Forcing with Language Fragments, Extending Namba Forcing, and Models of Theories with Constraints in Interpretation

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Abstract

We develop a forcing framework based on the idea of amalgamating language fragments into a theory with a canonical Henkin model. We then demonstrate the usefulness of this framework by applying it to both the extended Namba problem and the analysis of models of certain *theories with constraints in interpretation* (*TCIs*). The groundwork for a theory of TCIs and their models is laid in parallel to the development of our framework, and is of independent interest.

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1   Introduction

Forcing is a technique in mathematical logic, whereby a set is proven to exist through a non-constructive but coherent assembly of known components. It is most often used in relative consistency proofs, and has been so ever since Cohen invented syntactic forcing in 1963, when he proved that the continuum hypothesis is independent of ZFC in [1] and [2].

An application of forcing involves coming up with a partial order (called a *forcing notion*) and analysing a filter (called a *generic filter*) intersecting suitably many dense subsets of the aforementioned partial order. A set of which existence we want to show (called a *generic object*) typically manifests as an amalgamation of members of the generic filter. In set theoretic applications, a generic object generates an extension (called a *generic extension*) of the original universe.

Over the years, set theorists have discovered important relationships between forcing notions and their generic objects/extensions. These relationships usually associate higher-order properties of a partial order with its forcing consequences, to the extent that the former has become a good description of forcing theory. Now, in order to derive a generic extension we desire, we need to construct a forcing notion drawing from forcing theory with our generic object in mind, at the same time easing tensions between requirements. This balancing act can be extremely tough, in part because it sees little in the way of systematic support.

In this paper we develop a framework in which certain desiderata of a generic object can be naturally realised. Using this framework, all we need to do is to translate requirements on the generic extension into requirements on the generic object. We see via a non-trivial example, how this structured approach can make things more convenient and intuitive in practice.

It turns out that the bare-bones framework we initially envisioned can be adorned with additional layers of syntactic sugar to capture analogous properties of a specific type of first-order structures. These structures are incidentally models of what we term first-order *theories with constraints in interpretation* (*TCIs*). The concept of a first-order TCI, in some sense, generalises that of a first-order theory in logic, and can be useful in expressing different kinds of objects logicians care about. We devote some space in

this paper to the basic theory of TCIs and their models, before relating forcing to generic models of TCIs in a variety of ways that traverses set theory and computability theory. As a side effect, we characterise the expressiveness of forcing as a technique by the kind of truths it is able to impose on generic objects.

1.1 Sectional Content and Dependencies

Section 2 lays the technical foundation for the rest of the paper. Potential philosophical and meta-theoretic concerns are addressed, prerequisite knowledge highlighted, and background readings recommended. Important definitions and conventions are made explicit, especially those that are more niched, and those that lack community consensus. Consequentially, the topic of forcing and genericity (Subsection 2.4) gets a more comprehensive treatment than the others. All subsequent sections depend on the materials here.

Section 3 concerns itself with the development of our central framework for forcing. The technical machinery of this section concentrates around Lemma 3.28, which itself is a generalisation of Lemma 3.35. On the other hand, Lemma 3.35 is the accessible and more applicable backbone of the paper, a throughline tying subsequent sections together. Subsection 5.1 is notable for not depending on Lemma 3.35, nor in fact, on any of Section 3's technology.

The entirety of Section 4 is devoted to applying the forcing framework developed in Section 3 to a non-trivial problem in set theory: a variant of the extended Namba problem. Said framework is used to construct a specific forcing notion that gives rise to generic extensions satisfying a given set of requirements. This construction is wholly contained in the proof of the section's key theorem, Theorem 4.8. No other section is dependent on what transpires here.

Section 5 introduces the notions of (first-order) TCIs and models of TCIs, before relating them to forcing and genericity. In particular, Subsection 5.1 develops the basic theory of TCIs and their models, and can be read right after Section 2. The other subsections depend in part on the results of Section 3, and give applications of Section 3's forcing framework to more general contexts of genericity. Here, genericity is investigated in both set-theoretic and computability-theoretic senses of the word. Like Section 4, this section is not the object of any dependency.

2 Preliminaries

2.1 The Meta-theory

At the meta-level, it suffices to assume ZFC. We frame relative consistency proofs involving additional assumptions as proofs of statements of the form

$$\text{``Con}(\mathsf{ZFC} + \phi) \implies \text{Con}(\mathsf{ZFC} + \psi)\text{''}$$

over ZFC, where ϕ is the conjunction of the relevant assumptions. Implicit are the invocations of Gödel's completeness theorem at the meta-level, whenever we argue using models (a.k.a. *universes*) of set theory.

Our meta-theoretic approach to forcing is only slightly more complicated. Conventionally we start with a countable transitive model of ZFC, called a *ground model*, which is not guaranteed to exist by our meta-theory. There are many ways to sidestep this problem and treat the statement "countable transitive ground model" as a convenient abuse of notation, a few of which are noted by Kunen in IV.5 of [11]. We adopt the first approach detailed in IV.5.1 of [11], an approach that is pretty much standard in the community, and one we feel is most immediately and formally accessible.

2.2 Basic Mathematical Logic

Unless specified otherwise, we follow the standard definitions of concepts related to the syntax and semantics of first-order logic, as seen in e.g. [9].

Convention 2.1.

(1) We call any set of first order formulas a *first-order language*.

(2) We assume the first-order languages we consider to contain only $\neg, \wedge, \vee, \implies, \iff$ as their zeroth-order logical symbols, interpreted semantically in the usual sense.

(3) Each first-order logical symbol is identified with a unique member of $H(\omega) \setminus \omega$.

(4) Given a first-order language $\mathcal{L}$, let $\text{Ter}(\mathcal{L})$ denote the set of all terms occurring in (some formula in) $\mathcal{L}$.

(5) A first-order structure $\mathfrak{A}$ is presented in the form $(A; \mathcal{I})$, where A is the base set of $\mathfrak{A}$ and $\mathcal{I}$ is the interpretation of the signature of $\mathfrak{A}$. In this presentation, the signature of $\mathfrak{A}$ is simply $dom(\mathcal{I})$.

Sometimes, when the correspondence between a signature and its interpretation is clear, we might write $(A; \mathcal{I})$ as $(A; \vec{S})$, where $\vec{S}$ is some ordering of $ran(\mathcal{I})$.

(6) The signature of a first-order language or structure can contain only relation symbols of a non-zero finite arity, function symbols of a non-zero finite arity, and constant symbols, identified as follows:

- an n-ary relation symbol is a triple of the form $(X, 0, n)$,
- an n-ary function symbol is a triple of the form $(X, 1, n)$, and
- a constant symbol is a pair of the form $(X, 2)$.

We call any such symbol a *signature-related symbol*.

(7) A *form-preserving signature embedding* is an injective function from a set of signature-related symbols into the class of signature-related symbols, such that

- n-ary relation symbols are mapped to n-ary relation symbols,
- n-ary function symbols are mapped to n-ary function symbols, and
- constant symbols are mapped to constant symbols.

(8) We will assume that the class of signature-related symbols is disjoint from the set of first-order logical symbols.

(9) The variables occurring in any first-order formula must come from a fixed countably infinite set Var. We will assume that Var is disjoint from both the set of first-order logical symbols and the class of signature-related symbols.

(10) A string over a vocabulary set Σ is a member of $\Sigma^{<\omega}$.

We write "structure(s)" as the abbreviation of "first-order structure(s)" henceforth. There should be no confusion as these are the only type of structures we will be dealing with.

Definition 2.2. Given any set X and any signature σ, the *language associated with* $(X; \sigma)$ is the set of first-order formulas over σ with parameters from X. Similarly, given any structure $\mathfrak{A} = (A; \mathcal{I})$, the *language associated with* $\mathfrak{A}$ is the set of first-order formulas over the signature of $\mathfrak{A}$ with parameters from A.

Definition 2.3. For any structure $\mathfrak{A} = (A; \mathcal{I})$, an $\mathfrak{A}$-*valuation* is a function from Var into A.

Definition 2.4. The *size of a structure* $\mathfrak{A} = (A; \mathcal{I})$ is equal to

$$max\{|A|, |\mathcal{I}|\}.$$

We say $\mathfrak{A}$ is a *countable structure* iff its size is a countable cardinal.

2.3 Basic Set Theory

Unless specified otherwise, we follow the standard definitions of concepts typically encountered in a foundational set theory course, following e.g. [8].

Convention 2.5.

(1) Unless otherwise specified, V always refers to the universe we are currently working in. For all practical purposes, we can assume it is a countable transitive model of ZFC, so that it doubles as a ground model in case forcing arguments are to be run.

(2) We adopt the set-theoretic interpretation of functions as sets of ordered pairs satisfying certain properties. So when we say a function is definable, we actually mean its graph is definable as a set — usually a subset of an ambient structure that should be clear in context, if not explicitly mentioned.

(3) We say a structure $\mathfrak{A}$ is *a structure in the language of set theory* iff

- the signature of $\mathfrak{A}$ is a singleton set,

- the signature of $\mathfrak{A}$ contains a binary relation symbol that is interpreted in $\mathfrak{A}$ as the membership relation on V restricted to the base set of $\mathfrak{A}$.

More formally,
$$\mathfrak{A} = (A; \mathcal{I}) \text{ and } \mathcal{I} = \{(\dot{\in}, \in \cap A)\},$$
where $\dot{\in}$ is a binary relation symbol and $\in$ is the membership relation on V. In this case, we can just write $\mathfrak{A} = (A; \in)$.

(4) We say a structure $\mathfrak{A}$ is *a structure in a possibly expanded language of set theory* iff we omit the cardinality requirement on the signature of $\mathfrak{A}$ in (3). More formally,
$$\mathfrak{A} = (A; \mathcal{I}) \text{ and } \mathcal{I} = \{(\dot{\in}, \in \cap A)\} \cup \vec{X},$$
where $\dot{\in}$ is a binary relation symbol, $\vec{X}$ is a function and $\in$ is the membership relation on V. In this case, we can just write $\mathfrak{A} = (A; \in, \vec{X})$.

(5) We identify a universe of set theory W with the structure $(W; \in)$. This should not cause confusion in the circumstances we find ourselves in.

(6) A real is a subset of ω. We identify a real with its characteristic function on ω, as is standard in computability theory. As in (5), this ambiguity should not cause any confusion.

(7) When we say "a sufficiently strong set theory", we mean to emphasise the low strength of the theory. In such instances, well-known set theories far weaker than ZFC, such as Z^- and KP, typically work.

Definition 2.6. Let X be a set and $\mathfrak{A}$ be a structure in a possibly expanded language of set theory. We say X is *definable in the language associated with* $\mathfrak{A}$ iff
$$\mathfrak{A} \models \text{``}X = \{y : \phi(y)\}\text{''}$$
for some ϕ in one free variable in the language associated with $\mathfrak{A}$.

Definition 2.7. Given that V' and W are models of ZFC, we say V' is a *weak inner model* of W (or equivalently, W is a *weak outer model* of V') iff

- V' and W are both transitive, and

- $V' \subset W$.

If in addition, V' and W share the same ordinals i.e. $ORD^{V'} = ORD^W$, then V' is an *inner model* of W (or equivalently, W is an *outer model* of V').

Definition 2.8. Let Y be any set.
   We say X *codes* Y (or equivalently, X *is a code of* Y) iff

- X is a set of ordinals, and

- every transitive model of ZFC $-$ Powerset containing X also contains Y.

X *codes* Y *as a real* iff X codes Y and $X \subset \omega$.
   Y *has a real code* (or equivalently, Y *can be coded as a real*) iff X codes Y as a real for some X.

Lemma 2.9. *Let X be a set with $|trcl(X)| = \kappa$. Then there is $A \subset \kappa$ such that A codes X.*
   In particular, any set with a countable transitive closure has a real code.

Proof. Let $Y := trcl(X) \cup \{X\}$. Note that

$$X \text{ is the unique } \in\text{-maximal member of } Y \tag{2.1}$$

in any transitive model of ZFC $-$ Powerset containing Y. Choose any bijection f from Y into κ. Define

$$R := \{(f(x), f(y)) \in \kappa \times \kappa : (x, y) \in Y \times Y \text{ and } x \in y\}.$$

Now R can be thought of as a subset A of κ via the canonical Gödel numbering of pairs. If V' is a transitive model of ZFC $-$ Powerset containing A, then we can recover R in V'. The Mostowski collapse function works on R in V' to give us $Y \in V'$. This implies $X \in V'$ since X definable from Y via (2.1). $\qquad\square$

2.4 Forcing and Generic Objects

Following the meta-theoretic convention highlighted in Subsection 2.1, we start with a countable transitive ground model V. In the language of forcing, a forcing notion in V is just any partial order in V. If $\mathbb{P}$ is a forcing notion in V, then a $\mathbb{P}$-*generic filter over* V is a filter on $\mathbb{P}$ intersecting all dense subsets of $\mathbb{P}$ in V.

Given a forcing notion $\mathbb{P}$ in V, the class of $\mathbb{P}$-names in V — denoted $V^{\mathbb{P}}$ — and the forcing relation $\Vdash_{\mathbb{P}}^{V}$ (which relates elements of $\mathbb{P}$ with formulas parametrised by $\mathbb{P}$-names in V), are both essential to a forcing argument involving $\mathbb{P}$ carried out in V. These two classes are uniformly definable in V over the class of all forcing notions $\mathbb{P}$. $\mathbb{P}$-names in V are "evaluated at" a $\mathbb{P}$-generic filter g over V to obtain the $\mathbb{P}$-generic extension $V[g]$, which is necessarily countable and transitive. In other words, if g is a $\mathbb{P}$-generic filter over V, then

$$V \subset V[g] = \{\dot{x}[g] : \dot{x} \in V^{\mathbb{P}}\},$$

where $\dot{x}[g]$ means "x evaluated at g". The evaluation procedure is done outside V because g typically (in order to be of use at all) does not exist in V.

Convention 2.10. When it is clear that the background universe is V, we suppress mention of V when writing forcing relations in V. This means that given a forcing notion $\mathbb{P}$ in V, $\Vdash_{\mathbb{P}}$ is used interchangeably with $\Vdash_{\mathbb{P}}^{V}$.

Definition 2.11. We call W a *generic extension* (or a *forcing extension*) of V iff there exists a forcing notion $\mathbb{P}$ in V and a $\mathbb{P}$-generic filter g over V, such that $W = V[g]$.

Definition 2.12. We write "$\Vdash_{\mathbb{P}} \phi$" to mean

$$\text{``}\forall p \ (p \in \mathbb{P} \implies p \Vdash_{\mathbb{P}} \phi)\text{''}.$$

Remark 2.13. A theorem fundamental to the technique of forcing intricately connects the forcing relation $\Vdash_{\mathbb{P}}^{V}$ with truth in $\mathbb{P}$-generic extensions. It goes as follows:

> If $\mathbb{P}$ is a forcing notion in V, $p \in \mathbb{P}$, ϕ is a formula with n free variables, and $\dot{x}_1, ..., \dot{x}_n$ are $\mathbb{P}$-names in V, then
>
> - $p \Vdash_{\mathbb{P}} \phi(\dot{x}_1, \ldots, \dot{x}_n) \iff$
> $\forall g \ ((g$ is $\mathbb{P}$-generic over V and $p \in g)$
> $\implies V[g] \models \phi(\dot{x}_1[g], \ldots, \dot{x}_n[g]))$, and
> - $\forall g \ ((g$ is $\mathbb{P}$-generic over V and $V[g] \models \phi(\dot{x}_1[g], \ldots, \dot{x}_n[g]))$
> $\implies \exists q \ (q \Vdash_{\mathbb{P}} \phi(\dot{x}_1, \ldots, \dot{x}_n)$ and $q \in g))$.

This theorem, colloquially known as the *forcing theorem*, enables us to reason about truth in generic extensions from within the ground model, and often reduces the argument from one about semantic entailment to one pertaining to combinatorial properties of partial orders. For a technical lowdown of forcing terminology and the proof of the forcing theorem, the reader is encouraged to read Chapter IV of [11].

Definition 2.14. If $\mathbb{B} = (B, \mathbb{0}, \mathbb{1}, \leq_{\mathbb{B}}, \wedge_{\mathbb{B}}, \vee_{\mathbb{B}}, \sim_{\mathbb{B}})$ is a boolean algebra, then define

$$
\begin{aligned}
B^+ &:= B \setminus \{\mathbb{0}\} \\
\leq_{\mathbb{B}^+} &:= \ \leq_{\mathbb{B}} \cap (B^+ \times B^+) \\
\wedge_{\mathbb{B}^+} &:= \ \wedge_{\mathbb{B}} \cap (B^+ \times B^+) \\
\vee_{\mathbb{B}^+} &:= \ \vee_{\mathbb{B}} \cap (B^+ \times B^+) \\
\sim_{\mathbb{B}^+} &:= \ \sim_{\mathbb{B}} \cap (B^+ \times B^+) \\
\mathbb{B}^+ &:= (B^+, \mathbb{1}, \leq_{\mathbb{B}^+}, \wedge_{\mathbb{B}^+}, \vee_{\mathbb{B}^+}, \sim_{\mathbb{B}^+}) \\
\mathbb{B}_0^+ &:= (B^+, \leq_{\mathbb{B}^+}).
\end{aligned}
$$

Definition 2.15. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion and X be any set. The *upward closure of X in $\mathbb{P}$*, denoted $\mathrm{UC}(\mathbb{P}, X)$, is the set

$$
\{p \in P : \exists q \ (q \in X \text{ and } q \leq_{\mathbb{P}} p)\}.
$$

Definition 2.16. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion, $D \subset P$ and A be any set. We say a subset g of P *meets D in A* iff

$$
g \cap \{p \in P : p \in D \text{ or } \forall q \ (q \leq_{\mathbb{P}} p \implies q \notin D)\} \cap A \neq \emptyset.
$$

We say g *meets D* iff g meets D in V.

Definition 2.17. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion and $\mathfrak{A} = (A; \in, \vec{X})$ be a structure in a possibly expanded language of set theory. We say a subset g of P is $\mathbb{P}$-*generic over* $\mathfrak{A}$ (or g is a $\mathbb{P}$-*generic subset over* $\mathfrak{A}$) iff g meets D in A for all D such that

- $D \subset P$

- D is dense in $\mathbb{P}$, and

- D is definable in the language associated with $\mathfrak{A}$.

If in addition, g is a filter on $\mathbb{P}$, then we call g a $\mathbb{P}$-*generic filter over* $\mathfrak{A}$.

Definition 2.18. Let $\mathbb{P}$ be a forcing notion and $\mathfrak{A} = (A; \in, \vec{X})$ be a *transitive* structure in a possibly expanded language of set theory. A set x is a $(\mathbb{P}, \mathfrak{A})$-*generic object* iff there exists g a $\mathbb{P}$-generic filter over $\mathfrak{A}$ such that

- $x \in A[g]$, and

- g is definable in the language associated with $(A \cup \{x\}; \in, \vec{X})$,

in which case we say g *witnesses* x is a $(\mathbb{P}, \mathfrak{A})$-generic object.

Definition 2.19. Let x be any set. We call x a *generic object* iff there is a pair $(\mathbb{P}, \mathfrak{A})$ for which x is a $(\mathbb{P}, \mathfrak{A})$-generic object. We call x a V-*generic object* iff there is $\mathbb{P}$ for which x is a $(\mathbb{P}, V)$-generic object.

Observation 2.20. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion and X be any set. Then there is a structure $\mathfrak{A} = (A; \in) \in V$ such that in every weak outer model of V,

$$x \text{ is a } (\mathbb{P}, \mathfrak{A})\text{-generic object} \iff x \text{ is a } (\mathbb{P}, V)\text{-generic object}$$

for all $x \subset X$. In fact, we can choose A to be $H(\kappa)$ for any $\kappa > |trcl(\{P, X\})|$.

Observation 2.20 allows us to refer to $(\mathbb{P}, V)$-generic objects for any forcing notion $\mathbb{P}$, without needing to quantify over all formulas.

There are looser definitions of a generic filter or a generic object in the literature. For example, we can require the filter to only meet subsets with definitions belonging to a certain complexity class, as is commonly seen in computability theory. Informally then, the study of genericity boils down to observing the effects of a filter meeting a bunch of subsets.

Section 1 hinted at a key difference between forcing theory and the study of partial orders in order theory, and that is the nature of the properties studied apropos of their common subjects. In order theory, only first-order properties of partial orders are considered, whereas forcing theory concerns itself with their higher-order properties. Now, another such differentiating factor is the overwhelming focus on generic objects in forcing theory. In fact, so much attention is paid to generic objects in forcing theory that one might as well call it *genericity theory*.

Definition 2.21. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\mathbb{Q} = (Q, \leq_{\mathbb{Q}})$ be two forcing notions. We say π is an *embedding from $\mathbb{P}$ into $\mathbb{Q}$* iff π is an injective function from P into Q satisfying

- $p_1 \leq_{\mathbb{P}} p_2 \iff \pi(p_1) \leq_{\mathbb{Q}} \pi(p_2)$, and
- $p_1 \not\perp_{\mathbb{P}} p_2 \implies \pi(p_1) \not\perp_{\mathbb{Q}} \pi(p_2)$.

Definition 2.22. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\mathbb{Q} = (Q, \leq_{\mathbb{Q}})$ be two forcing notions. An embedding π from $\mathbb{P}$ into $\mathbb{Q}$ is *dense* iff $ran(\pi)$ is dense in $\mathbb{Q}$.

Fact 2.23. Let $\mathbb{P}$ and $\mathbb{Q}$ be two forcing notions. If there is a dense embedding from $\mathbb{P}$ into $\mathbb{Q}$, then $\mathbb{P}$ and $\mathbb{Q}$ are forcing equivalent.

Definition 2.24. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion. Define

$$w(\leq_{\mathbb{P}}) := \{(p, q) \in P \times P : \{q' : q' \leq_{\mathbb{P}} q\} \text{ is dense below } p\}, \text{ and}$$
$$w(\mathbb{P}) := (P, w(\leq_{\mathbb{P}})).$$

Fact 2.25. For any forcing notion $\mathbb{P}$,

(1) $w(w(\mathbb{P})) = w(\mathbb{P})$, and

(2) $\mathbb{P}$ is forcing equivalent to $w(\mathbb{P})$.

Definition 2.26. A forcing notion $\mathbb{P}$ is *separative* iff $w(\mathbb{P}) = \mathbb{P}$.

Fact 2.27. If $\mathbb{B}$ is a boolean algebra, then $\mathbb{B}_0^+$ is a separative forcing notion.

Fact 2.28. For any forcing notion $\mathbb{P}$, there is a unique (up to isomorphism) complete boolean algebra $\mathbb{B}$ such that a dense embedding exists from $\mathbb{P}$ into $\mathbb{B}_0^+$.

Remark 2.29. By Facts 2.23 and 2.28, every forcing extension of V is a $\mathbb{B}_0^+$-generic extension for some complete boolean algebra $\mathbb{B}$.

Definition 2.30. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\mathbb{Q} = (Q, \leq_{\mathbb{Q}})$ be two forcing notions. We say π is a *weak embedding from $\mathbb{P}$ into $\mathbb{Q}$* iff π is an embedding from $w(\mathbb{P})$ into $w(\mathbb{Q})$.

Definition 2.31. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\mathbb{Q} = (Q, \leq_{\mathbb{Q}})$ be two forcing notions. A weak embedding π from $\mathbb{P}$ into $\mathbb{Q}$ is *dense* iff π is dense as an embedding from $w(\mathbb{P})$ into $w(\mathbb{Q})$.

Remark 2.32. Let $\mathbb{P}$ and $\mathbb{Q}$ be two forcing notions. By Facts 2.23 and 2.25, if there is a dense weak embedding from $\mathbb{P}$ into $\mathbb{Q}$, then $\mathbb{P}$ and $\mathbb{Q}$ are forcing equivalent.

Definition 2.33. If $\mathbb{P} = (P, \leq_{\mathbb{P}})$ is a forcing notion and $p \in P$, we let $g_p(\mathbb{P})$ denote the set

$$\{q \in P : p \not\perp_{\mathbb{P}} q\}.$$

Definition 2.34. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ be a forcing notion. A member p of P is an *atom* of $\mathbb{P}$ iff

$$\forall q_1 \, \forall q_2 \, ((q_1 \leq_{\mathbb{P}} p \text{ and } q_2 \leq_{\mathbb{P}} p) \implies q_1 \not\perp_{\mathbb{P}} q_2).$$

Lemma 2.35. *If $\mathbb{P} = (P, \leq_{\mathbb{P}})$ is a forcing notion and p is an atom of $\mathbb{P}$, then $g_p(\mathbb{P})$ is a $\mathbb{P}$-generic filter over V.*

Proof. If D is dense in $\mathbb{P}$, then there is $q \in D$ with $q \leq_{\mathbb{P}} p$. Obviously, $q \in g_p(\mathbb{P})$. Therefore $g_p(\mathbb{P})$ is a $\mathbb{P}$-generic subset over V. To see that $g_p(\mathbb{P})$ is a filter, let q_1 and q_2 be members of $g_p(\mathbb{P})$. By the definition of $g_p(\mathbb{P})$, there are r_1 and r_2 such that

- $r_1 \leq_{\mathbb{P}} q_1$,
- $r_1 \leq_{\mathbb{P}} p$,
- $r_2 \leq_{\mathbb{P}} q_2$,
- $r_2 \leq_{\mathbb{P}} p$.

As p is an atom of $\mathbb{P}$, it must be the case that $r_1 \not\perp_{\mathbb{P}} r_2$, which means $q_1 \not\perp_{\mathbb{P}} q_2$. $\square$

Definition 2.36. A forcing notion $\mathbb{P}$ is *atomic* iff the set of atoms of $\mathbb{P}$ is dense in $\mathbb{P}$.

Definition 2.37. A forcing notion $\mathbb{P} = (P, \leq_{\mathbb{P}})$ is *atomless* iff no member of P is an atom of $\mathbb{P}$.

Definition 2.38. Define

$$C := 2^{<\omega}, \text{ and}$$
$$\leq_{\mathbb{C}} := \{(p, q) : q \subset p\}.$$

Call the forcing notion $\mathbb{C} := (C, \leq_{\mathbb{C}})$ *Cohen forcing*.

Definition 2.39. Let $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\mathbb{Q} = (Q, \leq_{\mathbb{Q}})$ be forcing notions. We say $\mathbb{P}$ is a *regular suborder* of $\mathbb{Q}$, denoted $\mathbb{P} \lessdot \mathbb{Q}$, iff

- $\mathbb{P}$ is a suborder of $\mathbb{Q}$, and
- for all $q \in Q$ there is $p \in P$ such that every $p' \leq_{\mathbb{P}} p$ is compatible with q in $\mathbb{Q}$.

Fact 2.40. If $\mathbb{P} = (P, \leq_{\mathbb{P}}) \lessdot \mathbb{Q}$, then for every $\mathbb{Q}$-generic filter g over V, $g \cap P$ is a $\mathbb{P}$-generic filter over V.

3 Forcing with Language Fragments

In model theory, a Henkin construction involves building a model of a theory over a language, from terms of that language. When such a construction is unequivocally guided by a given theory, we can safely identify the resulting Henkin model with said theory. As such, we have the following viable means of proving the existence of a object with property P:

1. Translate P into a specification S for a theory, such that the unique Henkin model of any theory satisfying S has property P.
2. Prove that a theory satisfying S exists.

In a similar fashion, we can force the existence of an object with property P by forcing the existence of a theory satisfying S. Naturally, this leads to forcing notions with conditions being fragments of the language over which a theory satisfying S is defined.

The idea of forcing models of a theory has been studied by model theorists — notably, Robinson and Barwise — in the 1970s (see e.g. [5]). More recently, set theorists have leveraged on this idea to solve problems of set-theoretic interest. Some examples include [10] and [12]. However, the forcing conditions used in these cases are highly complicated and specialised structures, and it is not immediately clear how much of the analysis of one forcing notion can be recycled in the analysis of another. Asperó and Schindler are perhaps the first to present a construction with language fragments as forcing conditions, in their seminal work [13]. It quickly became clear that similar presentations can be adopted for the cited past works. Streamlining and modularising the analysis and construction of forcing notions of this ilk thus seems like a useful proposition.

This section details a framework for constructing forcing notions with fragments of a language $\mathcal{L}$ as conditions, based on specifications of a theory over $\mathcal{L}$. To state such specifications, a "meta-language" dependent on $\mathcal{L}$ is required. Our goal is to ensure that the generic filters of each forcing notion produced indeed give rise to theories satisfying the given specifications. We will make precise the relevant technical terms and concepts as we build our framework over the subsequent three subsections.

The main result in this section is Lemma 3.28, which is stated and proven in greater generality than is needed for our framework. From Lemma 3.28 we derive Lemma 3.35, the primary workhorse of the entire paper.

3.1 General Languages and Meta-languages

The initial step in the development of our framework involves the ability to potentially interpret any set as a language.

Definition 3.1. The canonical negation function $\neg$ on V is defined as follows.

$$\neg x := \neg(x) = \begin{cases} y & \text{if } x = \ulcorner \neg y \urcorner \text{ for some } y \\ \ulcorner \neg x \urcorner & \text{otherwise.} \end{cases}$$

Basically, $\neg$ takes a member of $\mathcal{L}$ as input, and check whether it is a string with first (leftmost) character $\ulcorner \neg \urcorner$. If so, it removes the leading $\ulcorner \neg \urcorner$; otherwise, it "casts" the input as a string (mapping the input to a string of length 1 containing the input as the only character, if the input is not already a string) and prepend $\ulcorner \neg \urcorner$ to the result. For ease of argument, we identify the string containing a single character x with x itself.

Note that $\neg$ is Δ_0-definable with a single parameter $\ulcorner \neg \urcorner$, which we assume is in every of a sufficiently set theory. So the definition of $\neg$ is absolute for models of a sufficiently strong set theory.

We sometimes abuse notation and use $\neg$ the function and $\ulcorner \neg \urcorner$ the first-order logical symbol interchangeably. However, we take special care to distinguish them wherever is crucial in our definitions and arguments.

Definition 3.2. A set $\mathcal{L}$ is *closed under negation* iff for each $\phi \in \mathcal{L}$, $\neg\phi \in \mathcal{L}$.

Before we proceed, fix a set $\mathcal{L}$ that is closed under negation. We will stick to this $\mathcal{L}$ for the rest of this section.

Definition 3.3. A structure $\mathfrak{A} = (A; \in, \vec{R})$ is *$\mathcal{L}$-suitable* iff

- $\vec{R}$ is a set of relations on A,
- $(A; \in)$ is a model of a sufficiently strong set theory,
- $\mathcal{L} \subset A$, and

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- $\mathcal{L}$ is definable in the language associated with $\mathfrak{A}$.

Fix an $\mathcal{L}$-suitable $\mathfrak{A} = (A; \in, \vec{R})$ for the rest of this section. For notational convenience, we shall, without loss of generality, assume $\mathcal{L}$ occurs in $\vec{R}$.

Definition 3.4. Define $\mathcal{L}_{\mathfrak{A}}^*$ to be the language associated with $(A; \{\in, \vec{R}, E\})$, where E is a unary predicate symbol not occurring in $\vec{R}$.

We want to use $\mathcal{L}_{\mathfrak{A}}^*$ to reason about subsets of $\mathcal{L}$. Intuitively, a richer $\mathfrak{A}$ should allow us to formulate more statements about these subsets. Certain subsets of $\mathcal{L}$ are particularly interesting.

Definition 3.5. A set Σ is $\mathcal{L}$-*nice* iff

- $\Sigma \subset \mathcal{L}$,
- for all $\phi \in \mathcal{L}$,

  ○ $\{\phi, \neg\phi\} \not\subset \Sigma$, and
  ○ either $\phi \in \Sigma$ or $\neg\phi \in \Sigma$.

Definition 3.6. Let $\phi \in \mathcal{L}_{\mathfrak{A}}^*$. Define $\mathsf{pos}(\phi)$ to be the formula in $\mathcal{L}_{\mathfrak{A}}^*$ resulting from the following operation:

- for each subformula φ of ϕ, if $\varphi = \ulcorner \neg E(x) \urcorner$ for some $x \in \mathcal{L}$, then replace φ with $\ulcorner E(\neg x) \urcorner$.

Definition 3.7. Let

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, and
- ν be any subset of an $\mathfrak{A}$-valuation.

Then $\nu^*(\phi)$ is defined to be the sentence in $\mathcal{L}_{\mathfrak{A}}^*$ resulting from the following operation:

- for each $c \in dom(\nu)$, replace every free occurrence of c in ϕ with $\nu(c)$.

Definition 3.8. Let

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, and
- ν be an $\mathfrak{A}$-valuation.

We say $X \models_{\mathfrak{A}, \nu}^* \phi$ iff

$$(A; \in, \vec{R}, X \cap A) \models \nu^*(\phi)(\vec{x}; \in, \vec{R}, E).$$

We say $X \models_{\mathfrak{A}}^* \phi$ iff for every $\mathfrak{A}$-valuation ν, $X \models_{\mathfrak{A}, \nu}^* \phi$.

If $\Gamma \subset \mathcal{L}_{\mathfrak{A}}^*$ then we say

$$X \models_{\mathfrak{A}, \nu}^* \Gamma \text{ iff } X \models_{\mathfrak{A}, \nu}^* \phi \text{ for all } \phi \in \Gamma, \text{ and}$$
$$X \models_{\mathfrak{A}}^* \Gamma \text{ iff } X \models_{\mathfrak{A}}^* \phi \text{ for all } \phi \in \Gamma.$$

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Remark 3.9. Let

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$ be a sentence, and
- $C = A$ or C occur in $\vec{R}$.

Then
$$\{X \in C : X \models_{\mathfrak{A}}^* \phi\}$$

is a subset of A that is definable in the language associated with $\mathfrak{A}$. This definition is absolute for transitive models of $\mathsf{ZFC} - \mathsf{Powerset}$.

Note that for any $\mathcal{L}$-nice Σ and any $x \in \mathcal{L}$,

$$x \notin \Sigma \iff \neg x \in \Sigma,$$

so applying pos to a $\mathcal{L}_{\mathfrak{A}}^*$ formula does not alter its meaning with respect to $\mathcal{L}$-nice sets. More formally, we have the next lemma.

Lemma 3.10. *Let*

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, *and*
- ν *be an* $\mathfrak{A}$-*valuation.*

Then for every $\mathcal{L}$-*nice* Σ *in every weak outer model of* V,

$$\Sigma \models_{\mathfrak{A},\nu}^* \phi \iff \Sigma \models_{\mathfrak{A},\nu}^* \mathsf{pos}(\phi).$$

Proof. Σ being $\mathcal{L}$-nice means that for all $x \in \mathcal{L}$,

$$\Sigma \models_{\mathfrak{A}}^* \ulcorner \neg E(x) \urcorner \iff \Sigma \models_{\mathfrak{A}}^* \ulcorner E(\neg x) \urcorner.$$

The lemma follows immediately. $\qquad\qquad\square$

Definition 3.11. If $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, we say ϕ is $(\mathfrak{A}, \mathcal{L})$-satisfiable iff there are ν, W and X such that

- ν is an $\mathfrak{A}$-valuation,
- W is a weak outer model of V,
- $X \in W \cap \mathcal{P}(\mathcal{L})$, and
- $X \models_{\mathfrak{A},\nu}^* \phi$,

in which case the triple (ν, W, X) is said to *witness the* $(\mathfrak{A}, \mathcal{L})$-*satisfiability of* ϕ.

Definition 3.12. For any $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, define $\mathsf{set}(\phi)$ to be the pair (p, q) such that

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- $q = \{x \in \mathcal{L} : \ulcorner E(x) \urcorner \text{ is a subformula of } \phi\}$, and
- $p = \{x \in q : \ulcorner (\neg E(x)) \urcorner \text{ is not a subformula of } \phi\}$.

Definition 3.13. Let

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, and
- ν be an $\mathfrak{A}$-valuation.

For any sets p and q, we say ϕ *is* $\models_{\mathfrak{A},\nu}^*$*-true for* (p, q) iff

- $p \subset q \subset \mathcal{L}$, and
- for all

    ◦ weak outer models W of V, and
    ◦ $X \in W \cap \mathcal{P}(\mathcal{L})$,

$$X \cap q = p \implies X \models_{\mathfrak{A},\nu}^* \phi.$$

We say ϕ *is* $\models_{\mathfrak{A}}^*$ *-true for* (p, q) iff for every $\mathfrak{A}$-valuation ν, ϕ is $\models_{\mathfrak{A},\nu}^*$-true for (p, q). We say ϕ *is* $\models_{\mathfrak{A},\nu}^*$*-true for* p iff ϕ is $\models_{\mathfrak{A},\nu}^*$-true for (p, p).

For our purposes, being $\models_{\mathfrak{A},\nu}^*$-true (or $\models_{\mathfrak{A},\nu}^*$-true) can be too strong a requirement; it is often enough to narrow the scope of our "test models" to just $\mathcal{L}$-nice sets. This motivates the following definition.

Definition 3.14. Let

- $\phi \in \mathcal{L}_{\mathfrak{A}}^*$, and
- ν be an $\mathfrak{A}$-valuation.

For any sets p and q, we say ϕ *is* $\models_{\mathfrak{A},\nu}^*$*-nice for* (p, q) iff

- $p \subset q \subset \mathcal{L}$, and
- for all

    ◦ weak outer models W of V, and
    ◦ $X \in W \cap \mathcal{P}(\mathcal{L})$,

$$X \text{ is } \mathcal{L}\text{-nice and } X \cap q = p \implies X \models_{\mathfrak{A},\nu}^* \phi.$$

We say ϕ *is* $\models_{\mathfrak{A}}^*$ *-nice for* (p, q) iff for every $\mathfrak{A}$-valuation ν, ϕ is $\models_{\mathfrak{A},\nu}^*$-nice for (p, q). We say ϕ *is* $\models_{\mathfrak{A},\nu}^*$*-nice for* p iff ϕ is $\models_{\mathfrak{A},\nu}^*$-nice for (p, p).

Definition 3.15. Let $\mathcal{L}_{0,\mathfrak{A}}^*$ be the language associated with $\mathfrak{A}$. Obviously, $\mathcal{L}_{0,\mathfrak{A}}^* \subset \mathcal{L}_{\mathfrak{A}}^*$. Let $\mathcal{L}_{1,\mathfrak{A}}^*$ be the smallest $\mathcal{L}'$ satisfying the following conditions:

- $\mathcal{L}_{0,\mathfrak{A}}^* \subset \mathcal{L}'$,
- $\{\ulcorner E(x) \urcorner : x \in \mathrm{Ter}(\mathcal{L}_{0,\mathfrak{A}}^*)\} \subset \mathcal{L}'$,
- if

    ○ $\phi \in \mathcal{L}'$,

    ○ z is a variable not bound in ϕ,

    ○ $p \subset \mathcal{L}$, and

    ○ $|p| < \omega$,

    then
$$\ulcorner \forall z \, (z \in p \implies \phi) \urcorner \in \mathcal{L}',$$

- $\mathcal{L}'$ is closed under all zeroth-order logical operations.

Check that $\mathcal{L}_{1,\mathfrak{A}}^* \subset \mathcal{L}_{\mathfrak{A}}^*$. A first-order formula ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Delta_0$ iff $\phi \in \mathcal{L}_{1,\mathfrak{A}}^*$.

Definition 3.16. We define the subset $\mathcal{D}$ of $\mathcal{L}_{1,\mathfrak{A}}^*$ to contain formulas of the form

$$\bigvee_{i<m} (\bigwedge_{j<n_i} L_{ij}),$$

wherein for every $i < m$ and every $j < n_i$, there is P_{ij} such that

- $P_{ij} \in \mathcal{L}_{0,\mathfrak{A}}^*$ or $P_{ij} \in \mathcal{L}$ is of the form $\ulcorner E(x) \urcorner$, and
- $L_{ij} = P_{ij}$ or $L_{ij} = \ulcorner \neg P_{ij} \urcorner$.

Every formula in $\mathcal{L}_{1,\mathfrak{A}}^*$ is logically equivalent to a formula in $\mathcal{D}$; in fact, there is a function, definable in the language associated with $\mathfrak{A}$, that takes each formula in $\mathcal{L}_{1,\mathfrak{A}}^*$ to a formula in $\mathcal{D}$. Fix one such function and call it DNF.

Definition 3.17. A first-order formula ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Sigma_1$ iff it is of the form

$$\ulcorner \exists y_1 ... \exists y_j \, \phi \urcorner,$$

where

- $j < \omega$,
- ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Delta_0$,
- $y_1, ..., y_j$ are variables not bound in ϕ.

Definition 3.18. A first-order formula ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ iff it is of the form

$$\ulcorner \forall x_1 ... \forall x_i \ \phi \urcorner,$$

where

- $i < \omega$,
- ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Sigma_1$,
- $x_1, ..., x_i$ are variables not bound in ϕ.

Remark 3.19. Analogous to what the classification of general first-order formulas in prenex normal form, we can very naturally build on Definitions 3.15, 3.17 and 3.18, and inductively define $\mathcal{L}_{\mathfrak{A}}^*$-$\Sigma_n$ and $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_n$ sentences for all $n < \omega$. The only reason we did not is because our theorems and analyses neither mention nor require formulas outside of $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$.

Note also that in the definition of $\mathcal{L}_{\mathfrak{A}}^*$-$\Delta_0$ formulas, members of $\mathcal{L}_{0,\mathfrak{A}}^*$ are regarded, for all practical purposes, as atomic formulas. Further, bounded quantification is limited to finite sets — more in the spirit of arithmetical bounded quantification than the usual set-theoretic one. This is just as well, since we will be analysing these formulas in the context of arithmetic in the latter half of Subsection 5.2.

Lemma 3.20. *Let*

- *ϕ be a $\mathcal{L}_{\mathfrak{A}}^*$-$\Sigma_1$ formula, and*
- *(ν, W, X) witness the $(\mathfrak{A}, \mathcal{L})$-satisfiability of ϕ.*

Then there are finite sets $p \in A \cap \mathcal{P}(\mathcal{L})$ and $q \in A \cap \mathcal{P}(\mathcal{L})$ such that

- *$X \cap q = p$, and*
- *ϕ is $\models_{\mathfrak{A},\nu}^*$-true for (p, q).*

Proof. By induction on the length of ϕ.

Case 1: ϕ is $\mathcal{L}_{\mathfrak{A}}^*$-$\Delta_0$. Note first that

$$X \models_{\mathfrak{A},\nu}^* \phi \iff X \models_{\mathfrak{A},\nu}^* \nu^*(\mathsf{DNF}(\phi)),$$

so we have that $X \models_{\mathfrak{A},\nu}^* \nu^*(\mathsf{DNF}(\phi))$. Next, since $\nu^*(\mathsf{DNF}(\phi)) \in \mathcal{D}$, there is a disjunct φ of $\nu^*(\mathsf{DNF}(\phi))$ for which $X \models_{\mathfrak{A},\nu}^* \varphi$.

Let (p, q) be $\mathsf{set}(\varphi)$. Then p and q are finite sets with $p \subset q \subset \mathcal{L}$. Since $(A; \in)$ models enough set theory and $\mathcal{L} \subset A$, we too have $\{p, q\} \subset A$. As φ has all occurrences of literals over $\{E\}$ being conjuncts, we must have $X \cap q = p$. It is thus easy to see that for any $X' \subset \mathcal{L}$ in any weak outer model of V,

$$X' \cap q = p \implies X' \models_{\mathfrak{A},\nu}^* \varphi \implies X' \models_{\mathfrak{A},\nu}^* \phi.$$

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Case 2: $\phi = \ulcorner \exists y \, \phi' \urcorner$ for some y and ϕ'. There is an $\mathfrak{A}$-valuation ν' that agrees with ν on the free variables of ϕ, for which $X \models^*_{\mathfrak{A},\nu'} \phi'$. By the induction hypothesis, there are finite sets $p \in A \cap \mathcal{P}(\mathcal{L})$ and $q \in A \cap \mathcal{P}(\mathcal{L})$ for which $X \cap q = p$ and ϕ' is $\models^*_{\mathfrak{A},\nu'}$-true for (p,q). Since $(\nu')^*(\phi')$ logically implies $\nu^*(\phi)$, it must also be that ϕ is $\models^*_{\mathfrak{A},\nu}$-true for (p,q). $\qquad\square$

For any free variable x and any $\mathcal{L}^*_{\mathfrak{A}}$ formula ϕ, it is often a desideratum (if not an imperative) in practice to have ϕ explicitly "guarantee $E(\neg x)$" if $\ulcorner\neg E(x)\urcorner$ occurs in ϕ. Towards this end, we are incentivised to augment ϕ with a suitable gadget.

Definition 3.21. Let ϕ be a $\mathcal{L}^*_{\mathfrak{A}}$-$\Delta_0$ formula. Define

$$\mathrm{CS}(\phi) := \{\ulcorner(x \in \mathcal{L} \implies E(\neg x))\urcorner : x \text{ is free in } \phi \text{ and } \ulcorner\neg E(x)\urcorner \text{ occurs in } \phi\},$$

and

$$\mathsf{check}(\phi) := \ulcorner\phi \wedge \bigwedge \mathrm{CS}(\phi)\urcorner.$$

Remark 3.22. In the definition of $\mathrm{CS}(\phi)$, $E(\neg x)$ is a shorthand for

$$\exists y \, (y = \neg(x) \wedge E(y)),$$

where $\neg(\cdot)$ is the negation function on V, so that "$y = \neg(x)$" is expressible as a (Δ_0, no less) formula in $\mathcal{L}^*_{0,\mathfrak{A}}$. As a consequence, if $\ulcorner\exists x_1...\exists x_n \, \phi\urcorner$ is a $\mathcal{L}^*_{\mathfrak{A}}$-$\Sigma_1$ formula, then $\ulcorner\exists x_1...\exists x_n \, \mathsf{check}(\phi)\urcorner$ is logically equivalent to a $\mathcal{L}^*_{\mathfrak{A}}$-$\Sigma_1$ formula.

Lemma 3.23. *Let*

- $\phi = \bigwedge \mathrm{S} \wedge \bigwedge \mathrm{T}$ *for some*

  - S *is a finite subset of* $\mathcal{L}^*_{0,\mathfrak{A}}$, *and*
  - T *is a finite subset of* $\mathcal{L}$ *containing only formulas either of the form* $\ulcorner E(x)\urcorner$ *or of the form* $\ulcorner\neg E(x)\urcorner$,

- ν *be an* $\mathfrak{A}$-*valuation, and*
- $\Sigma \subset \mathcal{L}$.

If $\Sigma \models^*_{\mathfrak{A},\nu} \mathsf{check}(\mathsf{pos}(\phi))$, *then* $\mathsf{check}(\mathsf{pos}(\phi))$ *is* $\models^*_{\mathfrak{A},\nu}$-*nice for* Σ.

Proof. By the form of $\mathsf{check}(\mathsf{pos}(\phi))$, it is sufficient to prove that whenever

- $\Sigma' \supset \Sigma$,
- Σ' is $\mathcal{L}$-nice, and
- φ is a subformula of $\nu^*(\mathsf{check}(\mathsf{pos}(\phi)))$ of the form $\ulcorner E(x)\urcorner$,

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$\Sigma \models_{\mathfrak{A}}^* \varphi \iff \Sigma' \models_{\mathfrak{A}}^* \varphi$. That $\Sigma \subset \Sigma'$ means $\Sigma \models_{\mathfrak{A}}^* \varphi \implies \Sigma' \models_{\mathfrak{A}}^* \varphi$, so it is sufficient to prove $\Sigma \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner \implies \Sigma' \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner$. We examine the possible cases below.

Case 1: $x \notin \mathcal{L}$. Then $x \notin \Sigma'$, so $\Sigma' \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner$.

Case 2: $x \in L$ and φ occurs in $\mathsf{pos}(\phi)$. Since $\mathsf{check}(\mathsf{pos}(\phi))$ logically implies $\mathsf{pos}(\phi)$ and $\Sigma \models_{\mathfrak{A},\nu}^* \mathsf{check}(\mathsf{pos}(\phi))$, also $\Sigma \models_{\mathfrak{A},\nu}^* \mathsf{pos}(\phi)$. By the definition of pos, $\ulcorner \neg \varphi \urcorner$ must not occur in $\mathsf{pos}(\phi)$, so φ is a conjunct of $\mathsf{pos}(\phi)$ and $\Sigma \models_{\mathfrak{A}}^* \varphi$.

Case 3: $x \in L$, φ does not occur in $\mathsf{pos}(\phi)$ and $\ulcorner \neg \varphi \urcorner$ does not occur in $\nu^*(\mathsf{check}(\mathsf{pos}(\phi)))$. If φ occurs in $\nu^*(\mathsf{pos}(\phi))$ then by the same argument as in Case 2, $\Sigma \models_{\mathfrak{A}}^* \varphi$. Otherwise, φ is $\nu^*(\varphi')$ for some $\varphi' \in \mathrm{CS}(\mathsf{pos}(\phi))$. By the definition of CS and the fact that $\Sigma \models_{\mathfrak{A},\nu}^* \mathsf{check}(\mathsf{pos}(\phi))$, $\Sigma \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner$ means $x \notin \mathcal{L}$, in which case also $\Sigma' \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner$.

Case 4: $x \in L$, φ does not occur in $\mathsf{pos}(\phi)$ and $\ulcorner \neg \varphi \urcorner$ occurs in $\nu^*(\mathsf{check}(\mathsf{pos}(\phi)))$. Here $\ulcorner \neg \varphi \urcorner$ must occur in $\nu^*(\mathsf{pos}(\phi))$, so by the definition of CS, $\ulcorner (x \in \mathcal{L} \implies E(\neg x)) \urcorner$ occurs in $\nu^*(\mathsf{check}(\mathsf{pos}(\phi)))$. Now $\Sigma \models_{\mathfrak{A},\nu}^* \mathsf{check}(\mathsf{pos}(\phi))$ implies either $x \notin \mathcal{L}$ or $\neg x \in \Sigma$. In the former case $x \notin \Sigma'$; in the latter case $\neg x \in \Sigma'$ and Σ' being $\mathcal{L}$-nice tell us that $x \notin \Sigma'$. In both cases, we end up with $\Sigma' \models_{\mathfrak{A}}^* \ulcorner \neg \varphi \urcorner$. $\qquad\square$

3.2 Forcing Notions and Universal Sentences

Consider a forcing notion with conditions fragments of $\mathcal{L}$, ordered by reverse inclusion. Then genericity over $\mathbb{P}$ naturally gives us a subset of $\mathcal{L}$. We want to analyse this subset using $\mathcal{L}_{\mathfrak{A}}^*$.

Definition 3.24. A pair $(\mathfrak{A} = (A; \in, \vec{R}), \mathbb{P} = (P, \leq_P))$ is *good for $\mathcal{L}$* iff

- $\mathfrak{A}$ is $\mathcal{L}$-suitable,
- $\emptyset \neq P \subset \mathcal{P}(\mathcal{L}) \cap A$, and
- for all $\{p, q\} \subset P$, $p \leq_P q$ iff $q \subset p$.

  For the rest of this subsection, we fix a forcing notion $\mathbb{P} = (P, \leq_P))$ for which

- $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$, and
- $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}$.

As in the case for $\mathcal{L}$, we will assume $\mathbb{P}$ occurs in $\vec{R}$ without any loss of generality.

Definition 3.25. For $p \in P$, a *p-candidate for $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universality* is a set Σ for which

- $p \subset \Sigma$,

- for each $x \in [\Sigma]^{<\omega}$, there is $q \in P$ with $p \cup x \subset q$, and

- Σ is $\mathcal{L}$-nice.

Definition 3.26. Let $\phi \in \mathcal{L}^*_{\mathfrak{A}}$ be a sentence and $p \in P$. We say ϕ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universal for p* iff for all $q \leq_P p$ there is a set Σ in some weak outer model of V such that

- Σ is a q-candidate for $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universality, and

- $\Sigma \models^*_{\mathfrak{A}} \phi$.

We say ϕ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universal* iff for all $p \in P$, ϕ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal p.

For $\Gamma \subset \mathcal{L}^*_{\mathfrak{A}}$, we say Γ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universal* (*for p*) iff ϕ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal (for p) for all $\phi \in \Gamma$.

For notational convenience in the proof of Lemma 3.28 — indubitably our most important technical lemma — we introduce the following definition.

Definition 3.27. For $p \in \mathbb{P}$, let $\mathcal{F}_p$ denote the set

$$\{\Sigma : \Sigma \text{ is a } p\text{-candidate for } (\mathfrak{A}, \mathbb{P}, \mathcal{L})\text{-universality in some weak outer model of } V\}.$$

Lemma 3.28. *Let*

- *W be a weak outer model of V,*

- *$g \in W \cap \mathcal{P}(P)$ be a $\mathbb{P}$-generic filter over $\mathfrak{A}$,*

- *$p \in g$, and*

- *ϕ be a $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_2$ sentence which is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p.*

Then $\bigcup g$ is $\mathcal{L}$-nice and $\bigcup g \models^*_{\mathfrak{A}} \phi$.

Proof. We prove $\bigcup g \models^*_{\mathfrak{A}} \phi$ by induction on the length of ϕ. The proof that $\bigcup g$ is $\mathcal{L}$-nice will surface as a part of the induction argument.

Case 1: ϕ is $\mathcal{L}^*_{\mathfrak{A}}$-$\Delta_0$. By way of contradiction, assume $\bigcup g \models^*_{\mathfrak{A}} \neg\phi$. Since $\neg\phi$ is also a $\mathcal{L}^*_{\mathfrak{A}}$-$\Delta_0$ sentence, Lemma 3.20 tells us there are finite sets $p^\dagger \in A \cap \mathcal{P}(\mathcal{L})$ and $q^\dagger \in A \cap \mathcal{P}(\mathcal{L})$ such that $\neg\phi$ is $\models^*_{\mathfrak{A}}$-true for $(p^\dagger, q^\dagger)$. For each $z \in p^\dagger$, pick a $p_z \in g$ such that $z \in p_z$. Since g is a filter, there is some $p^* \in g$ for which $p^\dagger \subset (\bigcup_{z \in p^\dagger} p_z) \cup p \subset p^*$. If $p' \leq_P p^*$, then $p' \leq_P p$, so by ϕ being $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p, there is $\Sigma \in \mathcal{F}_{p'}$ with $\Sigma \models^*_{\mathfrak{A}} \phi$, whence $p^\dagger \subsetneq \Sigma \cap q^\dagger$. By our choice of Σ, we can find $q \in P$ with $p' \cup (\Sigma \cap q^\dagger) \subset q$. Necessarily, $q \leq_P p'$ and $p^\dagger \subsetneq q \cap q^\dagger$.

We have thus shown that the set

$$D_1 := \{q \in P : p^\dagger \subsetneq q \cap q^\dagger \text{ or } q \perp_{\mathbb{P}} p^*\}$$

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is dense in $\mathbb{P}$. D_1 is obviously definable in the language associated with $\mathfrak{A}$, so there is $q^* \in g \cap D_1$. That g is a filter containing p^* gives us $p^\dagger \subsetneq q^* \cap q^\dagger$, which implies

$$p^\dagger \subsetneq q^* \cap q^\dagger \subset \bigcup g \cap q^\dagger = p^\dagger,$$

a contradiction.

Case 2: ϕ is $\mathcal{L}_\mathfrak{A}^*$-$\Sigma_1$ but not $\mathcal{L}_\mathfrak{A}^*$-$\Delta_0$. Then ϕ is of the form $\ulcorner \exists x_1 ... \exists x_n \, \phi' \urcorner$ for some

- $\mathcal{L}_\mathfrak{A}^*$-$\Delta_0$ formula ϕ',
- n such that $1 \leq n < \omega$, and
- $\{x_1, ..., x_n\}$ the set of free variables of ϕ'.

We first show that $\bigcup g$ is $\mathcal{L}$-nice. Obviously, $\bigcup g \subset \mathcal{L}$, so we need only consider the other two conditions of being $\mathcal{L}$-nice. To that end, define

$$\varphi_x := \ulcorner (E(x) \vee E(\neg x)) \wedge (\neg E(x) \vee \neg E(\neg x)) \urcorner$$

for each $x \in \mathcal{L}$. Note that the φ_x's are $\mathcal{L}_\mathfrak{A}^*$-$\Delta_0$ sentences. Moreover, $X \models_\mathfrak{A}^* \varphi_x$ for every $\mathcal{L}$-nice set X and every $x \in \mathcal{L}$.

Let $x \in \mathcal{L}$ and $p' \leq_P p$. That ϕ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p means $\mathcal{F}_{p'} \neq \emptyset$, so choose any $\Sigma \in \mathcal{F}_{p'}$. We must have $\Sigma \models_\mathfrak{A}^* \varphi_x$ because Σ is $\mathcal{L}$-nice. This allows us to conclude that for all $x \in \mathcal{L}$, φ_x is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p. As Case 1 has been proven, we can apply it to yield

$$\bigcup g \models_\mathfrak{A}^* \varphi_x \text{ for all } x \in \mathcal{L},$$

which is just another way of saying $\bigcup g$ fulfils the last two conditions of Definition 3.5. We have thus shown that $\bigcup g$ is $\mathcal{L}$-nice.

Once more, let $p' \leq_P p$, so that there is $\Sigma \in \mathcal{F}_{p'}$ for which $\Sigma \models_\mathfrak{A}^* \phi$. Let ν be a $\mathfrak{A}$-valuation such that $\Sigma \models_{\mathfrak{A}, \nu}^* \phi'$. Then there is a disjunct φ of $\mathsf{DNF}(\phi')$ for which $\Sigma \models_\mathfrak{A}^* \nu^*(\varphi)$. Σ being $\mathcal{L}$-nice and Lemma 3.10 gives us $\Sigma \models_\mathfrak{A}^* \mathsf{pos}(\nu^*(\varphi))$. Since literals of the form

$$\ulcorner \neg E(x) \urcorner \text{ for some } x \in \mathcal{L}$$

do not occur in $\mathsf{pos}(\nu^*(\varphi))$, following the proof of Case 1 of Lemma 3.20, there must be a finite set $p^\dagger \in A \cap \mathcal{P}(\Sigma)$ such that $\mathsf{set}(\mathsf{pos}(\nu^*(\varphi))) = (p^\dagger, p^\dagger)$ and $\mathsf{pos}(\nu^*(\varphi))$ is $\models_\mathfrak{A}^*$-true for $p^\dagger$. By our choice of Σ, we can find $q \in P$ with $p' \cup p^\dagger \subset q$. In particular, $q \leq_P p'$ and $q \models_\mathfrak{A}^* \mathsf{pos}(\nu^*(\varphi))$.

The $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universality of p tells us that q can be extended to an $\mathcal{L}$-nice set. Necessarily, $q \models_\mathfrak{A}^* \ulcorner \neg E(x) \urcorner$ whenever $x \in \mathcal{L}$ and $\ulcorner \neg E(x) \urcorner$ occurs in $\nu^*(\varphi)$, so $q \models_\mathfrak{A}^* \nu^*(\mathsf{pos}(\varphi))$ too. If $\varphi' \in \mathsf{CS}(\mathsf{pos}(\varphi))$, then $q \models_\mathfrak{A}^*$

$\mathsf{pos}(\nu^*(\varphi))$ implies $q \models^*_{\mathfrak{A}} \nu^*(\varphi')$. So indeed, $q \models^*_{\mathfrak{A}} \nu^*(\mathsf{check}(\mathsf{pos}(\varphi)))$. As $\nu^*(\mathsf{check}(\mathsf{pos}(\varphi)))$ logically implies

$$\phi'' := \ulcorner\exists x_1...\exists x_n \left(\bigvee\{\mathsf{check}(\mathsf{pos}(\varphi)) : \varphi \text{ is a disjunct of } \mathsf{DNF}(\phi')\}\right)\urcorner,$$

we have $q \models^*_{\mathfrak{A}} \phi''$ as well.

This means the set

$$D_2 := \{q \in P : q \models^*_{\mathfrak{A}} \phi'' \text{ or } q \perp_{\mathbb{P}} p\}$$

is dense in $\mathbb{P}$. By Lemma 3.9, D_2 is definable in the language associated with $\mathfrak{A}$, so there is $q^* \in g \cap D_2$. That g is a filter containing p gives us $q^* \models^*_{\mathfrak{A}} \phi''$, which means we can find ν' a $\mathfrak{A}$-valuation and φ^* a disjunct of $\mathsf{DNF}(\phi')$ for which $q^* \models^*_{\mathfrak{A},\nu'} \mathsf{check}(\mathsf{pos}(\varphi^*))$. Now, seeing that $\bigcup g$ is $\mathcal{L}$-nice, and that φ^* takes the form required of ϕ in Lemma 3.23, we can invoke said lemma to get $\bigcup g \models^*_{\mathfrak{A},\nu'} \mathsf{check}(\mathsf{pos}(\varphi^*))$, or equivalently, $\bigcup g \models^*_{\mathfrak{A}} (\nu')^*(\mathsf{check}(\mathsf{pos}(\varphi^*)))$.

According to the definition of check in Definition 3.21, it is immediate that $(\nu')^*(\mathsf{check}(\mathsf{pos}(\varphi^*)))$ logically implies $(\nu')^*(\mathsf{pos}(\varphi^*))$, so we have $\bigcup g \models^*_{\mathfrak{A}} (\nu')^*(\mathsf{pos}(\varphi^*))$. By Lemma 3.10 and the fact that $\bigcup g$ is $\mathcal{L}$-nice, $\bigcup g \models^*_{\mathfrak{A}} (\nu')^*(\varphi^*)$. Finally, $(\nu')^*(\varphi^*)$ logically implies ϕ'', which in turn, logically implies ϕ — the latter implication being again, due to the definition of check. This nets us $\bigcup g \models^*_{\mathfrak{A}} \phi$.

Case 3: $\phi = \ulcorner\forall x\ \varphi(x, \vec{d})\urcorner$ for some x, φ and $\vec{d}$. Then for each $a \in A$, $\varphi(a, \vec{d})$ is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p. By the induction hypothesis, $\bigcup g \models^*_{\mathfrak{A}} \varphi(a, \vec{d})$ for all $a \in A$, so also $\bigcup g \models^*_{\mathfrak{A}} \phi$. $\qquad\square$

Remark 3.29.

(1) Let p be an arbitrary member of P and define

$$\mathbb{P}_{\leq p} := (\{q \in P : q \leq_{\mathbb{P}} p\}, \leq_{\mathbb{P}}).$$

Then $\mathbb{P}_{\leq p}$ is definable in the language associated with $\mathfrak{A}$ and $(\mathfrak{A}, \mathbb{P}_{\leq p})$ is good for $\mathcal{L}$. In addition, any $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_2$ sentence which is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p is also $(\mathfrak{A}, \mathbb{P}_{\leq p}, \mathcal{L})$-universal. Consequently, Lemma 3.28 is equivalent to, and can be restated as:

Lemma 3.28′. *Let*

- W *be a weak outer model of* V,
- $g \in W \cap \mathcal{P}(P)$ *be a* $\mathbb{P}$-*generic filter over* $\mathfrak{A}$, *and*

- ϕ *be a* $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_2$ *sentence which is* $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universal.*

Then $\bigcup g$ *is* $\mathcal{L}$-*nice and* $\bigcup g \models^*_{\mathfrak{A}} \phi$.

(2) Notice that we can substitute the requirement of g being $\mathbb{P}$-generic over $\mathfrak{A}$ with the weaker requirement

$$\phi_{wgen} := \text{``}g \cap D \neq \emptyset \text{ for all } D \text{ such that}$$

- $D \subset P$
- D is dense in $\mathbb{P}$, and
- D is definable in the language associated with $\mathfrak{A}$"

in Lemma 3.28, and still have the same proof go through. However, since we enforced $P \subset A$ from the get-go, g being $\mathbb{P}$-generic over $\mathfrak{A}$ is equivalent to ϕ_{wgen}, and the lemma is no stronger after the substitution.

(3) It is easy to see that the dense sets g is required to meet in the proof of Lemma 3.28 have reasonably simple definitions. Indeed, since we have assumed both $\mathbb{P}$ and $\mathcal{L}$ occur in $\vec{R}$, said dense sets are all Δ_2-definable in the language associated with $\mathfrak{A}$. In light of (2), we can substitute the requirement of g being $\mathbb{P}$-generic over $\mathfrak{A}$ with the weaker requirement

"$g \cap D \neq \emptyset$ for all D such that

- $D \subset P$
- D is dense in $\mathbb{P}$, and
- D is Δ_2-definable in the language associated with $\mathfrak{A}$"

in Lemma 3.28, without affecting its proof.

A natural strengthening of Lemma 3.28 is to have ϕ be an arbitrary $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_3$ sentence which is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal for p. As per Case 3 in the proof of Lemma 3.28, we can always get the outermost universal quantification for free, so we only have to prove the strengthened lemma assuming ϕ is $\mathcal{L}^*_{\mathfrak{A}}$-$\Sigma_2$ instead of $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_3$. However, the nice "characterisation" of $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_2$ sentences we will uncover in Section 5 (brought about by Theorems 5.32 and 5.33) seems to suggest that such a strengthening is impossible. With (1) of Remark 3.29 in mind, the following conjecture is then sufficiently motivated.

Conjecture 3.30. *There are sets* $\mathfrak{A}$, $\mathcal{L}$, $\mathbb{P} = (P, \leq_{\mathbb{P}})$, W, g, *and* ϕ *such that*

- $\mathcal{L}$ *is closed under negation.*
- $\mathbb{P}$ *is definable in the language associated with* $\mathfrak{A}$,

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- $(\mathfrak{A}, \mathbb{P})$ *is good for* $\mathcal{L}$,

- W *is a weak outer model of* V,

- $g \in W \cap \mathcal{P}(P)$ *is a* $\mathbb{P}$-*generic filter over* $\mathfrak{A}$,

- ϕ *is a* $\mathcal{L}_{\mathfrak{A}}^*$-$\Sigma_2$ *sentence which is* $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universal, and*

- $\bigcup g \not\models_{\mathfrak{A}}^* \phi$.

3.3 A Useful Framework

Fix a set of $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ sentences, Γ, for this subsection.

In the previous subsection, we saw how a forcing notion $\mathbb{P}$ can generate witnesses to certain $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ sentences when $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$ and $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}$. Leveraging on this fact, we shall develop a framework for defining forcing notions that generate witnesses to a given set of $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ sentences.

This framework both generalises and is inspired by the forcing construction Asperó and Schindler carried out in the proof of the main theorem of [13].

Definition 3.31. A set B is $\mathcal{L}$-*closed under finite extensions* iff

- $B \subset \mathcal{P}(\mathcal{L})$, and

- for all $x \in B$ and all $y \in [\mathcal{L}]^{<\omega}$, $x \cup y \in B$.

Definition 3.32. For any Σ and any p, we say Σ $\Gamma(\mathcal{L}, \mathfrak{A})$-*certifies* p iff

1. $p \subset \Sigma$,

2. Σ is $\mathcal{L}$-nice, and

3. $\Sigma \models_{\mathfrak{A}}^* \Gamma$.

It is easy to see that if $\mathcal{L}$, $\mathfrak{A}$, Γ, Σ and p are such that Σ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies p, then for all $q \subset \Sigma$, Σ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies q. This gives us the following proposition.

Proposition 3.33. *Let* $|\mathcal{L}| \leq \lambda$ *and* B *be* $\mathcal{L}$-*closed under finite extensions. If we define* $\mathbb{P} := (P, \leq_{\mathbb{P}})$, *where*

$$P := \{p \in B : \Vdash_{Col(\omega, \lambda)} \exists \Sigma \ (\text{``}\Sigma \ \Gamma(\mathcal{L}, \mathfrak{A})\text{-}certifies \ p\text{''})\}, \ and$$
$$\leq_{\mathbb{P}} := \{(p, q) \in P \times P : q \subset p\},$$

then as long as $P \neq \emptyset$,

- $(\mathfrak{A}, \mathbb{P})$ *is good for* $\mathcal{L}$, *and*

- *whenever* $p \in P$ *and* $\Sigma \ \Gamma(\mathcal{L}, \mathfrak{A})$-*certifies* p, Σ *is a* p-*candidate for* $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-*universality.*

Lemma 3.34. *Let $|trcl(\mathfrak{A})| \leq \lambda$ and $p \subset \mathcal{L}$. Assume there is Σ in a weak outer model W of V such that Σ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies p. Then*

$$\Vdash_{Col(\omega,\lambda)} \exists\Sigma \ (\text{``}\Sigma \ \Gamma(\mathcal{L}, \mathfrak{A})\text{-}certifies \ p\text{''}).$$

Proof. Suppose otherwise, so there is $q \in Col(\omega, \lambda)$ such that

$$q \Vdash_{Col(\omega,\lambda)} \neg\psi, \tag{3.1}$$

where
$$\psi := \exists\Sigma \ (\text{``}\Sigma \ \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''}).$$

Let g be $Col(\omega, \lambda)$-generic over W with $q \in g$, so that g is also $Col(\omega, \lambda)$-generic over V. First, that $W \models \psi$ means $W[g] \models \psi$. Next, notice that if $\varphi(\Sigma, y)$ is the conjunction of the statements

- $y = \{\phi \in \mathcal{L}^*_{\mathfrak{A}} : \Sigma \models^*_{\mathfrak{A}} \phi\}$,
- $\Gamma \subset y$,
- $p \subset \Sigma$, and
- "Σ is $\mathcal{L}$-nice",

then φ is a Σ_1 formula in the language of set theory, with parameters among $p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma$. This is because φ is equivalent to the statement of there being a function f with domain $\mathcal{L}^*_{\mathfrak{A}}$ — a Δ_0-definable subset of A — and codomain $\{0, 1\}$, such that

- f fulfils the inductive properties of Tarski's definition of the satisfaction relation, applied to the structure $(A; \in, \vec{R}, \Sigma)$,
- $y = \{\phi \in \mathcal{L}^*_{\mathfrak{A}} : f(\phi) = 1\}$,
- $\Gamma \subset y$,
- $p \subset \Sigma$, and
- "Σ is $\mathcal{L}$-nice",

every of which aforementioned points are Δ_0-expressible in the language of set theory. Note also that whenever $\varphi(\Sigma, y)$ holds, Σ, y and any witness f must have transitive closures of cardinalities no larger than λ. Moreover, we have

$$\psi \iff \exists\Sigma \ \exists y \ \varphi(\Sigma, y).$$

As $p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma$ are subsets of $trcl((p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma))$ and $|trcl((p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma))| \leq \lambda$, the structure
$$\mathfrak{B} := (trcl((p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma)); \in, p, \mathfrak{A}, \mathcal{L}, \mathcal{L}^*_{\mathfrak{A}}, \Gamma)$$

can be coded as a real in $V[g]$, by Lemma 2.9. This means that in all weak outer models of $V[g]$, ψ can be thought of as a Σ_1^1 sentence involving a real code of $\mathfrak{B}$ found in $V[g]$. In particular, ψ is absolute for $V[g]$ and $W[g]$. Now $W[g] \models \psi$ implies $V[g] \models \psi$, contradicting (3.1). $\qquad \square$

Note that the proof of Lemma 3.34 does not require that Γ contains only $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ sentences. Indeed, for a litany of properties K, the existence of an object satisfying K is absolute between V and its weak outer models and hence, between V and its forcing extensions. However, it is often useful — if not integral — to have a proper handle on such an object. It is towards this end that we are often interested in the existence of a V-generic object k such that k satisfies K in $V[k]$.

Specifying "k satisfies K" to be "$k \models_{\mathfrak{A}}^* \Gamma$", the following lemma is thus well-motivated.

Lemma 3.35. *Let*

- $|trcl(\mathfrak{A})| \leq \lambda$,
- B *be $\mathcal{L}$-closed under finite extensions,*
- $P := \{p \in B : \Vdash_{Col(\omega,\lambda)} \exists \Sigma\ ("\Sigma\ \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p")\}$,
- $\leq_{\mathbb{P}} := \{(p, q) \in P \times P : q \subset p\}$,
- $\mathbb{P} := (P, \leq_{\mathbb{P}})$,
- $\mathbb{P}$ *be definable in the language associated with $\mathfrak{A}$,*
- W *be a weak outer model of V,*
- $g \in W \cap \mathcal{P}(P)$ *be a filter on $\mathbb{P}$, and*
- g *is $\mathbb{P}$-generic over $\mathfrak{A}$.*

If there is Σ in a weak outer model W' of V such that $\Sigma\ \Gamma(\mathcal{L}, \mathfrak{A})$-certifies $\emptyset$, then $\bigcup g$ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies $\emptyset$.

In particular, if g is $\mathbb{P}$-generic over V, then $\bigcup g$ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies $\emptyset$ in $V[g] = V[\bigcup g]$.

Proof. The general statement is clear from Lemma 3.28, (1) of Remark 3.29, Proposition 3.33 and Lemma 3.34. That $\bigcup g$ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies $\emptyset$ is absolute for transitive models of $\mathsf{ZFC} - \mathsf{Powerset}$, so if g is $\mathbb{P}$-generic over V, then $\bigcup g$ $\Gamma(\mathcal{L}, \mathfrak{A})$-certifies $\emptyset$ in $V[g]$. Moreover, since $g = [\bigcup g]^{<\omega}$, we have $V[g] = V[\bigcup g]$. $\qquad \square$

We close this subsection with another absoluteness result. This time, instead of looking for witnesses in forcing extensions, we turn our focus to the forcing notions themselves.

Lemma 3.36. *The definition of $\mathbb{P}$ from parameters B, $\mathcal{L}$, $\mathfrak{A}$, Γ in Proposition 3.33, where λ is additionally specified to be $|trcl(\mathfrak{A})|$, is absolute for transitive models of ZFC.*

Proof. It suffices to show that the set

$$\{p \in B : \Vdash_{Col(\omega, |trcl(\mathfrak{A})|)} \exists \Sigma \, (\text{``}\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''})\} \tag{3.2}$$

is absolute for transitive models of ZFC.

Let V' and W be transitive models of ZFC such that $\{B, \mathcal{L}, \mathfrak{A}, \Gamma\} \subset V' \subset W$. Have $P^{V'}$ and P^W denote the versions of the set (3.2) defined in V' and W respectively. We want to prove $P^{V'} = P^W$.

First note that

$$Col(\omega, |trcl(\mathfrak{A})|)^{V'} = Col(\omega, |trcl(\mathfrak{A})|^{V'}) \cong Col(\omega, |trcl(\mathfrak{A})|^W) = Col(\omega, |trcl(\mathfrak{A})|)^W$$

in W, so

$$p \in P^W \iff \Vdash_{Col(\omega, |trcl(\mathfrak{A})|)^{V'}} \exists \Sigma \, (\text{``}\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''}) \tag{3.3}$$

in W. Since any forcing extension of W is a weak outer model of V, a direct application of Lemma 3.34 gives us $P^W \subset P^{V'}$. Next, fix any $p \in P^{V'}$ and any $Col(\omega, |trcl(\mathfrak{A})|)^{V'}$-generic filter g over W. Now g is also a $Col(\omega, |trcl(\mathfrak{A})|)^{V'}$-generic filter over V', so $V'[g]$ and $W[g]$ are transitive models of ZFC and moreover, $V'[g] \subset W[g]$. By the definition of $P^{V'}$ in V',

$$V'[g] \models \text{``}\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''}$$

for some $\Sigma \in V'[g]$. That $\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})$-certifying p is absolute for transitive models of ZFC implies

$$W[g] \models \text{``}\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''}.$$

We have thus shown

$$W \models (\Vdash_{Col(\omega, |trcl(\mathfrak{A})|)^{V'}} \exists \Sigma \, (\text{``}\Sigma \, \Gamma(\mathcal{L}, \mathfrak{A})\text{-certifies } p\text{''})),$$

whence $p \in P^W$ by (3.3). This allows us to conclude $P^{V'} = P^W$. $\qquad\square$

4 Extending Namba Forcing

This section illustrates how the framework introduced in Subsection 3.3 can be applied to resolve a problem in set theory.

4.1 An Extension Problem

Before stating our problem of interest, we feel obliged to present, at least in brief, the history surrounding it.

Fix a limit ordinal α and consider the chain of inequalities

$$cof(\alpha) \leq |\alpha| \leq \alpha, \tag{4.1}$$

which is provable in ZFC. Set theorists have long investigated the ability to change the signs in (4.1) via forcing. If $|\alpha| < \alpha$ in V, then the same must hold in any forcing extension . If $cof(\alpha) < |\alpha|$ in V, we can always force $cof(\alpha) = |\alpha|$ by collapsing both $cof(\alpha)$ and $|\alpha|$ to a regular cardinal in V no greater than $cof(\alpha)$.

On the flipside, if $|\alpha| = \alpha$ in V, then α is a cardinal there. As long as α is uncountable, a forcing notion that collapses α exists and necessarily forces $|\alpha| < \alpha$. We are left with the case where $cof(\alpha) = |\alpha|$ in V. Note that by swapping α with a smaller ordinal if necessary, we can assume α is regular in V without loss of generality. So assume α is an uncountable regular cardinal in V. Forcing $cof(\alpha) < |\alpha|$ under this assumption turned out to be highly non-trivial.

In his doctoral dissertation [3], Prikry assumed α is a measurable cardinal, and gave an example of a forcing notion that preserves all cardinalities, yet changes $cof(\alpha)$ to ω. A natural follow-up question to Prikry's result is thus:

> can we force the the separation of $cof(\alpha)$ and $|\alpha|$ on an uncountable regular α which provably exist over ZFC?

As successor cardinals are the only uncountable regular cardinals proven to exist over ZFC, a forcing notion separating $cof(\alpha)$ and $|\alpha|$ for any such α must collapse α. But can we ensure α is not collapsed "too far"? In other words, we want to force $cof(\alpha) < |\alpha|$ while preserving all cardinals below α.

The late 1960s saw two independent solutions to this problem in the affirmative, by Bukovský [6] and Namba [4]. Both solutions work with $\alpha = \omega_2$, which is the smallest possible value α can take in an affirmative answer. Simplifications were made to the presentation of Bukovský's and Namba's forcing notions over the years, without losing sight of the goal of their constructions. These simplifications culminated in what is now commonly known as *Namba forcing*. Since the focus of this section is on extending the key effects of Namba forcing, we feel obliged to define the forcing notion for the sake of completeness.

Definition 4.1. We say $(T, \leq)$ is a κ-splitting in A iff

- $(T, \leq)$ is a tree, and
- for every $s \in T \cap A$, s has κ many immediate $\leq$-successors in T.

Definition 4.2. If $(T, \leq)$ is a partial order and $s \in T$, we use $T_s^{\leq}$ to denote the set of $\leq$-successors of s in T. More formally,

$$T_s^{\leq} := \{t \in T : s \leq t\}.$$

Definition 4.3. Define the order $\leq^{\dagger}$ to be

$$\{(s, t) \in \omega_2^{<\omega} \times \omega_2^{<\omega} : dom(s) \subset dom(t) \text{ and } t {\restriction}_{dom(s)} = s\}.$$

Definition 4.4. A *Namba tree* is a subset T of $\omega_2^{<\omega}$ containing a *root* s such that

- $(T, \leq^{\dagger})$ is ω_2-splitting in $T_s^{\leq^{\dagger}}$, and
- whenever $t \in T$, either $s \leq^{\dagger} t$ or $t \leq^{\dagger} s$.

Definition 4.5 (Namba). Define

$$P_N := \{T \subset \omega_2^{<\omega} : T \text{ is a Namba tree}\}, \text{ and}$$
$$\leq_N := \{(p, q) \in P_N \times P_N : p \subset q\}.$$

We call the forcing notion $\mathbb{P}_N := (P_N, \leq_N)$ *Namba forcing*.

Namba forcing belongs to the class of *uniformly-splitting tree forcings*, one of which earliest-known members is Mathias forcing. A typical condition of a uniformly-splitting tree forcing is a tree, and it can be divided into two components — the stem and the crown. The stem is the main working part of a condition; stems in a generic filter combine to form a function that is the primary generic object we desire. The crowns work as *side conditions*, which in unity, endow the forcing notion with specific regularity properties. These properties are often crucial to the satisfaction of constraints placed on the forcing extension. If T is a Namba tree with root s, then its stem is

$$\{t \in T : t \subset s\}$$

and its crown is $T_s^{\leq^{\dagger}} \setminus \{s\}$.

By means of tree combinatorics, one can show that Namba forcing gives ω_2^V a cofinality of ω without collapsing ω_1^V. In fact, Namba forcing is a textbook example of such a forcing notion. It also has a stronger property than not collapsing ω_1^V — it is stationary-preserving. In the parlance of the preceding paragraph, the primary generic object here is a cofinal function from ω into ω_2^V, whereas the regularity property of pertinence is being stationary-preserving. We can then observe the following division of labour: the stems of Namba forcing are in charge of changing the cofinality of ω_2^V to ω, while the crowns of Namba forcing ensure all stationary subsets of ω_1 in V have their stationarity preserved.

The extended Namba problem, at its most rudimentary, asks (in V) for which regular cardinals λ above ω_2 is the statement

$Nb_0(\lambda) := $ 'there is a stationary-preserving forcing notion $\mathbb{P}$ such that

$\Vdash_{\mathbb{P}}$ "$cof(\alpha) = \omega$" for all regular cardinals α satisfying $\omega_2 \leq \alpha < \lambda$'

true. This formulation of the problem is already non-trivial, as the following fact tells us that iterating Namba forcing in the usual way is insufficient to get us further than $Nb_0(\omega_3^V)$, without assuming a strong failure of GCH.

Fact 4.6. Assume GCH holds below ω_2. Then

(1) $\Vdash_{\mathbb{P}_N}$ "$cof(\omega_3^V) = \omega_1$", and

(2) $\Vdash_{\mathbb{P}_N}$ "$cof(\beta) > \omega$" for all regular cardinals β satisfying $\omega_3 < \beta$.

Remark 4.7. That the naive two-step iteration of Namba forcing fails to be stationary-preserving assuming the hypothesis of Fact 4.6, suggests our natural conception of iteration is incompatible with the side conditions of Namba forcing. To overcome this incompatibility, it makes sense to consider either a new kind of iteration or an overhaul of the side conditions. The latter option could mean a departure from the intuition of uniformly-splitting tree forcings so radical, that the resultant forcing notion has conditions best presented as objects other than trees.

Indeed, a stronger variant of the extended Namba problem, asking for which regular cardinals λ above ω_2 is the statement

$Nb_1(\lambda) := $ 'there is a stationary-preserving forcing notion $\mathbb{P}$ such that

$\Vdash_{\mathbb{P}}$ "$cof(\alpha) = \omega$" for all regular cardinals α satisfying $\omega_2 \leq \alpha < \lambda$,

$\Vdash_{\mathbb{P}}$ "$cof(\lambda) = \omega_1$", and

$\Vdash_{\mathbb{P}}$ "$cof(\beta) > \omega$" for all regular cardinals β satisfying $\lambda < \beta$'

true, naturally arises from Fact 4.6.

Drawing from the deep and complex theories of subcomplete forcing and $\mathcal{L}$-forcing, Jensen showed in [12] that, modulo weak fragments of GCH, $Nb_1(\lambda)$ holds for all successor and strongly inaccessible cardinals above ω_2. Jensen used very different methods to construct the witnesses $\mathbb{P}$ for different categories of λ, but in doing so, he also ensured that $\mathbb{P}$ never adds reals.

But can we have $Nb_1(\lambda)$ hold for a bigger class of cardinals λ if we allow $\mathbb{P}$ to add reals?

4.2 A Conditional Solution

It turns out there is a relatively simple proof of

"$Nb_1(\lambda)$ for all λ above ω_2"

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if we assume something of greater consistency strength. This proof adopts a novel side-condition technique first employed in [13] (cf. Remark 4.7). It also demonstrates how amenable the forcing framework of Subsection 3.3 is in bolstering natural and obvious forcing conditions with said side conditions.

Theorem 4.8. *Assume the non-stationary ideal on ω_1 is precipitous. Let λ_f be a regular cardinal above 2^{ω_1}. Then there is a stationary-preserving forcing notion $\mathbb{P}$ such that*

(1) $\Vdash_{\mathbb{P}}$ "$cof(\alpha) = \omega$" for all regular cardinals α satisfying $\omega_2 \leq \alpha < \lambda_f$,

(2) $\Vdash_{\mathbb{P}}$ "$cof(\lambda_f) = \omega_1$", and

(3) $\Vdash_{\mathbb{P}}$ "$cof(\beta) > \omega$" for all regular cardinals β satisfying $\lambda_f < \beta$.

In other words, $Nb_1(\lambda_f)$ holds for all λ_f above 2^{ω_1}.

Proof. Fix a regular cardinal λ_f as specified in the statement of the theorem.

Let h be a generic filter on $Col(\lambda_f, \lambda_f)$, so that in $V[h]$,

- $cof(\alpha) = cof^V(\alpha)$ for all $\alpha \leq \lambda_f$,
- the non-stationary ideal on ω_1 is still precipitous,
- $|H(\lambda_f)| = \lambda_f$, and
- there is a a $\Diamond_{\lambda_f}$-sequence $(\bar{A}_\lambda : \lambda < \lambda_f)$.

Lemma 4.9. *If in $V[h]$ there is a stationary-preserving forcing notion $\mathbb{P}$ of size at most λ_f fulfilling (1) and (2) of Theorem 4.8, then $Col(\lambda_f, \lambda_f) * \dot{\mathbb{P}}$ is a stationary-preserving forcing notion fulfilling (1) to (3) of Theorem 4.8 in V, where $\dot{\mathbb{P}}$ is a $Col(\lambda_f, \lambda_f)$-name for $\mathbb{P}$.*

Proof. In V, $Col(\lambda_f, \lambda_f) * \dot{\mathbb{P}}$ is a stationary-preserving forcing notion fulfilling (1) and (2) of Theorem 4.8 because

- $Col(\lambda_f, \lambda_f)$ is a stationary-preserving forcing notion forcing $\dot{\mathbb{P}}$ to be such, and
- $Col(\lambda_f, \lambda_f)$ forces $cof(\alpha) = cof^V(\alpha)$ for all $\alpha \leq \lambda_f$.

To show $Col(\lambda_f, \lambda_f) * \dot{\mathbb{P}}$ fulfils (3) of Theorem 4.8, let $\beta > \lambda_f$ be regular in V. Then $cof^{V[h]}(\beta) \geq \lambda_f$ because $Col(\lambda_f, \lambda_f)$ is λ_f-closed. That $\mathbb{P}$ is of size at most λ_f implies it preserves cofinalities $\geq (\lambda_f^+)^{V[h]}$. That it also fulfils (2) of Theorem 4.8 implies it forces $cof(\beta) > \omega$. $\square$

Allow W to denote $V[h]$. Unless otherwise stated, we work in W from now on. For brevity, we write ω_1^V as just ω_1. Let

$$\kappa := (2^{\lambda_f})^+,$$
$$\mathfrak{A} := (H(\kappa); \in), \text{ and}$$
$$R := \{i < \lambda_f : \omega_2 \leq i \text{ and } i \text{ is regular}\}.$$

As $|H(\lambda_f)| = \lambda_f$, we can fix a bijection $c : \lambda_f \longrightarrow H(\lambda_f)$, and define

- $Q_\lambda := c"\lambda$ and
- $A_\lambda := c"(\bar{A}_\lambda)$

for each $\lambda < \lambda_f$.

Making use of straightforward closure arguments, we inductively define C such that

(Q1) C is a club in λ_f, and

(Q2) for all $\lambda \in C$,

- Q_λ is transitive,
- $(Q_\lambda; \in, c \cap Q_\lambda) \prec (H(\lambda_f); \in, c)$.

We will let $Q_{\lambda_f} = H(\lambda_f)$.

Now, given any $P, B \subset H(\lambda_f)$, the set

$$\{\lambda \in C : (Q_\lambda; \in, P, B) \prec (H(\lambda_f); \in, P, B)\}$$

is a club in λ_f. Moreover, we can derive from $(\bar{A}_\lambda : \lambda < \lambda_f)$ being a $\Diamond_{\lambda_f}$-sequence, that the set

$$\{\lambda \in C : B \cap Q_\lambda = A_\lambda\}$$

is stationary in λ_f. We thus obtain

($\diamond$) for all $P, B \subset H(\lambda_f)$, the set

$$\{\lambda \in C : (Q_\lambda; \in, P, A_\lambda) \prec (H(\lambda_f); \in, P, B)\}$$

is stationary in λ_f.

We want to define $\mathbb{P}$ as a forcing notion comprising finite fragments of some language $\mathcal{L} \subset H(\lambda_f)$, such that $\mathbb{P}$ satisfies the hypothesis of Lemma 4.9.

Let us first define $\mathcal{L}$.

Definition 4.10. The language $\mathcal{L}$ requires the following distinguished symbols:

- $\dot{F}_i$ for $i \in R$, and
- $\dot{X}_{\delta,\lambda}$ for $\delta < \omega_1$ and $\lambda \in C$.

Now fix $\mathcal{L}$ to be the smallest set closed under negation, that contains expressions (strings) of the following types:

(L1) $\ulcorner \dot{F}_i(n) = \alpha \urcorner$, for $i \in R$, $n < \omega$ and $\alpha < i$, and

(L2) $\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner$, for $\delta < \omega_1$, $\lambda \in C$ and $x \in Q_\lambda$.

Morally, each $\dot{F}_i$ labels an increasing and cofinal partial map from $f(i)$ into i, and each $\dot{X}_\delta$ labels a side condition. The side conditions will be used to preserve stationary subsets of ω_1. As $\mathcal{L} \subset H(\lambda_f)$, $\mathfrak{A}$ is $\mathcal{L}$-suitable.

Sometimes, we want to pass certain parameters of an expression in $\mathcal{L}$ via variables. In such circumstances we are formally passing the parameters through the functions

$$\chi_1 : (i, n, \alpha) \mapsto \ulcorner \dot{F}_i(n) = \alpha \urcorner,$$
$$\chi_2 : (\delta, \lambda, x) \mapsto \ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner,$$
$$\chi_3 : (i, n, \alpha) \mapsto \ulcorner \neg \dot{F}_i(n) = \alpha \urcorner, \text{ and}$$
$$\chi_4 : (\delta, \lambda, x) \mapsto \ulcorner \neg x \in \dot{X}_{\delta,\lambda} \urcorner,$$

with their domains restricted to $H(\lambda_f)$. Note that under this domain restriction, χ_1, χ_2, χ_3 and χ_4 are all members of $H(\kappa)$; in fact, they are all definable functions of $H(\lambda_f)$. For brevity's and clarity's sake, we will abuse notation and suppress mention of said functions, whenever it is clear that we are using variables as placeholders for parameters in our construction of $\mathcal{L}_\mathfrak{A}^*$ sentences involving the symbol $\ulcorner E \urcorner$ — see e.g. Definition 4.15. Since the language associated with $\mathfrak{A}$ is of the lowest complexity class ($\mathcal{L}_\mathfrak{A}^*$-$\Delta_0$) in our classification of $\mathcal{L}_\mathfrak{A}^*$ formulas, this method of passing parameters incurs no additional cost to the complexity of any $\mathcal{L}_\mathfrak{A}^*$ sentence thus abbreviated.

By our choice of C, if $\lambda \in C$, then $\mathcal{L} \cap Q_\lambda$ is precisely the smallest set closed under negation, that contains expressions of the following types:

(L1)$_\lambda$ $\ulcorner \dot{F}_i(n) = \alpha \urcorner$, for $i \in R \cap \lambda$, $n < \omega$ and $\alpha < i$, and

(L2)$_\lambda$ $\ulcorner x \in \dot{X}_{\delta,\lambda'} \urcorner$, for $\delta < \omega_1$, $\lambda' \in C \cap \lambda$ and $x \in Q_{\lambda'}$.

Definition 4.11. An object of the form

$$\langle \langle F_i : i \in R \rangle, \langle X_{\delta,\lambda} : \delta < \omega_1, \lambda \in C \rangle \rangle$$

interprets $\mathcal{L}$ iff

- each F_i is a partial function from ω into i, and

- each $X_{\delta,\lambda}$ is a subset of Q_λ.

Sometimes it is convenient to talk about interpretations of specific symbols occurring in $\mathcal{L}$.

Definition 4.12. For any pair (i, Σ), define $F_i(\Sigma)$ to be the set

$$\{(n, \alpha) : \ulcorner \dot{F}_i(n) = \alpha \urcorner \in \Sigma\}.$$

Definition 4.13. For any triple $(\delta, \lambda, \Sigma)$, define $X_{\delta,\lambda}(\Sigma)$ to be the set

$$\{x : \ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma\}.$$

Definition 4.14. Given

$$\mathfrak{C} := \langle\langle F_i : i \in R\rangle, \langle X_{\delta,\lambda} : \delta < \omega_1, \lambda \in C\rangle\rangle$$

interpreting $\mathcal{L}$ and $\mathcal{L}' \subset \mathcal{L}$, let $\Sigma(\mathfrak{C}, \mathcal{L}')$ denote the union of the following sets:

- $\{\ulcorner \dot{F}_i(n) = \alpha \urcorner \in \mathcal{L}' : F_i(n) = \alpha\}$,
- $\{\ulcorner \neg \dot{F}_i(n) = \alpha \urcorner \in \mathcal{L}' : F_i(n) \neq \alpha\}$,
- $\{\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in \mathcal{L}' : x \in X_{\delta,\lambda}\}$, and
- $\{\ulcorner \neg x \in \dot{X}_{\delta,\lambda} \urcorner \in \mathcal{L}' : x \notin X_{\delta,\lambda}\}$.

It is clear that $\Sigma(\mathfrak{C}, \mathcal{L}) \cap Q_\lambda = \Sigma(\mathfrak{C}, \mathcal{L} \cap Q_\lambda)$ is $\mathcal{L} \cap Q_\lambda$-nice for all $\lambda \in C \cup \{\lambda_f\}$.

We will define $\{\mathbb{P}_\lambda : \lambda \in C \cup \{\lambda_f\}\}$ by induction on λ. Assume that $\mathbb{P}_{\lambda'}$ has been defined for all $\lambda' \in \lambda \cap C$. Also, for $\lambda' \in \lambda \cap C$, allow

- $\mathcal{L}^*$ to denote the set of first order formulas over the signature $\{\in, P, B\}$,
- $Ef_0^\lambda(\lambda')$ to denote the set

$$\begin{aligned}
\{(\phi, \bar{s}, r) : \phi \in \mathcal{L}^* \text{ and} \\
r \in Q_{\lambda'} \text{ and} \\
\bar{s} \in (Q_{\lambda'})^{<\omega} \text{ and} \\
dom(\bar{s}) + 1 = arity(\phi) \text{ and} \\
(Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'}) \models \phi(r, \bar{s})\},
\end{aligned}$$

- $Ef_1^\lambda(\lambda')$ to denote the set

$$\{(\phi, \bar{s}) : \exists r \, ((\phi, \bar{s}, r) \in Ef_0(\lambda'))\},$$

  and

- $Df^\lambda(\lambda')$ to denote the set

$$\begin{aligned}
\{(\phi, \bar{s}) : \phi \in \mathcal{L}^* \text{ and} \\
\bar{s} \in (Q_{\lambda'})^{<\omega} \text{ and} \\
dom(\bar{s}) + 1 = arity(\phi) \text{ and} \\
\{y \in \mathbb{P}_{\lambda'} : (Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'}) \models \phi(y, \bar{s})\} \text{ is dense in } \mathbb{P}_{\lambda'}\}.
\end{aligned}$$

The functions

$$\begin{aligned}
Ef_0^\lambda &: \lambda \cap C \longrightarrow H(\lambda_f) \, [\lambda' \mapsto Ef_0(\lambda')], \\
Ef_1^\lambda &: \lambda \cap C \longrightarrow H(\lambda_f) \, [\lambda' \mapsto Ef_1(\lambda')], \text{ and} \\
Df^\lambda &: \lambda \cap C \longrightarrow H(\lambda_f) \, [\lambda' \mapsto Df(\lambda')]
\end{aligned}$$

are clearly members of $H(\kappa)$.

Definition 4.15. Let Γ_λ be the set of $\mathcal{L}_{\mathfrak{A}}^*$-$\Pi_2$ sentences enumerated below.

(C1)$_\lambda$ $\ulcorner\forall i \, \forall n \, \forall\alpha \, \forall\beta \, ((E(\ulcorner \dot{F}_i(n) = \alpha\urcorner) \wedge E(\ulcorner \dot{F}_i(n) = \beta\urcorner)) \implies \alpha = \beta)\urcorner$,

(C2)$_\lambda$ $\ulcorner\forall i \, \forall\alpha \, \forall m \, \forall\gamma \, \exists n \, \exists\beta$
$$((E(\ulcorner \dot{F}_i(m) = \gamma\urcorner) \vee E(\ulcorner \neg\dot{F}_i(m) = \gamma\urcorner)) \wedge \alpha \in i$$
$$\implies (\alpha \in \beta \wedge E(\ulcorner \dot{F}_i(n) = \beta\urcorner)))\urcorner,$$

(C3)$_\lambda$ $\ulcorner\forall\delta \, \forall\lambda_0 \, \forall\lambda_1 \, \forall x \, \forall y$
$$((E(\ulcorner x \in \dot{X}_{\delta,\lambda_0}\urcorner) \wedge E(\ulcorner y \in \dot{X}_{\delta,\lambda_1}\urcorner)) \implies \lambda_0 = \lambda_1)\urcorner,$$

(C4)$_\lambda$ $\ulcorner\forall\delta_0 \, \forall\delta_1 \, \forall\lambda_0 \, \forall\lambda_1 \, \forall x \, \forall y$
$$((E(\ulcorner x \in \dot{X}_{\delta_0,\lambda_0}\urcorner) \wedge E(\ulcorner y \in \dot{X}_{\delta_1,\lambda_1}\urcorner) \wedge \delta_0 \in \delta_1)$$
$$\implies \lambda_0 \in \lambda_1)\urcorner$$

(C5)$_\lambda$ $\ulcorner\forall\delta \, \forall\lambda' \, \forall\delta' \in \delta \, \forall x \, (E(\ulcorner x \in \dot{X}_{\delta,\lambda'}\urcorner) \implies E(\ulcorner \delta' \in \dot{X}_{\delta,\lambda'}\urcorner))\urcorner$,

(C6)$_\lambda$ $\ulcorner\forall\delta \, \forall\lambda' \, \forall\delta' \in \omega_1 \, \forall x \, ((E(\ulcorner x \in \dot{X}_{\delta,\lambda'}\urcorner) \wedge \delta \in \delta') \implies \neg E(\ulcorner \delta' \in \dot{X}_{\delta,\lambda'}\urcorner))\urcorner$,

(C7)$_\lambda$ $\ulcorner\forall\delta \, \forall\lambda' \, \forall\phi \, \forall\bar{s} \, \forall x \, \exists r$
$$((E(\ulcorner x \in \dot{X}_{\delta,\lambda'}\urcorner) \wedge$$
$$(\phi, \bar{s}) \in Ef_1^\lambda(\lambda') \wedge$$
$$\forall n \in dom(\bar{s}) \, (E(\ulcorner \bar{s}(n) \in \dot{X}_{\delta,\lambda'}\urcorner)))$$
$$\implies (E(\ulcorner r \in \dot{X}_{\delta,\lambda'}\urcorner) \wedge (\phi, \bar{s}, r) \in Ef_0^\lambda(\lambda')))\urcorner,$$

(C8)$_\lambda$ $\ulcorner\forall\delta \, \forall\lambda' \, \forall\phi \, \forall\bar{s} \, \forall x \, \exists p \in \mathbb{P}_{\lambda'}$
$$((E(\ulcorner x \in \dot{X}_{\delta,\lambda'}\urcorner) \wedge$$
$$(\phi, \bar{s}) \in Df^\lambda(\lambda') \wedge$$
$$\forall n \in dom(\bar{s}) \, (E(\ulcorner \bar{s}(n) \in \dot{X}_{\delta,\lambda'}\urcorner)))$$
$$\implies ((\phi, \bar{s}, p) \in Ef_0^\lambda(\lambda') \wedge$$
$$E(\ulcorner p \in \dot{X}_{\delta,\lambda'}\urcorner) \wedge$$
$$\forall e \in p \, (E(e))))\urcorner.$$

In Definition 4.15, we give a list of constraints on the $\dot{F}_i$'s and the $\dot{X}_{\delta,\lambda'}$'s, that are meant to dictate how the objects interpreting them behave. To be more formal, let Σ interpret the predicate symbol E occurring in $\mathcal{L}^*$. Here, we are using the term "interpret" in the conventional model-theoretic sense. Also, let

- $F_i := F_i(\Sigma)$, and
- $X_{\delta,\lambda'} := X_{\delta,\lambda'}(\Sigma)$,

as i, δ and λ' range over their appropriate domains. Then

- $(C1)_\lambda$ and $(C3)_\lambda$ mean to say that the F_i's and the set

$$\{(\delta, \lambda') : X_{\delta,\lambda'} \neq \emptyset\}$$

  are functions,

- $(C2)_\lambda$ means to say that the image of each F_i is cofinal in i,

- $(C4)_\lambda$ means to say that the function

$$\{(\delta, \lambda') : X_{\delta,\lambda'} \neq \emptyset\}$$

  is strictly increasing,

- $(C5)_\lambda$ and $(C6)_\lambda$ mean to say that $X_{\delta,\lambda'} \cap \omega_1 = \delta$ whenever $X_{\delta,\lambda'}$ is non-empty,

- $(C7)_\lambda$ means to tell us that if $X_{\delta,\lambda'}$ is non-empty, then $(X_{\delta,\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'})$ is an elementary submodel of $(Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'})$, and

- $(C8)_\lambda$ means to tell us that if $X_{\delta,\lambda'}$ is non-empty, then for every D a dense subset of $\mathbb{P}_{\lambda'}$ definable over $(Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'})$ with parameters from $X_{\delta,\lambda'}$,

$$[\Sigma]^{<\omega} \cap X_{\delta,\lambda'} \cap D \neq \emptyset.$$

Now we can define $\mathbb{P}_\lambda := (P_\lambda, \leq_\lambda)$, where

$P_\lambda := \{p \in [\mathcal{L} \cap Q_\lambda]^{<\omega} : \Vdash_{Col(\omega,|H(\kappa)|)} \exists \Sigma \, (\text{"} \Sigma \, \Gamma_\lambda(\mathcal{L} \cap Q_\lambda, \mathfrak{A})\text{-certifies } p\text{"})\}$, and
$\leq_\lambda := \{(p, q) \in P_\lambda \times P_\lambda : q \subset p\}$.

We let $\mathbb{P}$ denote $\mathbb{P}_{\lambda_f}$.

Lemma 4.16. $|\mathbb{P}| \leq \lambda_f$.

Proof. This follows immediately from the observation that

$$\mathbb{P} \subset [\mathcal{L}]^{<\omega} \subset H(\lambda_f), \text{ and}$$
$$|H(\lambda_f)| = \lambda_f. \qquad \square$$

By Proposition 3.33 and the lemma below, $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}$. Obviously, $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$ because $\mathbb{P} \in H(\kappa)$.

Lemma 4.17. *For all $\lambda \in C \cup \{\lambda_f\}$, $\emptyset \in P_\lambda$.*

Proof. Let g be $Col(\omega, |H(\kappa)|)$-generic over W. In $W[g]$, for every $i \in R$, choose a cofinal map from ω into i and call it F_i. For every $\delta < \omega_1$ and every $\lambda \in C$, let $X_{\delta,\lambda}$ be the empty set. Then

$$\mathfrak{C} := \langle \langle F_i : i \in R \rangle, \langle X_{\delta,\lambda} : \delta < \omega_1, \lambda \in C \rangle \rangle$$

interprets $\mathcal{L}$ and $\Sigma(\mathfrak{C}, \mathcal{L} \cap Q_\lambda) \, \Gamma_\lambda(\mathcal{L} \cap Q_\lambda, \mathfrak{A})$-certifies $\emptyset$ for all $\lambda \in C \cup \{\lambda_f\}$. $\qquad \square$

Using a argument similar to that in the proof of Lemma 4.17, we get the following.

Lemma 4.18. *If*

- $\lambda_0, \lambda_1 \in C \cup \{\lambda_f\}$,
- $\lambda_0 \leq \lambda_1$, *and*
- $\Sigma \, \Gamma_{\lambda_0}(\mathcal{L} \cap Q_{\lambda_0}, \mathfrak{A})$-*certify* p,

then there is $\Sigma' \supset \Sigma$ for which $\Sigma' \, \Gamma_{\lambda_1}(\mathcal{L} \cap Q_{\lambda_1}, \mathfrak{A})$-certify p.

It can be gleaned from Lemma 4.18 and the definition of the $\mathbb{P}_\lambda$'s that

(P1) $P_{\lambda_0} = P_{\lambda_1} \cap Q_{\lambda_0}$ whenever $\lambda_0, \lambda_1 \in C \cup \{\lambda_f\}$ and $\lambda_0 \leq \lambda_1$, and

(P2) $P_\lambda = \bigcup\{P_{\lambda'} : \lambda' \in C \cap \lambda\}$ whenever $\lambda \in C \cup \{\lambda_f\}$ and $sup(\lambda \cap C) = \lambda$.

Lemma 4.19. *Let*

- $\lambda' \in C \cup \{\lambda_f\}$, *and*
- g *be a $\mathbb{P}_{\lambda'}$-generic filter over W.*

Then $\bigcup g \models^*_{\mathfrak{A}} \Gamma_{\lambda'}$.

Proof. We apply Lemma 3.35 with

- $\mathfrak{A}$, $\mathcal{L}$, $\mathbb{P}_{\lambda'}$ and g as defined or given above,
- $\Gamma_{\lambda'}$ in place of Γ,
- $|H(\kappa)|$ in place of λ,
- $[\mathcal{L} \cap Q_{\lambda'}]^{<\omega}$ in place of B,
- W in place of V, and
- $W[g]$ in place of W,

noting that

- $|trcl(\mathfrak{A})| \leq |H(\kappa)|$,
- $[\mathcal{L} \cap Q_{\lambda'}]^{<\omega}$ is closed under finite extensions,
- the definition of $\mathbb{P}_{\lambda'}$ in relation to the other parameters is faithful to the hypotheses of Lemma 3.35,
- g satisfies the hypotheses of Lemma 3.35 with respect to the other parameters, and
- $\mathbb{P}_{\lambda'}$ being non-empty (per Lemma 4.17) implies there is Σ in some weak outer model W' of W such that $\Sigma \, \Gamma_{\lambda'}(\mathcal{L} \cap Q_{\lambda'}, \mathfrak{A})$-certifies $\emptyset$,

to arrive at $\bigcup g \models^*_{\mathfrak{A}} \Gamma_{\lambda'}$. $\qquad\qquad\square$

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The proof of Lemma 4.19 serves as an instructive example of the utility of Lemma 3.35. We shall omit details in subsequent applications of Lemma 3.35, wherever the use cases are deemed similarly straightforward.

Lemma 4.20. $\mathbb{P}$ *fulfils (1) of Theorem 4.8.*

Proof. For any $\mathbb{P}_{\lambda'}$-generic filter g over W, $\bigcup g \models^*_{\mathfrak{A}} \Gamma_{\lambda_f}$ by Lemma 4.19. In particular, $\bigcup g \models^*_{\mathfrak{A}} (C2)_{\lambda_f}$. That $\mathbb{P}$ fulfils (1) of Theorem 4.8 follows immediately. $\qquad\square$

Definition 4.21. Let

- $S(\omega_1)$ denote the set of all stationary subsets of ω_1, and
- $U(C, \lambda_f)$ denote the set of all subsets of C unbounded in λ_f.

Check that both $S(\omega_1)$ and $U(C, \lambda_f)$ are members of $H(\kappa)$.

Lemma 4.22. *The $\mathcal{L}^*_{\mathfrak{A}}$-$\Pi_2$ sentence*

$$\ulcorner \forall x \in H(\lambda_f) \, \forall S \in S(\omega_1) \, \forall U \in U(C, \lambda_f) \, \exists \mu \in S \, \exists \nu \in U \, (E(\ulcorner x \in \dot{X}_{\mu,\nu} \urcorner)) \urcorner \quad (4.2)$$

is $(\mathfrak{A}, \mathbb{P}, \mathcal{L})$-universal.

Proof. Fix arbitrary

- $p \in \mathbb{P}$,
- $x \in H(\lambda_f)$,
- $S \in S(\omega_1)$, and
- $U \in U(C, \lambda_f)$.

It suffices to show that there are

- $\mu \in S$, and
- $\nu \in U$

for which $p \cup \ulcorner x \in \dot{X}_{\mu,\nu} \urcorner \in \mathbb{P}$.

To that end, let $\nu \in U$ be such that $x \in Q_\nu$ and $p \in \mathbb{P}_\nu$. This is possible by (P2). Choose $g \times f$ a $\mathbb{P}_\nu \times Col(\omega, \nu)$-generic filter over W with $p \in g$, so that $g \in W[g \times f]$ is a $\mathbb{P}_\nu$-generic filter over W and $|\nu|^{W[g \times f]} = \omega$. By Lemma 4.19, $\bigcup g \models^*_{\mathfrak{A}} \Gamma_\nu$.

Label the non-stationary ideal on ω_1 in W, I. Since I is precipitous in W, $(W; \in, I)$ is generically iterable in $W[g \times f]$. Consider a one-step iteration $\mathfrak{I}_1 := \langle (W; \in, I), (W_1; \in, I_1) \rangle$ in $W[g \times f]$, where $(W_1; \in, I_1)$ is the generic ultrapower of $(W; \in, I)$ via a W-generic ultrafilter on I containing S. Extend $\mathfrak{I}_1$ to a generic iteration $\mathfrak{I}$ of length $\omega_1^{W[g \times f]}$ in $W[g \times f]$. Said iteration gives rise to an elementary embedding $j : W \longrightarrow M$,

where M, a transitive inner model of $W[g \times f]$, is the final iterate of $\mathfrak{I}$. Moreover, $crit(j) = \omega_1^W \in j(S)$.

Let

$$\Sigma' := j''(\bigcup g) \cup \{\ulcorner j(y) \in \dot{X}_{\omega_1^W, j(\nu)} \urcorner : y \in Q_\nu\} \subset j(\mathcal{L}),$$

$$e \text{ be } Col(\omega, j(\lambda_f))\text{-generic over } W[g \times f], \text{ and}$$

$$W^* := W[g \times f][e].$$

Working in W^*, define

$$\mathfrak{C} := \langle\langle F_i : i \in j(R)\rangle, \langle X_{\delta,\lambda} : \delta < \omega_1^M, \lambda \in j(C)\rangle\rangle$$

as follows:

- $F_i(n) := \alpha$ iff $\ulcorner \dot{F}_i(n) = \alpha \urcorner \in \Sigma'$ whenever $\dot{F}_i$ occurs in Σ',

- F_i is some (any) strictly increasing cofinal function from ω into i whenever $\dot{F}_i$ does not occur in Σ', and

- $x \in X_{\delta,\lambda}$ iff $\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma'$.

Then $\mathfrak{C}$ interprets $j(\mathcal{L})$, and $\Sigma^* := \Sigma(\mathfrak{C}, j(\mathcal{L}))$ is $j(\mathcal{L})$-nice. Obviously $j''\bigcup g \subset \Sigma' \subset \Sigma^*$, so $j(p) = j''p \subset \Sigma^*$. By the definition of Σ', we too have

$$q^* := j(p) \cup \{\ulcorner j(x) \in \dot{X}_{\omega_1^W, j(\nu)} \urcorner\} \subset \Sigma^*.$$

In order to conclude that Σ^* $j(\Gamma_{\lambda_f})(j(\mathcal{L}), j(\mathfrak{A}))$-certifies q^*, we are left with showing $\Sigma^* \models^*_{j(\mathfrak{A})} j(\Gamma_{\lambda_f})$. That

$$\Sigma^* \models^*_{j(\mathfrak{A})} j((\mathbf{C}\mathbf{<}k\mathbf{>})_{\lambda_f})$$

for $\mathbf{<}k\mathbf{>} \in \{1, 3, 4, 5, 6\}$ follows immediately from the construction of Σ^*, the elementarity of j, as well as the fact that $crit(j) = \omega_1^W$.

For $\mathbf{<}k\mathbf{>} \in \{2, 7, 8\}$, we check that $\Sigma^* \models^*_{j(\mathfrak{A})} j((\mathbf{C}\mathbf{<}k\mathbf{>})_{\lambda_f})$ in greater detail below.

- $\mathbf{<}k\mathbf{>} = 2$: Let $i \in j(R)$. If $\dot{F}_i$ does not occur in Σ', there is nothing to check, because the definition of $\mathfrak{C}$ guarantees $ran(F_i(\Sigma^*))$ is cofinal in i. Otherwise, $\dot{F}_i$ occurs in Σ', which means $\dot{F}_i$ occurs in $j''\bigcup g$. Then there is $i' \in R$ such that $\dot{F}_i = \dot{F}_{j(i')} = j(\dot{F}_{i'})$. That $\bigcup g \models^*_{\mathfrak{A}} (\mathbf{C}2)_\nu$ implies $ran(F_{i'}(\bigcup g))$ is cofinal in i'. By a basic property of elementary embeddings associated with generic iterations, we know that for any ordinal α satisfying the inequality $\omega_1^W < cof^W(\alpha)$, we must have $j(\alpha) = sup(j''\alpha)$. Thus, $ran(F_i(\Sigma^*)) = j''ran(F_{i'}(\bigcup g))$ is cofinal in $j(i') = i$, and we are done.

- $\mathbf{<}k\mathbf{>} = 7$: Let δ, λ, ϕ, $\bar{s}$ and x be such that

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(K7.1) $\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma^*$,

(K7.2) $(\phi, \bar{s}) \in j(Ef_1^{\lambda_f})(\lambda)$, and

(K7.3) $\ulcorner \bar{s}(n) \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma^*$ for all $n \in dom(\bar{s})$.

Combining (K7.2), (K7.3) and the definition of Σ^* gives us

- $\lambda = j(\lambda')$ for some $\lambda' \in C$,
- $j(\phi) = \phi \in j(\mathcal{L}^*) = \mathcal{L}^*$,
- $\bar{s} \in (j"Q_{\lambda'})^{<\omega} = j"(Q_{\lambda'})^{<\omega}$,
- $dom(\bar{s}) + 1 = arity(\phi)$, and
- $(j(Q_{\lambda'}); \in, j(\mathbb{P}_{\lambda'}), j(A_{\lambda'})) \models \exists r\, \phi(r, \bar{s})$.

We want to show that
$$(X_{\delta,\lambda}(\Sigma^*); \in, j(\mathbb{P}_{\lambda'}), j(A_{\lambda'})) \models \exists r\, \phi(r, \bar{s}).$$

By the elementarity of j,

- $j^{-1}(\bar{s}) \in (Q_{\lambda'})^{<\omega}$,
- $dom(j^{-1}(\bar{s})) + 1 = arity(\phi)$, and
- $(Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'}) \models \exists r\, \phi(r, j^{-1}(\bar{s}))\}$,

so $(\phi, j^{-1}(\bar{s})) \in Ef_1^{\lambda_f}(\lambda')$. Henceforth, there are two possible cases. We will analyse them with reference to the way Σ^* is constructed.

Case 1: $\delta = \omega_1^W$. Then $\lambda' = \nu$ and $X_{\delta,\lambda}(\Sigma^*) = j"Q_\nu$. As
$$X_{\delta,\lambda}(\Sigma^*) \cap j(\mathbb{P}_{\lambda'}) = j"\mathbb{P}_{\lambda'} \text{ and}$$
$$X_{\delta,\lambda}(\Sigma^*) \cap j(A_{\lambda'}) = j"A_{\lambda'},$$

we can conclude
$$(X_{\delta,\lambda}(\Sigma^*); \in, j(\mathbb{P}_{\lambda'}), j(A_{\lambda'})) \models \exists r\, \phi(r, \bar{s})$$

by invoking the elementarity of j once again.

Case 2: $\delta \neq \omega_1^W$. Then $\delta < \omega_1^W$, $\lambda' < \nu$, and
$$\ulcorner j^{-1}(\bar{s})(n) \in \dot{X}_{\delta,\lambda'} \urcorner \in \bigcup g$$

for all $n \in dom(j^{-1}(\bar{s}))$. Moreover, $X_{\delta,\lambda}(\Sigma^*) = j"X_{\delta,\lambda'}(\bigcup g)$, so $X_{\delta,\lambda}(\Sigma^*)$ being non-empty implies $X_{\delta,\lambda'}(\bigcup g)$ is non-empty as well. Since $\bigcup g \models_{\mathfrak{A}}^* (C7)_\nu$, we have
$$(X_{\delta,\lambda'}(\bigcup g); \in, \mathbb{P}_{\lambda'}, A_{\lambda'}) \models \exists r\, \phi(r, j^{-1}(\bar{s})).$$

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As

$$X_{\delta,\lambda}(\Sigma^*) \cap j(\mathbb{P}_{\lambda'}) = j"(X_{\delta,\lambda'}(\bigcup g) \cap \mathbb{P}_{\lambda'}) \text{ and}$$

$$X_{\delta,\lambda}(\Sigma^*) \cap j(A_{\lambda'}) = j"(X_{\delta,\lambda'}(\bigcup g) \cap A_{\lambda'}),$$

we can conclude

$$(X_{\delta,\lambda}(\Sigma^*); \in, j(\mathbb{P}_{\lambda'}), j(A_{\lambda'})) \models \exists r\, \phi(r, \bar{s})$$

by invoking the elementarity of j yet again.

$<k> = 8$: This is similar to the argument in the case of $<k> = 7$. We provide details for the sake of completeness, and to elucidate the ample similarity.

Let δ, λ, ϕ, $\bar{s}$ and x be such that

(K8.1) $\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma^*$,

(K8.2) $(\phi, \bar{s}) \in j(Df^{\lambda_f})(\lambda)$, and

(K8.3) $\ulcorner \bar{s}(n) \in \dot{X}_{\delta,\lambda} \urcorner \in \Sigma^*$ for all $n \in dom(\bar{s})$.

Combining (K8.2), (K8.3) and the definition of Σ^* gives us

- $\lambda = j(\lambda')$ for some $\lambda' \in C$,
- $j(\phi) = \phi \in j(\mathcal{L}^*) = \mathcal{L}^*$,
- $\bar{s} \in (j" Q_{\lambda'})^{<\omega} = j"(Q_{\lambda'})^{<\omega}$,
- $dom(\bar{s}) + 1 = arity(\phi)$, and
- $D := \{y \in j(\mathbb{P}_{\lambda'}) : (j(Q_{\lambda'}); \in, j(\mathbb{P}_{\lambda'}), j(A_{\lambda'})) \models \phi(y, \bar{s})\}$ is dense in $j(\mathbb{P}_{\lambda'})$.

We want to show that

$$[\Sigma^*]^{<\omega} \cap X_{\delta,\lambda}(\Sigma^*) \cap D \neq \emptyset.$$

By the elementarity of j,

- $j^{-1}(\bar{s}) \in (Q_{\lambda'})^{<\omega}$,
- $dom(j^{-1}(\bar{s})) + 1 = arity(\phi)$, and
- $j^{-1}(D) = \{y \in \mathbb{P}_{\lambda'} : (Q_{\lambda'}; \in, \mathbb{P}_{\lambda'}, A_{\lambda'}) \models \phi(y, j^{-1}(\bar{s}))\}$ is dense in $\mathbb{P}_{\lambda'}$,

so $(\phi, j^{-1}(\bar{s})) \in Df^{\lambda_f}(\lambda')$. Henceforth, there are two possible cases. We will analyse them with reference to the way Σ^* is constructed.

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Case 1: $\delta = \omega_1^W$. Then $\lambda' = \nu$ and $X_{\delta,\lambda}(\Sigma^*) = j''Q_\nu$. Clearly,

$$[\bigcup g]^{<\omega} \cap Q_\nu \cap j^{-1}(D) = g \cap j^{-1}(D) \neq \emptyset,$$

as g is $\mathbb{P}_\nu$-generic over W. That $j'' \bigcup g \subset \Sigma^*$ means

$$j(p) \in [\Sigma^*]^{<\omega} \cap X_{\delta,\lambda}(\Sigma^*) \cap D \neq \emptyset$$

for any $p \in g \cap j^{-1}(D)$.

Case 2: $\delta \neq \omega_1^W$. Then $\delta < \omega_1^W$, $\lambda' < \nu$, and

$$\ulcorner j^{-1}(\bar{s})(n) \in \dot{X}_{\delta,\lambda'} \urcorner \in \bigcup g$$

for all $n \in dom(j^{-1}(\bar{s}))$. Moreover, $X_{\delta,\lambda}(\Sigma^*) = j'' X_{\delta,\lambda'}(\bigcup g)$, so $X_{\delta,\lambda}(\Sigma^*)$ being non-empty implies $X_{\delta,\lambda'}(\bigcup g)$ is non-empty too. Since $\bigcup g \models_{\mathfrak{A}}^* (C8)_\nu$, we have

$$[\bigcup g]^{<\omega} \cap X_{\delta,\lambda'}(\bigcup g) \cap j^{-1}(D) \neq \emptyset.$$

As in Case 1, we can conclude

$$[\Sigma^*]^{<\omega} \cap X_{\delta,\lambda}(\Sigma^*) \cap D \neq \emptyset.$$

Now that

- $\Sigma^* \in W^*$,
- $\Sigma^* \ j(\Gamma_{\lambda_f})(j(\mathcal{L}), j(\mathfrak{A}))$-certifies q^*,
- W^* is a weak outer model of M, and
- $|H(\kappa)^W|^W = |trcl(\mathfrak{A})|^W$,

we can apply Lemma 3.34 with

- M in place of V,
- W^* in place of W,
- $j(|H(\kappa)^W|^W) = |H(j(\kappa))^M|^M$ in place of λ,
- Σ^* in place of Σ,
- $j(\Gamma_{\lambda_f})$ in place of Γ,
- $j(\mathcal{L})$ in place of $\mathcal{L}$,
- $j(\mathfrak{A})$ in place of $\mathfrak{A}$, and

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- q^* in place of p,

noting that in M,

$$Col(\omega, |H(j(\kappa))|) = j(Col(\omega, |H(\kappa)^W|^W)).$$

The application yields

$$(M; \in) \models \text{“} \Vdash_{j(Col(\omega,|H(\kappa)^W|^W)^W)} \exists\Sigma \text{ (“}\Sigma j(\Gamma_{\lambda_f})(j(\mathcal{L}), j(\mathfrak{A}))\text{-certifies } q^*\text{”})\text{”}.$$

But this means $q \in j(\mathbb{P})$, which implies

$$(M; \in) \models \exists\mu \in j(S) \, (j(p) \cup \{\ulcorner j(x) \in \dot{X}_{\mu,j(\nu)}\urcorner\} \in j(\mathbb{P})).$$

By the elementarity of j,

$$(W; \in) \models \exists\mu \in S \, (p \cup \{\ulcorner x \in \dot{X}_{\mu,\nu}\urcorner\} \in \mathbb{P}),$$

completing the proof. □

Lemma 4.23. $\mathbb{P}$ *is stationary-preserving.*

Proof. Let

- $S \in S(\omega_1)$,
- $p \in \mathbb{P}$,
- $\dot{C}$ be a $\mathbb{P}$-name in such that $p \Vdash_{\mathbb{P}} \text{“}\dot{C}$ is a club in $\omega_1^{W}\text{”}$,
- $D := \{(q, \eta) \in \mathbb{P} \times \omega_1 : q \Vdash_{\mathbb{P}} \eta \in \dot{C}\}$,
- g be a $\mathbb{P}$-generic filter over W with $p \in g$.

Applying ($\diamond$) with

- $\mathbb{P}$ in place of P, and
- D in place of B,

we get

$$U := \{\lambda \in C : (Q_\lambda; \in, \mathbb{P}, A_\lambda) \prec (H(\lambda_f); \in, \mathbb{P}, D)\}$$

is stationary in λ_f, so $U \in U(C, \lambda_f)$.

In $W[g]$, there are $\mu \in S$ and $\nu \in U$ such that

$$\emptyset \neq (X_{\mu,\nu}(\bigcup g); \in, \mathbb{P}, A_\nu) \prec (Q_\nu; \in, \mathbb{P}, A_\nu) \prec (H(\lambda_f); \in, \mathbb{P}, D),$$

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since

$$\bigcup g \models_{\mathfrak{A}}^{*} \Gamma_{\lambda_f} \text{ (in particular } \bigcup g \models_{\mathfrak{A}}^{*} (C7)_{\lambda_f}) \text{ and } \bigcup g \models_{\mathfrak{A}}^{*} (4.2)$$

by Lemmas 3.28, 4.19 and 4.22. Now, noting (P1), we have

$$\emptyset \neq (X_{\mu,\nu}(\bigcup g); \in, \mathbb{P}, D) \prec (Q_\nu; \in, \mathbb{P}_\nu, A_\nu) \prec (H(\lambda_f); \in, \mathbb{P}, D). \qquad (4.3)$$

It suffices to show that μ is a limit point of $\dot{C}[g]$. We fix $\zeta < \mu$ and seek some $\eta \in \dot{C}[g]$ with $\zeta < \eta < \mu$.

The set

$$E_\zeta := \{q \in \mathbb{P} : \exists \eta > \zeta \, ((q, \eta) \in D)\}$$

is dense in $\mathbb{P}$, so (4.3) tells us

$$E_\zeta \cap Q_\nu = \{q \in \mathbb{P}_\nu : \exists \eta > \zeta \, ((q, \eta) \in A_\nu)\}$$

is dense in $\mathbb{P}_\nu$. As $\bigcup g \models_{\mathfrak{A}}^{*} (C5)_{\lambda_f}$, we know $\zeta \in X_{\mu,\nu}(\bigcup g)$. Having $\bigcup g \models_{\mathfrak{A}}^{*} (C8)_{\lambda_f}$ then bestows us the existence of some

$$q \in [\bigcup g]^{<\omega} \cap X_{\mu,\nu}(\bigcup g) \cap E_\zeta \cap Q_\nu \neq \emptyset.$$

That $q \in X_{\mu,\nu}(\bigcup g)$ and (4.3) holds means

$$(X_{\mu,\nu}(\bigcup g); \in, \mathbb{P}, D) \models \exists \eta > \zeta \, ((q, \eta) \in D).$$

Invoking the fact that $\bigcup g \models_{\mathfrak{A}}^{*} (C6)_{\lambda_f}$ gives us some η such that

- $\zeta < \eta < \mu$, and
- $(q, \eta) \in D$.

Recalling the definition of D, we conclude $\eta \in \dot{C}[g]$ because $q \in [\bigcup g]^{<\omega} = g$, $\qquad \square$

Lemma 4.24. $\mathbb{P}$ *fulfils (2) of Theorem 4.8.*

Proof. Let g be $\mathbb{P}$-generic over W. By Lemma 4.19, $\bigcup g \models_{\mathfrak{A}}^{*} \Gamma_{\lambda_f}$, so that in $W[g]$,

$$K := \{(\delta, \lambda) : \exists x \, (\ulcorner x \in \dot{X}_{\delta,\lambda} \urcorner \in g)\}$$

is a strictly increasing function with domain contained in ω_1^W and range contained in λ_f. Lemma 4.22 tells us that $dom(K)$ is cofinal in ω_1^W and $ran(K)$ is cofinal in λ_f, hence

$$cof^{W[g]}(\lambda_f) = cof^{W[g]}(\omega_1^W).$$

By Lemma 4.23, we have $\omega_1^W = \omega_1^{W[g]}$, and consequently,

$$cof^{W[g]}(\lambda_f) = \omega_1^{W[g]}. \qquad \square$$

In view of Lemma 4.9, the theorem follows from Lemmas 4.16, 4.20, 4.23 and 4.24. $\qquad \square$

According to Theorem 4.8, if the non-stationary ideal on ω_1 is precipitous, then there is a uniform way of generating witnesses — in place of $\mathbb{P}_N$ — to analogues of Fact 4.6, that helps overcome the difficulty presented by the same fact to the extension of Namba forcing via iteration. To wit, we have the following corollary.

Corollary 4.25. *Assume*

- *the non-stationary ideal on ω_1 is precipitous, and*
- *GCH holds below ω_2.*

Then $Nb_1(\lambda)$ holds for all λ above ω_2.

4.3 Open Questions

In the proof of Theorem 4.8, fix any θ such that $[Q_\lambda]^{<\theta} \subset Q_\lambda$ for all $\lambda \in C$. Should we alter the definition of Γ_λ and P_λ for each $\lambda \in C \cup \{\lambda_f\}$ as follows:

$$\Gamma_\lambda := \Gamma_\lambda(\theta) := \text{the set defined as per Definition 4.15,}$$
$$\text{but with } \{\mathbb{P}_{\lambda'}(\theta) : \lambda' \in \lambda \cap C\} \text{ as the set of parameters,}$$
$$P_\lambda := P_\lambda(\theta) := \{p \in [\mathcal{L} \cap Q_\lambda]^{<\theta} :$$
$$\Vdash_{Col(\omega,|H(\kappa)|)} \exists \Sigma \left(\text{``}\Sigma\, \Gamma_\lambda(\theta)(\mathcal{L} \cap Q_\lambda, \mathfrak{A})\text{-certifies } p\text{''}\right)\}, \text{ and}$$
$$\mathbb{P}_\lambda := \mathbb{P}_\lambda(\theta) := (P_\lambda(\theta), \supset),$$

all subsequent lemmas in the proof still go through without a hitch. What then differentiates the $\mathbb{P}_{\lambda_f}(\theta)$'s that result from varying θ? More specifically, because we can always choose C such that $[Q_\lambda]^{<\omega_1} \subset Q_\lambda$ for all $\lambda \in C$, this next question is of particular interest.

Question 4.26. *Let λ_f, C and*

$$(C \longrightarrow \mathcal{P}(H(\lambda_f)))[\lambda \mapsto Q_\lambda]$$

be as in the proof of Theorem 4.8. Can C be chosen such that, furthermore,

(1) $[Q_\lambda]^{<\omega_1} \subset Q_\lambda$ for all $\lambda \in C$, and

(2) there is a first-order sentence φ in the language of set theory for which

$$\text{`` } \Vdash_{\mathbb{P}_{\lambda_f}(\omega)} \varphi \text{ and } \Vdash_{\mathbb{P}_{\lambda_f}(\omega_1)} \neg\varphi \text{'' ?}$$

We are optimistic that an answer to Question 4.26 can be generalised to associate forcing consequences of similar forcing notions with cardinality limitations placed upon their conditions. For example, an affirmative answer might allow us to separate the forcing consequences of two $\mathbb{P}$'s fulfilling the hypothesis of Lemma 3.35, according to differences in cardinality bounds on the members of their respective B's.

In the spirit of Jensen's results on the extended Namba problem, we are interested in whether the forcing notion $\mathbb{P}_{\lambda_f}$ constructed in the proof of Theorem 4.8 adds reals. More generally, we can ask the same about the parametrised versions of said forcing notion.

Question 4.27. *Let λ_f, C and*

$$(C \longrightarrow \mathcal{P}(H(\lambda_f)))[\lambda \mapsto Q_\lambda]$$

be as in the proof of Theorem 4.8. Can C and a cardinal θ be chosen such that, furthermore,

(1) $[Q_\lambda]^{<\theta} \subset Q_\lambda$ for all $\lambda \in C$, and

(2) $\mathbb{P}_{\lambda_f}(\theta)$ does not add reals?

In the likely event that the answer to Question 4.27 is in the negative, it makes sense to consider a more general question.

(Q1) Must $Nb_2(\lambda)$ hold for all λ above ω_2, where

$Nb_2(\lambda) :=$ 'there is a stationary-preserving forcing notion $\mathbb{P}$ such that

$\mathbb{P}$ does not add reals,

$\Vdash_{\mathbb{P}}$ "$cof(\alpha) = \omega$" for all regular cardinals α satisfying $\omega_2 \leq \alpha < \lambda$,

$\Vdash_{\mathbb{P}}$ "$cof(\lambda) = \omega_1$", and

$\Vdash_{\mathbb{P}}$ "$cof(\beta) > \omega$" for all regular cardinals β satisfying $\lambda < \beta$'?

As a consequence of Jensen's work, we need only consider the case of λ being a weakly inaccessible cardinal without further qualification.

Question 4.28. *Must $Nb_2(\lambda)$ hold for a weakly inaccessible λ, if λ is not strongly inaccessible?*

In another direction, we can ask about the possibility of eschewing the assumptions in Corollary 4.25.

Question 4.29. *Is it true that* ZFC $\vdash$ *"$Nb_1(\lambda)$ holds for all λ above ω_2"?*

Apparently, Question 4.29 has been resolved in the affirmative by Bondt and Velickovic in the former's forthcoming thesis. In fact, Bondt claims the duo have proven an even stronger theorem in ZFC alone, when answering a question of mine on MathOverflow (see [14]). However, as the finer details of their proof are not provided in his MathOverflow post, I shall conservatively mark Question 4.29 as open.

5 Theories with Constraints in Interpretation (TCIs) and their Models

In the previous section, we discussed a method of forcing the cofinality of regular cardinals within an interval to be ω. The idea of changing cofinalities via forcing involves extracting a cofinal function from an existing relation R on a structured set A. We can view A coupled with R as a new structure $\mathfrak{A}$ that sets the context of the problem. Then a subset of R being a cofinal function becomes a definable property over $\mathfrak{A}$. Compare and contrast this with the notion of a first-order theory, which defines a property over nothing more than a vocabulary; it makes sense that the addition of a structure interpreting said vocabulary would allow us to define more intricate properties.

Essentially, a structure can be used to provide additional constraints to a first-order theory, and forcing-related questions can often be framed as consistency questions that ask about the existence of models of first-order theories satisfying such constraints. This section is dedicated to studying the aforementioned models, with a focus on their relationship with genericity.

5.1 Definitions and Basic Properties

We first make precise the notion of first-order theories with added constraints, so that we can compare these mathematical objects with the first-order theories we are familiar with.

Definition 5.1. A *first-order theory with constraints in interpretation* (*first-order TCI*) — henceforth, just *theory with constraints in interpretation* (*TCI*) — is a tuple $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$, where

- T is a first order theory with signature σ,

- $\dot{\mathcal{U}}$ is a unary relation symbol not in σ,

- ϑ is a function (the *constraint function*) with domain $\sigma \cup \{\dot{\mathcal{U}}\}$,

- if $x \in ran(\vartheta)$, then there is y such that

    ○ either $x = (y, 0)$ or $x = (y, 1)$, and

    ○ if $\vartheta(\dot{\mathcal{U}}) = (z, a)$, then $y \subset z^n$ for some $n < \omega$, and

- if $\vartheta(\dot{\mathcal{U}}) = (z, a)$, then

    ○ $z \cap z^n = \emptyset$ whenever $1 < n < \omega$, and

    ○ $z^m \cap z^n = \emptyset$ whenever $1 < m < n < \omega$.

Definition 5.2. Let $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI. We say $\mathcal{M} := (U; \mathcal{I}) \models^* (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ — or $\mathcal{M}$ *models* $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ — iff all of the following holds:

- $\mathcal{M}$ is a structure,
- σ is the signature of $\mathcal{M}$,
- $\mathcal{M} \models T$,
- if $\vartheta(\dot{\mathcal{U}}) = (y, 0)$, then $U \subset y$,
- if $\vartheta(\dot{\mathcal{U}}) = (y, 1)$, then $U = y$, and
- for $\dot{X} \in \sigma$,

  ○ if $\dot{X}$ is a constant symbol and $\vartheta(\dot{X}) = (y, z)$, then $\mathcal{I}(\dot{X}) \in y \cap U$,

  ○ if $\dot{X}$ is an n-ary relation symbol and $\vartheta(\dot{X}) = (y, 0)$, then $\mathcal{I}(\dot{X}) \subset y \cap U^n$,

  ○ if $\dot{X}$ is an n-ary relation symbol and $\vartheta(\dot{X}) = (y, 1)$, then $\mathcal{I}(\dot{X}) = y \cap U^n$,

  ○ if $\dot{X}$ is an n-ary function symbol and $\vartheta(\dot{X}) = (y, 0)$, then

  $$\{z \in U^{n+1} : \mathcal{I}(\dot{X})(z \!\restriction_n) = z(n)\} \subset y \cap U^{n+1}, \text{ and}$$

  ○ if $\dot{X}$ is an n-ary function symbol and $\vartheta(\dot{X}) = (y, 1)$, then

  $$\{z \in U^{n+1} : \mathcal{I}(\dot{X})(z \!\restriction_n) = z(n)\} = y \cap U^{n+1}.$$

We say $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ has a model if there exists $\mathcal{M}$ for which $\mathcal{M} \models^* (T, \sigma, \dot{\mathcal{U}}, \vartheta)$.

Example 5.3. Let T be any first-order theory over the signature σ, $\dot{\mathcal{U}}$ be a unary relation symbol not in σ, and A be any set. Define ϑ on $\sigma \cup \{\dot{\mathcal{U}}\}$ such that

- $\vartheta(\dot{\mathcal{U}}) = (A, 1)$,
- $\vartheta(\dot{X}) = (A, 0)$ whenever $\dot{X} \in \sigma$ is a constant symbol,
- $\vartheta(\dot{X}) = (A^n, 0)$ whenever $\dot{X} \in \sigma$ is an n-ary relation symbol, and
- $\vartheta(\dot{X}) = (A^{n+1}, 0)$ whenever $\dot{X} \in \sigma$ is an n-ary function symbol.

If we set $\mathfrak{T} := (T, \sigma, \dot{\mathcal{U}}, \vartheta)$, then the models of $\mathfrak{T}$ are precisely the models of T with base set A.

Definition 5.4. Let $\mathfrak{A}$ and T be a structure and a first-order theory respectively, over the same signature σ. Define

$$\mathrm{Sub}(\mathfrak{A}, T) := \{\mathfrak{B} : \mathfrak{B} \text{ is a substructure of } \mathfrak{A} \text{ and } \mathfrak{B} \models T\}.$$

Members of $\mathrm{Sub}(\mathfrak{A}, T)$ are called T-substructures of $\mathfrak{A}$.

Example 5.5. Let T be any first-order theory over the signature σ, and $\mathfrak{A} = (A; \mathcal{I})$ be a structure interpreting σ. Define ϑ on $\sigma \cup \{\dot{\mathcal{U}}\}$ such that

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- $\vartheta(\dot{\mathcal{U}}) = (A, 0)$,
- $\vartheta(\dot{X}) = (\{\mathcal{I}(\dot{X})\}, 1)$ whenever $\dot{X} \in \sigma$ is a constant symbol,
- $\vartheta(\dot{X}) = (\mathcal{I}(\dot{X}), 1)$ whenever $\dot{X} \in \sigma$ is an n-ary relation symbol, and
- $\vartheta(\dot{X}) = (R_{\mathcal{I}}(\dot{X}), 1)$ whenever $\dot{X} \in \sigma$ is an n-ary function symbol, where

$$R_{\mathcal{I}}(\dot{X}) := \{z \in A^{n+1} : \mathcal{I}(\dot{X})(z\restriction_n) = z(n)\}.$$

If we set $\mathfrak{T} := (T, \sigma, \dot{\mathcal{U}}, \vartheta)$, then

$$\{\text{models of } \mathfrak{T}\} = \mathrm{Sub}(\mathfrak{A}, T).$$

In practice, we can view forcing as a technique to refine structures that provably exist in V. Often, such refinements cannot be carried out in V, for any successful attempt would result in objects that cannot exist in V. In each of these cases, forcing can be used to extend V to include an instance of the refined structure. The way we define TCIs allows them to specify — and act as blueprints for — refinements of this ilk. If $\mathfrak{T}$ is a TCI specifying a particular refinement, then models of $\mathfrak{T}$ correspond to the results of said refinement. We hope the next example can help illustrate our aforementioned idea of specification.

Example 5.6. Let $\dot{\mathcal{U}}$ be a unary relation and $\dot{R}$ be a binary relation. Define ϑ on $\{\dot{\mathcal{U}}, \dot{R}\}$ such that

- $\vartheta(\dot{\mathcal{U}}) = (\omega_1, 1)$, and
- $\vartheta(\dot{R}) = (\omega_1 \times \omega, 0)$.

Set T to contain exactly the sentences

$$\ulcorner \forall x\, \exists y\, (\dot{R}(x, y)) \urcorner,$$
$$\ulcorner \forall x\, \forall y\, \forall z\, ((\dot{R}(x, y) \wedge \dot{R}(x, z)) \implies y = z) \urcorner \text{ and}$$
$$\ulcorner \forall x\, \forall y\, \forall z\, ((\dot{R}(x, z) \wedge \dot{R}(y, z)) \implies x = y) \urcorner.$$

Now $\mathfrak{T} := (T, \{\dot{R}\}, \dot{\mathcal{U}}, \vartheta)$ is a TCI that specifies a refinement of the structure

$$\mathfrak{A} := (\omega_1; \{(\dot{R}, \omega_1 \times \omega)\})$$

to some

$$\mathfrak{A}' := (\omega_1; \{(\dot{R}, F)\}),$$

where $F \subset \omega_1 \times \omega$ is an injection from ω_1 into ω. As an implication, $\mathfrak{T}$ must not have any model in V. However, an outer model of V in which ω_1^V is collapsed to ω necessarily contains models of $\mathfrak{T}$.

Example 5.6 reminds us that the existence of models for a TCI is not absolute between V and its outer models. There is thus a fundamental difference between the model existence of a TCI and that of a first-order theory. This should reflect in our definition of what it means for a TCI to be consistent.

Definition 5.7. A TCI $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is *consistent* iff $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ has a model in some outer model of V.

Remark 5.8. It might seem at first glance, that the the consistency of a TCI is not a first-order property in the language of set theory, since it involves quantifying over outer models of V. This is not a real problem, as we shall see in the next subsection, because said definition is equivalent to a first-order property at the metalevel.

Definition 5.9. A TCI $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is *finitely consistent* iff for all finite $T' \subset T$, $(T', \sigma, \dot{\mathcal{U}}, \vartheta)$ is consistent.

Definition 5.10. A TCI $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is Π_n iff T contains only Π_n sentences.
  A TCI $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is Σ_n iff T contains only Σ_n sentences.
  A TCI $(T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is $\Sigma_n \cup \Pi_n$ iff every sentence in T is either Σ_n or Π_n.

TCIs allow natural constraints that are not first-order definable to be imposed on the models of a theory. However, they are not a "true" generalisation of first-order theories because their models have uppers bounds in size. In fact, we can show that the size limitation of models of TCIs is in some sense, the only shortcoming of TCIs *vis-a-vis* first-order theories.

Lemma 5.11. *Let T be a first-order theory over the signature σ, and $\dot{\mathcal{U}}$ be a unary relation symbol not in σ. For every cardinal κ, there is a TCI $\mathfrak{T}$ such that*

- $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ *for some ϑ, and*
- *every model $\mathfrak{A} = (A; \mathcal{I})$ of T with $|A| \leq \kappa$ is isomorphic to some model of $\mathfrak{T}$.*

Proof. Define

- $\vartheta(\dot{\mathcal{U}}) := (\kappa, 0)$,
- $\vartheta(\dot{X}) := (\kappa, 0)$ if $\dot{X}$ is a constant symbol,
- $\vartheta(\dot{X}) := (\kappa^n, 0)$ if $\dot{X}$ is an n-ary relation symbol,
- $\vartheta(\dot{X}) := (\kappa^{n+1}, 0)$ if $\dot{X}$ is an n-ary function symbol.

Then

$$\mathfrak{T} := (T, \sigma, \dot{\mathcal{U}}, \vartheta)$$

is as required. $\qquad\qquad\square$

It turns out that there is an analogue of the downward Lowenheim-Skolem theorem for TCIs.

Lemma 5.12. *Let*

- $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ *be a TCI,*

- y *be a set, and*

- α *be an infinite ordinal.*

Assume that $\vartheta(\dot{\mathcal{U}}) = (y, 0)$ and in some weak outer model W of V, there is a pair $(\mathcal{M}, f)$ such that

- $\mathcal{M} = (U; \mathcal{I}) \models^* \mathfrak{T}$, *and*

- $f : \alpha \longrightarrow U$ *is a bijection.*

Then for every β with $\omega \leq \beta \leq \alpha$, there is a pair $(\mathcal{M}', f')$ in W such that

- $\mathcal{M}' = (U'; \mathcal{I}') \models^* \mathfrak{T}$, *and*

- $f' : \beta \longrightarrow U'$ *is a bijection.*

Proof. Let β be such that $\omega \leq \beta \leq \alpha$. By the downward Lowenheim-Skolem theorem applied to $\mathcal{M}$ in W, there exists a structure $\mathcal{M}' := (U'; \mathcal{I}')$ for which $\mathcal{M}' \prec \mathcal{M}$ and $|U'| = |\beta|$. This means $\mathcal{M}' \models T$. Further, $U' \subset U \subset y$ and whenever $\dot{X} \in \sigma$ is a constant symbol, $\mathcal{I}(\dot{X}) = \mathcal{I}'(\dot{X}) \in U'$. The other criteria for $\mathcal{M}' \models^* \mathfrak{T}$ are easy to check. Fix f' to be any bijection from β into U', and we are done. $\qquad\square$

If we allow movement among outer models of V, we get the following (somewhat trivial) version of the general Lowenheim-Skolem theorem for TCIs.

Lemma 5.13. *Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI with an infinite model in some outer model of V. Then for every infinite ordinal β, there is a pair $(\mathcal{M}, f)$ in some outer model of V such that*

- $\mathcal{M} = (U; \mathcal{I}) \models^* \mathfrak{T}$, *and*

- $f : \beta \longrightarrow U$ *is a bijection.*

Proof. By our assumptions on $\mathfrak{T}$, it has a model $\mathcal{M} = (U, \mathcal{I})$ in some outer model W of V, such that U is infinite. Let g be $Col(\omega, |U \cup \beta|^W)$-generic over W. In $W[g]$, $\mathcal{M}$ is still a model of $\mathfrak{T}$; moreover, both U and β are countably infinite, so there is a bijection f from β into U. Obviously, $W[g]$ is an outer model of V, so $(\mathcal{M}, f)$ is as required. $\qquad\square$

On the other hand, we have no good analogue of the compactness theorem for TCIs. Indeed, there are simple examples in which compactness fails. We give one such example below.

Lemma 5.14. *There is a $\Sigma_1 \cup \Pi_1$ TCI $\mathfrak{T} := (T, \{\dot{R}\}, \dot{\mathcal{U}}, \vartheta)$ with a countable transitive closure, such that*

- $\dot{R}$ *is binary relation symbol,*
- *if $x \in ran(\vartheta)$, then $x = (y, 0)$ for some set y, and*
- $\mathfrak{T}$ *is finitely consistent but not consistent.*

Proof. Choose $\dot{R}$ and $\dot{\mathcal{U}}$ to be relation symbols of their appropriate arity in $H(\omega)$. We define ϑ on $\{\dot{\mathcal{U}}, \dot{R}\}$ as follows:

$$\vartheta(\dot{\mathcal{U}}) := (\omega, 0)$$
$$S_n := \{(k, l) : 2^n \leq k, l < 2^n + n \text{ and } k < l\}$$
$$S := \bigcup \{S_n : n < \omega\}$$
$$\vartheta(\dot{R}) := (S, 0).$$

Here, ϑ encodes a set of disjoint finite linear orders of unbounded lengths. Quite clearly, $\mathfrak{T}$ has a countable transitive closure, as all first-order sentences over the signature $\{\dot{R}\}$ are members of $H(\omega)$. Next, we want T to contain the first-order definition of a strict linear ordering, namely the conjunction of the three sentences (properties):

$$\varphi_1 \text{ (irreflexivity)} : \ulcorner \forall x \, (\neg \dot{R}(x, x)) \urcorner$$
$$\varphi_2 \text{ (transitivity)} : \ulcorner \forall x \, \forall y \, \forall z \, ((\dot{R}(x, y) \wedge \dot{R}(y, z)) \implies \dot{R}(x, z)) \urcorner$$
$$\varphi_3 \text{ (trichotomy)} : \ulcorner \forall x \, \forall y \, (\dot{R}(x, y) \vee \dot{R}(y, x) \vee y = x) \urcorner.$$

Finally, we define

$$T' := \{ \ulcorner \exists x_1 \, \exists x_2 \, ... \, \exists x_n \, (\bigwedge_{1 \leq k < n} \dot{R}(x_k, x_{k+1})) \urcorner : 1 < n < \omega \}$$

and let

$$T := T' \cup \{\varphi_1, \varphi_2, \varphi_3\}.$$

Note that any finite subset of T can be satisfied by a sufficiently long finite linear order, examples of which ϑ provides in abundance. However, a model of T is necessarily an infinite linear order, and our definition of ϑ precisely prohibits all infinite linear orders. We thus have $\mathfrak{T}$ being finitely consistent but not consistent. $\square$

Fix any infinite set X. By the Lowenheim-Skolem theorem for first-order logic, the compactness theorem for first-order logic holds even if we require the base set of the models in question to be subsets of X. As a result, the failure of compactness in a TCI of the form specified by Lemma 5.14 must come from restrictions imposed by ϑ. In a sense, then, Lemma 5.14 gives one of the simplest examples of such a ϑ, considering it has a singleton as its domain.

Definition 5.15. Given a TCI $\mathfrak{T}$ and any $\mathcal{M}$, we say $\mathcal{M}$ is a *finitely determined model of* $\mathfrak{T}$ iff $\mathcal{M} \models^* \mathfrak{T}$ and for some quantifier-free sentence φ in the language associated with $\mathcal{M}$,

$$\forall W \, \forall \mathcal{M}' \, ((W \text{ is an outer model of } V \text{ and } \mathcal{M}' \in W \text{ and } \mathcal{M}' \models^* \mathfrak{T} \text{ and } \mathcal{M}' \models \varphi)$$
$$\implies \mathcal{M}' = \mathcal{M}).$$

In this case, $\mathcal{M}$ is *finitely determined by φ.*

Naturally, all finite models of any TCI are finitely determined. As it turns out, if a TCI is consistent, then all its finitely determined models can be read off a forcing notion associated with it. We will prove this in the next subsection.

5.2   Forcing Extensions and Models of TCIs

In this subsection and the next one, we investigate how one could "force" the existence of models of TCIs, under different restrictions and in a variety of settings. As a starting point, we would like to frame the problem of constructing models of TCIs in the context of Section 3, just so we can utilise Lemma 3.35, among other things.

Lemma 5.16. *There is a formula ψ_{cert} in two free variables, such that in any model of* ZFC,

- $\psi_{cert}(\mathfrak{T}, (\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}}))$ *defines a function*

$$\mathfrak{T} \mapsto (\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}})$$

  on the class of all TCIs, wherewith

  - $\mathfrak{A}_{\mathfrak{T}} = (H(|trcl(\mathfrak{T})|^+); \in)$,
  - $\mathcal{L}_{\mathfrak{T}}$ *is a set closed under negation,*
  - $\mathfrak{A}_{\mathfrak{T}}$ *is $\mathcal{L}_{\mathfrak{T}}$-suitable, and*
  - $\Gamma_{\mathfrak{T}}$ *a set of $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sentences, and*

- *whenever*

  - $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$,
  - $\psi_{cert}(\mathfrak{T}, (\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}}))$, *and*
  - T *contains only Π_2 sentences,*

  $\Gamma_{\mathfrak{T}}$ *must contain only $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$-$\Pi_2$ sentences.*

Proof. Fix any TCI $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$. We will constructively define the tuple $(\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}})$ based on $\mathfrak{T}$ alone, and in the process, check that the requirements of the lemma are satisfied.

Of course, we have to set

$$\mathfrak{A}_{\mathfrak{T}} := (H(|trcl(\mathfrak{T})|^+); \in).$$

Note that $\mathfrak{A}_{\mathfrak{T}} \models \mathsf{ZFC} - \mathsf{Powerset}$, so $\mathfrak{A}_{\mathfrak{T}}$ is a model of a sufficiently strong set theory. Next, let

$$\sigma' := \sigma \cup \{\dot{\mathcal{U}}\}, \text{ and}$$
$$U := \text{the unique } y \text{ for which there exists } z \text{ such that } \vartheta(\dot{\mathcal{U}}) = (y, z).$$

For $\dot{X} \in \sigma'$, define $\mathcal{L}_{\mathfrak{T}}(\dot{X})$ as follows:

- if $\dot{X}$ is a constant symbol and $\vartheta(\dot{X}) = (y, z)$, then

$$\mathcal{L}_{\mathfrak{T}}(\dot{X}) := \{\ulcorner \dot{X} = x \urcorner : x \in y \cap U\},$$

- if $\dot{X}$ is an n-ary relation symbol and $\vartheta(\dot{X}) = (y, z)$, then

$$\mathcal{L}_{\mathfrak{T}}(\dot{X}) := \{\ulcorner \dot{X}(x) \urcorner : x \in y \cap U^n\}, \text{ and}$$

- if $\dot{X}$ is an n-ary function symbol and $\vartheta(\dot{X}) = (y, z)$, then

$$\mathcal{L}_{\mathfrak{T}}(\dot{X}) := \{\ulcorner \dot{X}(x \restriction_n) = x(n) \urcorner : x \in y \cap U^{n+1}\}.$$

Then

$$\mathcal{L}'_{\mathfrak{T}} := \bigcup \{\mathcal{L}_{\mathfrak{T}}(\dot{X}) : \dot{X} \in \sigma'\}, \text{ and}$$
$$\mathcal{L}_{\mathfrak{T}} := \text{the closure of } \mathcal{L}'_{\mathfrak{T}} \text{ under negation.}$$

Obviously, $\mathcal{L}_{\mathfrak{T}}$ is both a member and a subset of $H(|trcl(\mathfrak{T})|^+)$, so it is definable in the language associated with $H(|trcl(\mathfrak{T})|^+)$. We thus have that $\mathfrak{A}_{\mathfrak{T}}$ is $\mathcal{L}_{\mathfrak{T}}$-suitable.

Before we get to $\Gamma_{\mathfrak{T}}$, a remark (or rather, a reminder) is imperative.

Remark 5.17. In the same vein as what was elaborated after Definition 4.10, we will use functions to pass parameters of an expression in $\mathcal{L}_{\mathfrak{T}}$ via variables, whenever necessary in the construction of $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sentences involving the symbol $\ulcorner E \urcorner$. In fact, this can be done uniformly by the universal function $\chi_{\mathfrak{T}}$:

$$(S, x_1, ..., x_n, x_{n+1}) \mapsto \begin{cases} \ulcorner S(x_1, ..., x_n, x_{n+1}) \urcorner & \text{if } S \text{ is a } n+1\text{-ary relation symbol} \\ \ulcorner S(x_1, ..., x_n) = x_{n+1} \urcorner & \text{if } S \text{ is a } n\text{-ary function symbol} \\ \ulcorner S = x_{n+1} \urcorner & \text{if } S \text{ is a constant symbol and } n = 0 \\ \emptyset & \text{otherwise,} \end{cases}$$

which is defined on V by a Δ_0 formula.

As in the case of the proof of Lemma 4.8, we will abuse notation and abbreviate the use of $\chi_{\mathfrak{T}}$ with straightforward substitutions of variables for parameters in the writing of $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sentences. There are no intrinsic "hidden costs" in terms of complexity to such a presentation of $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sentences.

Now, we define $\Gamma_{\mathfrak{T}}$ as follows:

(1) For each constant symbol $\dot{X} \in \sigma'$,

$$\ulcorner \exists x \, (E(\ulcorner \dot{\mathcal{U}}(x) \urcorner) \wedge E(\ulcorner \dot{X} = x \urcorner)) \urcorner \in \Gamma_{\mathfrak{T}},$$

$$\ulcorner \forall x \, \forall y \, ((E(\ulcorner \dot{X} = x \urcorner) \wedge E(\ulcorner \dot{X} = y \urcorner)) \implies x = y) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(2) For each n-ary relation symbol $\dot{X} \in \sigma'$,

$$\ulcorner \forall x_1 ... \forall x_n \, (E(\ulcorner \dot{X}((x_1, ..., x_n)) \urcorner) \implies (\bigwedge_{1 \leq k \leq n} E(\ulcorner \dot{\mathcal{U}}(x_k) \urcorner))) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(3) If $\vartheta(\dot{\mathcal{U}}) = (y, 1)$, then

$$\ulcorner \forall x \, (x \in y \implies E(\ulcorner \dot{\mathcal{U}}(x) \urcorner)) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(4) For each n-ary relation symbol $\dot{X} \in \sigma$ such that $\vartheta(\dot{X}) = (y, 1)$,

$$\ulcorner \forall x_1 ... \forall x_n \, (((\bigwedge_{1 \leq k \leq n} E(\ulcorner \dot{\mathcal{U}}(x_k) \urcorner)) \wedge \ulcorner \dot{X}((x_1, ..., x_n)) \urcorner \in \mathcal{L}_{\mathfrak{T}})$$

$$\implies E(\ulcorner \dot{X}((x_1, ..., x_n)) \urcorner)) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(5) For each n-ary function symbol $\dot{X} \in \sigma'$,

$$\ulcorner \forall x_1 ... \forall x_{n+1} \, (E(\ulcorner \dot{X}(x_1, ..., x_n) = x_{n+1} \urcorner) \implies (\bigwedge_{1 \leq k \leq n+1} E(\ulcorner \dot{\mathcal{U}}(x_k) \urcorner))) \urcorner \in \Gamma_{\mathfrak{T}},$$

$$\ulcorner \forall x_1 ... \forall x_n \, \exists y \, ((\bigwedge_{1 \leq k \leq n} E(\ulcorner \dot{\mathcal{U}}(x_k) \urcorner)) \implies E(\ulcorner \dot{X}(x_1, ..., x_n) = y \urcorner)) \urcorner \in \Gamma_{\mathfrak{T}},$$

$$\ulcorner \forall x_1 ... \forall x_n \, \forall y \, \forall z \, ((E(\ulcorner \dot{X}(x_1, ..., x_n) = y \urcorner) \wedge E(\ulcorner \dot{X}(x_1, ..., x_n) = z \urcorner))$$

$$\implies y = z) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(6) For each n-ary function symbol $\dot{X} \in \sigma'$ such that $\vartheta(\dot{X}) = (y, 1)$,

$$\ulcorner \forall x_1 ... \forall x_{n+1} \, (((\bigwedge_{1 \leq k \leq n+1} E(\ulcorner \dot{\mathcal{U}}(x_k) \urcorner)) \wedge \ulcorner \dot{X}(x_1, ..., x_n) = x_{n+1} \urcorner \in \mathcal{L}_{\mathfrak{T}})$$

$$\implies E(\ulcorner \dot{X}(x_1, ..., x_n) = x_{n+1} \urcorner)) \urcorner \in \Gamma_{\mathfrak{T}}.$$

(7) We finally deal with members of T. So let $\phi \in T$. We first assume that for every atomic subformula φ of ϕ,

$$\varphi \text{ contains no more than one symbol from } \sigma \text{ (counting recurrences).} \quad (5.1)$$

To see why this assumption can be made without loss of generality, notice that there are canonical algorithms M_1 and M_2 such that, when given any atomic formula φ' over σ as input,

- M_1 returns a Σ_1 formula φ which is logically equivalent to φ' and satisfies (5.1), and

- M_2 returns a Π_1 formula φ which is logically equivalent to φ' and satisfies (5.1).

Consequently, by

(i) replacing atomic subformulas of ϕ via M_1 or M_2 according to their parities, and then

(ii) canonically converting the result of (i) to prenex normal form,

we can obtain a sentence that

- is logically equivalent to ϕ,

- has each of its atomic subformulas φ satisfy (5.1), and

- is Π_2 whenever ϕ is Π_2.

Next, we transform ϕ into ϕ' by first inductively relativising ϕ to "members of $\dot{\mathcal{U}}$", then simultaneously translating all its atomic subformulas to correspond to membership in $\mathcal{L}'_{\mathfrak{T}}$. In more detail, we carry out the procedure below.

(a) Cast ϕ as a string.

Given any string A, we can view A as a sequence of (possibly non-distinct) characters. The ordering of this sequence gives rise to the notion of (relative) *position*. Intuitively, the leftmost character of A marks its first position (position = 1), and for any k, the character at the $(k+1)$-th position of A necessarily lies to the immediate right of the character at position k. Therefore, the positions of A must range from 1 to the length of A.

(b) Initialise a pointer p at the first position of the ϕ.

The rationale of having p is to help us traverse the characters of ϕ as we modify it. Like any pointer, p occupies exactly one position at any point in time. Specifically, we want p to keep moving rightwards, even though ϕ as we now know it might change in length over the run of this procedure.

We will modify ϕ in steps, each step being a pass of a numbered stage in the enumeration of this procedure. For clarity of exposition, it is useful to distinguish ϕ pre- and post-modification. As we describe the procedure going forward, we shall let ϕ refer to the unaltered string: its state right after (a). At any particular step, *the current frame* denotes the modified form of ϕ at the beginning of said step.

(c) Let x be the current position occupied by p. If there is φ such that

- φ is a subformula of ϕ,
- the leading character of φ is a first-order quantifier, and
- φ is a substring of the current frame starting at position x,

then we let φ' be the shortest such string, and proceed according to the cases below. Otherwise, skip to (d).

Case 1: $\varphi' = \ulcorner \forall x\, \psi \urcorner$ for some x and ψ. Then we replace φ' starting at x of the current frame with the string

$$\ulcorner \forall x\ (E(\ulcorner \dot{\mathcal{U}}(x) \urcorner) \implies \psi) \urcorner.$$

Case 2: $\varphi' = \ulcorner \exists x\, \psi \urcorner$ for some x and ψ. Then we replace φ' starting at x of the current frame with the string

$$\ulcorner \exists x\ (E(\ulcorner \dot{\mathcal{U}}(x) \urcorner) \wedge \psi) \urcorner.$$

When we speak of replacing a substring Y starting at x of F with another string Z, we mean to produce the concatenated string $A^\frown Z^\frown B$, where A and B are the two unique strings for which

- either A is empty or the last character of A occupies position $x - 1$ of F, and
- $F = A^\frown Y^\frown B$.

After the replacement, the position occupied by p remains unchanged — it should still be at x relative to $A^\frown Z^\frown B$. Note that replacements of this kind make no changes to F at any position less than (to the left of) x.

(d) If p is not at the rightmost position of the current frame, increment the position it occupies by 1. Otherwise skip to (f).

(e) Go to (c).

(f) Substitute each atomic subformula ψ of ϕ occurring in the current frame with $\ulcorner E(\psi) \urcorner$, bearing in mind the abbreviations adopted in Remark 5.17. These substitutions can be done simultaneously because it is impossible to have two

distinct substitutable instances occupy overlapping positions of the current frame.

The aforementioned procedure produces a sentence $\phi' \in (\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sharing the *quantification structure* of ϕ. More precisely, this means the existence of a string A such that

- A contains only quantifiers,
- A is a subsequence of both ϕ and ϕ',
- if B is a subsequence of ϕ containing only quantifiers, then B is a subsequence of A, and
- if B' is a subsequence of ϕ' containing only quantifiers, then B' is a subsequence of A.

Now, convert ϕ' to a logically equivalent formula ϕ^* in prenex normal form, through an application of the standard conversion algorithm. This algorithm preserves the quantification structure of ϕ' — so that ϕ^* and ϕ have the same quantification structure — whenever ϕ is in prenex normal form.

   Enforce that $\phi^* \in \Gamma_{\mathfrak{T}}$.

(8) Nothing else is in $\Gamma_{\mathfrak{T}}$.

Assume T contains only Π_2 sentences. Then necessarily every member of T is in prenex normal form. As a consequence, the transformation

$$\varpi : \phi \mapsto \phi^*$$

described in (7) takes every member of T to a $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$ sentence in prenex normal form with the same quantification structure, making $\varpi''T$ a set of $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$-$\Pi_2$ sentences. Clearly, all additions to $\Gamma_{\mathfrak{T}}$ as per (1) to (6) are $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$-$\Pi_2$ sentences. By (8), $\Gamma_{\mathfrak{T}}$ contains only $(\mathcal{L}_{\mathfrak{T}})^*_{\mathfrak{A}_{\mathfrak{T}}}$-$\Pi_2$ sentences. $\qquad\square$

Fix ψ_{cert} to be as in Lemma 5.16. We are then justified in our next definition.

Definition 5.18. Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI. Define

$\mathfrak{A}_{\mathfrak{T}} :=$ the unique $\mathfrak{A}$ for which there are $\mathcal{L}$ and Γ satisfying $\psi_{cert}(\mathfrak{T}, (\mathfrak{A}, \mathcal{L}, \Gamma))$,

$\mathcal{L}_{\mathfrak{T}} :=$ the unique $\mathcal{L}$ for which there are $\mathfrak{A}$ and Γ satisfying $\psi_{cert}(\mathfrak{T}, (\mathfrak{A}, \mathcal{L}, \Gamma))$, and

$\Gamma_{\mathfrak{T}} :=$ the unique Γ for which there are $\mathfrak{A}$ and $\mathcal{L}$ satisfying $\psi_{cert}(\mathfrak{T}, (\mathfrak{A}, \mathcal{L}, \Gamma))$.

Definition 5.19. Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI. Define

$P(\mathfrak{T}) := \{p \in [\mathcal{L}_{\mathfrak{T}}]^{<\omega} : \Vdash_{Col(\omega, |trcl(\mathfrak{A}_{\mathfrak{T}})|)} \exists \Sigma \, (\text{``}\Sigma \, \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } p\text{''})\}$,

$\leq_{\mathbb{P}(\mathfrak{T})} := \{(p, q) \in P(\mathfrak{T}) \times P(\mathfrak{T}) : q \subset p\}$, and

$\mathbb{P}(\mathfrak{T}) := (P(\mathfrak{T}), \leq_{\mathbb{P}(\mathfrak{T})})$.

By Lemma 5.16, if $\mathfrak{T}$ is a TCI and A the base set of $\mathfrak{A}_{\mathfrak{T}}$, then

- $\mathbb{P}(\mathfrak{T}) \in A \cap \mathcal{P}(A)$, so $\mathbb{P}(\mathfrak{T})$ is definable in the language associated with $\mathfrak{A}_{\mathfrak{T}}$, and

- $(\mathfrak{A}_{\mathfrak{T}}, \mathbb{P}(\mathfrak{T}))$ is good for $\mathcal{L}_{\mathfrak{T}}$.

At this juncture, it is customary for us to revisit the main forcing construction of the previous section.

Remark 5.20. Consider the sequence of forcing notions

$$\{\mathbb{P}_\lambda : \lambda \in C \cup \{\lambda_f\}\}$$

constructed within W in the proof of Theorem 4.8. The adaptation of this inductive construction to the language of TCIs is straightforward: given $\lambda \in C \cup \{\lambda_f\}$ and $\vec{\mathbb{P}}_\lambda := \{\mathbb{P}_\theta : \theta \in \lambda \cap C\}$, define σ to contain

(A) the variable names of all subsets of $H(\lambda_f)$ germane to the definition of $\mathbb{P}_\lambda$ in the original proof (this set includes $\vec{\mathbb{P}}_\lambda$), along with

(B) two other ternary relation symbols, $\dot{F}$ and $\dot{X}$,

and nothing else. If $\dot{Z} \in \sigma$ is of type (A), then we interpret $\dot{Z}$ invariably as whichever set it is defined to be in the original proof; for example, when $\dot{Z} = \ulcorner C \urcorner$, we set $\vartheta(\dot{Z}) = (C, 1)$. On the other hand, if $\dot{Z} \in \sigma$ is of type (B), then we interpret $\dot{Z}$ according to either (L1)$_\lambda$ or (L2)$_\lambda$ based on the identity of $\dot{Z}$; for example, when $\dot{Z} = \dot{F}$, we set

$$\vartheta(\dot{Z}) = (\{(i, n, \alpha) : i \in R \cap \lambda, n < \omega \text{ and } \alpha < i\}, 0).$$

Also set $\vartheta(\dot{\mathcal{U}}) = (H(\lambda_f), 1)$.

Next, modify (C1)$_\lambda$ to (C8)$_\lambda$ such that each subformula of the form $\ulcorner E(\ulcorner \phi \urcorner) \urcorner$ is replaced by $\ulcorner \phi \urcorner$, and let T contain only these formulas (noting and adjusting for the abuse of notation in the original presentation). Finally, letting $\mathfrak{T}_\lambda = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$, it takes no more than a routine unfurling and checking of definitions to ascertain that $\mathbb{P}(\mathfrak{T}_\lambda)$ is isomorphic to the forcing notion $\mathbb{P}_\lambda$ defined in the original proof.

We see in Remark 5.20 that our procedure of associating a partial order with each TCI can be used to generate forcing notions as complex as the ones constructed to solve a difficult problem in set theory. More formal declarations of the power of this procedure will appear — and be proven — in the later parts of this section. But before that, let us return to the setting of the ground.

Lemma 5.21. *There is a formula ψ_{trans} in two free variables, absolute for transitive models of* ZFC $-$ Powerset, *such that $\psi_{trans}(\mathcal{M}, \Sigma)$ defines a bijection from*

$$\{\mathcal{M} : \mathcal{M} \models^* \mathfrak{T}\}$$

into

$$\{\Sigma : \Sigma \; \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } \emptyset\}$$

for every TCI $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$.

Proof. Let U be the unique y for which there exists z such that $\vartheta(\dot{\mathcal{U}}) = (y, z)$. Given a model $\mathcal{M} = (M; \mathcal{I})$ of $\mathfrak{T}$, define

$$U(\mathcal{M}) := \{\ulcorner \dot{\mathcal{U}}(x) \urcorner : x \in M\} \cup \{\ulcorner \neg \dot{\mathcal{U}}(x) \urcorner : x \in U \setminus M\}.$$

Now define ψ_{trans} as follows:

$$\psi_{trans}(\mathcal{M}, \Sigma) \iff \Sigma = (U(\mathcal{M}) \cup Diag(\mathcal{M})) \cap \mathcal{L}_{\mathfrak{T}},$$

where $Diag(\mathcal{M})$ is the atomic diagram of $\mathcal{M}$. Verily, ψ_{trans} is a Δ_0 formula with parameters $\mathfrak{T}$ and $\mathcal{L}_{\mathfrak{T}}$, while $\mathcal{L}_{\mathfrak{T}}$ is Δ_1-definable with parameter $\mathfrak{T}$. As such, ψ_{trans} is absolute for transitive models of $\mathsf{ZFC} - \mathsf{Powerset}$. We can then straightforwardly check that ψ_{trans} defines a bijection as required by the lemma, based on how the triple $(\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}})$ is constructed from $\mathfrak{T}$. $\qquad\square$

Remark 5.22. For any TCI $\mathfrak{T}$ and any structure $\mathfrak{A}$, if $\mathfrak{A}$ is $\mathcal{L}_{\mathfrak{T}}$-suitable and $\Gamma_{\mathfrak{T}}$ is a set of $(\mathcal{L}_{\mathfrak{T}})_{\mathfrak{A}}^*$ sentences, then for all Σ and p,

$$\Sigma\,\Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } p \iff \Sigma\,\Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A})\text{-certifies } p.$$

We can therefore replace $\mathfrak{A}_{\mathfrak{T}}$ in Lemma 5.21 with any appropriate $\mathfrak{A}$ and still have the lemma hold true for the same ψ_{trans}.

Fix ψ_{trans} to be as in Lemma 5.21 for the next definition.

Definition 5.23. Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI in V and $\mathcal{M}$ be a model of $\mathfrak{T}$ in some outer model of V. Define

$$\Sigma(\mathcal{M}) := \text{the unique } \Sigma \text{ for which } \psi_{trans}(\mathcal{M}, \Sigma).$$

It is time to cash the cheque issued in Remark 5.8.

Lemma 5.24. *Let* $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ *be a TCI. Then* $\mathfrak{T}$ *is consistent iff*

$$\Vdash_{Col(\omega, \lambda)} \exists \mathcal{M}\ (\text{``}\mathcal{M} \models^* \mathfrak{T}\text{''}),$$

where $\lambda := |H(|trcl(\mathfrak{T})|^+)|$.

Proof. By Lemmas 5.16 and 5.21, we can find a triple $(\mathfrak{A}_{\mathfrak{T}}, \mathcal{L}_{\mathfrak{T}}, \Gamma_{\mathfrak{T}})$ such that

- $\mathfrak{A}_{\mathfrak{T}} = (H(|trcl(\mathfrak{T})|^+); \in)$,
- $\mathcal{L}_{\mathfrak{T}}$ is a set closed under negation,
- $\mathfrak{A}_{\mathfrak{T}}$ is $\mathcal{L}_{\mathfrak{T}}$-suitable,
- $\Gamma_{\mathfrak{T}}$ a set of $(\mathcal{L}_{\mathfrak{T}})_{\mathfrak{A}_{\mathfrak{T}}}^*$ sentences, and

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- $\mathfrak{T}$ is consistent iff for some outer model W of V,

$$\exists \Sigma \in W \; (\text{``}\Sigma \, \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } \emptyset\text{''}).$$

Then the conjunction of

- Lemma 3.34,
- the fact that every forcing extension of V is an outer model of V, and
- the fact that every outer model of V is a weak outer model of V

tells us that $\mathfrak{T}$ is consistent iff

$$\Vdash_{Col(\omega, \lambda)} \exists \Sigma \; (\text{``}\Sigma \, \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } \emptyset\text{''}),$$

where $\lambda := |H(|trcl(\mathfrak{T})|^+)|$.

As every $Col(\omega, \lambda)$-generic extension of V is a transitive model of ZFC $-$ Powerset, we can apply Lemma 5.21 again to complete the proof. $\qquad\square$

Intuitively, the consistency of a theory — however it is defined — should be absolute in a sufficiently strong sense. This is the case for first-order theories, any of which consistency is absolute for transitive models of set theory. The following Lemma establishes a similar absoluteness property with regards to the consistency of a TCI.

Lemma 5.25. *Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI. Then $\mathfrak{T}$ being consistent is absolute for transitive models of* ZFC *sharing the same ordinals.*

Proof. This is very much similar to the proof of Lemma 3.36. Nevertheless, we shall provide details.

Let V' and W be transitive models of ZFC with $ORD^{V'} = ORD^W$ and $\mathfrak{T} \in V' \subset W$. If $\mathfrak{T}$ is consistent in W, then $\mathfrak{T}$ has a model in some outer model of W. Said outer model is also an outer model of V', so $\mathfrak{T}$ is consistent in V' as well.

Now assume $\mathfrak{T}$ is consistent in V'. Letting

$$\lambda := |H(((|trcl(\mathfrak{T})|^{V'})^+)^{V'})^{V'}|^{V'},$$

Lemma 5.24 gives us

$$\Vdash_{Col(\omega, \lambda)} \exists \mathcal{M} \; (\text{``}\mathcal{M} \models^* \mathfrak{T}\text{''})$$

in V'. Note that

$$\mathbb{P} := Col(\omega, \lambda)^{V'}$$

remains a forcing notion in W, so consider g a $\mathbb{P}$-generic filter over W. Necessarily, g is also $\mathbb{P}$-generic over V', and further, $V'[g] \subset W[g]$. In $V'[g]$, $\mathfrak{T}$ is forced to have a model — call it $\mathcal{M}$. Being a model of $\mathfrak{T}$ is absolute for transitive models of ZFC, so $\mathcal{M} \models^* \mathfrak{T}$ holds in $W[g]$ too. Since $W[g]$ is an outer model of W, $\mathfrak{T}$ must be consistent in W. $\quad\square$

We now define a class of generic objects that manifest as models of TCIs.

Definition 5.26. Let $\mathfrak{T}$ be a consistent Π_2 TCI.

If $\mathfrak{A}$ and $\mathbb{P}$ are such that $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$ and $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}_{\mathfrak{T}}$, then a $(\mathbb{P}, \mathfrak{A})$-*generic model* of $\mathfrak{T}$ is a model $\mathcal{M}$ of $\mathfrak{T}$ satisfying

$$\Sigma(\mathcal{M}) = \left(\bigcup g\right) \cap \mathcal{L}_{\mathfrak{T}}$$

for some $\mathbb{P}$-generic filter g over $\mathfrak{A}$. In this case, we say g *witnesses $\mathcal{M}$ is a $(\mathbb{P}, \mathfrak{A})$-generic model of* $\mathfrak{T}$. We say g *witnesses a $(\mathbb{P}, \mathfrak{A})$-generic model of* $\mathfrak{T}$ iff for some $\mathcal{M}$, g witnesses $\mathcal{M}$ is a $(\mathbb{P}, \mathfrak{A})$-generic model of $\mathfrak{T}$.

We call $\mathcal{M}$ a $\mathfrak{A}$-*generic model* of $\mathfrak{T}$ iff for some $\mathbb{P}$ definable in the language associated with $\mathfrak{A}$ such that $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}_{\mathfrak{T}}$, $\mathcal{M}$ is a $(\mathbb{P}, \mathfrak{A})$-generic model of $\mathfrak{T}$.

We call $\mathcal{M}$ a *generic model* of $\mathfrak{T}$ iff for some $\mathfrak{A}$ and $\mathbb{P}$ such that $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$ and $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}_{\mathfrak{T}}$, $\mathcal{M}$ is a $(\mathbb{P}, \mathfrak{A})$-generic model of $\mathfrak{T}$.

Observation 5.27.

(1) If $\mathfrak{T}$ is a consistent Π_2 TCI, and $\mathfrak{A}$ and $\mathbb{P}$ are such that $\mathbb{P}$ is definable in the language associated with $\mathfrak{A}$ and $(\mathfrak{A}, \mathbb{P})$ is good for $\mathcal{L}_{\mathfrak{T}}$, then every $(\mathbb{P}, \mathfrak{A})$-generic model of $\mathfrak{T}$ is a $(\mathbb{P}, \mathfrak{A})$-generic object.

(2) In the same vein as Observation 2.20, we see that given any consistent Π_2 TCI $\mathfrak{T}$,

$$\forall x \ (x \text{ is a } (\mathbb{P}(\mathfrak{T}), \mathfrak{A}_{\mathfrak{T}})\text{-generic model of } \mathfrak{T}$$
$$\iff \ x \text{ is a } (\mathbb{P}(\mathfrak{T}), V)\text{-generic model of } \mathfrak{T})$$

in every outer model of V. As a result, we can thus safely talk about $(\mathbb{P}(\mathfrak{T}), V)$-generic models of $\mathfrak{T}$ without the need to quantify over all sets.

Our definition of generic models might seem overly restrictive at first glance. The next lemma provides justification that it is not so.

Lemma 5.28. *Let $\mathfrak{T}$ be a TCI. If $\mathcal{M}$ is a model of $\mathfrak{T}$ in some forcing extension of V, then $\mathcal{M}$ is a V-generic model of $\mathfrak{T}$.*

Proof. Let $\mathcal{M}$ be a model of $\mathfrak{T}$ in a forcing extension W of V. Without loss of generality, we can assume the existence of $\mathbb{P} = (P, \leq_{\mathbb{P}})$ and $\dot{\Sigma}$ such that

- $\mathbb{P}$ is a forcing notion,
- $P \cap \mathcal{L}_{\mathfrak{T}} = \emptyset$,
- $\dot{\Sigma}$ is a $\mathbb{P}$-name,
- $\Vdash_{\mathbb{P}} \text{``}\exists \mathcal{M}' \ (\mathcal{M}' \models^* \mathfrak{T} \text{ and } \Sigma(\mathcal{M}') = \dot{\Sigma})\text{''}$, and

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- for some $\mathbb{P}$-generic filter g_0 over V, $\Sigma(\mathcal{M}) = \dot{\Sigma}[g_0]$.

Define

$$P^* := \{x \in [P \cup \mathcal{L}_{\mathfrak{I}}]^{<\omega} : x \cap P \text{ has a common extension in } \mathbb{P} \text{ and}$$
$$\forall y \, \exists p \, (y \in x \cap \mathcal{L}_{\mathfrak{I}} \implies (p \in x \cap P \text{ and } p \Vdash_{\mathbb{P}} y \in \dot{\Sigma}))\}$$

and have $\mathbb{P}^* := (P^*, \supset)$.

Fix $x \in P^*$, and let p be a common extension in $\mathbb{P}$ of the members of $x \cap P$. Then any extension of $\{p\}$ in $\mathbb{P}^*$ is compatible with x in $\mathbb{P}^*$. This means that

$$(\pi : P \longrightarrow P^*) \, [p \mapsto \{p\}]$$

is a dense weak embedding from $\mathbb{P}$ into $\mathbb{P}^*$. As a result,

$$\sigma : g \mapsto \mathrm{UC}(w(\mathbb{P}^*), \pi"g)$$

is a bijection from

$$\{g : g \text{ is a } \mathbb{P}\text{-generic filter over } V\} \cap W$$

into

$$\{h : h \text{ is a } \mathbb{P}^*\text{-generic filter over } V\} \cap W$$

with inverse

$$\tau : h \mapsto \pi^{-1}h,$$

in every weak outer model W of V. The following fact is easy to see.

Fact 5.29. If h is a $\mathbb{P}^*$-generic filter over V, then $[(\bigcup h) \cap P]^{<\omega} \subset h$.

Proposition 5.30. *Let g be a $\mathbb{P}$-generic filter over V. Then $(\bigcup \sigma(g)) \cap P = g$.*

Proof. Denote $(\bigcup \sigma(g)) \cap P$ as g'. By the definitions of π and σ, $g \subset g'$ clearly. Choose $p \in g'$. Since $\sigma(g)$ is a $\mathbb{P}^*$-generic filter over V, Fact 5.29 tells us that $\{p\} \in \sigma(g)$. As $\tau = \sigma^{-1}$, $p \in g$, and we are done. $\qquad\square$

Proposition 5.31. *Let g be a $\mathbb{P}$-generic filter over V. Then*

$$(\bigcup \sigma(g)) \cap \mathcal{L}_{\mathfrak{I}} = \{y \in \mathcal{L}_{\mathfrak{I}} : \exists p \, (p \in (\bigcup \sigma(g)) \cap P \text{ and } p \Vdash_{\mathbb{P}}^V y \in \dot{\Sigma})\}.$$

Proof. Denote $(\bigcup \sigma(g)) \cap \mathcal{L}_{\mathfrak{I}}$ as Σ and

$$\{y \in \mathcal{L}_{\mathfrak{I}} : \exists p \, (p \in (\bigcup \sigma(g)) \cap P \text{ and } p \Vdash_{\mathbb{P}}^V y \in \dot{\Sigma})\}$$

as Σ'. By the definition of P^* and the fact that $\sigma(g) \subset P^*$, $\Sigma \subset \Sigma'$. Choose $y \in \Sigma'$, so that there is $p \in (\bigcup \sigma(g)) \cap P$ with $p \Vdash_{\mathbb{P}}^V y \in \dot{\Sigma}$. But that entails the density of

$$D_p := \{x \in P^* : y \in x\}$$

below $\{p\}$ in $\mathbb{P}^*$. Since $\{p\} \in \sigma(g)$ by Fact 5.29, $D_p \cap \sigma(g) \neq \emptyset$, whence $y \in \Sigma$. $\qquad\square$

Combining Propositions 5.30 and 5.31, we know that whenever g is a $\mathbb{P}$-generic filter over V, $(\bigcup \sigma(g)) \cap \mathcal{L}_{\mathfrak{T}} = \dot{\Sigma}[g]$. In particular, $(\bigcup \sigma(g_0)) \cap \mathcal{L}_{\mathfrak{T}} = \Sigma(\mathcal{M})$. As $\sigma(g_0)$ is a $\mathbb{P}^*$-generic filter over V, $\mathcal{M}$ is a $(\mathbb{P}^*, V)$-generic model of $\mathfrak{T}$. $\square$

Henceforth, we will look into the extent to which generic models of Π_2 TCIs are abundant. We begin with two of the main theorems in this subsection, thereby establishing the connection between Π_2 TCIs and forcing extensions.

Theorem 5.32. *Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a Π_2 TCI. If $\mathfrak{T}$ is consistent, then every $\mathbb{P}(\mathfrak{T})$-generic filter over V witnesses a $(\mathbb{P}(\mathfrak{T}), V)$-generic model.*

Proof. The lemma follows directly from Lemmas 3.35, 5.16 and 5.21, noting that the hypotheses of Lemma 3.35 are satisfied with

- $\mathfrak{A}_{\mathfrak{T}}$ in place of $\mathfrak{A}$,
- $\mathcal{L}_{\mathfrak{T}}$ in place of $\mathcal{L}$,
- $[\mathcal{L}_{\mathfrak{T}}]^{<\omega}$ in place of B,
- $\mathbb{P}(\mathfrak{T})$ in place of $\mathbb{P}$, and
- $\Gamma_{\mathfrak{T}}$ in place of Γ. $\square$

Theorem 5.33. *Let $\mathbb{B} = (B, \mathbb{0}, \mathbb{1}, \leq_{\mathbb{B}}, \wedge_{\mathbb{B}}, \vee_{\mathbb{B}}, \sim_{\mathbb{B}})$ be a boolean algebra. Then there is a consistent Π_2 TCI $\mathfrak{T}$ and a dense weak embedding from $\mathbb{B}_0^+$ into $\mathbb{P}(\mathfrak{T})$.*

Proof. Choose

- $\dot{\mathbb{1}}$ to be a constant symbol, and
- $\dot{\leq}, \dot{\sim}, \dot{G}$ and $\dot{\mathcal{U}}$ to be distinct relation symbols of arities $2, 2, 1$ and 1 respectively.

Set σ to be $\{\dot{\mathbb{1}}, \dot{\leq}, \dot{\sim}, \dot{G}\}$. We define ϑ on $\{\dot{\mathcal{U}}, \dot{\mathbb{1}}, \dot{\leq}, \dot{\sim}, \dot{G}\}$ as follows:

$$\vartheta(\dot{\mathcal{U}}) := (B^+, 1)$$
$$\vartheta(\dot{\mathbb{1}}) := (\{\mathbb{1}\}, 1)$$
$$\vartheta(\dot{\sim}) := (\sim_{\mathbb{B}^+}, 1)$$
$$\vartheta(\dot{\leq}) := (\leq_{\mathbb{B}^+}, 1)$$
$$\vartheta(\dot{G}) := (B^+, 0).$$

Now, have T contain only the sentences

$$\ulcorner \forall p \, \forall q \, \exists r \, ((\dot{G}(p) \wedge \dot{G}(q)) \implies (\dot{G}(r) \wedge \dot{\leq}(r,p) \wedge \dot{\leq}(r,q))) \urcorner,$$
$$\ulcorner \forall p \, \forall q \, ((\dot{\leq}(p,q) \wedge \dot{G}(p)) \implies \dot{G}(q)) \urcorner, \text{ and}$$
$$\ulcorner \forall p \, \exists q \, (p = \dot{\mathbb{1}} \vee \dot{G}(p) \vee (\dot{G}(q) \wedge \dot{\sim}(p,q))) \urcorner.$$

Let $\mathfrak{T} := (T, \sigma, \dot{\mathcal{U}}, \vartheta)$. Then $\mathfrak{T}$ is clearly a consistent Π_2 TCI, for any ultrafilter on $\mathbb{B}$ is an interpretation of $\dot{G}$ satisfying $\mathfrak{T}$. Further,

(1) for each $p \in B^+$, $\{\ulcorner \dot{G}(p) \urcorner\} \in P(\mathfrak{T})$, and

(2) for each $x \in P(\mathfrak{T})$,

$$\exists p \, \forall q \, (x \cup \{\ulcorner \dot{G}(p) \urcorner\} \in P(\mathfrak{T}) \wedge (\ulcorner \dot{G}(q) \urcorner \in x \implies p \leq_{\mathbb{B}^+} q)$$
$$\wedge \, (\ulcorner \neg \dot{G}(q) \urcorner \in x \implies p \leq_{\mathbb{B}^+} (\sim_{\mathbb{B}^+} (q)))).$$

Define $\pi : B^+ \longrightarrow P(\mathfrak{T})$ to be

$$p \mapsto \{\ulcorner \dot{G}(p) \urcorner\},$$

which is possible by (1). We argue that π is a dense weak embedding from $\mathbb{B}_0^+$ into $\mathbb{P}(\mathfrak{T})$.

Denote $w(\leq_{\mathbb{P}(\mathfrak{T})})$ as $\leq^*$, and $w(\mathbb{P}(\mathfrak{T}))$ as $\mathbb{P}^*$.

We first show that π is a weak embedding. Since $\mathbb{B}_0^+$ is separative by Fact 2.27, $\leq_{\mathbb{B}^+} = w(\leq_{\mathbb{B}^+})$. Assume $p \leq_{\mathbb{B}^+} q$ and let $x \leq_{\mathbb{P}(\mathfrak{T})} \pi(p)$. By the definition of T, $x \cup \pi(q) \in P(\mathfrak{T})$, so $x \not\perp_{\mathbb{P}(\mathfrak{T})} \pi(q)$. We thus have $\pi(p) \leq^* \pi(q)$. Next, assume $p \not\leq_{\mathbb{B}^+} q$. Then there is $r \leq_{\mathbb{B}^+} p$ such that $r \perp_{\mathbb{B}_0^+} q$, because $\mathbb{B}_0^+$ is separative. This means $(\pi(p) \cup \pi(r)) \perp_{\mathbb{P}(\mathfrak{T})} \pi(q)$, and $\pi(p) \not\leq^* \pi(q)$. Lastly, if $\pi(p) \not\perp_{\mathbb{P}^*} \pi(q)$, then $\pi(p) \not\perp_{\mathbb{P}(\mathfrak{T})} \pi(q)$, which guarantees $p \not\perp_{\mathbb{B}_0^+} q$.

To see that $ran(\pi)$ is dense in $\mathbb{P}^*$, fix any $x \in P(\mathfrak{T})$. By (2), there is $p \in B^+$ for which

- $x \cup \{\ulcorner \dot{G}(p) \urcorner\} \in P(\mathfrak{T})$,
- $\forall q \, (\ulcorner \dot{G}(q) \urcorner \in x \implies p \leq_{\mathbb{B}^+} q)$, and
- $\forall q \, (\ulcorner \neg \dot{G}(q) \urcorner \in x \implies p \leq_{\mathbb{B}^+} (\sim_{\mathbb{B}^+} (q)))$.

What this entails is, whenever $\mathcal{M} \models^* \mathfrak{T}$ and $\pi(p) \subset \Sigma(\mathcal{M})$, we must have $x \subset \Sigma(\mathcal{M})$. By the definition of $\mathbb{P}(\mathfrak{T})$ and the bijectivity of the function $\mathcal{M} \mapsto \Sigma(\mathcal{M})$, $\pi(p) \leq^* x$. $\quad \square$

From the perspective of forcing extensions, Theorems 5.32 and 5.33, in view of Remarks 2.29 and 2.32, tell us that

1. every consistent Π_2 TCI has a model that "generates" a forcing extension, and

2. every forcing extension is "generated" by a model of a consistent Π_2 TCI.

This lends credence and weight to the informal thesis (a slogan, rather),

"Forcing is exactly Π_2."

Further, it suggests one way to use a complexity class defined on TCIs as a measure of — or a proxy for — accessibility within the context of the set-theoretic multiverse.

5.3 More Generic Models of Π_2 TCIs

The remainder of this section concerns itself with finer details regarding the existence of generic models of Π_2 TCIs. First up is a generic version of Lemma 5.13.

Lemma 5.34. *Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a Π_2 TCI with an infinite model in some outer model of V. Then for every infinite ordinal β, there is a forcing notion $\mathbb{P}$ such that whenever g is a $\mathbb{P}$-generic filter over V, there are sets $\mathcal{M} = (U; \mathcal{I})$ and f in some outer model of V for which*

(a) g witnesses $(\mathcal{M}, f)$ is a $(\mathbb{P}, V)$-generic object,

(b) g witnesses $\mathcal{M}$ is a $(\mathbb{P}, V)$-generic model of $\mathfrak{T}$, and

(c) $f : \beta \longrightarrow U$ is a bijection.

Proof. Fix an infinite ordinal β. We want to modify $\mathfrak{T}$ to get another consistent Π_2 TCI $\mathfrak{T}^*$ such that from every model $\mathcal{M}^*$ of $\mathfrak{T}^*$ we can read off a structure $\mathcal{M} = (U; \mathcal{I})$ and a function f satisfying both

(b)' $\mathcal{M} \models^* \mathfrak{T}$ and $\Sigma(\mathcal{M}^*) \cap \mathcal{L}_{\mathfrak{T}} = \Sigma(\mathcal{M})$

as well as (c) of the lemma.

Note that we can, without loss of generality, assume σ contains only relation symbols and constant symbols. This is because for any function symbol $\dot{X}$ and any $n < \omega$, $\dot{X}$ being an n-ary function is definable in a $(n+1)$-ary relation symbol $\dot{Y}$ via the conjunction of the Π_2 sentences

$$\ulcorner \forall x_1 ... \forall x_n \, \exists y \, (\dot{Y}(x_1, ..., x_n, y)) \urcorner \text{ and}$$
$$\ulcorner \forall x_1 ... \forall x_n \, \forall y \, \forall z \, ((\dot{Y}(x_1, ..., x_n, y) \wedge \dot{Y}(x_1, ..., x_n, z)) \implies y = z) \urcorner,$$

if we interpret formulas of the form $\dot{Y}(x_1, ..., x_n, x_{n+1})$ as $\dot{X}(x_1, ..., x_n) = x_{n+1}$.

Have (y, z) be $\vartheta(\dot{\mathcal{U}})$ and σ' be $\sigma \cup \{\dot{\mathcal{U}}\}$. Choose

- $\dot{F}$ to be a unary function symbol not in σ',

- $\dot{\mathcal{U}}^*$ and $\dot{V}$ to be distinct unary relation symbols not in σ', and

- $\dot{c}$ to be a constant symbol not in σ', for each $c \in y$, such that $\dot{c} \neq \dot{d}$ if $\{c, d\} \subset y$ and $c \neq d$.

Let

$$\sigma^* := \sigma \cup \{\dot{F}, \dot{\mathcal{U}}, \dot{V}\} \cup \{\dot{c} : c \in y\}.$$

We specify ϑ^* by how it acts on members of its domain. Pick a set b of cardinality $|\beta|$ that is disjoint from y, and set $\vartheta^*(\dot{\mathcal{U}}^*) := (y \cup b, 1)$. Make the assignments

$$\vartheta^*(\dot{F}) := (b \times y, 0)$$
$$\vartheta^*(\dot{\mathcal{U}}) := (y, z)$$
$$\vartheta^*(\dot{V}) := (b, 1)$$
$$\vartheta^*(\dot{c}) := (\{c\}, 0) \text{ for each } c \in y.$$

Whenever $\dot{X} \in \sigma$ and $\vartheta(\dot{X}) = (y', z')$, we define $\vartheta^*(\dot{X}) := (y', min\{z, z'\})$.

Now, we modify $\Gamma_{\mathfrak{T}}$ by first removing members of the type described in (3) and (4) of Lemma 5.16, and then for each remaining member φ of $\Gamma_{\mathfrak{T}}$, replacing every subformula of φ of the form $\ulcorner E(\ulcorner x \urcorner) \urcorner$ with $\ulcorner x \urcorner$.

Call the result of said modification T'. Whenever $\dot{X} \in \sigma$ is an n-ary relation symbol with $\vartheta(\dot{X}) = (y', 1)$, define

$$T(\dot{X}) := \{\ulcorner \bigwedge_{1 \leq k \leq n} \dot{\mathcal{U}}(\dot{c}_k) \implies \dot{X}(\dot{c}_1, ..., \dot{c}_n) \urcorner : (c_1, ..., c_n) \in y' \cap y^n\}.$$

Finally, define T^* to be the union of T',

$$\bigcup \{T(\dot{X}) : \dot{X} \in \sigma \text{ and } \exists y' \, (\vartheta(\dot{X}) = (y', 1))\},$$

and the finite set of sentences

$$T_0^* := \{\ulcorner \forall x \, \exists y \, (\dot{V}(x) \implies (\dot{\mathcal{U}}(y) \wedge \dot{F}(x) = y)) \urcorner,$$
$$\ulcorner \forall x \, \exists y \, (\dot{\mathcal{U}}(x) \implies (\dot{V}(y) \wedge \dot{F}(y) = x)) \urcorner,$$
$$\ulcorner \forall x \, \forall y \, (\dot{F}(x) = \dot{F}(y) \implies x = y) \urcorner\}.$$

Clearly T^* is a set of Π_2 sentences over the vocabulary σ^*.

A routine verification should enable the reader to see that

$$T_1^* := T' \cup (\bigcup \{T(\dot{X}) : \dot{X} \in \sigma \text{ and } \exists y' \, (\vartheta(\dot{X}) = (y', 1))\})$$

is basically a translation of $\Gamma_{\mathfrak{T}}$ in our expanded vocabulary σ^*, with the set of constants $\{\dot{c} : c \in y\}$ fulfilling a role similar to that of the parameter $\mathcal{L}_{\mathfrak{T}}$ (say, in the context of (4) in the proof of Lemma 5.16). On the other hand, T_0^* expresses precisely the requirement that a bijection from b (and thus from β, in any outer model of V) into $\mathcal{I}(\dot{\mathcal{U}})$ exists for every $\mathcal{I}$ satisfying $(y \cup b; \mathcal{I}) \models^* (T_0^*, \sigma^*, \dot{\mathcal{U}}^*, \vartheta^*)$ — said bijection is just $\mathcal{I}(\dot{F})$. In fact, it does so in a manner independent of truths over the vocabulary σ', so that whenever $\mathcal{M}^*$ is a model of

$$\mathfrak{T}^* := (T^*, \sigma^*, \dot{\mathcal{U}}^*, \vartheta^*) = (T_0^* \cup T_1^*, \sigma^*, \dot{\mathcal{U}}^*, \vartheta^*),$$

we have

$$\Sigma(\mathcal{M}^*) \cap \mathcal{L}_{\mathfrak{T}} = \Sigma(\mathcal{M})$$

for some model $\mathcal{M}$ of $\mathfrak{T}$. By Lemma 5.13 and our assumptions on $\mathfrak{T}$, $\mathfrak{T}^*$ is consistent.

We have thus checked that $\mathfrak{T}^*$ possesses the properties we want: it is a consistent Π_2 TCI, and from every model $\mathcal{M}^*$ of $\mathfrak{T}^*$ we can read off a structure $\mathcal{M} = (U; \mathcal{I})$ and a function f satisfying both (b)' defined at the beginning of the proof as well as (c) of the lemma. An invocation of Theorem 5.32 with $\mathfrak{T}^*$ in place of $\mathfrak{T}$ then completes the proof. $\qquad\square$

By strengthening the hypotheses on $\mathfrak{T}$ in Lemma 5.34, we can derive more from our witnesses.

Lemma 5.35. *Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a Π_2 TCI with only infinite model(s) across all outer models of V. Then for every infinite ordinal β, there is a forcing notion $\mathbb{P}$ such that whenever g is a $\mathbb{P}$-generic filter over V, there are sets $\mathcal{M} = (U; \mathcal{I})$ and f in some outer model of V for which*

   (a) g witnesses $(\mathcal{M}, f)$ is a $(\mathbb{P}, V)$-generic object,

   (b) g witnesses $\mathcal{M}$ is a $(\mathbb{P}, V)$-generic model of $\mathfrak{T}$,

   (c) $g \cap P(\mathfrak{T})$ witnesses $\mathcal{M}$ is a $(\mathbb{P}(\mathfrak{T}), V)$-generic model of $\mathfrak{T}$, and

   (d) $f : \beta \longrightarrow U$ is a bijection.

Proof. Construct $\mathfrak{T}^*$ from $\mathfrak{T}$ as per the proof of Lemma 5.34.

Proposition 5.36. $\mathbb{P}(\mathfrak{T}) \lessdot \mathbb{P}(\mathfrak{T}^*)$.

Proof. Observe that, if $\mathcal{M}$ is a model of $\mathfrak{T}$ is some outer model W of V, then $\mathcal{M}$ extends to a model of $\mathfrak{T}^*$ in an outer model W' of W. As a result, $\mathbb{P}(\mathfrak{T})$ is a suborder of $\mathbb{P}(\mathfrak{T}^*)$. To show the regularity of $\mathbb{P}(\mathfrak{T})$ as a suborder of $\mathbb{P}(\mathfrak{T}^*)$, let $p \in \mathbb{P}(\mathfrak{T}^*)$. Define

$$q_0 := \{\ulcorner \dot{\mathcal{U}}(j) \urcorner : \ulcorner \dot{F}(i) = j \urcorner \in p\}$$

and let

$$q := (p \cap \mathcal{L}_{\mathfrak{T}}) \cup q_0.$$

Obviously $q \in \mathbb{P}(\mathfrak{T})$. Consider any $\Sigma \, \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})$-certifying q in some outer model of V. Since the set U defined by $\dot{\mathcal{U}}$ in Σ is guaranteed to be infinite following our assumptions on $\mathfrak{T}$, the finitely many restrictions imposed by p on the relationship between (the function interpreting) $\dot{F}$ and U can be circumvented with ease. In other words, Σ can be extended to some $\Sigma^* \, \Gamma_{\mathfrak{T}^*}(\mathcal{L}_{\mathfrak{T}^*}, \mathfrak{A}_{\mathfrak{T}^*})$-certifying p. But this means every $q' \leq_{\mathbb{P}(\mathfrak{T})} q$ is compatible with p in $\mathbb{P}(\mathfrak{T}^*)$. $\qquad\square$

By Fact 2.40, Proposition 5.36, and the identity

$$\bigcup(g \cap P(\mathfrak{T})) = (\bigcup g) \cap \mathcal{L}_{\mathfrak{T}}$$

which holds for every $\mathbb{P}(\mathfrak{T}^*)$-generic filter g over V, we are done. $\qquad\square$

Models of a TCI $\mathfrak{T}$ across all outer models of V can be very complicated. However, when a model of $\mathfrak{T}$ is finitely determined, its atomic diagram can be easily read off $\mathbb{P}(\mathfrak{T})$.

Lemma 5.37. *Let $\mathfrak{T}$ be a TCI and $\mathcal{M}$ be a finitely determined model of $\mathfrak{T}$ in some outer model of V. Then for some atom p of $\mathbb{P}(\mathfrak{T})$, $\Sigma(\mathcal{M}) = g_p(\mathbb{P}(\mathfrak{T}))$. In particular, $\mathcal{M} \in V$.*

Proof. Let $\mathcal{M}$ be finitely determined by φ. Without loss of generality, we can assume φ is the conjunction of a set of literals $\{l_i : i < n\}$ for some $n < \omega$. This means

$$p := \{\ulcorner E(l_i) \urcorner : i < n\}$$

is an atom of $\mathbb{P}(\mathfrak{T})$. Lemma 2.35 tells us that $g_p(\mathbb{P}(\mathfrak{T}))$ is $\mathbb{P}(\mathfrak{T})$-generic over V, so necessarily $\Sigma(\mathcal{M}) = g_p(\mathbb{P}(\mathfrak{T}))$ by Theorem 5.32. Then according to Lemma 5.21, $\mathcal{M} \in V$ because $g_p(\mathbb{P}(\mathfrak{T})) \in V$. $\qquad\square$

It is possible to have an analogue of Lemma 5.37 for models that are "close to being finitely determined".

Definition 5.38. Let $\mathfrak{T}$ be a TCI. Inductively define $\Gamma_{\mathfrak{T}}^{(\alpha)}$, $P(\mathfrak{T})^{(\alpha)}$ and $\mathbb{P}(\mathfrak{T})^{(\alpha)}$ for all ordinals $\alpha \leq |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$ as follows:

$$\Gamma_{\mathfrak{T}}^{(0)} := \Gamma_{\mathfrak{T}},$$
$$P(\mathfrak{T})^{(0)} := P(\mathfrak{T}),$$
$$\Gamma_{\mathfrak{T}}^{(\alpha)} := \Gamma_{\mathfrak{T}}^{(\alpha-1)} \cup \{\ulcorner \bigvee_{x \in p}(\neg E(x)) \urcorner : p \text{ is an atom of } \mathbb{P}(\mathfrak{T})^{(\alpha-1)}\}$$

$$\text{if } \alpha \text{ is a successor ordinal,}$$
$$\Gamma_{\mathfrak{T}}^{(\alpha)} := \bigcup_{\beta < \alpha} \Gamma_{\mathfrak{T}}^{(\beta)}$$

$$\text{if } \alpha \text{ is a limit ordinal,}$$
$$P(\mathfrak{T})^{(\alpha)} := \{p \in [\mathcal{L}_{\mathfrak{T}}]^{<\omega} : \Vdash_{Col(\omega, |trcl(\mathfrak{A}_{\mathfrak{T}})|)} \exists \Sigma \ (\text{``}\Sigma \ \Gamma_{\mathfrak{T}}^{(\alpha)}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}})\text{-certifies } p\text{''})\},$$
$$\mathbb{P}(\mathfrak{T})^{(\alpha)} := (P(\mathfrak{T})^{(\alpha)}, \leq_{\mathbb{P}(\mathfrak{T})}).$$

By a simple cardinality argument, there must exist some $\alpha < |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$ for which $\Gamma_{\mathfrak{T}}^{(\alpha)} = \Gamma_{\mathfrak{T}}^{(\alpha+1)}$, whence $\mathbb{P}(\mathfrak{T})^{(\alpha)} = \mathbb{P}(\mathfrak{T})^{(\alpha+1)}$.

Definition 5.39. Let $\Gamma_{\mathfrak{T}}^{\top}$ denote the unique Γ such that $\Gamma = \Gamma_{\mathfrak{T}}^{(\alpha)} = \Gamma_{\mathfrak{T}}^{(\alpha+1)}$ for some $\alpha < |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$. Similarly, $\mathbb{P}(\mathfrak{T})^{\top}$ shall denote the unique $\mathbb{P}$ such that $\mathbb{P} = \mathbb{P}(\mathfrak{T})^{(\alpha)} = \mathbb{P}(\mathfrak{T})^{(\alpha+1)}$ for some $\alpha < |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$.

It is not hard to see that $P(\mathfrak{T})^{\top}$ is an atomless upward closed subset of $\mathbb{P}(\mathfrak{T})$ and $\Gamma_{\mathfrak{T}} \subset \Gamma_{\mathfrak{T}}^{\top}$.

Remark 5.40. In constructing the $\mathbb{P}(\mathfrak{T})^{(\alpha)}$'s, we are inductively removing atoms of $\mathbb{P}(\mathfrak{T})$. These atoms are representatives of isolated models of a TCI. By looking at Definition 5.38 in this way, we can draw obvious parallels between $\mathbb{P}(\mathfrak{T})^{(\alpha)}$ and the α-th-order Cantor-Bendixson derivative of a set. Such parallels culminate in $\mathbb{P}(\mathfrak{T})^{\top}$ being analogous to the "perfect core" of $\mathbb{P}(\mathfrak{T})$.

Definition 5.41. Given a TCI $\mathfrak{T}$ and any $\mathcal{M}$, we say $\mathcal{M}$ is an *almost finitely determined model of* $\mathfrak{T}$ iff $\mathcal{M} \models^* \mathfrak{T}$ and for some $\alpha < |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$ and an atom p of $\mathbb{P}(\mathfrak{T})^{(\alpha)}$,

$$p \subset \Sigma(\mathcal{M}).$$

We have as our next lemma, the promised analogue of Lemma 5.37.

Lemma 5.42. *Let $\mathfrak{T}$ be a TCI and $\mathcal{M}$ be an almost finitely determined model of $\mathfrak{T}$ in some outer model of V. Then for some $\alpha < |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$ and some atom p of $\mathbb{P}(\mathfrak{T})^{(\alpha)}$, $\Sigma(\mathcal{M}) = g_p(\mathbb{P}(\mathfrak{T})^{(\alpha)})$. In particular, $\mathcal{M} \in V$.*

Proof. Choose any model $\mathcal{M}$ of $\mathfrak{T}$ in an outer model of V. It suffices to prove by induction on $\alpha \leq |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$ that

$$\forall q \,\exists\beta \leq \alpha \,\exists p \,((q \text{ is an atom of } \mathbb{P}(\mathfrak{T})^{(\alpha)} \text{ and } q \subset \Sigma(\mathcal{M}))$$
$$\implies (p \text{ is an atom of } \mathbb{P}(\mathfrak{T})^{(\beta)} \text{ and } \Sigma(\mathcal{M}) = g_p(\mathbb{P}(\mathfrak{T})^{(\beta)}))).$$

The base case where $\alpha = 0$ is just Lemma 5.37. For the inductive case, assume $0 < \alpha \leq |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|^+$. and let q be an atom of $\mathbb{P}(\mathfrak{T})^{(\alpha)}$ with $q \subset \Sigma(\mathcal{M})$. Then by Lemma 2.35 and the definition of $\mathbb{P}(\mathfrak{T})^{(\alpha)}$, either $\Sigma(\mathcal{M}) = g_q(\mathbb{P}(\mathfrak{T})^{(\alpha)})$ or there is $\beta' < \alpha$ and an atom q' of $\mathbb{P}(\mathfrak{T})^{(\beta')}$ such that $q' \subset \Sigma(\mathcal{M})$. In the latter case, the inductive hypothesis gives us $\beta \leq \beta'$ and an atom p of $\mathbb{P}(\mathfrak{T})^{(\beta)}$ for which $\Sigma(\mathcal{M}) = g_p(\mathbb{P}(\mathfrak{T})^{(\beta)})$. Either way we are done. $\square$

The way $\mathbb{P}(\mathfrak{T})$ and $\mathbb{P}(\mathfrak{T})^{\top}$ are defined from a TCI $\mathfrak{T}$ allows us to establish a nice dichotomy on the $(\mathbb{P}(\mathfrak{T}), V)$-generic models of $\mathfrak{T}$ when $\mathfrak{T}$ is Π_2.

Lemma 5.43. *Let $\mathfrak{T}$ be a Π_2 TCI and $\mathcal{M}$ be a $(\mathbb{P}(\mathfrak{T}), V)$-generic model of $\mathfrak{T}$. Then one of the following must hold:*

(1) $\mathcal{M}$ is almost finitely determined.

(2) $\mathcal{M}$ *is a* $(\mathbb{P}(\mathfrak{T})^\top, V)$-*generic model of* $\mathfrak{T}$.

Proof. Let g be a $\mathbb{P}(\mathfrak{T})$-generic filter over V and assume $\mathcal{A} \cap g = \emptyset$, where

$$\mathcal{A} := \{p : \exists \alpha \ (\alpha < |[\mathcal{L}_\mathfrak{T}]^{<\omega}|^+ \text{ and } p \text{ is an atom of } \mathbb{P}(\mathfrak{T})^{(\alpha)})\}.$$

This latter assumption is equivalent to saying that the unique model $\mathcal{M}$ of $\mathfrak{T}$ for which $\bigcup g = \Sigma(\mathcal{M})$ is not almost finitely determined. By Theorem 5.32, it suffices to show that g is a $\mathbb{P}(\mathfrak{T})^\top$-generic filter over V. Clearly, $\bigcup g \ \Gamma_\mathfrak{T}^\top(\mathcal{L}_\mathfrak{T}, \mathfrak{A}_\mathfrak{T})$-certifies p, so $g \subset \mathbb{P}(\mathfrak{T})^\top$. That $\mathbb{P}(\mathfrak{T})^\top$ is a suborder of $\mathbb{P}(\mathfrak{T})$ means g is a filter on $\mathbb{P}(\mathfrak{T})^\top$.

To see g is $\mathbb{P}(\mathfrak{T})^\top$-generic over V, let E be predense in $\mathbb{P}(\mathfrak{T})^\top$. Note that if $p \in \mathbb{P}(\mathfrak{T})$ is incompatible in $\mathbb{P}(\mathfrak{T})$ with every member of $\mathcal{A}$, then $p \in \mathbb{P}(\mathfrak{T})^\top$. As such, $E \cup \mathcal{A}$ must be predense in $\mathbb{P}(\mathfrak{T})$. But this implies $E \cap g \neq \emptyset$ because g is $\mathbb{P}(\mathfrak{T})$-generic and $\mathcal{A} \cap g = \emptyset$. $\qquad\square$

The following is a stronger version of Theorem 5.32.

Theorem 5.44. *Let* $\mathfrak{T}$ *be a* Π_2 *TCI. If not all models of* $\mathfrak{T}$ *are almost finitely determined, then* $\mathbb{P}(\mathfrak{T})^\top$ *is non-empty and every* $\mathbb{P}(\mathfrak{T})^\top$-*generic filter over* $\mathfrak{A}_\mathfrak{T}$ *witnesses* $\mathcal{M}$ *is a* $(\mathbb{P}(\mathfrak{T})^\top, \mathfrak{A}_\mathfrak{T})$-*generic model of* $\mathfrak{T}$ *for some* $\mathcal{M}$.

Proof. Assume not all models of $\mathfrak{T}$ are almost finitely determined, and let $\mathcal{M}$ be a model of $\mathfrak{T}$ not almost finitely determined in some outer model of V. Then $\Sigma(\mathcal{M}) \ \Gamma_\mathfrak{T}^\top(\mathcal{L}_\mathfrak{T}, \mathfrak{A}_\mathfrak{T})$-certifies $\emptyset$, so $\mathbb{P}(\mathfrak{T})^\top$ is non-empty.

Check that the hypotheses of Lemma 3.35 are satisfied when we have

- $\mathfrak{A}_\mathfrak{T}$ in place of $\mathfrak{A}$,
- $\mathcal{L}_\mathfrak{T}$ in place of $\mathcal{L}$,
- $[\mathcal{L}_\mathfrak{T}]^{<\omega}$ in place of B,
- $\mathbb{P}(\mathfrak{T})^\top$ in place of $\mathbb{P}$, and
- $\Gamma_\mathfrak{T}^\top$ in place of Γ.

A direct application of said lemma, coupled with the knowledge that $\Gamma_\mathfrak{T} \subset \Gamma_\mathfrak{T}^\top$, completes the proof. $\qquad\square$

For a countable TCI $\mathfrak{T}$, the consistency of $\mathfrak{T}$ implies the existence of a model of $\mathfrak{T}$ in V.

Lemma 5.45. *Let* $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ *be a* Π_2 *TCI such that*

$$|\sigma \cup y| \leq \aleph_0$$

whenever $\vartheta(\dot{\mathcal{U}}) = (y, z)$ *for some* z. *If* $\mathfrak{T}$ *is consistent then* $\mathfrak{T}$ *has a model in* V.

Proof. Fix f_1 a bijection between y and $|y|$, and f_2 a form-preserving signature embedding from $\sigma \cup \{\dot{\mathcal{U}}\}$ into $H(\omega)$. Then f_1 and f_2 naturally induce

- a TCI $\mathfrak{T}' := (T', \sigma', \dot{\mathcal{U}}, \vartheta')$ in V such that $\mathfrak{T}'$ has a countable transitive closure, and
- a bijection between $\{\mathcal{M} : \mathcal{M} \models^* \mathfrak{T}\}$ and $\{\mathcal{M} : \mathcal{M} \models^* \mathfrak{T}'\}$ in every weak outer model of V.

As a consequence, we can assume $\mathfrak{T}$ has a countable transitive closure without loss of generality. By Lemma 2.9, $\mathcal{T}$ can be coded as a real. Besides, if $\mathfrak{T}$ has a model $\mathcal{M}$ in an outer model W of V, then M has a real code. By a routine check while unfurling the definition of $\models^*$ (see e.g. the proof of Lemma 3.34 for an argument of the satisfaction relation being Σ_1), we get that the statement

$$\exists \mathcal{M} \, (\mathcal{M} \models^* \mathfrak{T})$$

is equivalent to a Σ_1^1 sentence involving a real code of $\mathfrak{T}$ found in V, so it is absolute for V and any of its weak outer models. If $\mathfrak{T}$ is consistent, it must have a model in some outer model of V, whence it has a model in V. $\qquad\square$

Ideally, in the spirit of Lemmas 5.34 and 5.35, we want to prove a generic version of Lemma 5.45. This can be done through a relatively effective version of Theorem 5.44 for a certain class of countable TCIs, so as to kill two birds with one stone. Some definitions and facts are prerequisites.

For the rest of this subsection, fix a bijection

$$f^\dagger : \mathrm{Var} \cup \{x : x \text{ is a first-order logical symbol}\} \cup \{\ulcorner \in \urcorner\} \longrightarrow \{n < \omega : n \text{ is odd}\}.$$

Definition 5.46. For any countable set X, we say (r, f) *witnesses* $(X; \in)$ *is computable* iff

- f is a bijection from X into $\{n < \omega : n \text{ is even}\}$,
- $r = \{\langle (f \cup f^\dagger)^*(\varphi) \rangle : \varphi \text{ is a member of the } \Delta_0\text{-elementary diagram of } (X; \in)\}$,

  where

    - $\langle \cdot \rangle$ is the standard computable Gödel numbering of strings over ω, and
    - $(f \cup f^\dagger)^*$ is the canonical bijection from the set of finite strings over $dom(f \cup f^\dagger)$ into the set of finite strings over ω, induced by $f \cup f^\dagger$, and

- r is computable.

We say r is a *nicely computable code of* $(X; \in)$ iff there is f for which (r, f) witnesses $(X; \in)$ is computable.

Fact 5.47. There is a nicely computable code of $(H(\omega); \in)$.

Fact 5.48. If r is a nicely computable code of $(H(\omega); \in)$, then there is a unique f for which (r, f) witnesses $(H(\omega); \in)$ is computable.

Let $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ be a TCI, and y be such that $\vartheta(\dot{\mathcal{U}}) = (y, z)$ for some z. Assume $|\sigma \cup y| \leq \aleph_0$. Then we can find f_1 and f_2 such that

- f_1 is a bijection from y into $|y|$, and

- f_2 is a form-preserving signature embedding from $\sigma \cup \{\dot{\mathcal{U}}\}$ into $H(\omega)$.

Together, f_1 and f_2 naturally induce a TCI $\mathfrak{T}'$ with its associated $\mathcal{L}_{\mathfrak{T}'}$ being a subset of $H(\omega)$. Moreover, in every model $\mathfrak{A}$ of ZFC − Powerset containing $\{\mathfrak{T}, f_1, f_2\}$, f_1 and f_2 also induce a bijection $h^{\mathfrak{A}}$ from

$$\{\mathcal{M} : \mathcal{M} \models^* \mathfrak{T}\}$$

into

$$\{\mathcal{M}' : \mathcal{M}' \models^* \mathfrak{T}'\},$$

such that for all $\mathcal{M} \in dom(h^{\mathfrak{A}})$, $\mathcal{M} \cong h^{\mathfrak{A}}(\mathcal{M})$.

Hence, if we only care about models of $\mathfrak{T}$ up to isomorphism, we can without loss of generality, assume y is an ordinal at most ω and $\mathcal{L}_{\mathfrak{T}}$ is a subset of $H(\omega)$.

Definition 5.49. A TCI $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ is *code-friendly* iff

- $\vartheta(\dot{\mathcal{U}}) = (y, z) \in (\omega + 1) \times 2$, and

- $\mathcal{L}_{\mathfrak{T}} \subset H(\omega)$.

Code-friendly TCIs are relatively well-behaved and easy to reason about, especially when it comes to things like absoluteness. Notice that given any code-friendly TCI $\mathfrak{T} = (T, \sigma, \dot{\mathcal{U}}, \vartheta)$ and any ordinal α,

- $\mathcal{L}_{\mathfrak{T}}$ and $\Gamma_{\mathfrak{T}}$ are definable subsets of $H(\omega)$ over the structure

$$\mathfrak{A}_{\mathfrak{T}}^* := (H(\omega); \in, T, \sigma, \dot{\mathcal{U}}, \vartheta),$$

- $P(\mathfrak{T})^{(\alpha)} = \{p \in [\mathcal{L}_{\mathfrak{T}}]^{<\omega} : \exists \Sigma \, (\text{``}\Sigma \, \Gamma_{\mathfrak{T}}^{(\alpha)}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A}_{\mathfrak{T}}^*)\text{-certifies } p\text{''})\}$ by straightforward induction incorporating an argument similar to that which proved Lemma 3.34, and hence

- $P(\mathfrak{T})^{\top}$ is Δ_1-definable in $\mathfrak{A}_{\mathfrak{T}}^*$.

This means that the definition of $\mathbb{P}(\mathfrak{T})^{\top}$ from a code-friendly TCI $\mathfrak{T}$ is absolute for transitive models of ZFC − Powerset.

Recall Cohen forcing $\mathbb{C} = (C; \leq_{\mathbb{C}})$. We will use this labelling in the statements and proofs of the subsequent lemmas.

Lemma 5.50. *Let $\mathfrak{T}$ be a code-friendly Π_2 TCI, and (r, f) witness $(H(\omega); \in)$ is computable.*

Then either all models of $\mathfrak{T}$ are almost finitely determined, or else there is an oracle machine Ψ and a structure $\mathfrak{A}$ in the language of set theory, such that whenever g is a $\mathbb{C}$-generic filter over $\mathfrak{A}$, there is a unique model $\mathcal{M}_g$ of $\mathfrak{T}$ satisfying

$$\Psi^{(f"g) \oplus (f"(P(\mathfrak{T})^\top))} = f"(\Sigma(\mathcal{M}_g)).$$

Moreover, the function $g \mapsto \mathcal{M}_g$ defined as such is injective.

Proof. Assume not all models of $\mathfrak{T}$ are almost finitely determined. For brevity, let us write

$$s := f"(P(\mathfrak{T})^\top) \text{ and}$$
$$\leq_s := f"(\leq_{\mathbb{P}(\mathfrak{T})^\top}).$$

We shall identify s with $(s, \leq_s)$ whenever contextually necessary. This can be done without loss of generality because $\leq_s$ is computable in s.

Going forward, even beyond this proof, we would often argue about things in $H(\omega)$ even though our intended domain of discourse is the set of natural numbers. This is because first-order truths about $(H(\omega); \in)$ are uniformly propagated by f onto its range, so that specific versions of them hold there as well. If one such truth is sufficiently simple, then r knows the version of it on $ran(f)$ and can then relay that to the appropriate machines for further processing.

Let $\mathfrak{A} = (A; \in)$ be any countable elementary substructure of $\mathfrak{A}_\mathfrak{T}$ with $P(\mathfrak{T})^\top \in A$. Since $\mathcal{L}_\mathfrak{T}$ is just the closure of $\bigcup P(\mathfrak{T})^\top$ under negation and $\mathfrak{A} \models \mathsf{ZFC} - \mathsf{Powerset}$, we have $\mathcal{L}_\mathfrak{T} \in A$. That $\mathfrak{T}$ is code-friendly and $H(\omega) \subset A$ implies $\mathcal{L}_\mathfrak{T} \subset A$ too, so $\mathfrak{A}$ is $\mathcal{L}_\mathfrak{T}$-suitable. The members of $\Gamma_\mathfrak{T}$ are $(\mathcal{L}_\mathfrak{T})^*_{\mathfrak{A}_\mathfrak{T}}$-$\Pi_2$ sentences with a single parameter $P(\mathfrak{T})^\top$ and quantification exclusively over $H(\omega)$, entailing that $\Gamma_\mathfrak{T}$ is also a set of $(\mathcal{L}_\mathfrak{T})^*_\mathfrak{A}$-$\Pi_2$ sentences. Similar to what we did in the proof of Theorem 5.44, check that the hypotheses of Lemma 3.35 are satisfied with

- $\mathfrak{A}$ as defined,
- $\mathcal{L}_\mathfrak{T}$ in place of $\mathcal{L}$,
- $[\mathcal{L}_\mathfrak{T}]^{<\omega}$ in place of B,
- $\mathbb{P}(\mathfrak{T})^\top$ in place of $\mathbb{P}$, and
- $\Gamma_\mathfrak{T}^\top$ in place of Γ.

Following the proof of Theorem 5.44, while bearing in mind

- $\Gamma_\mathfrak{T} \subset \Gamma_\mathfrak{T}^\top$,

- Remark 5.22 and how its invocation is justified by the preceding paragraph, as well as

- the injectivity of the function $\mathcal{M} \mapsto \Sigma(\mathcal{M})$,

we apply Lemma 3.35 to give us

$$\forall \bar{g} \; \exists! \mathcal{M}_{\bar{g}} \; (\bar{g} \text{ is a } \mathbb{P}(\mathfrak{T})^\top\text{-generic filter over } \mathfrak{A}$$
$$\implies (\mathcal{M}_{\bar{g}} \models^* \mathfrak{T} \text{ and } \bigcup \bar{g} = \Sigma(\mathcal{M}_{\bar{g}}))). \tag{5.2}$$

Passing (5.2) through f leads us to the presence of an oracle machine $\bar{\Phi}$ fulfilling

$$\forall \bar{g} \; \exists! \mathcal{M}_{\bar{g}} \; (\bar{g} \text{ is a } s\text{-generic filter over } \mathfrak{A}$$
$$\implies (\mathcal{M}_{\bar{g}} \models^* \mathfrak{T} \text{ and } \bar{\Phi}^{\bar{g}} = f"(\Sigma(\mathcal{M}_{\bar{g}})))). \tag{5.3}$$

Next, note that

$$u := f" C \text{ and}$$
$$\leq_u := f"(\leq_\mathbb{C})$$

are computable subsets of ω. We shall, without loss of generality, identify u with $(u, \leq_u)$ whenever contextually necessary.

Proposition 5.51. *There is a dense embedding π of $\mathbb{C}$ into $\mathbb{P}(\mathfrak{T})^\top$ such that π is Δ_1-definable over the structure $(H(\omega); \in, P(\mathfrak{T})^\top)$.*

Proof. First, C is a Δ_0 subset of $H(\omega)$. That (r, f) witnesses $H(\omega)$ is computable means f and $g := f^{-1}$ are functions of which graphs are Δ_1-definable over

$$\mathfrak{B} := (H(\omega); \in, P(\mathfrak{T})^\top).$$

Also, $\mathcal{L}_\mathfrak{T}$ is Δ_1-definable over $\mathfrak{B}$ because

$$x \in \mathcal{L}_\mathfrak{T} \iff \{x\} \in P(\mathfrak{T})^\top \text{ or } \{\neg x\} \in P(\mathfrak{T})^\top.$$

Inductively define sequences $\{a_n : n < \omega\}$ and $\{k_n : n < \omega\}$ as follows:

$$a_0 := \emptyset,$$
$$k_n := min((f"\mathcal{L}_\mathfrak{T}) \setminus (f"a_n)), \text{ and}$$
$$a_{n+1} := a_n \cup \{g(k_n)\} \cup \{\neg(g(k_n))\}.$$

Note that $\{a_n\}$ and $\{k_n\}$ are Δ_1-definable over $\mathfrak{B}$. Next, let P^* be such that

$$x \in P^* \iff x \in P(\mathfrak{T})^\top \text{ and } x \subset a_{|x|},$$

and have x *split in P^** iff

$$x \in P^* \text{ and } \forall y \, (y \in a_{|x|+1} \setminus a_{|x|} \implies x \cup \{y\} \in P^*),$$

so that both P^* and the set of all its members that split in P^* are Δ_1-definable over $\mathfrak{B}$. We say x is a *P^*-least split above z* iff

$$z \subset x \text{ and } x \text{ splits in } P^* \text{ and } \forall y \, (z \subset y \subsetneq x \implies y \text{ does not split in } P^*).$$

Clearly, P^* is dense in $P(\mathfrak{T})^\top$, so $P(\mathfrak{T})^\top$ being atomless entails P^* is too. This yields the existence of a — necessarily unique — P^*-least split above z for every $z \in P^*$.

Finally, we can inductively define π on C as such:

$$\pi(\emptyset) := \text{the } P^*\text{-least split above } \emptyset,$$
$$\pi(x ^\frown \langle 0 \rangle) := \text{the } P^*\text{-least split above } \pi(x) \cup \{\varphi_{|x|,0}\}, \text{ and}$$
$$\pi(x ^\frown \langle 1 \rangle) := \text{the } P^*\text{-least split above } \pi(x) \cup \{\varphi_{|x|,1}\},$$

where

$$\varphi_{n,0} := \text{the unique member of } a_{n+1} \setminus a_n \text{ with leading symbol } \ulcorner \neg \urcorner, \text{ and}$$
$$\varphi_{n,1} := \text{the unique member of } a_{n+1} \setminus a_n \text{ with leading symbol not } \ulcorner \neg \urcorner.$$

It in not difficult to see that $ran(\pi)$ is dense in P^*, and thus in $P(\mathfrak{T})^\top$. Moreover, since for each $x \in C$, the definition of $\pi(x)$ depends only on the finite set

$$\{\pi(x \restriction_n) : n < |x|\}$$

and finitely many parameters which are Δ_1-definable over $\mathfrak{B}$, π must be Δ_1-definable over $\mathfrak{B}$ as well. $\qquad\square$

Proposition 5.51, via f, implies the existence of a dense embedding π of $\mathbb{C}$ into $\mathbb{P}(\mathfrak{T})^\top$ with $f"\pi$ computable in s, which is all we need to proceed. Fix any such π. It is not difficult to verify that $f \circ \pi = (f"\pi) \circ f$ on domain C and taking upward closure of a set in a forcing notion commutes with f. As such,

$$\forall g \, (g \text{ is a } \mathbb{C}\text{-generic filter over } \mathfrak{A} \implies (\mathrm{UC}(s, (f"\pi)"(f"g)) \text{ is a filter on } s \text{ and}$$
$$\Phi^{(f"g) \oplus s} = \mathrm{UC}(s, (f"\pi)"(f"g)))) \qquad (5.4)$$

for some oracle machine Φ.

Proposition 5.52. *Let g be a $\mathbb{C}$-generic subset over $\mathfrak{A}$. Then $\mathrm{UC}(\mathbb{P}(\mathfrak{T})^\top, \pi"g)$ is a $\mathbb{P}(\mathfrak{T})^\top$-generic subset over $\mathfrak{A}$.*

Proof. Let h denote $\mathrm{UC}(\mathbb{P}(\mathfrak{T})^\top, \pi"g)$, and D be a dense subset of $\mathbb{P}(\mathfrak{T})^\top$ definable in the language associated with $\mathfrak{A}$. Then the set

$$D' := \{p \in C : \exists q\ (q \in D \text{ and } \pi(p) \leq_{\mathbb{P}(\mathfrak{T})^\top} q)\}$$

is also definable in the language associated with $\mathfrak{A}$. Choose any $p_0 \in C$. By the density of D in $\mathbb{P}(\mathfrak{T})^\top$, there is $q_0 \in D$ such that $q_0 \leq_{\mathbb{P}(\mathfrak{T})^\top} \pi(p_0)$. That π is a dense embedding tells us there exists $q_1 \in ran(\pi)$ with $q_1 \leq_{\mathbb{P}(\mathfrak{T})^\top} q_0$. Now, for some $p_1 \leq_{\mathbb{C}} p_0$, $q_1 = \pi(p_1)$ and $p_1 \in D'$. We can therefore conclude that D' is dense in $\mathbb{C}$.

As g is $\mathbb{C}$-generic over $\mathfrak{A}$, we can find $p \in g \cap D'$. Seeing that $\pi(p) \in h$ and h is upward closed, we have by the definition of D', $h \cap D \neq \emptyset$. $\qquad\square$

Passing Proposition 5.52 through f strengthens (5.4) to

$$\forall g\ (g \text{ is a } \mathbb{C}\text{-generic filter over } \mathfrak{A} \implies (\mathrm{UC}(s, (f"\pi)"(f"g)) \text{ is a } s\text{-generic}$$
$$\text{filter over } \mathfrak{A} \text{ and}$$
$$\Phi^{(f"g)\oplus s} = \mathrm{UC}(s, (f"\pi)"(f"g)))). \quad (5.5)$$

Now (5.3) and (5.5) in conjunction tells us that we can combine $\bar{\Phi}$ and Φ into an oracle machine Ψ such that

$$\forall g\ \exists! \mathcal{M}_g\ (g \text{ is a } \mathbb{C}\text{-generic filter over } \mathfrak{A}$$
$$\implies (\mathcal{M}_g \models^* \mathfrak{T} \text{ and } \Psi^{(f"g)\oplus s} = f"(\Sigma(\mathcal{M}_g)))) \quad (5.6)$$

and

$$F_\Psi := ((g \text{ a } \mathbb{C}\text{-generic filter over } \mathfrak{A}) \mapsto \mathcal{M}_g \text{ as per } (5.6)) = F_{\bar{\Phi}} \circ F_\Phi,$$

where

$$F_{\bar{\Phi}} := (\bar{g} \text{ a } s\text{-generic filter over } \mathfrak{A}) \mapsto \mathcal{M}_{\bar{g}} \text{ as per } (5.3), \text{ and}$$
$$F_\Phi := (g \text{ a } \mathbb{C}\text{-generic filter over } \mathfrak{A}) \mapsto \mathrm{UC}(s, (f"\pi)"(f"g)) \text{ as per } (5.5)$$

are both injective. $\qquad\square$

Remark 5.53. Observe that we derived π in a uniform way from the parameters given in Proposition 5.51. Turning our attention to the proof of Lemma 5.50, said observation passes through f to imply F_Φ is derivable uniformly in $\mathbb{P}(\mathfrak{T})^\top$, and thus in $\mathfrak{T}$. Obviously, $F_{\bar{\Phi}}$ is derivable uniformly in $\mathfrak{T}$, so F_Ψ is too. As a result, this same Ψ works uniformly in $\mathfrak{T}$ to witness Lemma 5.50 for all $\mathfrak{T}$ and g as given in said lemma.

In addition to Remark 5.53, we can also strengthen Lemma 5.50 by lowering the requirement on the genericity of g and omitting $(A; \in)$ altogether. We formulate a strengthened version below in the nomenclature of computability theory.

Lemma 5.54. *Let (r, f) witness $(H(\omega); \in)$ is computable. Then there is an oracle machine Ψ such that whenever $\mathfrak{T}$ is a code-friendly Π_2 TCI, one of the following must hold:*

(1) All models of $\mathfrak{T}$ are almost finitely determined.

(2) For every $(f"(P(\mathfrak{T})^\top))$-1-generic real t, there is a unique model $\mathcal{M}_t$ of $\mathfrak{T}$ satisfying

$$\Psi^{t \oplus (f"(P(\mathfrak{T})^\top))} = f"(\Sigma(\mathcal{M}_t)).$$

Moreover for every $\mathfrak{T}$, the function $t \mapsto \mathcal{M}_t$ defined as such is injective.

Proof. Choose an arbitrary a code-friendly Π_2 TCI $\mathfrak{T}$ with not all models almost finitely determined. Adopt the abbreviations

$$s := f"(P(\mathfrak{T})^\top) \text{ and}$$
$$\leq_s := f"(\leq_{\mathbb{P}(\mathfrak{T})^\top}),$$

and identify s with $(s, \leq_s)$ whenever contextually necessary. We will modify the proof of Lemma 5.50 to get an oracle machine Ψ witnessing (2), before checking that a very slightly modified version of Remark 5.53 applies to Ψ.

First, retrace the relevant parts of the proof of Lemma 5.50 to arrive at (5.2).

By the same argument we used to conclude (3) of Remark 3.29 from the proof of Lemma 3.28, having a filter g on $\mathbb{P}(\mathfrak{T})^\top$ satisfy

- $\bigcup g \models_{\mathfrak{A}}^* \Gamma_{\mathfrak{T}}$ and
- $\bigcup g$ is $\mathcal{L}_{\mathfrak{T}}$-nice,

so that $\bigcup g \ \Gamma_{\mathfrak{T}}(\mathcal{L}_{\mathfrak{T}}, \mathfrak{A})$-certifies $\emptyset$, does not require g to meet all dense sets definable in the language associated with $\mathfrak{A}$. Indeed, it suffices for g to meet the dense sets "associated with" $(\mathcal{L}_{\mathfrak{T}})_{\mathfrak{A}}^*$-$\Sigma_1$ subformulas of sentences in $\Gamma_{\mathfrak{T}}$ — or rather, with $(\mathcal{L}_{\mathfrak{T}})_{\mathfrak{A}}^*$-$\Sigma_1$ sentences obtained by substituting constants (in the language associated with $\mathfrak{A}$) for variables in the aforementioned subformulas — sets defined by the types D_1 and D_2 in the proof of Lemma 3.28. One should keep in mind Remark 3.22 as one ascertains that — following the notation of the proof of Lemma 3.28 — every instance of ϕ'' appearing in some definition of a set of type D_2 is always logically equivalent to a $(\mathcal{L}_{\mathfrak{T}})_{\mathfrak{A}}^*$-$\Sigma_1$ sentence.

Per the definition of $\Gamma_{\mathfrak{T}}$ through the proof of Lemma 5.16, these are sets with Δ_2 definitions merely involving the predicates $\mathbb{P}(\mathfrak{T})^\top$ and $\mathcal{L}_{\mathfrak{T}}$ (the latter being a Δ_1-definable subset of $H(\omega)$ in the former over the structure $(H(\omega); \in)$), parameters in $H(\omega)$ and quantification over $H(\omega)$. Noting that

- $H(\omega)$ is closed under the function $\chi_{\mathfrak{T}}$ as defined in Remark 5.17, and
- the f-image of $\chi_{\mathfrak{T}}$ is Δ_0^0 (i.e. computable),

the f-image of each of said dense sets is thus $\Delta_2^{0,s}$. The upshot of all this is the justification that, by passing through f, every filter on s meeting all $\Delta_2^{0,s}$ dense subsets of s is appropriate for our purposes.

In fact, we can consider more filters on s.

Respectively, let D_1^* and D_2^* be dense sets of $\mathbb{P}(\mathfrak{T})^\top$ of the forms D_1 and D_2 as defined in the proof of Lemma 3.28. Then there are sets $D_{1,1}^*$, $D_{1,2}^*$, $D_{2,1}^*$, $D_{2,2}^*$ such that

(a) $D_1^* = D_{1,1}^* \cup D_{1,2}^*$ and $D_2^* = D_{2,1}^* \cup D_{2,2}^*$,

(b) for some p_1, $D_{1,1}^*$ is dense below p_1 in $\mathbb{P}(\mathfrak{T})^\top$ and $D_{1,2}^* = \{q \in \mathbb{P}(\mathfrak{T})^\top : q \perp_{\mathbb{P}(\mathfrak{T})^\top} p_1\}$,

(c) for some p_2, $D_{2,1}^*$ is dense below p_2 in $\mathbb{P}(\mathfrak{T})^\top$ and $D_{2,2}^* = \{q \in \mathbb{P}(\mathfrak{T})^\top : q \perp_{\mathbb{P}(\mathfrak{T})^\top} p_2\}$, and

(d) $D_{1,1}^*$ and $D_{2,1}^*$ are Σ_1-definable with a single parameter $\mathbb{P}(\mathfrak{T})^\top$ and quantification over $H(\omega)$.

Proposition 5.55.

(1) Every filter on $\mathbb{P}(\mathfrak{T})^\top$ that meets $D_{1,1}^*$ also meets D_1^*.

(2) Every filter on $\mathbb{P}(\mathfrak{T})^\top$ that meets $D_{2,1}^*$ also meets D_2^*.

Proof. We only prove (1), as the proof of (2) is the same up to labelling.

Let g be a filter on $\mathbb{P}(\mathfrak{T})^\top$ meeting $D_{1,1}^*$. Since $D_{1,1}^* \subset D_1^*$ by (a), we are done if $g \cap D_{1,1}^* \neq \emptyset$. Otherwise, let $p \in g$ be such that no $q \leq_{\mathbb{P}(\mathfrak{T})^\top} p$ is a member of $D_{1,1}^*$. In this case, according to (a) and (b), we are done if we can show $p \perp_{\mathbb{P}(\mathfrak{T})^\top} p_1$. To that end, assume the converse that there is a common extension p' of p and p_1 in $\mathbb{P}(\mathfrak{T})^\top$. By (b), $D_{1,1}^*$ is dense below p_1, so it is also dense below p'. Choose $q \leq_{\mathbb{P}(\mathfrak{T})^\top} p'$ with $q \in D_{1,1}^*$. But $q \leq_{\mathbb{P}(\mathfrak{T})^\top} p$, contradicting our choice of p. $\qquad\square$

Passing what we have hitherto shown through f the way (5.2) was in the proof of Lemma 5.50, for some oracle machine $\bar{\Phi}$ we have

$$\forall \bar{g} \; \exists! \mathcal{M}_{\bar{g}} \; (\bar{g} \text{ is a filter on } s \text{ meeting all } \Sigma_1^{0,s} \text{ subsets of } s$$
$$\implies (\mathcal{M}_{\bar{g}} \models^* \mathfrak{T} \text{ and } \bar{\Phi}^{\bar{g}} = f"(\Sigma(\mathcal{M}_{\bar{g}})))). \tag{5.7}$$

Adopt the abbreviations

$$u := f"C \text{ and}$$
$$\leq_u := f"(\leq_{\mathbb{C}}),$$

and identify u with $(u, \leq_u)$ whenever contextually necessary.

To deal with generic reals instead of generic filter, we first fix an oracle machine $\ddot{\Phi}$ that computes the set

$$f"\{t \restriction_n : n < \omega\}$$

when given any real t as oracle. Next, we follow an argument similar to the one used to derive (5.4) in the proof of Lemma 5.50, so that for some

- dense embedding π of $\mathbb{C}$ into $\mathbb{P}(\mathfrak{T})^\top$ with $f"\pi$ computable in s, and

- oracle machine Φ,

we have

$$\forall t \ \exists! c \ (t \text{ is a } s\text{-}1\text{-generic real} \implies (\mathrm{UC}(s, (f"\pi)"c) \text{ is a filter on } s \text{ and}$$
$$\ddot{\Phi}^t = c \text{ and } \Phi^{c \oplus s} = \mathrm{UC}(s, (f"\pi)"c))). \quad (5.8)$$

The upcoming proposition is an analogue of Proposition 5.52, formulated to restrict the universe of discourse to the (even) natural numbers.

Proposition 5.56. *Let c be a subset of u meeting all $\Sigma_1^{0,s}$ subsets of u. Then $\mathrm{UC}(s, (f"\pi)"c)$ meets all $\Sigma_1^{0,s}$ subsets of s.*

Proof. Let z be the subset of s defined by a $\Sigma_1^{0,s}$ formula $\varphi(x)$ in one free variable. Define

$$a_z := \{p \in u : \exists q \ (\varphi(q) \text{ and } (f"\pi)(p) \leq_s q)\}.$$

That u, $f"\pi$ and $\leq_s$ are all computable in s (u is even outright computable) gives us the $\Sigma_1^{0,s}$-definability of a_z as a subset of u. Consequently, c must meet a_z.

If $c \cap a_z \neq \emptyset$, then by the definition of a_z, there are conditions $p \in c$ and $q \in z$ for which $(f"\pi)(p) \leq_s q$, so $\mathrm{UC}(s, (f"\pi)"c)$ meets z. Otherwise, there is $p \in c$ that cannot be extended in u to a member of a_z. Consider any $q \leq_s (f"\pi)(p)$. As $f"\pi$ densely embeds u into s, we can find conditions $p' \in u$ and $q' \in s$ for which $q' \leq_s q$ and $(f"\pi)(p') = q'$. Now $p' \leq_u p$, which according to our choice of p, means $p' \notin a_z$. Unfurling the definition of a_z gives us $q \notin z$. Having thus shown that $(f"\pi)(p)$ cannot be extended in s to a member of z, we are done. $\qquad\square$

We can passing the definition of a s-1-generic real through f to conclude that for every such real t, $\ddot{\Phi}^t$ is a subset of u meeting all $\Sigma_1^{0,s}$ subsets of u. With (5.7) and (5.8) in mind, Proposition 5.56 then tells us that we can combine $\bar{\Phi}$, $\ddot{\Phi}$ and Φ to get an oracle machine Ψ fulfilling the requirements

$$\forall t \ \exists! \mathcal{M}_t \ (t \text{ is a } s\text{-}1\text{-generic} \implies (\mathcal{M}_t \models^* \mathfrak{T} \text{ and } \Psi^{t \oplus s} = f"(\Sigma(\mathcal{M}_t)))) \quad (5.9)$$

and

$$F_\Psi := ((t \text{ a } s\text{-}1\text{-generic real}) \mapsto \mathcal{M}_t \text{ as per (5.9)}) = F_{\bar{\Phi}} \circ F_\Phi \circ F_{\ddot{\Phi}},$$

where all of

$$F_{\bar{\Phi}} := (\bar{g} \text{ a filter on } s \text{ meeting all } \Sigma_1^{0,s} \text{ subsets of } s) \mapsto \mathcal{M}_{\bar{g}} \text{ as per (5.7)},$$
$$F_\Phi := (c \text{ a filter on } u \text{ meeting all } \Sigma_1^{0,s} \text{ subsets of } u)$$
$$\mapsto \mathrm{UC}(s, (f"\pi)"c) \text{ as per (5.8), and}$$
$$F_{\ddot{\Phi}} := (t \text{ a } s\text{-}1\text{-generic real}) \mapsto f"\{t{\upharpoonright}_n : n < \omega\}$$

are injective.

The argument in Remark 5.53 applies here to net us the uniformity of deriving $F_{\bar{\Phi}}$ and F_Φ from $\mathfrak{T}$, and clearly $F_{\ddot{\Phi}}$ does not depend on $\mathfrak{T}$ at all. As in Remark 5.53, we can then conclude that Ψ is the required witness to the lemma. $\qquad\square$

Fix any nicely computable code r of $H(\omega)$. Check that

- the function f given in Lemma 5.54 is a definable subset of $H(\omega)$ over the structure $(H(\omega); \in, r)$, in light of Fact 5.48,

- the oracle machine Ψ constructed in the proof of Lemma 5.54 is a definable element of $H(\omega)$ over the structure $(H(\omega); \in, r, f)$, and

- the injective function $t \mapsto f"\Sigma(\mathcal{M}_t)$ defined in the proof of Lemma 5.54 always has a left inverse computable using only $f_r"(P(\mathfrak{T})^\top)$ as parameter. Further, said left inverse is uniformly computable over all relevant code-friendly Π_2 TCIs $\mathfrak{T}$.

As such, we have actually proven a more general version of Lemma 5.54, which we formally present as our final theorem below.

Theorem 5.57. *There is a formula ψ_{gmc} in two free variables, absolute for transitive models of* $\mathsf{ZFC} - \mathsf{Powerset}$*, such that $\psi_{gmc}(r, (f_r, \Psi_r, \bar{\Psi}_r))$ defines a function*

$$r \mapsto (f_r, \Psi_r, \bar{\Psi}_r)$$

on the set of all nicely computable codes of $H(\omega)$, wherewith

(1) (r, f_r) witnesses $(H(\omega); \in)$ is computable,

(2) Ψ_r and $\bar{\Psi}_r$ are oracle machines, and

(3) whenever $\mathfrak{T}$ is a code-friendly Π_2 TCI, one of the following must hold:

(a) All models of $\mathfrak{T}$ are almost finitely determined.

(b) For every $(f_r"(P(\mathfrak{T})^\top))$-1-generic real t, there is a unique model $\mathcal{M}_t$ of $\mathfrak{T}$ satisfying

$$\Psi_r^{t \oplus (f_r"(P(\mathfrak{T})^\top))} = f_r"(\Sigma(\mathcal{M}_t)) \text{ and}$$
$$\bar{\Psi}_r^{(f_r"(\Sigma(\mathcal{M}_t))) \oplus (f_r"(P(\mathfrak{T})^\top))} = t.$$

In particular, for each such pair $(t, \mathcal{M}_t)$,

$$t \oplus (f_r"(P(\mathfrak{T})^\top)) \equiv_T (f_r"(\Sigma(\mathcal{M}_t))) \oplus (f_r"(P(\mathfrak{T})^\top)).$$

From Theorem 5.57, we can prove that certain T-substructures of a countable structure have the prefect set property.

Corollary 5.58. *Let $\mathfrak{A} = (A; \mathcal{I})$ and T be a countable structure and a first-order Π_2 theory respectively, over the same signature σ. Then $|\mathrm{Sub}(\mathfrak{A}, T)| \leq \aleph_0$ or $|\mathrm{Sub}(\mathfrak{A}, T)| = 2^{\aleph_0}$.*

Proof. Clearly $\mathrm{Sub}(\mathfrak{A}, T)$ is invariant under isomorphisms, so without loss of generality, we can assume A is some ordinal α with $\alpha \leq \omega$, and $dom(\mathcal{I}) \subset H(\omega) \setminus \omega$. By Example 5.5, there is a Π_2 TCI $\mathfrak{T}$ such that

$$\{\text{models of } \mathfrak{T}\} = \mathrm{Sub}(\mathfrak{A}, T).$$

Our assumptions on $\mathfrak{A}$ allow us to choose $\mathfrak{T}$ satisfying $\mathcal{L}_{\mathfrak{T}} \subset H(\omega)$, so that $\mathfrak{T}$ is also code-friendly. Note that $|\mathrm{Sub}(\mathfrak{A}, T)| \leq 2^{|A|} \leq 2^{\aleph_0}$.

If all models of $\mathfrak{T}$ are almost finitely generated, then Lemma 5.42 tells us that the number of models of $\mathfrak{T}$ is bounded above by

$$max\{|P(\mathfrak{T})|, |[\mathcal{L}_{\mathfrak{T}}]^{<\omega}|\} \leq \aleph_0,$$

which means $|\mathrm{Sub}(\mathfrak{A}, T)| \leq \aleph_0$. Otherwise, by Theorem 5.57, for some real X there is an injection from the set of X-1-generic reals into the set of models of $\mathfrak{T}$. As there are continuum many X-1-generic reals, $|\mathrm{Sub}(\mathfrak{A}, T)| = 2^{\aleph_0}$. $\qquad\square$

Corollary 5.58 also follows from a well-known fact in descriptive set theory (see e.g. [7]), via the Cantor-Bendixson theorem.

Fact 5.59. Let $\mathfrak{A} = (A; \mathcal{I})$ and T be a countable structure and a first-order theory respectively, over the same signature σ. Then $\mathrm{Sub}(\mathfrak{A}, T)$ is a closed set in some Polish space.

Precisely because Theorem 5.57 establishes a very strong version of perfect set property, it makes sense to treat it like a souped-up variant of the Cantor-Bendixson theorem, subject to restrictions on use cases. Instead of being applicable to arbitrary closed subsets of some Polish space, Theorem 5.57 only applies to sets which contain exactly all the models of some Π_2 TCI. This treatment comes off as a natural extension of the parallels we drew in Remark 5.40.

5.4   Open Questions

The study of how abundant generic models of a TCI are, can be approached from another direction: by comparing them with arbitrary models of the same TCI. As such, the propensity for a model of a TCI to be isomorphic to a generic model becomes of fundamental interest. In view of much of the work done in this section, the following is a most natural question.

(Q2) Is there a consistent Π_2 TCI $\mathfrak{T}$ such that every model of $\mathfrak{T}$ found in some outer model of V is isomorphic to a V-generic model of $\mathfrak{T}$?

Consider any consistent first-order Π_2 theory T with only finite models (there are many such theories with the empty signature). T can be used to define a Π_2 TCI $\mathfrak{T}$ such that every model of T is isomorphic to some model of $\mathfrak{T}$, and vice versa. This relation between T and $\mathfrak{T}$ remains true in all outer models of V. Hence, (Q2) can be answered in the affirmative.

Even if we require $\mathfrak{T}$ to have an infinite model in each of these questions, the same answers apply when we choose $\mathfrak{T}$ to be a Π_2 TCI such that all models of $\mathfrak{T}$ are isomorphic to the unique (up to isomorphism) $\aleph_0$-sized model of a $\aleph_0$-categorical first-order Π_2 theory T (say, the theory of dense linear orders without endpoints). As the $\aleph_0$-categoricity of a theory is absolute for transitive models of ZFC with the same ordinals, said relation between T and $\mathfrak{T}$ is preserved across outer models of V.

However, the question dual to (Q2) appears much more difficult.

Question 5.60. *Is there a Π_2 TCI $\mathfrak{T}$ with a model $\mathcal{M}$ in some outer model of V such that $\mathcal{M}$ is not isomorphic to any V-generic model of $\mathfrak{T}$?*

Since our impetus for studying TCIs stems from our interest in uncovering links between forcing/genericity and the semantics of first-order logic, it is perhaps fitting that we ask for a similar example of a first-order theory.

Question 5.61. *Is there a first-order Π_2 theory T with a model M in some outer model of V such that for no Π_2 TCI $\mathfrak{T}$ is M isomorphic to a V-generic model of $\mathfrak{T}$?*

The technique of forcing is apparently the only way we know how to extend an arbitrary model of set theory. By Lemma 5.28, Questions 5.60 and 5.61 would be intractable using current technology if they were to be answered in the affirmative. We can modify them into possibly easier questions by only asking for consistency — instead of outright existence — over the base set theory.

Question 5.62.

(1) *Is it consistent that there exists a Π_2 TCI $\mathfrak{T}$ with a model $\mathcal{M}$ in some outer model of V such that $\mathcal{M}$ is not isomorphic to any V-generic model of $\mathfrak{T}$?*

(2) *Is is consistent that there exists a first-order Π_2 theory T with a model M in some outer model of V such that for no Π_2 TCI $\mathfrak{T}$ is M isomorphic to a V-generic model of $\mathfrak{T}$?*

Our line of questioning can be extended to the paradigm of relative effectiveness.

Fix ψ_{gmc} to be as in Theorem 5.57. Let r be a nicely computable code of $H(\omega)$. Define

$f_r :=$ the unique f for which there are Ψ and $\bar{\Psi}$ satisfying $\psi_{gmc}(r, (f, \Psi, \bar{\Psi}))$,

$\Psi_r :=$ the unique Ψ for which there are f and $\bar{\Psi}$ satisfying $\psi_{gmc}(r, (f, \Psi, \bar{\Psi}))$, and

$\bar{\Psi}_r :=$ the unique $\bar{\Psi}$ for which there are f and Ψ satisfying $\psi_{gmc}(r, (f, \Psi, \bar{\Psi}))$.

Analogous to Question 5.60, we want to know if there is a code-friendly TCI $\mathfrak{T}$ with a model M that neither almost finitely determined nor isomorphic to any $\mathcal{M}_t$ born from a $(f_r"(P(\mathfrak{T})^\top))$-1-generic real t à la Theorem 5.57. As turns out, because $f_r \circ f_s^{-1}$ is computable for any other nicely computable code s of $H(\omega)$, the answer to this question is independent of the choice of r.

Question 5.63. *Is there a code-friendly Π_2 TCI $\mathfrak{T}$ with a model $\mathcal{M}$ such that*

(a) $\mathcal{M}$ is not almost finitely determined, and

(b) for every nicely computable code r of $H(\omega)$ and every $(f_r"(P(\mathfrak{T})^\top))$-1-generic real t, if $\mathcal{M}'$ is a model of $\mathfrak{T}$ satisfying

$$\Psi_r^{t \oplus (f_r"(P(\mathfrak{T})^\top))} = f_r"(\Sigma(\mathcal{M}')),$$

then $\mathcal{M} \not\cong \mathcal{M}'$?

We can modify Question 5.63 to get an analogue of Question 5.61 in the same spirit of relative effectiveness.

Question 5.64. *Is there a first-order Π_2 theory T with a model M such that whenever $\mathfrak{T}$ is a code-friendly Π_2 TCI and $\mathcal{M} \models^* \mathfrak{T}$,*

(a) if $\mathcal{M}$ is almost finitely determined then $M \not\cong \mathcal{M}$, and

(b) if r is a nicely computable code of $H(\omega)$ and t is a $(f_r"(P(\mathfrak{T})^\top))$-1-generic real for which

$$\Psi_r^{t \oplus (f_r"(P(\mathfrak{T})^\top))} = f_r"(\Sigma(\mathcal{M})),$$

then $M \not\cong \mathcal{M}$?

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