

## 8.1

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$r_1 \cap r_2 = A$  and  $A \rightarrow r_1$ , which means  $r_1 \cup r_2 \rightarrow r_1$ , and this decomposition is lossless.

## 8.13

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There is no such a dependency  $B \rightarrow \alpha$  in  $F_1^+$  and a dependency  $\alpha \rightarrow D$  in  $F_2^+$ , so the dependency  $B \rightarrow D$  cannot be preserved by this decomposition.

## 8.19

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$(A, B, C, E)$

$(B, D)$

## 8.20

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$(A, B, C, E)$

$(B, D)$

A relation in BCNF is also in 3NF.

## 8.29

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**a**

$B \rightarrow D \rightarrow AD \rightarrow ABCD \rightarrow ABCDE$

That is,  $B^+ = ABCDE$

**b**

$A \rightarrow ABCD$  (Augmentation with A)

$ABCD \rightarrow ABCDE$  (Augmentation with ABCD)

$A \rightarrow ABCDE$  (Transitivity)

$AF \rightarrow ABCDEF$  (Augmentation)

**c**

Remove extraneous  $D$ :

$A \rightarrow BC$

$$BC \rightarrow E$$

$$B \rightarrow D$$

$$D \rightarrow A$$

Then remove extraneous  $C$  and combine:

$$A \rightarrow BC$$

$$B \rightarrow DE$$

$$D \rightarrow A$$

**d**

Use the algorithm:

$$(A, B, C)$$

$$(B, D, E)$$

$$(D, A)$$

$$(B, F)$$

**e**

Use the algorithm:

$$(A, B, C, D)$$

$$(B, E)$$

$$(B, F) \text{ (primary key)}$$

**f**

No. Unless we can use  $F^{+}$  inferred from the canonical cover.