  $q_i \geq q_{\text{fetch}}$  (default 0.5)
3. Verify Merkle proof against on-chain root
4. Merge into local LanceDB with CRDT semantics

### 8.2 Incentive Alignment

Nodes that propagate high-quality priors earn fee rebates:

$$\text{rebate}_i \propto \sum_{j:\text{fetched from } i} q_j \cdot \text{size}_j. \quad (6)$$

## 9 Integration with PoAI

### 9.1 Attestation-Based Rewards

Each prior submission includes PoAI attestation (TEE report + task metrics):

- $\Delta I$ : Information gain from experience
- $\Delta U$ : Utility improvement on held-out tasks
- Cost metrics: compute, bandwidth, energy

Emissions formula (see HMM paper):

$$R_i = \gamma \Delta I + \beta \Delta U - \lambda_c \cdot \text{cost}_i, \quad (7)$$

where  $\gamma = 1.0$ ,  $\beta = 0.5$ ,  $\lambda_c = 0.1$  (defaults).

## 10 Experimental Evaluation

### 10.1 Multi-Agent Code Generation

Configuration	Resolved Rate	Avg. Prior Reuse
Isolated agents (no DSO)	16.3%	0.0
DSO (Byzantine honest)	18.8%	3.2 priors/task
DSO (20% Byzantine)	17.9%	2.8 priors/task
<b>DSO (median voting)</b>	<b>18.7%</b>	<b>3.1 priors/task</b>

Table 1: SWE-bench Verified (500 issues), 10 agents. Byzantine nodes submit random priors.

### 10.2 Storage Efficiency

- Full-precision priors: 2.4 GB/agent
- 1-bit compressed: 82 MB/agent (29.3× savings)
- IPFS overhead: +12% (metadata, proofs)
- Effective:  $\approx 26\times$  savings

## 11 Related Work

**Federated learning:** FedAvg, FedProx, byzantine-robust aggregation. **Decentralized ML:** Swarm learning, peer-to-peer training. **Parameter-efficient adaptation:** LoRA, adapters, prompt tuning. **Blockchain + ML:** Federated learning on blockchain, decentralized model markets.

## 12 Conclusion

Hanzo DSO enables decentralized, zero-training model adaptation by sharing compressed experiential priors. Byzantine-robust aggregation ensures resilience against adversarial nodes, while integration with PoAI provides attestation-based incentives. Future work includes cross-domain transfer (code  $\rightarrow$  data science) and hierarchical aggregation (specialized sub-groups).

## A Security Analysis

### A.1 Sybil Resistance

Stake-weighting limits influence of Sybil identities. With  $N$  honest nodes and  $M < N/2$  Sybil identities (each with stake  $s$ ), median voting ensures honest aggregate if total honest stake  $\geq$  total malicious stake.

### A.2 Data Poisoning

1-bit quantization limits information content per prior, reducing attack surface. Median voting filters extreme values. Quality scoring enables post-hoc detection.

*Disclaimer.* This document describes a proposed protocol. Security properties require formal verification.