## Generalized Turán Problems for Trees and More

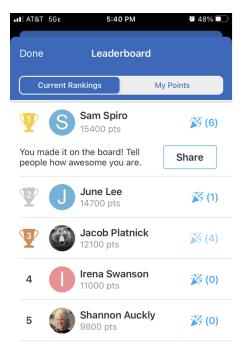
Sam Spiro, Georgia State University

Based on Joint Work with Sean English



Feedback can be given at binatorics.com







The Turán number ex(n, F) is defined to be the maximum number of edges that an F-free graph on n vertices can have.

The Turán number ex(n, F) is defined to be the maximum number of edges that an F-free graph on n vertices can have.

#### Theorem

$$ex(n, K_2) = 0.$$

# Theorem (Turán 1941)

$$\operatorname{ex}(n,K_r) = \left\lfloor \binom{r-1}{2} \frac{n^2}{(r-1)^2} \right\rfloor$$

## Theorem (Turán 1941)

$$\operatorname{ex}(n,K_r) = \left\lfloor \binom{r-1}{2} \frac{n^2}{(r-1)^2} \right\rfloor$$

## Theorem (Kővari-Sós-Turán 1954)

$$ex(n, K_{s,t}) = O(n^{2-1/s}).$$

## Theorem (Turán 1941)

$$\operatorname{ex}(n,K_r) = \left\lfloor \binom{r-1}{2} \frac{n^2}{(r-1)^2} \right\rfloor$$

## Theorem (Kővari-Sós-Turán 1954)

$$ex(n, K_{s,t}) = O(n^{2-1/s}).$$

## Theorem (Folklore)

If T is a tree with  $e(T) \ge 2$ , then

$$ex(n, T) = \Theta(n)$$
.

Given a family of graphs  $\mathcal{F}$ , we say G is  $\mathcal{F}$ -free if G is F-free for all  $F \in \mathcal{F}$ . We define  $\mathrm{ex}(n,\mathcal{F})$  to be the maximum number of edges in an n-vertex  $\mathcal{F}$ -free graph.

Given a family of graphs  $\mathcal{F}$ , we say G is  $\mathcal{F}$ -free if G is F-free for all  $F \in \mathcal{F}$ . We define  $ex(n, \mathcal{F})$  to be the maximum number of edges in an n-vertex  $\mathcal{F}$ -free graph.

# Conjecture (Compactness Conjecture)

If  $\mathcal{F}$  is a finite family of graphs, then  $ex(n, \mathcal{F}) = \Theta(ex(n, F))$  for some  $F \in \mathcal{F}$ .

## Definition (Alon-Shikhelman 2016)

Given a graph H and a family of graphs  $\mathcal{F}$ , we define the generalized Turán number  $\mathrm{ex}(n,H,\mathcal{F})$  to be the maximum number of copies of H in an n-vertex  $\mathcal{F}$ -free graph.

## Definition (Alon-Shikhelman 2016)

Given a graph H and a family of graphs  $\mathcal{F}$ , we define the generalized Turán number  $\mathrm{ex}(n,H,\mathcal{F})$  to be the maximum number of copies of H in an n-vertex  $\mathcal{F}$ -free graph.

## **Proposition**

$$ex(n, K_2, \mathcal{F}) = ex(n, \mathcal{F}).$$

## Definition (Alon-Shikhelman 2016)

Given a graph H and a family of graphs  $\mathcal{F}$ , we define the generalized Turán number  $\mathrm{ex}(n,H,\mathcal{F})$  to be the maximum number of copies of H in an n-vertex  $\mathcal{F}$ -free graph.

### Proposition

$$ex(n, K_2, \mathcal{F}) = ex(n, \mathcal{F}).$$

See the recent survey by Gerbner and Palmer for way more history than I'm going to give here.

# Theorem (Gerbner-Palmer 2019)

For all  $r \geq 2$  and families  $\mathcal{F}$  we have

$$\operatorname{ex}(n, K_r, \mathcal{F}) \leq \operatorname{ex}(n, \mathcal{F})^{r/2}.$$

## Theorem (Gerbner-Palmer 2019)

For all  $r \geq 2$  and families  $\mathcal{F}$  we have

$$\operatorname{ex}(n, K_r, \mathcal{F}) \leq \operatorname{ex}(n, \mathcal{F})^{r/2}.$$

## Theorem (Gerbner 2023)

Determines  $ex(n, K_{2,t}, F)$  asymptotically for any non-bipartite graph F.

## Theorem (Gerbner-Palmer 2019)

For all  $r \geq 2$  and families  $\mathcal{F}$  we have

$$\operatorname{ex}(n, K_r, \mathcal{F}) \leq \operatorname{ex}(n, \mathcal{F})^{r/2}.$$

## Theorem (Gerbner 2023)

Determines  $ex(n, K_{2,t}, F)$  asymptotically for any non-bipartite graph F.

## Theorem (Füredi-Kündgen 2006)

If  $\operatorname{ex}(n,\mathcal{F}) = \Theta(n^{2-\beta})$  and  $\mathcal{F}$  does not contain a star, then

$$ex(n, K_{1,t}, \mathcal{F}) = \tilde{\Theta}(max\{n^t, n^{t+1-t\beta}\}).$$

#### Question

Can we prove bounds on ex(n, T, F) for arbitrary trees T (for some possibly fixed F)?

#### Question

Can we prove bounds on ex(n, T, F) for arbitrary trees T (for some possibly fixed F)?

## Theorem (Gerbner 2023)

Determines the order of magnitude of  $ex(n, T, K_{2,t})$  for every tree T.

#### Question

Can we prove bounds on ex(n, T, F) for arbitrary trees T (for some possibly fixed F)?

## Theorem (Gerbner 2023)

Determines the order of magnitude of  $ex(n, T, K_{2,t})$  for every tree T.

## Theorem (Alon-Shikhelman 2016)

If T is a tree and F is a tree, then  $ex(n, T, F) = \Theta(n^k)$  for some integer k.

#### Question

Can we prove bounds on ex(n, T, F) for arbitrary trees T (for some possibly fixed F)?

## Theorem (Gerbner 2023)

Determines the order of magnitude of  $ex(n, T, K_{2,t})$  for every tree T.

## Theorem (Alon-Shikhelman 2016)

If T is a tree and F is a tree, then  $ex(n, T, F) = \Theta(n^k)$  for some integer k.

## Theorem (Letzter 2019)

If H is any graph and F is a tree, then  $ex(n, H, F) = \Theta(n^k)$  for some integer k.

#### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, T, F) = O(ex(n, F)^{k-1}).$$

Moreover, we can determine which of these two cases happen for a given  $T, k, \mathcal{F}$ .

### Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

## Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

## Corollary

If T is a tree and  $\mathcal{F}$  contains a forest, then  $ex(n, T, \mathcal{F}) = \Theta(n^k)$  for some integer k.

## Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

# Corollary

If T is a tree and  $\mathcal{F}$  contains a forest, then  $ex(n, T, \mathcal{F}) = \Theta(n^k)$  for some integer k.

#### Proof.

Let k be the smallest integer such that  $ex(n, T, \mathcal{F}) = O(n^k)$ .

## Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

# Corollary

If T is a tree and  $\mathcal{F}$  contains a forest, then  $ex(n, T, \mathcal{F}) = \Theta(n^k)$  for some integer k.

#### Proof.

Let k be the smallest integer such that  $ex(n, T, \mathcal{F}) = O(n^k)$ . If  $ex(n, T, \mathcal{F}) = \Omega(n^k)$  then we're done

## Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

## Corollary

If T is a tree and  $\mathcal{F}$  contains a forest, then  $ex(n, T, \mathcal{F}) = \Theta(n^k)$  for some integer k.

#### Proof.

Let k be the smallest integer such that  $ex(n, T, \mathcal{F}) = O(n^k)$ . If  $ex(n, T, \mathcal{F}) = \Omega(n^k)$  then we're done, otherwise our theorem implies that

$$ex(n, T, F) = O(ex(n, F)^{k-1})$$

## Theorem (English-S. 2025+)

Either  $ex(n, T, F) = \Omega(n^k)$  or  $ex(n, T, F) = O(ex(n, F)^{k-1})$ .

# Corollary

If T is a tree and  $\mathcal{F}$  contains a forest, then  $ex(n, T, \mathcal{F}) = \Theta(n^k)$  for some integer k.

#### Proof.

Let k be the smallest integer such that  $\mathrm{ex}(n,T,\mathcal{F})=O(n^k)$ . If  $\mathrm{ex}(n,T,\mathcal{F})=\Omega(n^k)$  then we're done, otherwise our theorem implies that

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}) = O(n^{k-1}).$$



#### Corollary

If  $T \neq K_2$  is a tree with  $\ell \geq 2$  leaves, then any family  $\mathcal{F}$  with  $\operatorname{ex}(n,T,\mathcal{F}) = O(n^\ell)$  has  $\operatorname{ex}(n,T,\mathcal{F}) = \Theta(n^k)$  for some integer k.

In particular, if  $ex(n, T, \mathcal{F}) = o(n^{\ell})$  then  $ex(n, T, \mathcal{F}) = O(n^{\ell-1})$ .

### Corollary

If  $T \neq K_2$  is a tree with  $\ell \geq 2$  leaves, then any family  $\mathcal F$  with  $\operatorname{ex}(n,T,\mathcal F) = O(n^\ell)$  has  $\operatorname{ex}(n,T,\mathcal F) = \Theta(n^k)$  for some integer k.

In particular, if  $ex(n, T, F) = o(n^{\ell})$  then  $ex(n, T, F) = O(n^{\ell-1})$ .

#### Proof.

If  ${\mathcal F}$  contains a forest then we are done by the previous result.

## Corollary

If  $T \neq K_2$  is a tree with  $\ell \geq 2$  leaves, then any family  $\mathcal{F}$  with  $\operatorname{ex}(n,T,\mathcal{F}) = O(n^\ell)$  has  $\operatorname{ex}(n,T,\mathcal{F}) = \Theta(n^k)$  for some integer k.

In particular, if  $ex(n, T, F) = o(n^{\ell})$  then  $ex(n, T, F) = O(n^{\ell-1})$ .

#### Proof.

If  $\mathcal F$  contains a forest then we are done by the previous result. Otherwise, the graph G obtained by duplicating each leaf of T a total of  $n/\ell$  times is  $\mathcal F$ -free

## Corollary

If  $T \neq K_2$  is a tree with  $\ell \geq 2$  leaves, then any family  $\mathcal F$  with  $\operatorname{ex}(n,T,\mathcal F) = O(n^\ell)$  has  $\operatorname{ex}(n,T,\mathcal F) = \Theta(n^k)$  for some integer k.

In particular, if  $ex(n, T, F) = o(n^{\ell})$  then  $ex(n, T, F) = O(n^{\ell-1})$ .

#### Proof.

If  $\mathcal F$  contains a forest then we are done by the previous result. Otherwise, the graph G obtained by duplicating each leaf of T a total of  $n/\ell$  times is  $\mathcal F$ -free and contains  $\Omega(n^\ell)$  copies of T.

# Theorem (English-S. 2025++)

If T is a tree with  $\ell \geq 2$  leaves, then every family  $\mathcal F$  with  $\operatorname{ex}(n,T,\mathcal F) = o(n^{\ell+1})$  satisfies

$$\operatorname{ex}(n,T,\mathcal{F}) = O\left(n^{\ell + \frac{\ell^2 - \ell}{e(T) - 1}}\right).$$

# Theorem (English-S. 2025++)

If T is a tree with  $\ell \geq 2$  leaves, then every family  $\mathcal F$  with  $\operatorname{ex}(n,T,\mathcal F)=o(n^{\ell+1})$  satisfies

$$\operatorname{ex}(n,T,\mathcal{F}) = O\left(n^{\ell + \frac{\ell^2 - \ell}{e(T) - 1}}\right).$$

Moreover, this bound is best possible for essentially all values of  $\ell$ , e(T).

# Theorem (English-S. 2025++)

If T is a tree with  $\ell \geq 2$  leaves, then every family  $\mathcal F$  with  $\operatorname{ex}(n,T,\mathcal F)=o(n^{\ell+1})$  satisfies

$$\operatorname{ex}(n,T,\mathcal{F}) = O\left(n^{\ell + \frac{\ell^2 - \ell}{e(T) - 1}}\right).$$

Moreover, this bound is best possible for essentially all values of  $\ell$ , e(T).

#### Theorem (English-S. 20??++)

For the path graph  $P_t$ , every graph F with  $ex(n, P_t, F) = O(n^{\alpha(P_t)})$  has  $ex(n, P_t, F) = \Theta(n^k)$  for some integer k.

# Stability for Generalized Turán Problems

# Stability for Generalized Turán Problems

# Proposition (English-Halfpap-Krueger 2024)

For stars  $K_{1,t}$ , every family of graphs  $\mathcal{F}$  either satisfies  $ex(n, K_{1,t}, \mathcal{F}) = \Omega(n^t)$  or  $ex(n, K_{1,t}, \mathcal{F}) = O(n)$ .

# Proposition (English-Halfpap-Krueger 2024)

For stars  $K_{1,t}$ , every family of graphs  $\mathcal{F}$  either satisfies  $ex(n, K_{1,t}, \mathcal{F}) = \Omega(n^t)$  or  $ex(n, K_{1,t}, \mathcal{F}) = O(n)$ .

### Proof.

If  $\mathcal F$  does not contain a subgraph of a star, then  $G=K_{1,n-1}$  is  $\mathcal F$ -free and shows  $\operatorname{ex}(n,K_{1,t},\mathcal F)=\Omega(n^t)$ .

## Proposition (English-Halfpap-Krueger 2024)

For stars  $K_{1,t}$ , every family of graphs  $\mathcal{F}$  either satisfies  $ex(n, K_{1,t}, \mathcal{F}) = \Omega(n^t)$  or  $ex(n, K_{1,t}, \mathcal{F}) = O(n)$ .

#### Proof.

If  $\mathcal{F}$  does not contain a subgraph of a star, then  $G=K_{1,n-1}$  is  $\mathcal{F}$ -free and shows  $\operatorname{ex}(n,K_{1,t},\mathcal{F})=\Omega(n^t)$ . Otherwise,  $\mathcal{F}$  must contain some subgraph of some star  $K_{1,r}$ , which means  $\operatorname{ex}(n,K_{1,t},\mathcal{F})=O_r(n)$ .

### Proposition (English-Halfpap-Krueger 2024)

For stars  $K_{1,t}$ , every family of graphs  $\mathcal{F}$  either satisfies  $ex(n, K_{1,t}, \mathcal{F}) = \Omega(n^t)$  or  $ex(n, K_{1,t}, \mathcal{F}) = O(n)$ .

#### Proof.

If  $\mathcal F$  does not contain a subgraph of a star, then  $G=K_{1,n-1}$  is  $\mathcal F$ -free and shows  $\operatorname{ex}(n,K_{1,t},\mathcal F)=\Omega(n^t)$ . Otherwise,  $\mathcal F$  must contain some subgraph of some star  $K_{1,r}$ , which means  $\operatorname{ex}(n,K_{1,t},\mathcal F)=O_r(n)$ .

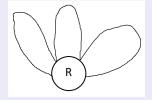
We want to generalize this idea by saying that for any graph H, there exists some "simple" family  $\mathcal{F}_H$  such that the behavior of  $\mathrm{ex}(n,K_{1,t},\mathcal{F})$  depends on how  $\mathcal{F}$  "interacts" with  $\mathcal{F}_H$ .

#### **Definition**

Given a graph H, a subset  $R \subseteq V(H)$ , and an integer q, we define the sunflower-power  $H_R^q$  to be the graph obtained by taking q copies of H which all agree on R and which are otherwise disjoint.

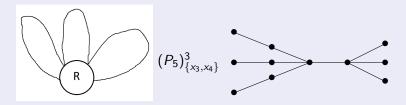
#### **Definition**

Given a graph H, a subset  $R \subseteq V(H)$ , and an integer q, we define the sunflower-power  $H_R^q$  to be the graph obtained by taking q copies of H which all agree on R and which are otherwise disjoint.



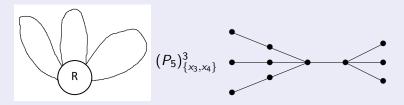
#### **Definition**

Given a graph H, a subset  $R \subseteq V(H)$ , and an integer q, we define the sunflower-power  $H_R^q$  to be the graph obtained by taking q copies of H which all agree on R and which are otherwise disjoint.



#### **Definition**

Given a graph H, a subset  $R \subseteq V(H)$ , and an integer q, we define the sunflower-power  $H_R^q$  to be the graph obtained by taking q copies of H which all agree on R and which are otherwise disjoint.

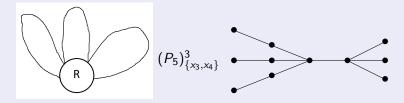


#### We define

$$\mathcal{F}_{H,k}^q = \{H_R^q : H - R \text{ has at least } k \text{ connected components}\}.$$

#### **Definition**

Given a graph H, a subset  $R \subseteq V(H)$ , and an integer q, we define the sunflower-power  $H_R^q$  to be the graph obtained by taking q copies of H which all agree on R and which are otherwise disjoint.



We define

$$\mathcal{F}_{H,k}^q = \{H_R^q : H - R \text{ has at least } k \text{ connected components}\}.$$

#### Claim

Every graph in  $\mathcal{F}_{H,k}^q$  has at least  $q^k$  copies of H.

### Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}_{H,k}^q$ -free.

## Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{H,k}$ -free.

#### Proof.

If some  $H_R^q \in \mathcal{F}_{H,k}^q$  is always  $\mathcal{F}$ -free, then the previous claim with  $q \approx n/v(H)$  shows  $\mathrm{ex}(n,H,\mathcal{F}) = \Omega(n^k)$ .

## Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{H,k}$ -free.

#### Proof.

If some  $H_R^q \in \mathcal{F}_{H,k}^q$  is always  $\mathcal{F}$ -free, then the previous claim with  $q \approx n/v(H)$  shows  $\operatorname{ex}(n,H,\mathcal{F}) = \Omega(n^k)$ . Otherwise there exists some q so that every  $H_R^q \in \mathcal{F}_{H,k}^q$  contains an element of  $\mathcal{F}$  as a subgraph.  $\square$ 

## Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{H,k}$ -free.

#### Proof.

If some  $H_R^q \in \mathcal{F}_{H,k}^q$  is always  $\mathcal{F}$ -free, then the previous claim with  $q \approx n/v(H)$  shows  $\operatorname{ex}(n,H,\mathcal{F}) = \Omega(n^k)$ . Otherwise there exists some q so that every  $H_R^q \in \mathcal{F}_{H,k}^q$  contains an element of  $\mathcal{F}$  as a subgraph.  $\square$ 

### Corollary (General Stability for Generalized Turán)

If  $ex(n, H, \mathcal{F}_{H,k}^q) = O_q(n^{\beta})$ , then every family  $\mathcal{F}$  either satisfies  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or  $ex(n, H, \mathcal{F}) = O(n^{\beta})$ .

#### **Proposition**

For every graph H and family  $\mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n)$  or  $ex(n, H, \mathcal{F}) = O(1)$ .

#### Proposition

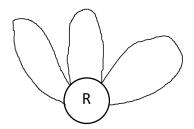
For every graph H and family  $\mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n)$  or  $ex(n, H, \mathcal{F}) = O(1)$ .

By the previous corollary, it suffices to prove  $ex(n, H, \mathcal{F}_{H,1}^q) = O_q(1)$ .

#### Proposition

For every graph H and family  $\mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n)$  or  $ex(n, H, \mathcal{F}) = O(1)$ .

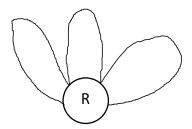
By the previous corollary, it suffices to prove  $ex(n, H, \mathcal{F}_{H,1}^q) = O_q(1)$ .



### **Proposition**

For every graph H and family  $\mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n)$  or  $ex(n, H, \mathcal{F}) = O(1)$ .

By the previous corollary, it suffices to prove  $ex(n, H, \mathcal{F}_{H,1}^q) = O_q(1)$ .



This is essentially equivalent to the Erdős-Rado Sunflower lemma.



### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

### Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{H,k}$ -free.

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

### Proposition (Key Observation)

For every  $H, k, \mathcal{F}$ , either  $ex(n, H, \mathcal{F}) = \Omega(n^k)$  or there exists some q such that every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{H,k}$ -free.

### Theorem (English-S. 2025+)

If T is a tree and if G is  $\mathcal{F}^q_{T,k}$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

## **Proof Sketch**

### Theorem (English-S. 2025+)

If T is a tree and if G is  $\mathcal{F}_{T,k}^q$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

### Theorem (English-S. 2025+)

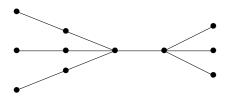
If T is a tree and if G is  $\mathcal{F}^q_{T,k}$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

Strategy: identify each copy of T in G by a set of k-1 edges E such that each E identifies at most O(1) copies of T.

### Theorem (English-S. 2025+)

If T is a tree and if G is  $\mathcal{F}^q_{T,k}$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

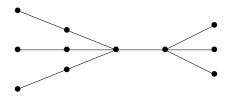
Strategy: identify each copy of T in G by a set of k-1 edges E such that each E identifies at most O(1) copies of T.



### Theorem (English-S. 2025+)

If T is a tree and if G is  $\mathcal{F}^q_{T,k}$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

Strategy: identify each copy of T in G by a set of k-1 edges E such that each E identifies at most O(1) copies of T.

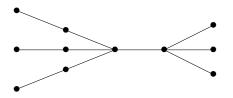


For this example, to specify a given copy of  $P_5$ , we **must** identify its last edge.

### Theorem (English-S. 2025+)

If T is a tree and if G is  $\mathcal{F}^q_{T,k}$ -free, then the number of copies of T in G is at most  $O(e(G)^{k-1})$ .

Strategy: identify each copy of T in G by a set of k-1 edges E such that each E identifies at most O(1) copies of T.



For this example, to specify a given copy of  $P_5$ , we **must** identify its last edge. More generally, whatever set of edges E we choose to identify a given copy K in a graph, the edges of E **must** intersect every subtree  $K' \subseteq K$  which has "many extensions."

#### Question

Given a set of subtrees  $\mathcal{T}$  of a tree  $\mathcal{T}$ , when can we guarantee that there exist a set of k-1 edges E which intersect all of these subtrees?

#### Question

Given a set of subtrees  $\mathcal{T}$  of a tree  $\mathcal{T}$ , when can we guarantee that there exist a set of k-1 edges  $\mathcal{E}$  which intersect all of these subtrees?

### Theorem (Helly Property for Trees)

If  $\mathcal T$  is a set of subtrees of a tree T such that the vertex sets of the subtrees pairwise intersect, then there exists a vertex  $v \in V(T)$  which intersects the vertex set of every subtree of  $\mathcal T$ .





#### **Definition**

We call a subtree  $T' \subseteq T$  leaf-cuttable if  $e(T) \ge 1$  and if every edge in  $E(T) \setminus E(T')$  which intersects T' intersects a leaf of T'.

### Theorem (English-S. 2025+)

If  $\mathcal T$  is a set of <u>leaf-cuttable</u> subtrees of a tree  $\mathcal T$  such that the edge sets of the subtrees pairwise intersect, then there exists an edge  $e \in E(\mathcal T)$  which intersects the edge set of every subtree of  $\mathcal T$ .

### Theorem (English-S. 2025+)

If  $\mathcal T$  is a set of <u>leaf-cuttable</u> subtrees of a tree  $\mathcal T$  such that the edge sets of the subtrees pairwise intersect, then there exists an edge  $e \in E(\mathcal T)$  which intersects the edge set of every subtree of  $\mathcal T$ .

### Theorem (English-S. 2025+)

If  $\mathcal T$  is a set of <u>leaf-cuttable</u> subtrees of a tree  $\mathcal T$  such that every k subtrees from  $\mathcal F$  contains two subtrees with intersecting edge sets, then there exists a set  $E\subseteq E(\mathcal T)$  of at most k-1 edges which intersects the edge set of every subtree of  $\mathcal T$ .

### Theorem (English-S. 2025+)

If  $\mathcal T$  is a set of <u>leaf-cuttable</u> subtrees of a tree  $\mathcal T$  such that the edge sets of the subtrees pairwise intersect, then there exists an edge  $e \in E(\mathcal T)$  which intersects the edge set of every subtree of  $\mathcal T$ .

### Theorem (English-S. 2025+)

If  $\mathcal{T}$  is a set of <u>leaf-cuttable</u> subtrees of a tree  $\mathcal{T}$  such that every k subtrees from  $\mathcal{F}$  contains two subtrees with intersecting edge sets, then there exists a set  $E \subseteq E(\mathcal{T})$  of at most k-1 edges which intersects the edge set of every subtree of  $\mathcal{T}$ .

This result implies the vertex-Helly result (and also König's Theorem for trees).

We are now at the "starting line" of the proof, the rest of which requires us to do a few more things:

We are now at the "starting line" of the proof, the rest of which requires us to do a few more things:

• Verify that the "highly extendable" subtrees  $\mathcal{T}_{\mathcal{K}}$  are leaf-cuttable (easy),

We are now at the "starting line" of the proof, the rest of which requires us to do a few more things:

- Verify that the "highly extendable" subtrees  $\mathcal{T}_K$  are leaf-cuttable (easy),
- Verify that any k subtrees from  $\mathcal{T}_K$  have a pair with intersecting edge sets (true because  $\mathcal{F}^q_{T_k}$ -free, but not so easy to prove),

#### **Proof Vibes**

We are now at the "starting line" of the proof, the rest of which requires us to do a few more things:

- Verify that the "highly extendable" subtrees  $\mathcal{T}_K$  are leaf-cuttable (easy),
- Verify that any k subtrees from  $\mathcal{T}_K$  have a pair with intersecting edge sets (true because  $\mathcal{F}^q_{T_k}$ -free, but not so easy to prove),
- Verify that finding this set of k-1 edges for each  $\mathcal{T}_K$  is enough to give the desired bound (annoying but doable).

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{T,k}$ -free and

$$ex(n, T, F) = O(ex(n, F)^{k-1}).$$

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{T,k}$ -free and

$$ex(n, T, F) = O(ex(n, F)^{k-1}).$$

#### Question

Does this hold with T replaced by some more general type of graph?

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{T,k}$ -free and

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

#### Question

Does this hold with T replaced by some more general type of graph?

If this result holds at  $k = \alpha(H) + 1$ , then

$$ex(n, H, \mathcal{F}) = O(ex(n, \mathcal{F})^{\alpha(H)})$$

### Theorem (English-S. 2025+)

For any tree T, integer  $k \geq 1$ , and family of graphs  $\mathcal{F}$ , either

$$ex(n, T, \mathcal{F}) = \Omega(n^k),$$

or every  $\mathcal{F}$ -free graph is  $\mathcal{F}^q_{T,k}$ -free and

$$ex(n, T, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

#### Question

Does this hold with T replaced by some more general type of graph?

If this result holds at  $k = \alpha(H) + 1$ , then

$$ex(n, H, \mathcal{F}) = O(ex(n, \mathcal{F})^{\alpha(H)}),$$

and this holds for bipartite graphs without isolated edges by König's Theorem.



#### Question

Is it the case that for every bipartite graph H without isolated vertices, integer k, and family of graphs  $\mathcal{F}$  that either

$$ex(n, H, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, H, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

#### Question

Is it the case that for every bipartite graph H without isolated vertices, integer k, and family of graphs  $\mathcal F$  that either

$$ex(n, H, \mathcal{F}) = \Omega(n^k),$$

or

$$ex(n, H, \mathcal{F}) = O(ex(n, \mathcal{F})^{k-1}).$$

#### **Fact**

This is false for  $C_4$  and k = 2:

We've seen a lot of ways that  $ex(n, H, \mathcal{F})$  can't behave like; what about the ways it can behave?

We've seen a lot of ways that  $ex(n, H, \mathcal{F})$  can't behave like; what about the ways it can behave?

#### **Definition**

We say that a rational number r is *realizable* for a graph H if there exists a (finite) family of graphs  $\mathcal{F}$  such that  $ex(n, H, \mathcal{F}) = \Theta(n^r)$ .

We've seen a lot of ways that  $ex(n, H, \mathcal{F})$  can't behave like; what about the ways it can behave?

#### **Definition**

We say that a rational number r is *realizable* for a graph H if there exists a (finite) family of graphs  $\mathcal{F}$  such that  $ex(n, H, \mathcal{F}) = \Theta(n^r)$ .

### Theorem (Bukh-Conlon 2018)

Every rational in [1,2] is realizable for  $H = K_2$ .

### Theorem (English-Halfpap-Krueger 2024)

Every rational in [1, t] is realizable for  $H = K_t$ .

### Theorem (English-Halfpap-Krueger 2024)

Every rational in [1, t] is realizable for  $H = K_t$ .

### Theorem (English-Halfpap-Krueger 2024)

Every rational in [t, t+1] is realizable for  $H = K_{1,t}$  and no rational in (1,t) is.

#### Theorem (English-S. 2025+)

For every graph H of maximum degree  $\Delta$ , every rational in the range

$$\left[v(H)-\frac{e(H)}{2\Delta^2},v(H)\right]$$

is realizable.

### Theorem (English-S. 2025+)

For every graph H of maximum degree  $\Delta$ , every rational in the range

$$\left[v(H)-\frac{e(H)}{2\Delta^2},v(H)\right]$$

is realizable.

### Theorem (English-S. 2025+)

For every tree T of maximum degree  $\Delta$ , every rational in the range

$$\left[v(H)-\frac{e(H)}{\Delta},v(H)\right]$$

is realizable.

This is best possible for stars.



### Question

Given a graph H, what rationals r are realizable?

#### Question

Given a graph H, what rationals r are realizable?

From our general stability theorem, we know that if  $ex(n, H, \mathcal{F}^q_{H,k}) = o(n^\beta)$ , then no rational in  $(\beta, k)$  is realizable.

#### Question

Given a graph H, what rationals r are realizable?

From our general stability theorem, we know that if  $ex(n, H, \mathcal{F}_{H,k}^q) = o(n^\beta)$ , then no rational in  $(\beta, k)$  is realizable.

#### Question

Is this the only way an exponent can fail to be rational?

#### Question

Given a graph H, what rationals r are realizable?

From our general stability theorem, we know that if  $ex(n, H, \mathcal{F}_{H,k}^q) = o(n^\beta)$ , then no rational in  $(\beta, k)$  is realizable.

#### Question

Is this the only way an exponent can fail to be rational?

We believe we can verify that this is true for all graphs on at most 4 vertices through various ad hoc techniques.

#### Question

Is it the case that every rational in  $[\alpha(H), \nu(H)]$  is realizable?

This is true and best possible for cliques and stars.

### Conjecture

For every graph H, every rational in

$$\left[v(H)-\frac{e(H)}{\Delta},v(H)\right]$$

is realizable.

Say that a rational r is strongly realizable if there exists a single graph F such that  $ex(n, H, F) = \Theta(n^r)$ .

Say that a rational r is strongly realizable if there exists a single graph F such that  $ex(n, H, F) = \Theta(n^r)$ .

#### Question

Is it the case that for all H, the number of strongly realizable exponents below  $\alpha(H)$  is  $O_H(1)$ ?

Say that a rational r is strongly realizable if there exists a single graph F such that  $ex(n, H, F) = \Theta(n^r)$ .

#### Question

Is it the case that for all H, the number of strongly realizable exponents below  $\alpha(H)$  is  $O_H(1)$ ?

The idea here is that any F of this form must be a subgraph of every member of  $\mathcal{F}^q_{H,\alpha(H)} \neq \emptyset$  for some q. For any fixed q this is a finite number of possibilities, and I don't think arbitrarily large q should give new behaviors (analogous to  $F \subseteq K_{2,t}$  implying  $\operatorname{ex}(n,F) = \Theta(n^r)$  for some  $r \in \{0,1,3/2\}$ ).

#### Question

Can one say anything about generalized Turán numbers of trees T satisfying

$$n^{\ell+1} \ll \operatorname{ex}(n, T, \mathcal{F}) \ll n^{\ell+2}$$
?

For this it would suffice to prove upper bounds onex $(n, T, \mathcal{F}^q_{T,\ell+2})$ .

Our new Helly theorem for edge sets of trees turns out to imply the usual Helly theorem for vertex sets.

#### Question

Does the usual vertex Helly Theorem imply the edge Helly Theorem?

Our new Helly theorem for edge sets of trees turns out to imply the usual Helly theorem for vertex sets.

#### Question

Does the usual vertex Helly Theorem imply the edge Helly Theorem? In particular, can you derive König's Theorem for trees?

Our new Helly theorem for edge sets of trees turns out to imply the usual Helly theorem for vertex sets.

#### Question

Does the usual vertex Helly Theorem imply the edge Helly Theorem? In particular, can you derive König's Theorem for trees?

I'm willing to offer up to 1 for a proof or disproof of either of these questions.

### Proposition

If G contains  $kn^{2-1/s}$  edges with  $k = \omega(1)$ , then it contains at least  $\Omega(k^{st}n^s)$  copies of  $K_{s,t}$ .

Note that this is best possible by considering a random graph with  $kn^{2-1/s}$  edges.

### Proposition

If G contains  $kn^{2-1/s}$  edges with  $k = \omega(1)$ , then it contains at least  $\Omega(k^{st}n^s)$  copies of  $K_{s,t}$ .

Note that this is best possible by considering a random graph with  $kn^{2-1/s}$  edges.

#### Question

Does the same phenomenon hold if G contains many copies of  $K_r$ ?

### Conjecture (Dubroff-Gunby-Narayanan-S.)

If  $2 \le r \le s \le t$  and if G contains at least  $kn^{r-\binom{r}{2}/s}$  copies of  $K_r$ , then it contains at least  $k^{st/\binom{r}{2}}n^{s-o(1)}$  copies of  $K_{s,t}$ .

### Conjecture (Dubroff-Gunby-Narayanan-S.)

If  $2 \le r \le s \le t$  and if G contains at least  $kn^{r-\binom{r}{2}/s}$  copies of  $K_r$ , then it contains at least  $k^{st/\binom{r}{2}}n^{s-o(1)}$  copies of  $K_{s,t}$ .

Things seem to behave differently if r > s and k is small.

### Conjecture (Dubroff-Gunby-Narayanan-S.)

If  $2 \le r \le s \le t$  and if G contains at least  $kn^{r-\binom{r}{2}/s}$  copies of  $K_r$ , then it contains at least  $k^{st/\binom{r}{2}}n^{s-o(1)}$  copies of  $K_{s,t}$ .

Things seem to behave differently if r > s and k is small.

### Conjecture (Dubroff-Gunby-Narayanan-S.)

For all  $1 \le k \le n^{1/2t}$ , there exist n-vertex graphs with  $kn^{3/2}$  triangles and at most  $k^t n^{3/2+o(1)}$  copies of  $K_{2,t}$ .