```
S(n) = 2^{i-1} + S(n-F_i) - S(F_{i+1} - n - 3) with: F_i \le n < F_{i+1}
Let L(n) = n - F_i and R(n) = F_{i+1} - n - 3
The recursion is : (1) S(n) = 2^{i-1} + S(L(n)) - S(R(n))
If we represent n by its binary Zeckendorf decomposition the function L
is equivalent to remove the hightest bit
Ex: 74 = F_9 + F_6 + F_4 + F_1 = 100101001b
L(100101001b) = 1010001b = F_6 + F_4 + F_1 = 19 = 74 - 55
R(100101001b) = 89 - 74 - 3 = 12 = F_5 + F_3 + F_1 = 10101b
Let E_L = \{i \ge 0; L^{(i)}(n)\} and E_R = \{i \ge 0; L^{(i)}(R(n))\}
E_L = \{100101001b, 101001b, 1001b, 1b\}; E_R = \{10101b, 101b, 1b\}
It can be proved that R(E_L) \subset E_R and R(E_R) \subset E_L
if F_i \le n < F_{i+1} and F_j \le n - F_i < F_{j+1}
we have j \leq i - 2 (no consecutive Fibonacci numbers)
R(n) = F_{i+1} - n - 3 and R(L(n)) = F_{j+1} - (n - F_i) - 3
so R(n) = R(L(n)) + F_{i+1} - F_i - F_{j+1} = R(L(n)) + F_{i-1} - F_{j+1}
the Zeckendorf decomposition of F_{i-1} - F_{j+1} uses Fibonacci with index \geq j
the Zeckendorf decomposition of R(L(n)) uses Fibonacci with index \leq j-2
R(L(n)) is a low part of Zeckendorf decomposition of R(n) \Rightarrow R(L(n)) \in E_R
the method is simple: we compute E_{LR} = E_L \bigcup E_R
as foreach element k \in E_{LR} L(k), R(k) \in E_{LR}
we can compute S(k) using (1) for k in ascending order
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