

P510 7

Let $f(n)$ denote the number of bit strings of length n that contain a pair of consecutive 0s.

a) By looking at the last two digits of bit strings of length $n-1$ we can conclude that $f(n) = f(n-1) + f(n-2) + 2^{n-2}$ for $n \geq 2$.

b) $f(0)=0$ and $f(1)=0$.

c) From the recurrence relation we can have $f(4)=8$, $f(5)=19$, $f(6)=43$, and finally $f(7)=94$.

So there are 94 bit strings of length 7 that contain a pair of consecutive 0s.

P511 17

Let $f(n)$ denote the number of ternary strings of length n that do not contain consecutive symbols that are the same.

a) By looking at the last two digits of bit strings of length $n-1$ we can conclude that $f(n)=2*f(n-1)$.

b) $f(1)=3$.

c) We can know that $f(n)=3 \cdot 2^{n-1}$ and so $f(6)=96$.

So there are 96 ternary strings of length 6 that do not contain consecutive symbols that are the same.

P525 13

We can have the characteristic equation $r^3-7r-6=0$ and we can have the roots $r_1=-1$, $r_2=3$ and $r_3=-2$.

Assume that $a_n=a \cdot (-1)^n + b \cdot 3^n + c \cdot (-2)^n$.

With the initial conditions we can have: $a+b+c=9$, $-a+3b-2c=10$ and $a+9b+4c=32$. Hence we can have $a=8$, $b=4$ and $c=-3$.

So $a_n=8 \cdot (-1)^n + 4 \cdot 3^n - 3 \cdot (-2)^n$.

P526 35

The roots of the characteristic equation $r^2=4r-3$ are $r_1=1$, $r_2=3$.

For 2^n from $k \cdot 2^n = 4k \cdot 2^{n-1} - 3k \cdot 2^n - 2 + 2^n$ we have $k=-4$. For $n+3$ from $a \cdot n^2 + b \cdot n = 4a \cdot (n-1)^2 + 4b \cdot (n-1) - 3a \cdot (n-2)^2 - 3b \cdot (n-2) + n + 3$ we have $a=-0.25$, $b=-2.5$. So the particular solution is $-4 \cdot 2^n - \frac{n^2}{4} - \frac{5n}{2}$.

Then we have $1=a+b-4$ and $4=a+3b-43/4$, and $a=1/8$, $b=39/8$.

So the solution is $\frac{1}{8} + \frac{39}{8} \cdot 3^n - 4 \cdot 2^n - \frac{n^2}{4} - \frac{5n}{2}$.

P535 21

a) We simply have $f(4)=3$, and then $f(16)=7$.

So $f(16)=7$.

b) Given that n is a perfect square and $f(2)=1$, we can assume that $n = 2^{2^m}$. From the recurrence relation we can have $f(n) = 2^m f\left(2^{\frac{n}{2^m}}\right) + 2^m - 1 = 2^m f(2) + 2^m - 1$, and $2^m = \log n$.

So $f(n) = O(\log n)$.