

## Practical No: 4

**Problem statement: Write a program for error detection and correction for 7/8 bits ASCII codes using Hamming Codes or CRC. Demonstrate the packets captured traces using Wireshark Packet Analyzer Tool for peer to peer mode.( 50% students will perform Hamming Code and others will perform CRC)**

### Objectives:

1. To understand how CRC/FCS is calculated
2. To Learn Error Detection and Correction.
3. To Learn Hamming Code method.

### Procedure for Error Correcting Codes

#### Cyclic Redundancy Check (CRC) $M = K$

bit messages

$F = n$  bits frame check sequence (FCS)

$T = (k + n)$  bits frame where  $n < k$

$P = n+1$  is the predetermined divisor

We want  $\text{mod}(T, P) = 0$

Rewrite  $T = M 2^n + F$   $F =$

$M 2^n / P$

**Ex:**  $M = 1010001101$  ( 10 bits)  $P =$

$110101$  (6 bits)

$P = 6$  hence  $n=5$  bits  $F = 5$  bits to becalculated  $M 2^n$

$= M 2^5 = 101000110100000$

$$2^5 M / P = 101000110100000 / 110101$$

Method –I : Modulo-2 Arithmetic

<b>110101</b>	1	0	1	0	0	0	1	1	0	1	0	0	0	0	0	Q
	1	1	0	1	0	1										1
	x	1	1	1	0	1										
		1	1	1	0	1	1									
		1	1	0	1	0	1									1
		X	0	1	1	1	0									
			0	1	1	1	0	1	nd							
				0	0	0	0	0								0
				1	1	1	0	1	0							
				1	1	0	1	0	1							1
				x	0	1	1	1	1	1						
					0	0	0	0	0	0						0
					1	1	1	1	1	0						
					1	1	0	1	0	1						1
					x	0	1	0	1	1	0					
						0	0	0	0	0	0					0
						x	1	0	1	1	0	0				
							1	1	0	1	0	1				1
							x	1	1	0	0	1	0			
								1	1	0	1	0	1			1
									x	x	0	1	1	1	0	0

F= Remainder =  
R

Therefore F =

$$01110 \text{ T} = M 2^n + F$$

101000110100000

+ 01110

101000110101110

At the receiver , it receives T=101000110101110

At receiving end, it divides T by same P = **110101**

<b>110101</b>	1	0	1	0	0	0	1	1	0	1	0	1	1	1	0	Q
<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>											1
x	1	1	1	0	1											
	1	1	1	0	1	1										
<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>											1
	X	0	1	1	1	0										
		0	1	1	1	0	1	nd								
			<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>									0
			1	1	1	0	1	0								
			<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>								1
			x	0	1	1	1	1	1							
				<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>							0
					1	1	1	1	1	0						
					<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>						1
					x	0	1	0	1	1	1					
						<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>					0
						x	1	0	1	1	1	1				

								1	1	0	1	0	1			1
								x	1	1	0	1	0	1		
									1	1	0	1	0	1		1
									0	0	0	0	0	0		
								F= Remainder = R								
								x	x	0	0	0	0	0	0	0

If remainder is 0 then frame T is received correctly

**Method –II : Polynomial Method** Message =  $M(x) = (k+n)$  -bits Divisor =  $G(x) = P$

=  $n+1$  bits

$F = g(x) = x^c = n$ -bits

$G(x) \mid M(x) * x^c \text{ (Q(x) T(x) = (k+n) bits}$

6 4 3

Ex.  $M(x) = x^6 + x^4 + x^3 + 1$   $g(x) = x^3$

Therefore  $M(x) * x^c$

9 7 6 3  
=  $x^9 + x^7 + x^6 + x^3$

-----  
 $(x^3 + 1) \mid x^9 + x^7 + x^6 + x^3 \text{ (} x^6 + x^4 + x + 1x^9 +$

$+ x^6$

-----

$$X^7 + \quad + X^3$$

$$X^7 + \quad + X^4$$

-----

$$X^4 + X^3$$

$$X^4 + X$$

-----

$$X^3 + X$$

$$X^3 + 1$$

-----

Remainder R=F  $x + 1$   $T(x) = x^9 + x^7 + x^6 + x^3 + x + 1$  At receiver

$$x^3 + 1) x^9 + x^7 + x^6 + x^3 + x + 1 (x^6 + x^4 + x + 1 x^9 +$$

$$+ x^6$$

$$x^7 + \quad + x^3$$

$$x^7 + \quad + x^4$$

$$x^4 + x^3 + x$$

$$x^4 + x$$

$$x^3 + 1 \text{ and finally } 0$$

### 1] MODULO 2 METHOD CODE: Sample Input and Output

Enter Message M: 1 1 1 0 0 1 1 0

M: 11100110

Enter Predetermined Divisor P: 1 1 0 0 1

P: 11001

-----T: 111001100000

R after the fist

time : 00101 1R=

001011

1S= 00000

2R= 010111

2S= 11001

3R= 011100

3S= 11001

4R= 001010

4S= 00000

5R= 010100

5S= 11001

6R= 011010

6S= 11001

7R= 000110

7S= 00000

CRC = 0110

## 2] POLYNOMIAL METHOD CODE: Sample Input and Output

Enter no of terms for M 5 : 1 2 5 6 7

Enter no of terms for G 3 : 0 3 4

M:  $x^7 + x^6 + x^5 + x^2 + x^1 + G$ :

$x^4 + x^3 + x^0 +$

T:  $x^{11} + x^{10} + x^9 + x^6 + x^5 + S$ :

$x^{11} + x^{10} + x^7 +$

11 cut off

10 cut off

T after putting -1s:  $x^{-1} + x^{-1} + x^9 + x^6 + x^5 +$  Now T

after sorting 1 :  $x^{-1} + x^{-1} + x^5 + x^6 + x^9 +$

T after appending elements of S:  $x^7 + x^5 + x^6 + x^9 +$

Final T:  $x^9 + x^7 + x^6 + x^5 +$

S:  $x^9 + x^8 + x^5 + x^{-1} + x^{-1} + x^{-1} +$

=====

9 cut off

5 cut off

T after putting -1s:  $x-1 + x7 + x6 + x-$

1 + Now T after sorting 1 :  $x-1 + x-1 +$

$x6 + x7 +$

T after appending elements of S:  $x8 + x6 + x7$

+ Final T:  $x8 + x7 + x6 +$

S:  $x8 + x7 + x4 + x-1 + x-1 + x-1 + x-1 + x-1 + x-1 +$

=====

8 cut off

7 cut off

T after putting -1s:  $x-1 + x-1 +$

$x6 +$  Now T after sorting 1 :  $x-1$

+  $x-1 + x6 +$

T after appending elements of S:  $x4$

+  $x6 +$  Final T:  $x6 + x4 +$

S:  $x6 + x5 + x2 + x-1 + x-1 + x-1 + x-1 + x-1 + x-1 + x-1 + x-1 + x-1 +$

=====

6 cut off





T after putting -1s:  $x^{-1} +$

$x^4 +$  Now T after sorting

1 :  $x^{-1} + x^4 +$

T after appending elements of S:  $x^2 + x^5$

$+ x^4 +$  Final T:  $x^5 + x^4 + x^2 +$

S:  $x^5 + x^4 + x^1 + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} +$

=====

5 cut off

4 cut off

T after putting -1s:  $x^{-1} + x^{-1} + x^2 +$

Now T after sorting 1 :  $x^{-1} + x^{-1} + x^2 +$

T after appending elements of S:  $x^1 +$

$x^2 +$  Final T:  $x^2 + x^1 +$

S:  $x^2 + x^1 + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} + x^{-1} +$   
 $x^{-1} + x^{-1} + x^{-2} +$

=====

-----FINAL ANSWER: CRC =  $x^2 + x^1 +$

**Conclusion:** students understand of error detection and correction for 7/8 bits ASCII codes using Hamming Codes or CRC. Demonstrate the packets captured traces using Wireshark Packet Analyzer Tool for peer to peer mode

**Signature with Date**