

Q2: Part 1

B-F (bArray)

$n = \text{bArray.size}()$

element = [0,1], final\_result = Null;

$(2^n) \rightarrow$  for (i from 1 to  $2^n$ ) {

res = list()

for (j from 1 to n) {

if element[i % 2] == 1 {

res.append(bArray[i])

}

~~for (i from 1 to n) {~~ valid = 0

k = res.size()

for (c from 1 to k) {

if (res[c] - res[c-1]) > c {

valid = 1  
break

}

} if (final\_result == Null) { and valid == 0

final\_result = res;

}

else {

if (final\_result.size() > res.size()) { and valid == 0

final\_result = res

}

} }

☆ ANALYSIS: ☆

The run-time of this code is  $O(2^n \cdot n)$ .

Generate one possible result

Check validation

Assign the result.

Part 2. Greedy (bArray) {

$n = \text{bArray.size}()$

$A = \{b[n-1]\}$

$x = n-1$

while ( $x \neq 0$ ) {

Let  $p$  be the largest break point,  $\text{bArray}[p] + C \leq \text{bArray}[x]$

$x = p$

$A = A \cup \{bArray[p]\}$

}

Q3: sub-optimality:

If  $A = \{a_1, a_2, \dots, a_k\}$  is an optimal solution, then  $A' = A - \{a_i\}$  is the optimal solution of  $B' = \{b_1, b_2, \dots, b_p\}$  where  $b_p$  is the largest break point,  $b_p + c \leq b_n$ .

Proof: If there exist an optimal solution  $O'$  for  $B'$  that  $|O'| < |A'|$ , then let  $O = O' + \{a_i\}$  which means  $O$  is the global optimal solution and

$$|O| < |A|$$

this means  $A$  is not an global optimal solution, which is a contradiction. With the given information, our assumption is not right. So, the  $A'$  is the optimal solution for  $B'$ .