### ECS 32B - Priority Queues

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UC Davis - Summer Session #2 2020

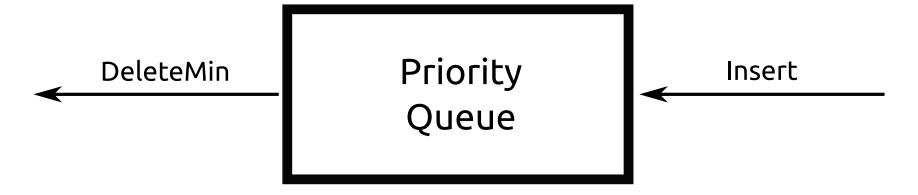


### Motivation

- Printer jobs.
- Operating system job scheduling.

### **Priority Queue**

- **priority queue**: ADT requiring these two operations:
  - o Insert.
  - o DeleteMin: retrieve and delete element with smallest key.



## Simple Implementations

#### **Unordered Linked List**

- Insert: constant time (at front).
- DeleteMin: linear time.

#### Ordered Linked List

- Insert: linear time.
- DeleteMin: constant time<sup>1</sup> (at front).

1. This *does* take constant time. During lecture, I accidentally talked as if it were an ordered Python list instead of an ordered linked list.

## Other Implementations

#### Binary Search Tree

- Insert: linear time.
  - Average: logarithmic time.
- DeleteMin: linear time.
  - Average: logarithmic time.

#### Self-Balancing BST

- Insert: logarithmic time.
- DeleteMin: logarithmic time.
  - Peeking at min: also logarithmic time.

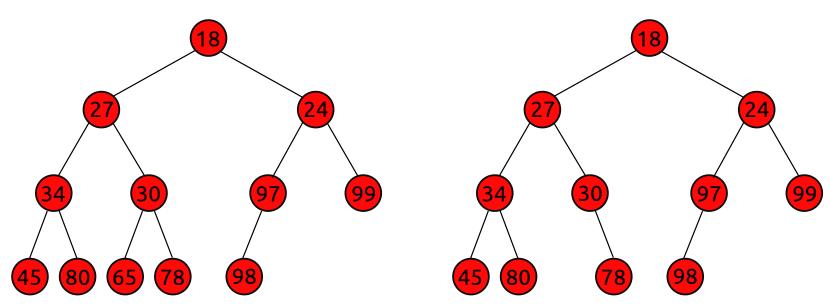
#### Binary Heap

- Preferred implementation for priority queue.
- Insert: logarithmic time.
  - Average: constant time.
- DeleteMin: logarithmic time.
  - Peeking at min: constant time.

- **binary heap**: binary tree with two properties:
  - 1. *structure property*: complete binary tree (completely filled, except possibly the last level, which is filled left to right).
  - 2. *heap-order property*: each node's key is less than its children's keys.
    - min heap: root is smallest element. (max heap reverses heap-order property.)
- Basic operations: insert, delete root (min if min heap).

#### Example

#### Nonexample

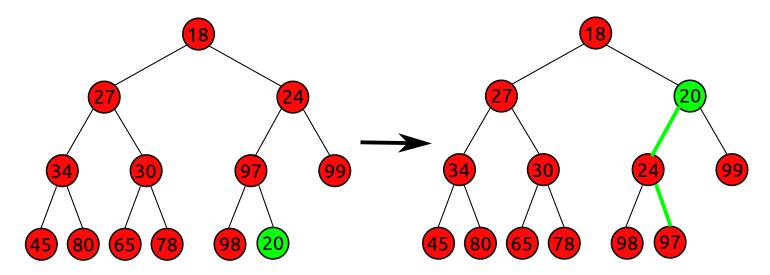


#### **Insert**

- 1. Insert new element in correct spot to maintain *structure property*.
- 2. Move the element up until *heap-order property* restored (**percolate up**).

#### Example

• Insert 20.



#### **Analysis**

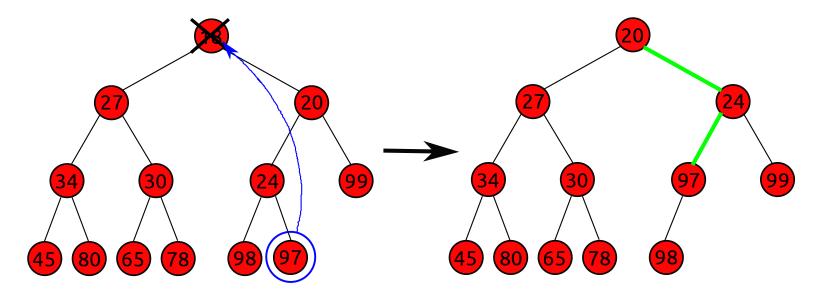
- ullet d + 1 assignments, where inserted node is percolated up d times.
- Worst-case time complexity:  $\Theta(\lg n)$  time (inserted node becomes root).
- Average-case time complexity: constant time<sup>1</sup>.

1. From p.251 of Weiss' book: "[I]t has been shown that 2.607 comparisons are required on average to perform an insert, so the average insert moves an element up 1.607 levels." 7/26

#### Delete

- 1. Remove root.
- 2. Move rightmost leaf to be root, to maintain *structure property*.
- 3. Move this new root down (by swapping) until *heap-order property* restored (**percolate down**).

#### Example

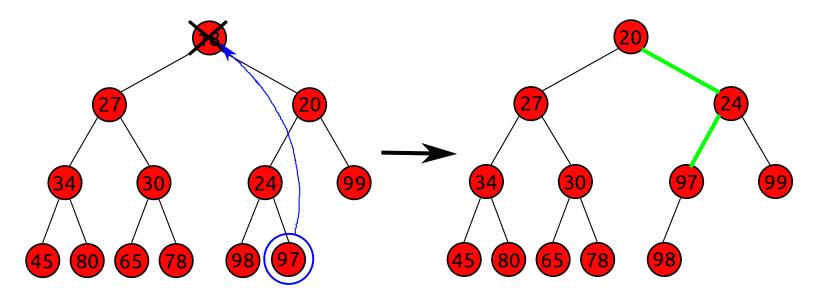


• For the first swap, could 97 be swapped with 27 instead of with 20?

#### Delete

- 1. Remove root.
- 2. Move rightmost leaf to be root, to maintain *structure property*.
- 3. Move this new root down (by swapping) until *heap-order property* restored (**percolate down**).

#### Example



• No; would violate *heap-order property*. Always swap with child that has smaller key.

#### Analysis

- Worst-case time complexity:  $\Theta(\lg n)$ .
- Average-case time complexity:  $\Theta(\lg n)$ .
  - Percolating down a node that was previously at bottom.

#### Representation: Python List

• Works because of rigid structure of binary heap (complete binary tree).

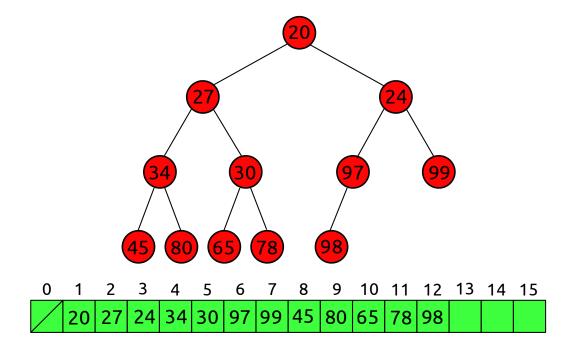
#### Pros/Cons

- Space-efficient (needn't waste space on the links as would with a typical tree implementation).
- Requires estimate of maximum heap size (can resize).

#### Parent/Children

- For node at index *i*:
  - Parent is at index  $\frac{i}{2}$  (truncate).
  - $\circ$  Left child is at index 2i.
  - Right child is at index 2i + 1.

#### Example



## Priority Queue ADT

#### Extended API (for Priority Queue *P* )

- If have additional data structure (e.g. hash table) to track position of each key/node in the heap (the underlying Python list), these extended operations take  $\Theta(\lg N)$  time:
  - DecreaseKey(k, change)
  - ∘ *IncreaseKey*(*k*, *change*)
  - $\circ$  *Remove*(k)

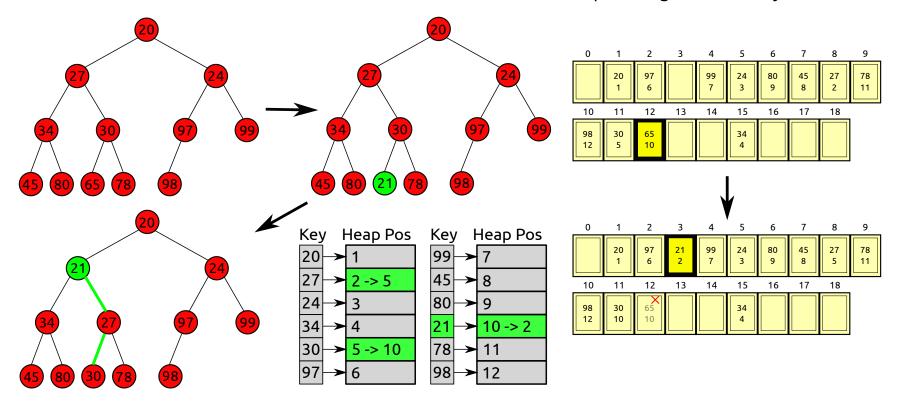
## Priority Queue ADT

#### $DecreaseKey(k, change)^{1}$

• Decreases the key of the target node by a specified amount.

#### Example<sup>2</sup>

- 1. Find the node using the additional data structure.
- 2. Decrease the key as desired.
- 3. Percolate up as long as necessary.



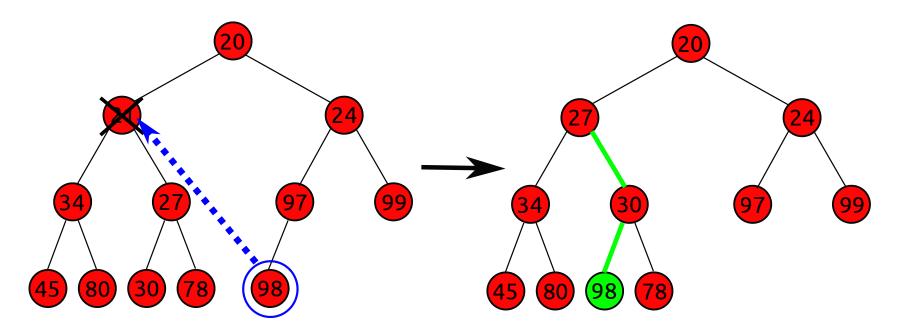
- 1. IncreaseKey is very similar.
- 2. The order in which the keys appear in the bottom-right of the image is unimportant (it should be a hash table anyways, npt a map).

## Priority Queue ADT

#### Remove(k) (any element, not necessarily minimum)<sup>1</sup>

- 1. Find the node using the additional data structure.
- 2. Remove the node.
- 3. Replace the node with the rightmost leaf, to maintain *structure property*.
- 4. Percolate up or down.

#### Example



1. In the book by Weiss, they implement Remove(P, k) by doing  $DecreaseKey(P, k, \infty)$  (i.e. forcing k to become the root) and then doing ExtractMin(P). This needlessly takes twice the amount of time but doesn't change the big-O.

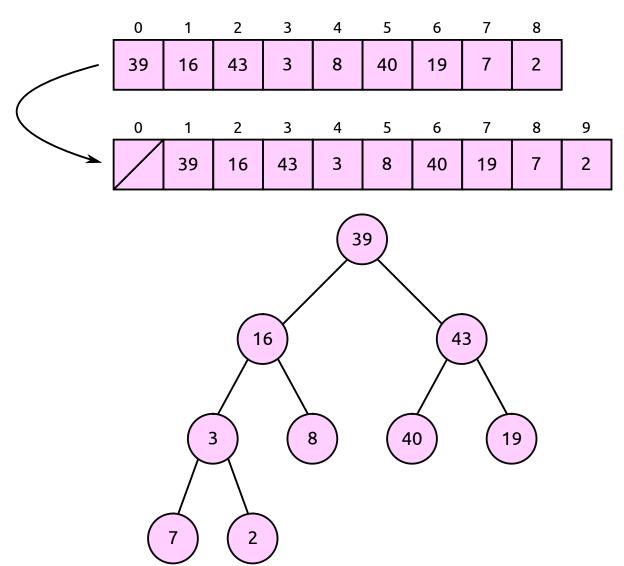
• **Goal**: Create a heap out of *n* given items (given in a Python list).

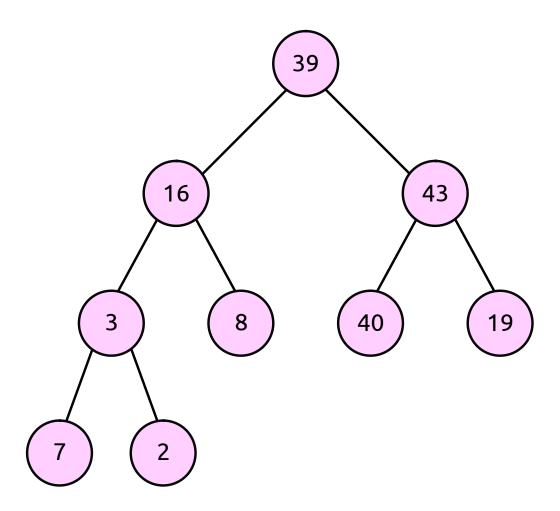
#### Naive Method

- Insert each item.
- Worst-case time complexity:  $\Theta(n \lg n)$ .
- Average-case time complexity:  $\Theta(n)$ .

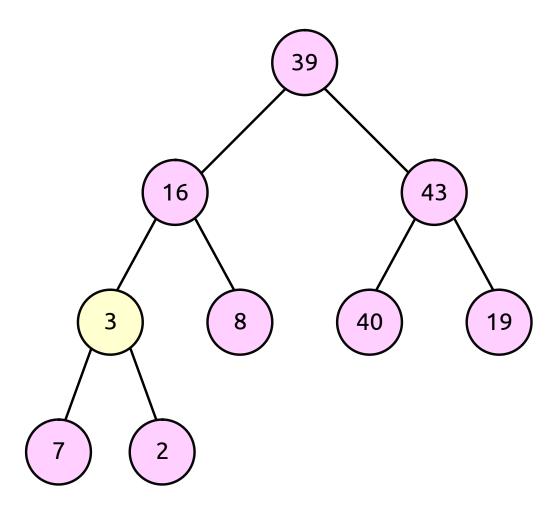
#### Build Heap

- Worst-case time complexity:  $\Theta(n)$  time.
  - We won't prove here.
- Steps:
  - 1. Treat list as a heap.
  - 2. For each level upwards (starting at second-to-lowest level), percolate each node down.

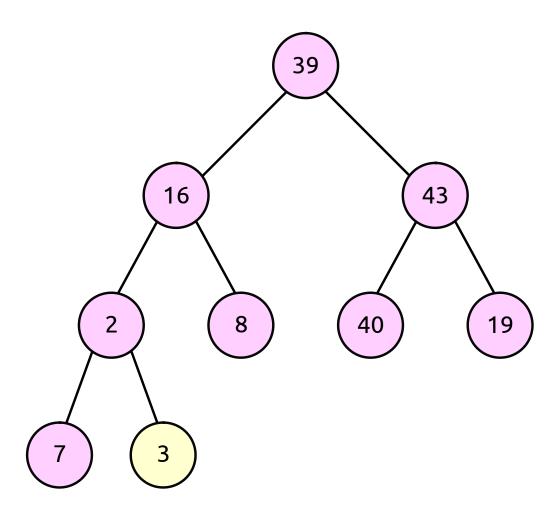


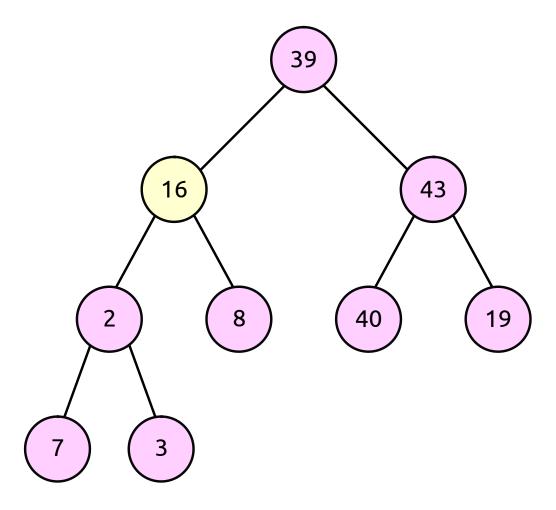


Build Heap Example

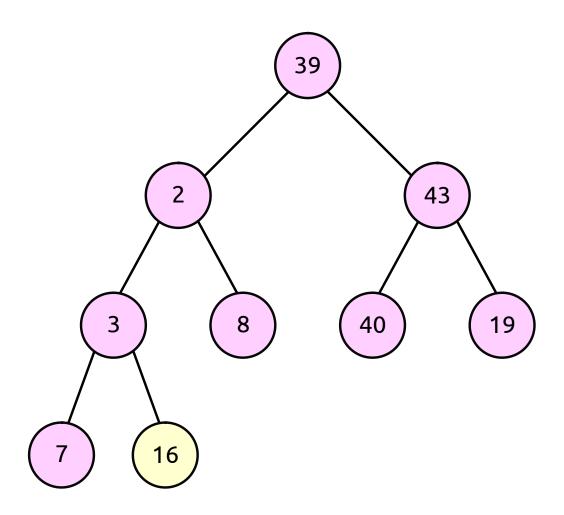


• Need to percolate 3 down.

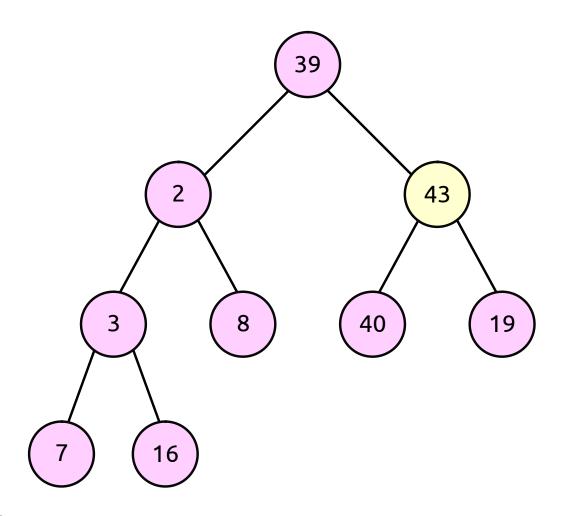




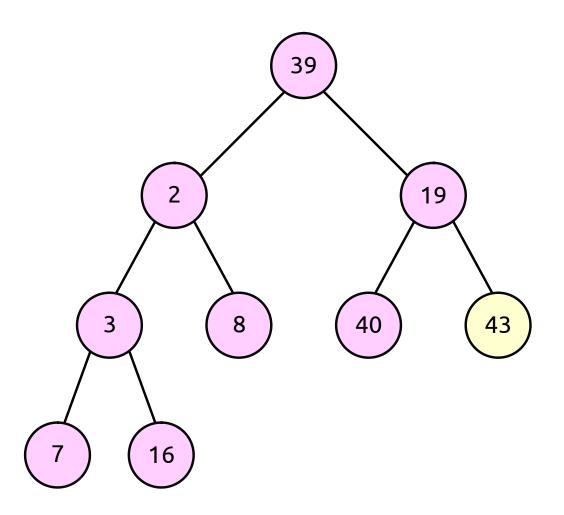
- 8, 40, and 19 are leaves, so skip them.
- Percolate 16 down.



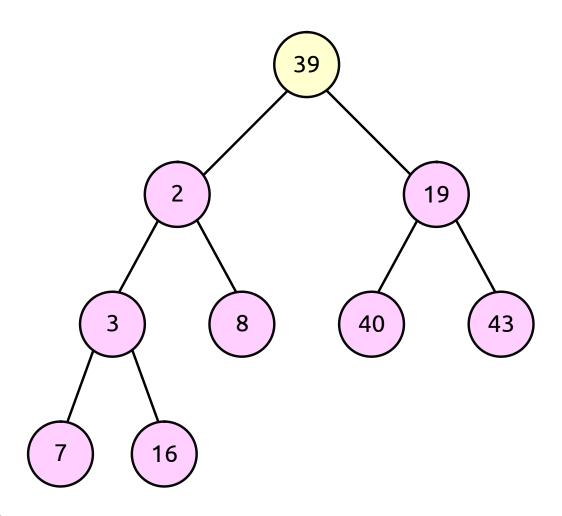
Build Heap Example



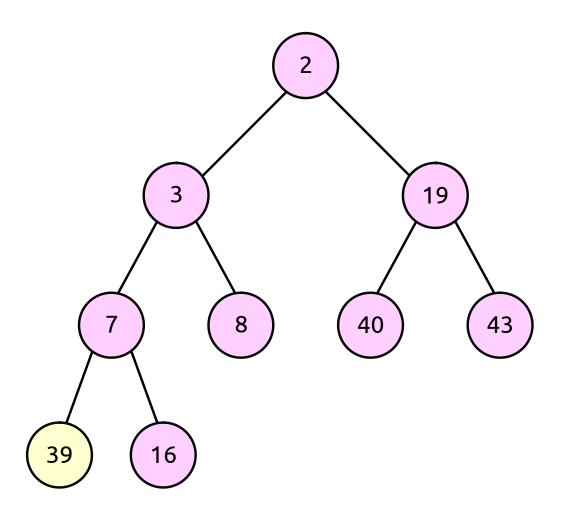
• Percolate 43 down.



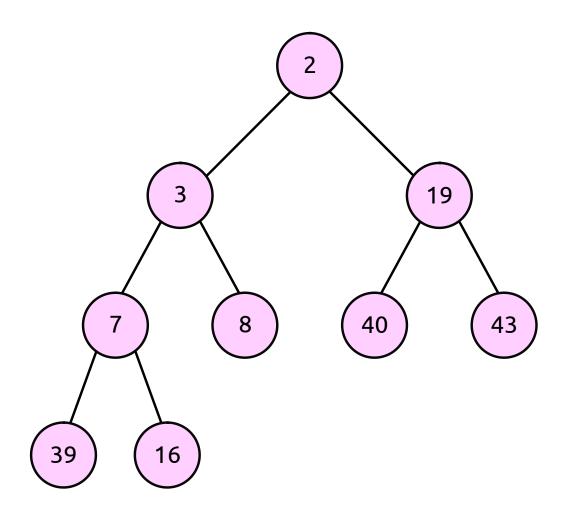
Build Heap Example



• Percolate 39 down.



Build Heap Example



• Done.

### References / Further Reading

- Chapter 6 of *Data Structures and Algorithm Analysis in C++* by Mark Allen Weiss (Fourth Edition).
- Sections 7.8-7.10 of *Problem Solving with Algorithms and Data Structures using Python* by Brad Miller and David Ranum.