1. for循环的程序步的计算方法:

中文书24页有介绍for循环程序步的计算方法:

- 一般情况下一个for循环计一个程序步,注意循环语句块内的程序执行n次时,for语句执行n+1次,在退出for循环时会多计算一次。
- 2. 有同学将count++写在for循环内,如

for(int 
$$i = 0$$
;  $i < m$ ,count++;  $i + +$ )

上课的时候说错了,这样是计算不出来count的值的。for的判断语句是*i<m,count++* ,c++中逗号表达式的值返回最右边的式子的值,所以会返回count++的值,因为count初始值为0,count++返回的值为0,所以执行第一次for时就会跳出for循环,压根不会进入循环体。

或

$$for(int i = 0; i < m; i ++,count++)$$

根据for循环的执行顺序,需要在最后加 last for的count++;

- 3. for循环执行的顺序。
- 4. 程序简化:

目的是为了计算步程数,所以简化掉计算语句,留下一些计数语句如i++,并且将count++合并,最外层的count++也可以合并。

- 5. 是求程序步而不是用大O法表示时间复杂度。
- 6.计算程序步的表格还是按书中频度表的格式写。
  - Determine the frequency counts for all statements in the following two program segments:

```
1 for (i = 1; i <= n; i++)

2 for (j = 1; j <= i; j++)

3 for (k = 1; k <= j; k++)

4 x++;

(a)

1 i = 1;

2 while (i \le n)

4 x++;

5 i++;

6 }
```

a:

程序	程序步		
forli=1; i <= n; i +t)	n+	n+1	
forljelijæisjtt)	$\frac{n^2}{2} + \frac{2n}{2}$	<u>n(n+d)</u> 2	\frac{\lambda}{\sum_{i=1}} (i+1)
for(K=1; K=j;j++)	$\frac{n^3}{b} + n^2 + \frac{5n}{5}$	<u>n(n+1)(n+5)</u> 3	N ≥ (j+1)
カサナ;	$\frac{n^3}{6} + \frac{n^2}{2} + \frac{n}{3}$	<u>n(n+1)(n+2)</u>	Y W Z Z
$5 : \frac{n^3}{3} + 2n^2 + \frac{11n}{3} + 1$			

<b>行</b> ↔ 1↔	s/e₽	频度₽	行 <u>步程数</u> 。
	1.0	n+1₽	n+1.
2.₽	1.0	$\frac{n^2}{2} + \frac{3n}{2}$	$\frac{n^2}{2} + \frac{3n}{2} e$
3.₽	1.0	$\frac{n^3}{6} + n^2 + \frac{5n}{6}$	$\frac{n^3}{6} + n^2 + \frac{5n}{6}$
4.	1.0	$\frac{n^3}{6} + \frac{n^2}{2} + \frac{n}{3} \varphi$	$\frac{n^3}{6} + \frac{n^2}{2} + \frac{n}{3} $
θ	ρ	<u>总步程</u> 数:	$\frac{n^3}{3} + 2n^2 + \frac{11n}{3} + 1$

b:

行 <sub>4</sub> 1 <sub>4</sub>	s/e₽	頻度₽	行 <u>步程数</u> 。
1 ₽	1.0	1.0	14
2.₽	1.0	N+1 &	N+1 ¢
3₽	0 &	1.0	0 0
4 ₽	10	n₽	n↔
5 ₽	1.0	n₽	n₽
4	٠	<u>总步程</u> 数: →	3n + 2 @

```
Introduce statements to increment count at all appropriate points in
  4. (a)
            Program 1.32.
 void D(int *x, int n)
     int i = 1;
     do {
         x[i] += 2;
         i += 2;
     while (i \le n);
     i = 1:
     while (i <= (n/2))
         x[i] += x[i+1];
         i++;
Program 1.32: Example program
          Simplify the resulting program by eliminating statements. The
          simplified program should compute the same value for count as com-
          What is the exact value of count when the program terminates? You
     (c)
    (d) Obtain the step count for Program 1.32 using the frequency method.
```

## 注意:

1. 取整

#### 向下取整与向上取整

对任意实数 x,我们用 $\lfloor x \rfloor$ 表示小于或等于 x 的最大整数(读作"x 的向下取整"),并用 $\lceil x \rceil$ 表示大于或等于 x 的最小整数(读作"x 的向上取整")。对所有实数 x,

$$x - 1 < \lfloor x \rfloor \leqslant x \leqslant \lceil x \rceil < x + 1 \tag{3.3}$$

## 2. do while循环的程序步:

do内程序块和while语句执行的次数一样

# a: 添加程序步骤计数语句

```
void D(int * x, int n) {
2 int i = 1;
3 count++; //i赋值
4 do {
5 x[i] += 2;
6 i += 2;
7 count += 2; //循环体内语句程序步
8 count++; //while条件程序步
9 } while (i <= n);</pre>
10 //count++; 没有这一步
11 i = 1;
12 count++; //i赋值
13 while (i <= (n / 2)) {
14 count++; //while条件
15
16 x[i] += x[i + 1];
17 i++;
18
19 count += 2;//循环体内语句程序步
20
21 count++; //出循环的最后一个while条件
22 }
```

#### b: 简化计数程序

```
void D(int * x, int n) {
  int i = 1;
  do {
  i += 2;
  count += 3;
  } while (i <= n);
  i = 1;
  while (i <= (n / 2)) {
  i++;
}</pre>
```

```
10 count += 3;

11 }

12 count += 3;

13 }
```

# c: 程序的准确程序步

3n + 3

# d: 给程序的行编号如下:

```
1 void D(int * x, int n) {
2  int i = 1;
3  do {
4  x[i] += 2;
5  i += 2;
6  } while (i <= n);
7  i = 1;
8  while (i <= (n / 2)) {
9  x[i] += x[i + 1];
10  i++;
11  }
12 }</pre>
```

	レング	20 4 20	12. A. K. D
	71为奇	力为偶	<b>任意情况</b>
int 1=1;	1	1	1
dos	o o	7	0
X[i] +2 );	<u>nt </u>	7/2	127
1+=25	<u>n+ </u>	1/2	T <u>4</u> 7
Inohile (1) <=n);	1+1	4	TAT .
るープ	1		1
while ( i <= Lr/z))	<u>M-1</u> +1	1/2+	Ln/21+/
}	0	0	0
かしう)ナニ かしけり	<u>n - 1</u>	4	<u>n</u> LŽI
à tt)	<u>n- </u>	7	N L <u>S</u> J
	0	Ö	0
	3n+3	3n+3	37=7+31=31+3

# 5. Do Exercise 4 for function Transpose (Program 1.33).

#### a: 添加计算步程数语句

```
1 void Transpose(int ** a, int n) {
2    for (int i = 0; i < n - 1; i++) {
3        count++;
4    for (int j = i + 1; j < n; j++) {
5        count++;
6        swap(a[i][j], a[j][i]);
7        count++;
8    }
9    count++;
10    }
11    count++;
12 }</pre>
```

#### b:简化的计算步程数程序:

```
void Transpose(int ** a, int n) {
for (int i = 0; i < n - 1; i++) {
  count += 2;

for (int j = i + 1; j < n; j++) {
  count += 2;
  }
}

count += 2;
}

count += 2;
}
</pre>
```

## c: 总步程数

```
n^2 + n - 1
```

# d: 给程序行编号

```
1 void Transpose(int ** a, int n) {
2  for (int i = 0; i < n - 1; i++) {
3  for (int j = i + 1; j < n; j++) {
4  swap(a[i][j], a[j][i]);</pre>
```

```
5 }6 }7 }
```

程店	单层循环 步数	程序步
for ( ant a = 0; a < n - 1; a + + )	$i = 0 \sim n - 2$ ; $(n - 2 - 0 + 1) + 1 = n$	Λ
for lint j = i+1; G <n; o+r)<="" td=""><td>j=1+1~n-1; [n-1-1+1)+1]+1=n-1</td><td><math display="block">\sum_{i=0}^{n-2} (n-i) = \frac{n^2}{2} + \frac{n}{2} -  </math></td></n;>	j=1+1~n-1; [n-1-1+1)+1]+1=n-1	$\sum_{i=0}^{n-2} (n-i) = \frac{n^2}{2} + \frac{n}{2} -  $
swap(ati][j], atj][i]),	n-i-	$\frac{12}{12}(n-1) = \frac{n^2}{2} - \frac{1}{2}$
$\sqrt{n^2+n-1}$		

行步程数。 s/e↵ 頻度↵ 行↵ 2 ₽ 1₽ n₽ 3₽  $\frac{n^2}{2} + \frac{n}{2} - 1 =$  $\frac{n^2}{2} + \frac{n}{2} - 1$ 1₽ 4₽  $\frac{n^2}{2} - \frac{n}{2}$  $\frac{n^2}{2} - \frac{n}{2}$ 1₽  $n^2 + n - 1$ 总步程数: ↩

6. Do Exercise 4 for Program 1.34. This program multiplies a = a + b two  $a \times b$ 

```
void Multiply( int **a, int **b, int **c, int n) 

{
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
        {
            c[i][j] = 0;
            for (int k = 0; k < n; k++)
            c[i][j] += a[i][k] * b[k][j];
        }
}

Program 1.34: Square matrix multiplication
```

# a: 添加计算步程数语句

```
1 void Multiply(int ** a, int **b, int **c, int n) {
2 for (int i = 0; i < n; i++) {
3 count++; //for循环
4 for (int j = 0; j < n; j++) {
5 count++; //for循环
```

```
6 c[i][j] = 0;
7 count++; //赋值
8 for (int k = 0; k < n; k++) {
9 count++; //for循环
10 c[i][j] += a[i][k] * b[k][j];
11 count++; //计算
12 }
13 count++; //for循环最后一步
14 }
15 count++; //for循环最后一步
16 }
17 count++; //for循环最后一步
18 }
```

## b:简化的计算步程数程序:

```
void Multiply(int ** a, int **b, int **c, int n) {
for (int i = 0; i < n; i++) {
  count += 2;
  for (int j = 0; j < n; j++) {
  count += 3;
  for (int k = 0; k < n; k++) {
  count += 2;
  }
}

count += 2;
}</pre>
```

#### c: 总步程数

# $2n^3 + 3n^2 + 2n + 1$

#### d: 给程序行编号

```
void Multiply(int ** a, int **b, int **c, int n) {
for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++) {
    c[i][j] = 0;
  for (int k = 0; k < n; k++) {
    c[i][j] += a[i][k] * b[k][j];
  }
}
}
</pre>
```

行↵	s/e₽	频度↩	行 <u>步程数</u> 。
2 &	1.₽	n + 1 &	n + 1 ø
3 &	1.₽	<u>n(</u> n + 1)₽	<u>n(</u> n + 1)₽
4.	1.	$n^2$ $\varphi$	$n^2$ $\varphi$
5.	1₽	$n^2(\mathbf{n+1})$ $\varphi$	$n^2(\boldsymbol{n+1})$
6₽	1.0	$n^3$ $_{arphi}$	$n^{oldsymbol{3}_{arphi}}$
<b>.</b>	₽3	<u>总步程</u> 数: ↩	$2n^3 + 3n^2 + 2n + 1$

- 7. (a) Do Exercise 4 for Program 1.35. This program multiplies two matrices a and b where a is an  $m \times n$  matrix and b is an  $n \times p$  matrix.
  - (b) Under what conditions will it be profitable to interchange the two outermost for loops?

```
void Multiply(int **a, int **b, int **c, int m, int n, int p)

{

for (int i = 0; i < m; i++)

for (int j = 0; j < p; j++)

{

c[i][j] = 0;

for (int k = 0; k < n; k++)

c[i][j] += a[i][k] * b[k][j];
}

Program 1.35: Matrix multiplication
```

#### 问题一:

## a: 添加计算步程数语句

```
void Multiply(int **a, int **b, int **c, int m, int n, int p) {
for (int i = 0; i < m; i++) {
  count++;

for (int j = 0; j < p; j++) {
  count++;

  c[i][j] = 0;
  count++;

for (int k = 0; k < n; k++) {</pre>
```

```
9  count++;
10  c[i][j] += a[i][k] * b[k][j];
11  count++;
12  }
13  count++;
14  }
15  count++;
16  }
17  count++;
18 }
```

# b:简化的计算步程数程序:

```
void Multiply(int **a, int **b, int **c, int m, int n, int p) {
for (int i = 0; i < m; i++) {
  count += 2;

  for (int j = 0; j < p; j++) {
  count += 3;

  for (int k = 0; k < n; k++) {
  count += 2;

  }
}

count++;
}</pre>
```

#### c: 总步程数

# 2mpn+3mp+2m+1

#### d: 给程序行编号

```
void Multiply(int **a, int **b, int **c, int m, int n, int p) {

for (int i = 0; i < m; i++) {

for (int j = 0; j < p; j++) {

c[i][j] = 0;

for (int k = 0; k < n; k++) {

c[i][j] += a[i][k] * b[k][j];

}

}

}

}

}

}
</pre>
```

行₽	s/e₽	频度₽	行 <u>步程数</u> ₽
2 &	10	m + 1₽	m + 10
3 ₽	1.0	m(p + 1) $\varphi$	m(p + 1) $\varphi$
4.₽	1.0	mp₽	mp₽
5₽	1.0	mp(n + 1).	mp(n + 1) $_{\circ}$
6 ₽	1.0	mpn₽	mpn₽
÷.	P	<u>总步程</u> 数: ↩	2mpn+3mp+2m+1₽

#### 问题二:

计算得出程序的总步程数是2mpn+3mp+2m+1,外两层循环涉及到的变量为m和p,交换外两次循环不会影响2mpn+3mp的大小,但是会影响2m的大小,因此当m > p时,交换外两层循环的步程数为2mpn+3mp+2p+1 < 2mpn+3mp+2m+1,因此此时交换能提高效率。

#### 程序题:

2. Given n Boolean variables  $x_1, \dots, x_n$  we wish to print all possible combinations of truth values they can assume. For instance, if n = 2, there are four possibilities: true, true; true, false; false, true; false, false. Write a C++ program to accomplish this and do a frequency count.

#### 方法:

1: 递归

2: 位运算

以下的程序没有用上面的任何一种方法,自己随便写的,同学们之后写程序的时候,可以多用STL库中包装好的东西,比如<vector>,<stack>,<queue>,<string>,<unordered\_map>,<list>等等代替使用char\*, int[],int\*等东西。

```
1 #include <vector>
2 #include <iostream>
3 using namespace std;
4
5 vector<vector<string>> solution(int n) {
```

```
if (n <= 0) {
   cout << "n必须大于等于1。\n";
7
   exit(0);
8
9
10
    //当n等于1时的结果数组
11
12
    vector<vector<string>> res;
   res.push_back({ "false" });
13
    res.push_back({ "true" });
14
15
    if (n == 1) {
16
17
    return res;
18
    }
    //将数组中的每一个值分别追加true和false,没循环一次,数组扩大一倍
19
    //例如{ {"false"} ,{ "true"}}会变为
20
    //{ {"false", "true"} ,{ "true", "true", "false"},{ "true",
21
"false"}}
    for (int i = 2; i <= n; i++) {
    int size = res.size();
   for (int j = 0; j < size; j++) {
24
25
   vector<string> a = res[j];
26
    a.push_back("true");
   res.push_back(a);
27
    res[j].push_back("false");
28
   }
29
30
   }
   return res;
31
  }
32
33
34 //int main() {
35 // vector<vector<string>> temp = solution(9);
36 // int raw = temp.size();
37 // int col = temp[0].size();
38 // for (int i = 0; i < raw; i++) {
39 // cout << "( " << temp[i][0].c_str();</pre>
40 // for (int j = 1; j < col; j++) {
41 // cout << " , " << temp[i][j].c_str();
42 // }
43 // cout << " ) " << endl;
44 // }
```

```
45 // cout << "有" << temp.size() << "中组合。\n";
46 //}
```

14. If S is a set of n elements, the powerset of S is the set of all possible subsets of S. For example, if S = (a,b,c), then powerset  $(S) = \{(), (a), (b), (c), (a,b), (a,c), (b,c), (a,b,c)\}$ . Write a recursive function to compute

# powerset (S).

```
1 #include <vector>
2 #include <iostream>
3 using namespace std;
5 vector<vector<string>> util(vector<string> S, int start, int end) {
  vector<vector<string>> res;
   if (start > end) {
   res.push_back({});
   return res;
10
   }
   res = util(S, start, end - 1);
11
    int size = res.size();
   for (int i = 0; i < size; i++) {
    vector<string> temp = res[i];
14
    temp.push_back(S[end]);
15
16
    res.push_back(temp);
    }
17
    return res;
18
   }
19
20
21
   vector<vector<string>> solution(vector<string> S) {
    return util(S, 0, S.size() - 1);
  }
24
25
26 //
27 //int main() {
28 // vector<string> S{ "a","b","c" ,"d"};
29 // vector<vector<string>> res = solution(S);
30 // int size = res.size();
31 // cout << "{ ";
```

```
32 //
33 // for (int i = 0; i < size; i++) {
34 //
35 // int size1 = res[i].size();
36 // if (size1 == 0)
37 // cout << "( )";
38 // else {
39 // cout << "( " << res[i][0].c_str();</pre>
40 // for (int j = 1; j < res[i].size(); j++) {
41 // cout << " , " << res[i][j].c_str();
42 // }
43 // cout << " )";
44 // }
45 // if( i != size - 1)
46 // cout << " , ";
47 // }
48 // cout << " }";
49 // cout << endl;
50 // cout << "超集中共有" << size << "个元素" << endl;
51 //}
```