Program Transformation and Analysis Assignment 1

Benjamin Brandt Ohrt, zpn492

April 28, 2019

1 Introduction

This the first of five weekly assignment in the course Program Transformation and Analysis (PAT) at Copenhagen University. The course professor is Robert Glück. The course is held in block 4, 2019.

In this assignment the focus is on Programs as Data Objects and Reversible Computing.

Before we start on the assignment we introduce an inverse interpreter *invint* and a program inverter *invtrans*. For this introduction we need to define the following: Let p be a program written in a language L, x an input into p and y an output from p. We define p^{-1} as the inverse of p such that givin y as input to p^{-1} produces x.

Inverse Interpreter

An inverse interpreter can be defined with the following definition[1, p. 271]:

$$[[invint]][p, y] = x \tag{1}$$

Program Inverter

An program inverter can defined with the following definition[1, p. 271]:

$$[[invtrans]][p] = p^{-1}$$

 $[[p^{-1}]][y] = x$ (2)

2 Assignment

Exercise 1

Define a trivial program inverter givin an inverse interpreter.

$$[[invtrans']] = \lambda p.\lambda y.[[invint]][p, y]$$
(3)

Exercise 2

Define a trivial inverse interpreter givin a program inverter.

$$[[invint']] = \lambda[p, y].[[Lint]][[[[invtrans]]p], y] \tag{4}$$

Where Lint is an L-interpreter.

Exercise 3

Run the Janus-program fib forward section 3 by hand, where n=4, x1=0 and x2=0. Trace the store (n, x1, x2) after each statement.

| Recursion | State | n | x1 | x2 | branch |
|-----------|-------------|---|----|----|----------|
| 0 | initial | 4 | 0 | 0 | else |
| 1 | initial | 3 | 0 | 0 | else |
| 2 | initial | 2 | 0 | 0 | else |
| 3 | initial | 1 | 0 | 0 | else |
| 4 | initial | 0 | 0 | 0 | then |
| 4 | terminating | 0 | 1 | 1 | fi:true |
| 3 | terminating | 0 | 1 | 2 | fi:false |
| 2 | terminating | 0 | 2 | 3 | fi:false |
| 1 | terminating | 0 | 3 | 5 | fi:false |
| 0 | terminating | 0 | 5 | 8 | fi:false |

Exercise 4

Run the Janus-program fib backward section 3 by hand where n = 0, x1 = 5 and x2 = 8. Trace the store (n, x1, x2) after each statement.

| Recursion | State | n | x 1 | x2 | branch |
|-----------|-------------|---|------------|----|----------|
| 0 | initial | 0 | 5 | 8 | else |
| 1 | initial | 0 | 3 | 5 | else |
| 2 | initial | 0 | 2 | 3 | else |
| 3 | initial | 0 | 1 | 2 | else |
| 4 | initial | 0 | 1 | 1 | then |
| 4 | terminating | 0 | 0 | 0 | fi:true |
| 3 | terminating | 1 | 0 | 0 | fi:false |
| 2 | terminating | 2 | 0 | 0 | fi:false |
| 1 | terminating | 3 | 0 | 0 | fi:false |
| 0 | terminating | 4 | 0 | 0 | fi:false |

Exercise 5

Describe formally but concisely, how you interpreted each statement (+=, -=, <=>, if, call) in the Janus-program fib in the fwd and bwd direction.

Switching sourcecode between a Janus-program p and its inverse p^{-1} can be done by interpreting statements using above table. After the interpretation you just swap if and fi and reverse the order of operations within the then and else statement. Ex. a list l of operations in p where $l := \{x1+=1; x2+=1\}$ is a list l' in p^{-1} where $l' := \{x2+=1; x1+=1\}$

3 Janus-programs

Fibonacci forward

Fibonacci backward

fi n=0

References

[1] Mogensen, Torben, Schmidt, David, Sudborough, I. Hal (Eds.). The Essence of Computation - Complexity, Analysis, Transformation. Springer-Verlag, Berlin Heidelberg, 2002.