

$$\begin{aligned}
S(n) &= \sum_{\substack{i,j,k,l \\ 0 < i < j < k < l \leq n}} (a_i + a_j + a_k + a_l) \times 1_{a_i < a_j < a_k < a_l} = \sum_{\substack{i,j,k,l \\ 0 < i < j < k < l \leq n}} ((a_i + a_j) + (a_k + a_l)) \times 1_{a_i < a_j < a_k} \times 1_{a_k < a_l} \\
&= \sum_{\substack{i,j,k,l \\ 0 < i < j < k < l \leq n}} ((a_i + a_j) \times 1_{a_i < a_j < a_k} \times 1_{a_k < a_l} + (a_k + a_l) \times 1_{a_i < a_j < a_k} \times 1_{a_k < a_l}) \\
&= \sum_{0 < k \leq n} \left\{ \sum_{\substack{i,j \\ 0 < i < j < k}} ((a_i + a_j) \times 1_{a_i < a_j < a_k}) \times \left( \sum_{\substack{l \\ k < l \leq n}} 1_{a_k < a_l} \right) \right\} + \sum_{0 < j \leq n} \left\{ \sum_{\substack{i \\ 0 < i < j}} 1_{a_i < a_j} \times \left( \sum_{\substack{k,l \\ j < k < l \leq n}} ((a_k + a_l) \times 1_{a_k < a_l}) \right) \right\} \\
&= \sum_{0 < k \leq n} S_{L2}(k) \times W_H(k) + \sum_{0 < j \leq n} S_{H2}(j) \times W_L(j)
\end{aligned}$$

Where:

$$\begin{aligned}
W_L(j) &= \sum_{\substack{i \\ 0 < i < j}} 1_{a_i < a_j} \quad (resp.) \quad W_H(k) = \sum_{\substack{l \\ k < l \leq n}} 1_{a_k < a_l} \\
S_{L2}(k) &= \sum_{\substack{i,j \\ 0 < i < j < k}} ((a_i + a_j) \times 1_{a_i < a_j < a_k}) \quad (resp.) \quad S_{H2}(j) = \sum_{\substack{k,l \\ j < k < l \leq n}} ((a_k + a_l) \times 1_{a_k < a_l}) \\
S_{L2}(k) &\text{ can be computed by :} \\
S_{L2}(k) &= \sum_{\substack{j \\ 0 < j < k}} 1_{a_j < a_k} (S_{L1}(j) + a_j \times W_L(j)) \quad \text{with: } S_{L1}(j) = \sum_{\substack{i \\ 0 < i < j}} a_i \times 1_{a_i < a_j}
\end{aligned}$$

We can compute by Fenwick tree for any function  $f(i)$   $F(i \rightarrow f(i), j) = \sum_{\substack{i \\ 0 < i < j}} f(i) \times 1_{a_i < a_j}$

$$(resp) \quad F^T(i \rightarrow f(i), j) = \sum_{\substack{j \\ j < i \leq n}} f(i) \times 1_{a_j < a_i}$$

So the algorithm is:

$$\begin{aligned}
W_L(j) &= F(i \rightarrow 1, j) \\
S_{L1}(j) &= F(i \rightarrow a_i, j) ; \quad S_{L1*}(j) = S_{L1}(j) + a_j W_L(j) ; \quad S_{L2}(j) = F(i \rightarrow S_{L1*}(j), j) \\
W_H(j) &= F^T(i \rightarrow 1, j) \\
S_{H1}(j) &= F^T(i \rightarrow a_i, j) ; \quad S_{H1*}(j) = S_{H1}(j) + a_j W_H(j) ; \quad S_{H2}(j) = F^T(i \rightarrow S_{H1*}(j), j) \\
S(n) &= \sum_{0 < j \leq n} S_{L2}(j) W_H(j) + S_{H2}(j) W_L(j)
\end{aligned}$$

$$\begin{aligned}
S_{L2}(k) &= \sum_{0 < i < \overset{i,j}{j} < k} ((a_i + a_j) \times 1_{a_i < a_j} \times 1_{a_j < a_k}) = \sum_{0 < \overset{j}{j} < k} 1_{a_j < a_k} \times \sum_{0 < \overset{i}{i} < j} ((a_i + a_j) \times 1_{a_i < a_j}) \\
&= \sum_{0 < \overset{j}{j} < k} 1_{a_j < a_k} \times S_{L1*}(j) \\
\text{with: } S_{L1*}(j) &= \sum_{0 < \overset{i}{i} < j} (a_i + a_j) \times 1_{a_i < a_j} = \sum_{0 < \overset{i}{i} < j} ((a_i \times 1_{a_i < a_j}) + a_j \times \sum_{0 < \overset{i}{i} < j} 1_{a_i < a_j}) \\
&= S_{L1}(j) + a_j \times W_L(j) \text{ with: } S_{L1}(j) = \sum_{0 < \overset{i}{i} < j} a_i \times 1_{a_i < a_j}
\end{aligned}$$

*Remark : to implement the division by 2, we compute  $R_{\Delta-1}$*