

Theory of Computation, Fall 2021

Assignment 7 (Due December 17 Friday 9:35am)

1. Σ^* is countable
2. subset of a countable set is also countable
3. every language is a subset of Σ^*

Q1. Prove that every language is countable. You can use any theorem that we've proved in class.

Q2. Prove that there is an undecidable subset of $\{1\}^*$. $\{0,1\}^* \not\rightarrow \{1\}^*$

Q3. We have already proved that the following language is not recursive. Show that it is recursively enumerable.

1. for $i=1,2,3,\dots$
run M on each for i steps
 2. halt if M halts within i steps
- $A = \{ \langle M \rangle : M \text{ is a Turing machine that halts on some string} \}$

(Hint: You may use the fact that Turing machines are able to enumerate all the strings in Σ^* as s_1, s_2, s_3, \dots)

Q4. Let A and B be two languages. Let f be a reduction from A to B . $w \in A$ iff $f(w) \in B$

- (a) Suppose that you have a Turing machine M_B that semidecides B . Try to construct a Turing machine M_A that semidecides A .
 1. run M_B on $f(w)$
 2. halt if halt
- (b) What conclusion can you draw from (a)?

Q5. Consider the following language. Show that the following language is not recursive by showing a reduction from H to L .

$A = \{ \langle M \rangle : M \text{ is a TM such that for any string } w, M \text{ halts on a string } w^R \text{ whenever it halts on } w \}.$

- (a) Show that the A is not recursive by showing a reduction from H to A .
- (b) Use Rice's theorem to prove that A is not recursive.

L is RE and both w^R and w are in L

- on input u
1. if 10, accept
 2. if 01, reject
 3. if 01, run M on w

References

- [1] Sipser M.. Introduction to the Theory of Computation. CENGAGE Learning (2013)
- [2] Lewis H., Papadimitriou C.. Elements of the Theory of Computation. Prentice-Hall (1998)