

On Artin's Conjecture: Pairs of Additive Forms

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Let $f(x_1, \dots, x_s) \in \mathbb{Z}[x_1, \dots, x_s]$ be a form (homogeneous polynomial) of degree k . The equation $f(\mathbf{x}) = 0$ has a non-trivial solution $\mathbf{x} \in \mathbb{Q}_p^s$ for all primes p provided that

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- $k = 1$,
- $k = 2$ (Meyer),
- $k = 3$ (Lewis).

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Let $f_1, \dots, f_r \in \mathbb{Z}[x_1, \dots, x_s]$ be forms of degree (k_1, \dots, k_r) . The equations $f_1 = \dots = f_r = 0$ have a non-trivial solution $\mathbf{x} \in \mathbb{Q}_p^s$ for all primes p provided that

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- $r = 2, k_1 = k_2 = 2$ (Dem'yanov).

How close is Artin's conjecture to the truth?

Counter examples

- 1966: Terjanian: There is a form of degree 4 in 18 variables with no non-trivial 2-adic solution.
- 1966: Browkin: For all p there exist forms without a non-trivial p -adic solution violating Artin's conjecture.
- 1981: Arkhipov and Karatsuba: Not even true if $s > k^n$ for any fixed $n \in \mathbb{N}$.

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Positive Indications

- Degree: All counter examples are of even degree k .
- Primes: 1965: Ax and Kochen: For a fixed degree k one has a non-trivial p -adic solution for all but finitely many p .
- Forms: 1963: Davenport and Lewis: Artin's conjecture holds for all additive forms.

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A form $f \in \mathbb{Z}[x_1, \dots, x_s]$ is called additive if $f(x_1, \dots, x_s) = a_1 x_1^k + a_2 x_2^k + \dots + a_s x_s^k$.

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Generalisation to system of additive forms:

- 1983: Lewis and Montgomery: Not true for all r -tuples (k_1, k_2, \dots, k_r) .
- 2015: Wooley: For $r = 2$ there are already counterexamples.

$$r = 2 \text{ and } k_1 = k_2$$

Let f, g be two additive forms with integer coefficients in s variables of degree k .

Question

How big does s have to be to ensure a non-trivial p -adic solution $f(\mathbf{x}) = g(\mathbf{x}) = 0$ for all primes p ?

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for $\tau \geq 1$. Else $s \geq 8k^2$, $s \geq \frac{8}{3}k^2$ and $s \geq 4k^2$ variables, respectively, are sufficient.

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- 2011, 2013: Godinho, de Souza Neto: For $k = p^\tau (p - 1)$ it is sufficient if $s > 2 \frac{p}{p-1} k^2 - 2k$ holds and either $p \in \{3, 5\}$ or $p \geq 7$ and $\tau \geq \frac{p-1}{2}$

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- 2013: Godinho, Knapp and Rodrigues: Expected bound holds for $k = 6$.

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- 2017: Godinho, Ventura: Expected bound hold for $k = 3^\tau \cdot 2$.

Theorem (K.)

For $p \geq 5$ prime, $\tau \geq 1$ and $k = p^\tau(p-1)$ the pair of additive forms with integers coefficients

$$\sum_{j=1}^s a_j x_j^k = \sum_{j=1}^s b_j x_j^k = 0$$

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Missing values of k :

- 2^τ with $2 \leq \tau \leq 15$
- $3 \cdot 2^\tau$ with $2 \leq \tau (\leq 15)$

Finding p -adic solutions

Let $f(x_1, \dots, x_s) = \sum_{i=1}^s a_i x_i^k$ and $g(x_1, \dots, x_s) = \sum_{i=1}^s b_i x_i^k$.

For $k = p^\tau k_0$ with $\gcd(p, k_0) = 1$ define

$$\gamma := \begin{cases} 1, & \text{if } \tau = 0 \\ \tau + 1, & \text{if } \tau > 0 \text{ and } p > 2 \\ \tau + 2, & \text{if } \tau > 0 \text{ and } p = 2. \end{cases}$$

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Hensel's Lemma

If the congruences

$$\sum_{i=1}^s a_i x_i^k \equiv 0 \pmod{p^\gamma}, \quad \sum_{i=1}^s b_i x_i^k \equiv 0 \pmod{p^\gamma}$$

have a solution in the integers for which the matrix

$$\begin{pmatrix} a_1 x_1 & a_2 x_2 & \dots & a_s x_s \\ b_1 x_1 & b_2 x_2 & \dots & b_s x_s \end{pmatrix}$$

has rank 2 modulo p , then the pair of forms f, g has a non-trivial p -adic solution.

Equivalence relation defined via the operations:

- $f' = f(p^{\nu_1}x_1, \dots, p^{\nu_s}x_s)$, $g' = g(p^{\nu_1}x_1, \dots, p^{\nu_s}x_s)$ for integers ν_i
- $f'' = \lambda_1 f + \lambda_2 g$, $g'' = \mu_1 f + \mu_2 g$ for $\lambda_1, \lambda_2, \mu_1, \mu_2 \in \mathbb{Q}$ with $\lambda_1\mu_2 - \lambda_2\mu_1 \neq 0$.

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Definition

A p -normalised pair f, g is a pair of forms with integer coefficients and $\vartheta(f, g) \neq 0$, where the power of p dividing $\vartheta(f, g)$ is as small as possible among all pairs of forms in the same p -equivalent class.

Definition

A variable x_i is called **at level l** if $\begin{pmatrix} a_i \\ b_i \end{pmatrix} = p^l \begin{pmatrix} \tilde{a}_i \\ \tilde{b}_i \end{pmatrix}$ and $p \nmid \begin{pmatrix} \tilde{a}_i \\ \tilde{b}_i \end{pmatrix}$. The vector $\begin{pmatrix} \tilde{a}_i \\ \tilde{b}_i \end{pmatrix}$ is called the **level coefficient vector** of x_i .

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Define

$$\mathcal{L}_0 := \left\{ c \begin{pmatrix} 1 \\ 0 \end{pmatrix} \mid 1 \leq c \leq p-1 \right\} \quad \text{and} \quad \mathcal{L}_\nu := \left\{ c \begin{pmatrix} \nu \\ 1 \end{pmatrix} \mid 1 \leq c \leq p-1 \right\}$$

for all $1 \leq \nu \leq p$. One says that the variable x_i **is of colour** ν , if $\begin{pmatrix} \tilde{a}_i \\ \tilde{b}_i \end{pmatrix} \in \mathcal{L}_\nu \pmod{p}$.

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- Let q_i be the number of variables at level i which are not in the biggest colour.

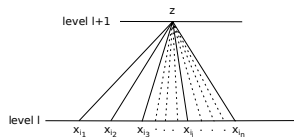
$$m_0 + \cdots + m_{j-1} + q_j \geq \frac{(j + \frac{1}{2})s}{k} \quad \text{for} \quad j = 0, 1, \dots, k-1$$

Contraction

Goal:

$$\sum_{i=1}^s a_i x_i^k \equiv 0 \pmod{p^\gamma}, \quad \sum_{i=1}^s b_i x_i^k \equiv 0 \pmod{p^\gamma} \quad (\text{solve})$$

$$\begin{pmatrix} a_1 x_1 & a_2 x_2 & \dots & a_s x_s \\ b_1 x_1 & b_2 x_2 & \dots & b_s x_s \end{pmatrix} \quad (\text{rank 2 modulo } p)$$



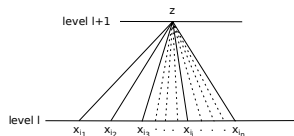
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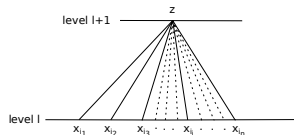
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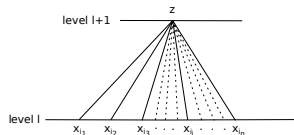


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Definition

Let x_{i_1}, \dots, x_{i_n} variables at level l and $y_1, \dots, y_n \in \mathbb{Z} \setminus p\mathbb{Z}$ such that

$$\sum_{j=1}^n a_{ij} y_j^k \equiv \sum_{j=1}^n b_{ij} y_j^k \equiv 0 \pmod{p^{l+1}},$$

then the variables x_{i_j} **contract** to a variable at level at least $l + 1$.

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Davenport and Lewis: One can construct at least $\min \left(\left\lfloor \frac{m_0}{2^{p-1}} \right\rfloor, \left\lfloor \frac{q_0}{p} \right\rfloor \right)$ primary variables at level at least 1.

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- Lift the primary variables to higher levels.

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→ Gives at best the bound $s > 2 \frac{p}{p-1} k^2 - 2k$.

Lemma

A set of $2p - 2$ variables at level l ($p - 1$ of some colour ν and $p - 1$ not of colour ν) contract together with a primary variable at level at least l to a primary variable at level at least $l + 1$.

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→ Take unused variables at lower levels to create helpful variables along the way.

Define

$$\mathcal{L}_{0\mu} := \left\{ c \left(\begin{pmatrix} 1 \\ 0 \end{pmatrix} + \mu \begin{pmatrix} 0 \\ p \end{pmatrix} \right) \mid c \in (\mathbb{Z}/p^2\mathbb{Z})^* \right\}$$
$$\mathcal{L}_{\nu\mu} := \left\{ c \left(\begin{pmatrix} \nu \\ 1 \end{pmatrix} + \mu \begin{pmatrix} p \\ 0 \end{pmatrix} \right) \mid c \in (\mathbb{Z}/p^2\mathbb{Z})^* \right\}$$

for all $0 \leq \nu \leq p$ and $0 \leq \mu \leq p-1$. One says that the variable x_i is of colour nuance (ν, μ) , if $\begin{pmatrix} \tilde{a}_i \\ \tilde{b}_i \end{pmatrix} \in \mathcal{L}_{\nu\mu} \pmod{p^2}$.

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Alternative way if $m_0 - q_0$ is big:

- The variables at level 0 not in the biggest colour take the place of primary variables.
→ Goal: Create a variable at level γ which traces back to one of those.
- Use the (many) variables at level 0 in the biggest colour to create helpful variables.

$$m_0 + \cdots + m_j \geq \frac{(j+1)s}{k} \quad \text{for } j = 0, 1, \dots, k-1$$

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- Produce primary variables (or work with the alternative, if $m_0 - q_0$ is big enough).
- Lift them to higher level, trying to minimise the factor which is lost with each lifting.
- Either reach at least level γ or gather information about the distribution of the variables at low levels.
- If necessary use level rotation to expand the knowledge to higher levels.

Thanks for your attention!