# Extremal Problems for Random Objects

Sam Spiro, Rutgers University



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## Theorem (Erdős-Stone 1946)

$$ex(n,F) = \left(1 - \frac{1}{\chi(F) - 1} + o(1)\right) \binom{n}{2}.$$

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The lower bound is tight when p = 1. The upper bound is tight if p is "small."

$$\frac{1}{2} p \binom{n}{2} \lesssim \text{ex}(\textit{G}_{n,p},\textit{K}_{3}) \lesssim p \binom{n}{2},$$

with the lower bound tight for p=1 and the upper bound tight for  $p\ll n^{-1/2}$ .

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# Theorem (Frankl-Rödl 1986)

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$$\operatorname{ex}(G_{n,p},K_3)\sim \frac{1}{2}p\binom{n}{2} \qquad p\gg n^{-1/2}.$$

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Theorem (Conlon-Gowers, Schacht 2010)

 $p \gg n^{-1/m_2(F)}$ 

$$\operatorname{ex}(G_{n,p},F) = p \cdot \left(1 - \frac{1}{\chi(F) - 1} + o(1)\right) \binom{n}{2} \qquad p \gg$$

 $\operatorname{ex}(G_{n,p},K_3)\sim \frac{1}{2}p\binom{n}{2}$ 

where  $m_2(F) = \max\{\frac{e(F')-1}{v(F')-2}: F' \subseteq F\}.$ 

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#### Conjecture

If F is a bipartite graph which is not a forest, then whp

$$\operatorname{ex}(G_{n,p},F) = \begin{cases} \Theta(p \cdot \operatorname{ex}(n,F)) & p \gg n^{-1/m_2(F)}, \\ (1+o(1))p\binom{n}{2} & p \ll n^{-1/m_2(F)}. \end{cases}$$

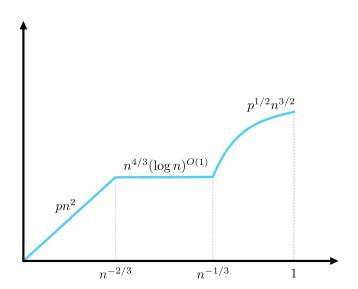
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This conjecture turns out to be completely false!



Plot of  $ex(G_{n,p}, C_4)$  (Füredi 1991)

## Conjecture (McKinley-S.)

If F is a graph with  $\operatorname{ex}(\mathsf{n},\mathsf{F}) = \Theta(\mathsf{n}^{\alpha})$  for some  $\alpha \in (1,2]$ , then whp

$$\exp(G_{n,p},F) = \max\{\Theta(p^{\alpha-1}n^{\alpha}), n^{2-1/m_2(F)}(\log n)^{O(1)}\},$$

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## Theorem (Nie-S. 2023 (Informal))

This conjecture (essentially) implies Sidorenko's conjecture.

## Theorem (Kővari-Sós-Turán 1954)

$$\operatorname{ex}(n,K_{s,t})=O(n^{2-1/s}).$$



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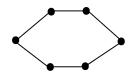
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Moreover, this bound is tight whenever  $ex(n, K_{s,t}) = \Theta(n^{2-1/s})$ .

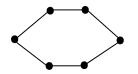
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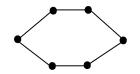


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## Theorem (Morris-Saxton 2013)

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Moreover, this is tight whenever  $ex(n, \{C_3, C_4, \dots, C_{2b}\}) = \Theta(n^{1+1/b})$ .

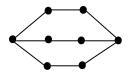
# Theorem (Jiang-Longbrake 2022)

If F satisfies "mild conditions", then

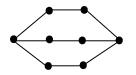
$$\operatorname{ex}(G_{n,p},F) = O(p^{1-m_2^*(F)(2-\alpha)}n^{\alpha}) \text{ for } p \text{ large},$$

where 
$$m_2^*(F) = \max\{\frac{e(F')-1}{v(F')-2} : F' \subsetneq F, \ e(F') \ge 2\}.$$

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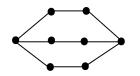


#### Theorem (Corsten-Tran 2021)

$$\operatorname{ex}(G_{n,p}, \theta_{a,b}) = O(p^{\frac{2}{ab}} n^{1+1/b}) \text{ for } p \text{ large.}$$

Note: our conjecture predicts  $p^{\frac{1}{b}}n^{1+1/b}$ .

$$\operatorname{ex}(n,\theta_{a,b}) = O(n^{1+1/b}).$$

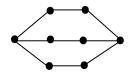


#### Theorem (McKinley-S. 2023)

For  $a \ge 100$ ,

$$ex(G_{n,p}, \theta_{a,b}) = O(p^{\frac{1}{b}}n^{1+1/b}) \text{ for } p \text{ large.}$$

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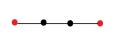
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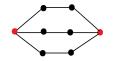
Moreover, this bound is tight whenever a is sufficiently large in terms of b.

# Theorem (Bukh-Conlon 2015)

If  $T^\ell$  is the " $\ell$ th power of a balanced tree" and  $\ell$  is sufficiently large, then

$$\operatorname{ex}(n, T^{\ell}) = \Omega(n^{2-\rho(T)}).$$





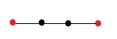


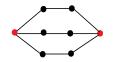


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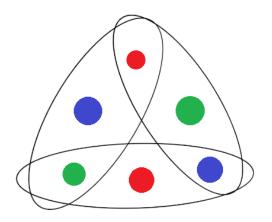


## Theorem (S. 2022)

$$ex(G_{n,p}, T^{\ell}) = \Omega(p^{1-\rho(T)}n^{2-\rho(T)}),$$

provided  $\ell$  is sufficiently large.

# Hypergraphs



## Theorem (S.-Verstraëte 2021)

Let  $K_{s_1,\ldots,s_r}^r$  denote the complete r-partite r-graph with parts of sizes  $s_1,\ldots,s_r$ . There exist constants  $\beta_1,\beta_2,\beta_3,\gamma$  depending on  $s_1,\ldots,s_r$  such that, for  $s_r$  sufficiently large in terms of  $s_1,\ldots,s_{r-1}$ , we have whp

$$\mathrm{ex}(\textit{G}^{r}_{\textit{n},\textit{p}},\textit{K}^{r}_{\textit{s}_{1},...,\textit{s}_{r}}) = \begin{cases} \Theta\left(\textit{p}\textit{n}^{r}\right) & \textit{n}^{-r} \ll \textit{p} \leq \textit{n}^{-\beta_{1}}, \\ \textit{n}^{r-\beta_{1}+o(1)} & \textit{n}^{-\beta_{1}} \leq \textit{p} \leq \textit{n}^{-\beta_{2}}(\log \textit{n})^{\gamma}, \\ \Theta\left(\textit{p}^{1-\beta_{3}}\textit{n}^{r-\beta_{3}}\right) & \textit{n}^{-\beta_{2}}(\log \textit{n})^{\gamma} \leq \textit{p} \leq 1. \end{cases}$$

#### Question

Does the McKinley-Spiro conjecture extend to hypergraphs?

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# Theorem (Nie-S. 2023 (Informal))

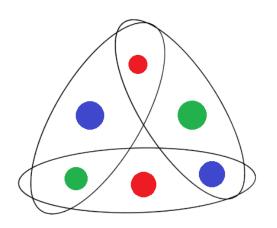
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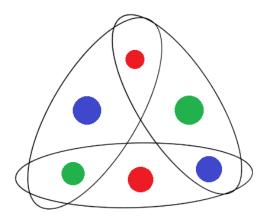
#### Question

Does the McKinley-Spiro conjecture extend to hypergraphs?

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Many hypergraphs fail to have a flat middle range. More precisely, any hypergraph which isn't Sidorenko fails to have a flat middle range.



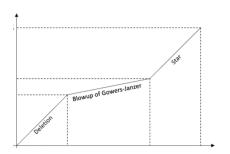


We define the *loose cycle*  $C_{\ell}^{r}$  to be the *r*-uniform hypergraph obtained by inserting r-2 distinct vertices into each edge of the graph cycle  $C_{\ell}$ .

### Theorem (Nie-S.-Verstaëte 2020; Nie 2023)

For  $r \ge 3$ , if  $p \gg n^{-r+3/2}$  then whp

$$ex(G_{n,p}^r, C_3^r) = max\{p^{\frac{1}{2r-3}}n^{2+o(1)}, pn^{r-1+o(1)}\}.$$



Picture due to Jiaxi Nie.

## Theorem (Mubayi-Yepremyan 2020; Nie 2023)

For  $r \ge 4$ , if  $p \gg n^{-r+1+\frac{1}{2\ell-1}}$  then whp

$$\operatorname{ex}(G_{n,p}^r, C_{2\ell}^r) = \max\{n^{1+\frac{1}{2\ell-1}}, pn^{r-1}\}.$$

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Bounds also are known for Berge cycles, but the bounds are significantly weaker (S.-Verstraëte; Nie).

## Theorem (Nie-S. 20XX (Informal))

If F is a graph and one has upper bounds for  $ex(G_{n,p}, F)$ , then one can prove corresponding bounds for  $ex(G_{n,p}^r, F^{+r})$ .

Here  $F^{+r}$  is the r-graph obtained by inserting r-2 new vertices inside each edge.

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Here  $F^{+r}$  is the r-graph obtained by inserting r-2 new vertices inside each edge.

#### Corollary

We have tight bounds for  $ex(G_{n,p}^r, K_{s,t}^{+r})$  if  $r \ge s + 2$ .

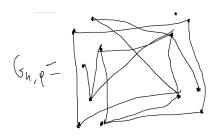
# **Upper Bound Techniques**

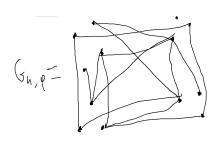
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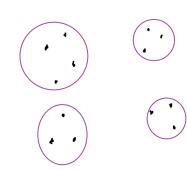
# Proof. Containers.

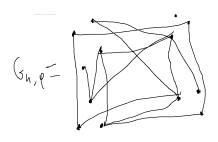
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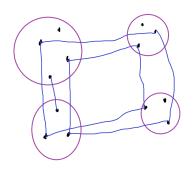
Proof.	
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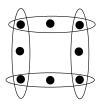


# Future Problems

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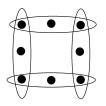
Prove tight bounds for the 3-uniform loose 4-cycle.



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#### **Problem**

Prove tight bounds for the 3-uniform loose 4-cycle.



#### **Problem**

Prove tight bounds for subdivisions of complete bipartite graphs.