

A NEW CLASS OF ASYNCHRONOUS ITERATIVE ALGORITHMS WITH ORDER INTERVALS

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ABSTRACT. This paper deals with a new class of parallel asynchronous iterative algorithms for the solution of nonlinear systems of equations. The main feature of the new class of methods presented here is the possibility of flexible communication between processors. In particular partial updates can be exchanged. Approximation of the associated fixed point mapping is also considered. A detailed convergence study is presented. A connection with the Schwarz alternating method is made for the solution of nonlinear boundary value problems. Computational results on a shared memory multiprocessor IBM 3090 are briefly presented.

1. INTRODUCTION

In the past few years a large number of parallel computing methods have been proposed for the solution of large scale numerical problems (see Bertsekas and Tsitsiklis [7], Hockney and Jesshope [22], Ortega [30], and Schendel [34]). Among parallel algorithms asynchronous iterative methods have received a considerable amount of attention. Asynchronous algorithmic models were introduced by Chazan and Miranker (see [8]) for the solution of linear systems of equations. A necessary and sufficient condition of convergence is given in [8] (see also [2]).

In the nonlinear case, the convergence properties of parallel asynchronous iterative algorithms are now well understood. Many authors have concentrated on fixed point problems. In particular Miellou (see [24] and [26]) and Baudet (see [3]) have shown contraction properties using a vectorial norm (see also [33]). El Tarazi [15] has shown a contraction property using an appropriate scalar norm.

A complementary approach dealing with the properties of the nonlinear equations operator was considered by Miellou and Spiteri [29], Giraud Spiteri [20], and Spiteri [35] (see also [36]); the authors quoted above have proposed sufficient conditions of convergence for asynchronous iterations. The reader is referred to Giraud and Spiteri [21] for an implementation of asynchronous algorithms.

Some authors have also made use of the discrete maximum principle. Monotone sequences of vectors are generated in this case. Miellou [25] has given a sufficient condition of convergence for asynchronous iterations under partial ordering. Other contributions are due to El Tarazi [16], Miellou [27], El Baz (see [10], [12], and

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[13]), Frommer (see [18] and [19]). Reference is also made to Cousot [9] for a study related to the proof of programs using fixed point techniques.

Finally, we note that Bertsekas has proposed a different approach based on the definition of a sequence of nested level sets (see [5] and [7], see also [4], [6], and [11]). Perturbation of fixed point iterative algorithms has been studied in this context by Miellou, Cortey-Dumont, and Boulbrachène (see [28]).

In this paper we consider partial ordering techniques and propose new parallel asynchronous iterative methods which are more general than the ones quoted above. An original feature of the new parallel iterative methods presented here is flexible communication between processors. More precisely, updating phases can use data which are issued from computations in progress. This leads to a better coupling between communication and computation. In contrast with previous studies (see [3], [8], and [25]), the methods presented here permit one to consider block iterative schemes with delayed access to block subvectors, the components of which are relative to different iteration numbers or intermediate computations not explicitly labelled by an iteration number. Another important feature of this paper is the use of mappings which approximate the solution. A detailed convergence study is presented. The new model of parallel asynchronous iterations proposed in this paper can be applied with success to subdomain methods for the solution of partial differential equations and gradient type methods for network flow problems. Computational studies using shared memory and distributed memory multiprocessors are presented in detail in two complementary papers (see [14] and [37]). Computational experiences have shown the practical interest of the proposed parallel asynchronous iterative methods. A brief presentation of numerical results for the solution of partial differential equations via subdomain methods is made at the end of the paper.

The paper contains five sections. In Section 2 we present the main problem and give some background material. Section 3 deals with the new class of parallel asynchronous iterative methods. Two kinds of approximations of the fixed point mapping are presented in Section 4. In Section 5 we consider the application of the theoretical study to the Schwarz alternating method for the numerical solution of boundary value problems; computational experience using a shared memory multiprocessor IBM 3090 is also presented in this section.

2. PROBLEM STATEMENT

2.1. Notations, background material. Let n be a positive integer and assume that

(2.1) a is a continuous mapping of R^n into R^n ,

(2.2) a is a surjective M -function.

We recall that the mapping a is an M -function (see [31] and [32]) if a is off-diagonally antitone i.e. for any $x \in R^n$, the functions defined as follows:

$$(2.3) \quad \begin{cases} a_{lk} : \{t \in R \mid x + te_k \in R^n\} \rightarrow R, \\ a_{lk}(t) = a_l(x + te_k), \quad l \neq k, l, k = 1, \dots, n, \end{cases}$$

are monotone decreasing, where $e_k \in R^n$ $k = 1, \dots, n$, are the unit canonical basis vectors and furthermore a is inverse isotone i.e.

$$(2.4) \quad a(x) \leq a(y) \quad \text{for any } x, y \in R^n \text{ implies that } x \leq y.$$

In the sequel we will make use of the following notation: for any sequence $\{x^k\} \subset R^n$, $\lim_{k \rightarrow \infty} x^k = +\infty$ ($\lim_{k \rightarrow \infty} x^k = -\infty$) if $\lim_{k \rightarrow \infty} x_i^k = +\infty$ ($\lim_{k \rightarrow \infty} x_i^k = -\infty$) for at least one index i .

Definition 1. The mapping $a: R^n \rightarrow R^n$ is order-coercive if for any sequence $\{x^k\} \subset R^n$, such that $x^k \leq x^{k+1}$, $k = 0, 1, \dots$, and $\lim_{k \rightarrow \infty} x^k = +\infty$, we have $\lim_{k \rightarrow \infty} a(x^k) = +\infty$ and for any sequence $\{x^k\} \subset R^n$, such that $x^k \geq x^{k+1}$, $k = 0, 1, \dots$, and $\lim_{k \rightarrow \infty} x^k = -\infty$, we have $\lim_{k \rightarrow \infty} a(x^k) = -\infty$.

We can characterize the surjectivity of M -functions by means of the order-coercivity concept (see [32, Theorem 3.7]).

Theorem 1. Let $a: R^n \rightarrow R^n$ be a continuous M -function. Then a is surjective if and only if a is order-coercive.

We consider now the solution of the following system of equations

$$(2.5) \quad a(y) = 0.$$

Under the above assumptions problem (2.5) has a unique solution.

2.2. Subproblem decomposition and associated fixed point mapping. Let $E = R^n$ and consider the following splitting of E : $E = \prod_{i=1}^{\alpha} E_i$, where α is a positive integer, $E_i = R^{n_i}$, and $\sum_{i=1}^{\alpha} n_i = n$. Each subspace E_i is endowed with the natural (or componentwise) partial ordering associated with the cone $K_i = R_+^{n_i}$ of vectors with non-negative components in R^{n_i} . Let $w \in E$. Consider the following block decomposition of w :

$$w = \{w_1, \dots, w_i, \dots, w_{\alpha}\} \in \prod_{i=1}^{\alpha} E_i,$$

and the following block decomposition of a :

$$a(w) = \{a_1(w), \dots, a_i(w), \dots, a_{\alpha}(w)\} \in \prod_{i=1}^{\alpha} E_i.$$

In the sequel $a_i(w_1, \dots, w_{i-1}, y_i, w_{i+1}, \dots, w_{\alpha})$ will also be denoted by $a_i(y_i; w)$. Let assumptions (2.1) and (2.2) hold. Then

$$(2.6) \quad \begin{cases} \text{for all } i \in \{1, \dots, \alpha\} \text{ and all } w \in E, \text{ the mapping: } y_i \rightarrow a_i(y_i, w), \\ \text{is a continuous surjective } M\text{-function of } R^{n_i} \text{ onto } R^{n_i} \end{cases}$$

(see [32, Theorem 3.5]). Moreover it follows from the assumption (2.2) that

$$(2.7) \quad \begin{cases} \text{for all } w \in E \text{ and } i \in \{1, \dots, \alpha\}, \text{ the problem } a_i(z_i; w) = 0, \\ \text{has a unique solution } z_i. \end{cases}$$

Thus the mapping F from R^n onto R^n such that

$$(2.8) \quad F(w) = z = \{z_1, \dots, z_i, \dots, z_{\alpha}\},$$

with z_i from (2.7) is well-defined. It is a fixed point formulation for problem (2.5) since

$$F(y) = y \Leftrightarrow a(y) = 0.$$

Proposition 1. F is isotone on E .

Proof. Let w and $w' \in E$ such that $w \leq w'$. Let $z = F(w)$ and $z' = F(w')$, where z'_i satisfies $a_i(z'_i; w') = 0$, $i \in \{1, \dots, \alpha\}$. By the off-diagonal antitonicity of a we have

$$a_i(z'_i; w) \geq 0, \quad i \in \{1, \dots, \alpha\}.$$

Thus for all $i \in \{1, \dots, \alpha\}$, we have

$$a_i(z'_i; w) \geq a_i(z_i; w) = 0,$$

and by the inverse isotonicity of $a_i(z_i; w)$, we have for all $i \in \{1, \dots, \alpha\}$

$$z_i \leq z'_i. \quad \square$$

2.3. **Supersolutions and a -supersolutions.**

Definition 2. A vector $y \in E$ is a supersolution (a subsolution) if $F(y) \leq y$ ($F(y) \geq y$).

Definition 3. A vector $y \in R^n_+$ is an a -supersolution (an a -subsolution) if $a(y) \geq 0$ ($a(y) \leq 0$).

Proposition 2. *If y is an a -supersolution, then y is a supersolution.*

Proof. It follows from (2.7) and (2.8) that for all $i \in \{1, \dots, \alpha\}$, we have

$$a_i(F_i(y), y) = 0 \leq a_i(y).$$

By the inverse isotonicity of a_i we have

$$F_i(y) \leq y_i. \quad \square$$

Remark 1. In general the converse is false. For example, consider the case where a is a linear mapping. Let A be the associated matrix, then it follows from assumption (2.2) that A is an M -matrix. We have

$$a_i(y) = \sum_{j=1}^{\alpha} A_{ij}y_j,$$

where $\{A_{ij}\}$ results from the block decomposition of A . The block diagonal matrices A_{ii} , are necessarily M -matrices. We have also

$$y_i - F_i(y) = y_i + \sum_{j \neq i} A_{ii}^{-1} A_{ij}y_j.$$

In particular, we note that $y_i - F_i(y) > 0$, does not imply $A_{ii}(y_i - F_i(y)) = a_i(y) > 0$. Indeed, consider the case where $E_i = R^2$ and $A_{ii} = \begin{bmatrix} 1 & -0.5 \\ 0.5 & 1 \end{bmatrix}$; then for $y_i - F_i(y) = \{\frac{1}{10}, 1\}^T$, we find that

$$A_{ii}(y_i - F_i(y)) = a_i(y) = \left\{ -\frac{2}{5}, \frac{19}{20} \right\}^T.$$

However in the linear case if F corresponds to the Jacobi mapping associated with a point decomposition of the M -matrix A (i.e. $\alpha = n$ and $E_j = R$, $j \in \{1, \dots, n\}$), then $w \geq F(w)$ is equivalent to $Aw \geq 0$.

Proposition 3. *Let a be an M -function and $w \in R^n$ an a -supersolution. Let $s \subset \{1, \dots, n\}$ and $v \in R^n$, be such that $v \leq w$, $a_i(v_i; w) \geq 0$, for all $i \in s$ and $v_i = w_i$, for all $i \notin s$, then v is an a -supersolution.*

Proof. For all $i \in \{1, \dots, n\}$ we have $a_i(v_i; w) \geq 0$, since w is an a -supersolution. Thus by the off-diagonal antitonicity of a we find that

$$a_i(v_i; v) \geq a_i(v_i; w) \geq 0. \quad \square$$

Proposition 4. *Let a be an M -function, v and w two a -supersolutions. Then $v \wedge w = \{\min\{v_1, w_1\}, \dots, \min\{v_r, w_r\}, \dots, \min\{v_n, w_n\}\}$ is an a -supersolution.*

Proof. We use the following notation: $a(v \wedge w) = \{a_1(v \wedge w), \dots, a_r(v \wedge w), \dots, a_n(v \wedge w)\}$, where a_r corresponds here to a point decomposition of the mapping a . For all $r \in \{1, \dots, n\}$, we have

$$a_r(v \wedge w) = a_r(v_r \wedge w_r; v \wedge w).$$

By the off-diagonal antitonicity of a , we have for all $r \in \{1, \dots, n\}$,

$$a_r(w_r; v \wedge w) \geq a_r(w) \geq 0, \text{ and } a_r(v_r; v \wedge w) \geq a_r(v) \geq 0.$$

Thus, $a_r(v \wedge w) \geq 0$. \square

3. ASYNCHRONOUS ITERATIONS WITH ORDER INTERVALS

In Section 2 we have considered a fixed point mapping associated with the exact solution of the subproblems (2.7). In this section we study fixed point mappings associated with approximate solutions. Such mappings are introduced for the solution of problem (2.5) via parallel asynchronous iterative methods with flexible communication.

3.1. Supermappings and a -supermappings associated with F . Throughout the paper we adopt the following notation for order intervals. Let $x_i, y_i \in E_i$, be such that $x_i \leq y_i$; then an order interval in E_i is denoted by

$$\langle x_i, y_i \rangle_i = \{z_i \in E_i \mid x_i \leq z_i \leq y_i\}.$$

Similarly, let $x, y \in E$, be such that $x \leq y$; then an order interval in E is denoted by

$$\langle x, y \rangle = \{z \in E \mid x \leq z \leq y\}.$$

Definition 4. Let a be an M -function. Then the mapping F^a with components F_i^a is an a -supermapping associated with F if for all $i \in \{1, \dots, \alpha\}$ and y elements of the domain of definition of $F_i^a: \{y \in E \mid a_i(y) \geq 0\}$, we have $F_i^a(y) \leq y_i$, $a_i(F_i^a(y), y) \geq 0$, and $F_i^a(y) \neq y_i$ if $F_i(y) \neq y_i$.

Definition 5. Let a be an M -function. Then the mapping F^a with components F_i^a is a supermapping associated with F if for all $i \in \{1, \dots, \alpha\}$ and y element of the domain of definition of $F_i^a: \{y \in E \mid a_i(y) \geq 0\}$, we have $F_i^a(y) \in \langle F_i(y), y_i \rangle_i$ and $F_i^a(y) \neq y_i$ if $F_i(y) \neq y_i$.

Proposition 5. *If F^a is an a -supermapping associated with F , then F^a is a supermapping associated with F .*

Proof. From (2.7), (2.8), and Definition 4 it follows that

$$a_i(F_i(y); y) = 0 \leq a_i(F_i^a(y); y).$$

Then

$$F_i(y) \leq F_i^a(y),$$

since a_i is inverse isotone. \square

Remark 2. The converse of Proposition 5 is in general false. This result can be shown using an argument similar to the one developed in Remark 1.

Remark 3. We can analogously define submappings and α -submappings associated with F by reversing the inequalities.

3.2. A new class of parallel asynchronous iterations. In this section we define a general class of iterative methods which contains in particular sequential relaxation methods like Jacobi or Gauss-Seidel and parallel iterative methods whereby components are updated simultaneously by concurrent processors without any order nor synchronization. The new class of asynchronous iterations presented here allows also flexible communication between processors. We introduce first the concepts of steering and sequence of delays.

Definition 6. A steering \mathcal{S} of components of the iteration vector, is a sequence $\{s(p)\}$ where $p \in N$, the set of natural integers, such that

$$(3.1) \quad \text{for all } p \in N, \quad s(p) \in \{1, \dots, \alpha\}.$$

Definition 7. A sequence $\{r(p)\}$ (or \mathcal{R}) of maximal delays in the access to the updated components of the iteration vector, is defined by

$$r(p) = \{r_1(p), \dots, r_i(p), \dots, r_\alpha(p)\} \in N^\alpha \quad \text{for all } p \in N,$$

where for all $i \in \{1, \dots, \alpha\}$ and $p \in N$, $0 \leq r_i(p) \leq p$ and $r_i(p) = 0$ if $i = s(p)$.

We note that the latest condition has been introduced in order to simplify the analysis of the algorithms. This assumption is more restrictive than the one used in the classical model, however it is natural since it implies that each processor has access to its own working space without delay. In particular, this condition has been used in the framework of interval arithmetic (see [19]).

We introduce the functions $\rho_i: N \rightarrow N$, $i \in \{1, \dots, \alpha\}$, defined by $\rho_i(p) = p - r_i(p)$ and satisfying for all $i \in \{1, \dots, \alpha\}$ and $p \in N$:

$$(3.2) \quad 0 \leq \rho_i(p) \leq p,$$

$$(3.3) \quad \rho_i(p) = p, \quad \text{if } i = s(p).$$

In the sequel, we will denote by $\rho(p)$ the mapping from N^α into itself, with components $\rho_i(p)$, $i \in \{1, \dots, \alpha\}$. We introduce now the sequences $\{K_i^p\}$, $p \in N$, defined by

$$(3.4) \quad K_i^p = \{k \in N \mid s(k) = i, 0 \leq k < p\}, \quad i \in \{1, \dots, \alpha\}, p \in N.$$

The elements of the set K_i^p correspond to the iteration numbers between 0 and p at which the i th block component of the iteration vector is updated. We note that

$$(3.5) \quad \{K_i^p\} \text{ is a denumerable sequence of finite elements of the set of parts of } N.$$

Moreover $\{K_i^p\}$ is nested:

$$(3.6) \quad K_i^p \subseteq K_i^{p+1}.$$

Definition 8. Asynchronous iterations with order intervals associated with the α -supermapping F^α are sequences $\{y^p\}$ defined recursively as follows. For all $p \in N$,

$$(3.7) \quad \begin{cases} y_i^{p+1} = F_i^\alpha(\tilde{y}^p), & \text{if } i = s(p), \\ y_i^{p+1} = y_i^p, & \text{if } i \neq s(p), \end{cases}$$

where \tilde{y}^p is implicitly defined by

$$(3.8) \quad \tilde{y}^0 = y^0, \text{ an } a\text{-supersolution}$$

and

$$(3.9) \quad \tilde{y}^p \in \langle y^p, y^{\rho(p)} \wedge \tilde{y}^q \rangle, \quad \text{if } p \geq 1,$$

where the block-components of the vector $y^{\rho(p)} \in E$ are the subvectors $y_i^{\rho(p)}$, $i \in \{1, \dots, \alpha\}$, and $q = \max_{k \in K_{s(p)}^p} k$.

We note that $\tilde{y}^p \in \langle y^p, y^{\rho(p)} \rangle$ if $K_i^p = \emptyset$, $i = 1, \dots, \alpha$, where \emptyset denotes the empty set.

Remark 4. Asynchronous iterations defined by (3.7)–(3.9) are general iterative methods whereby iterations are carried out in parallel by α processors in arbitrary order and possibly without any synchronization. Flexible communication between processors is the main feature of this class of algorithms. Indeed we note that the block components \tilde{y}_j^p of the iteration vector used at the time of the updating of y_i^{p+1} are taken anywhere in the order interval $\langle y_j^p, y_j^{\rho(p)} \wedge \tilde{y}_j^q \rangle_j$, where \tilde{y}_j^q denotes the value used at the time of the last updating of the block-component y_i and $y_j^{\rho(p)}$ is related to the non-deterministic behaviour of the iterative scheme. The introduction of order intervals permits us to take into account data coming from computations which are in progress and which are not explicitly labelled by an iteration number. So, there is a better coupling between communication and computation and we may expect a faster convergence. Moreover each communication of a block subvector is not necessarily associated with a given iteration number. The exchanged values of the components of each block subvector can be relative to different iteration numbers. We note that the model presented in this study permits us to analyse asynchronous block iterative methods directly in the block decomposition framework. So, it is not necessary to introduce a model defined with respect to individual components that would be artificial and cumbersome in the context of this study. The use of a -supermappings which approximate the solution of subproblems is another important feature of the parallel methods presented here.

Proposition 6. *Let a be a continuous surjective M -function, F the fixed point mapping defined by (2.7) and (2.8), F^a an a -supermapping associated with F , $y^0 \in E$ an a -supersolution, \mathcal{S} , and \mathcal{R} a steering and a sequence of delays, respectively. Then the asynchronous iteration $\{y^p\}$ given by (3.7)–(3.9) is well defined and satisfies*

$$y^* \leq \dots \leq y^{p+1} \leq y^p \leq \dots \leq y^0,$$

and $\lim_{p \rightarrow \infty} y^p = y^*$, where y^* is an a -supersolution of problem (2.5).

Proof. We proceed by induction. It follows from the equality (3.8) that

$$(3.10) \quad a_{s(0)}(\tilde{y}^0) = a_{s(0)}(y^0) \geq 0.$$

Thus, the relations (3.7) and (3.10) imply the existence of a unique $y^1 \in E$. By the definition of the a -supermapping F^a , it follows from (3.7) and (3.8) that

$$(3.11) \quad y_{s(0)}^1 \leq \tilde{y}_{s(0)}^0 = y_{s(0)}^0.$$

Moreover, it follows from (3.7) and (3.11) that

$$(3.12) \quad y^1 \leq y^0.$$

It follows from the definition of the a -supermapping F^a , (3.7) and (3.8) that

$$(3.13) \quad a_i(y^1) = a_i(y_i^1; y^0) \geq 0, \quad \text{if } i = s(0).$$

Moreover, by the off-diagonal antitonicity of a it follows from (3.7), (3.12), and (3.8) that

$$(3.14) \quad a_i(y^1) = a_i(y_i^0; y^1) \geq a_i(y^0) \geq 0, \quad \text{if } i \neq s(0).$$

The inequalities (3.13) and (3.14) imply

$$(3.15) \quad a(y^1) \geq 0.$$

The relations (3.10), (3.12), and (3.15) give the first step of the induction. We assume now that for some p there exists a unique $y^m \in E$, for all m , $0 < m \leq p$ and a unique $\tilde{y}^m \in E$, for all m , $0 \leq m < p$, such that

$$(3.16) \quad a_{s(m)}(\tilde{y}^m) \geq 0,$$

$$(3.17) \quad y^p \leq \dots \leq y^m \leq y^{m-1} \leq \dots \leq y^0,$$

$$(3.18) \quad a(y^m) \geq 0.$$

It follows from (3.17), (3.3), and (3.9) that

$$(3.19) \quad \tilde{y}_{s(m)}^m = y_{s(m)}^m, \quad \text{for all } m \leq p.$$

From the relations (3.4) and (3.7) we have

$$(3.20) \quad y_{s(p)}^p = y_{s(p)}^0, \quad \text{if } K_{s(p)}^p = \emptyset.$$

It follows from the relations (3.17), (3.9), and (3.2) that

$$(3.21) \quad \tilde{y}^p \leq y^0, \quad \text{if } K_{s(p)}^p = \emptyset.$$

By the off-diagonal antitonicity of a it follows from the relations (3.19)–(3.21) that

$$(3.22) \quad a_{s(p)}(\tilde{y}^p) = a_{s(p)}(y_{s(p)}^0; \tilde{y}^p) \geq a_{s(p)}(y^0) \geq 0, \quad \text{if } K_{s(p)}^p = \emptyset.$$

Consider now the case where $K_{s(p)}^p \neq \emptyset$. From (3.7) and (3.4) we have

$$(3.23) \quad y_{s(p)}^p = y_{s(q)}^{q+1}, \quad \text{if } K_{s(p)}^p \neq \emptyset,$$

where

$$(3.24) \quad q = \max_{k \in K_{s(p)}^p} k.$$

Thus, from the equalities (3.19) and (3.23) we have

$$(3.25) \quad \tilde{y}_{s(p)}^p = y_{s(p)}^p = y_{s(q)}^{q+1}, \quad \text{if } K_{s(p)}^p \neq \emptyset.$$

Moreover, from the relation (3.9) we find that

$$(3.26) \quad \tilde{y}^p \leq \tilde{y}^q, \quad \text{if } K_{s(p)}^p \neq \emptyset.$$

By the off-diagonal antitonicity of a and the definition of the a -supermapping F^a , it follows from the relations (3.25), (3.26), and (3.16) that

$$(3.27) \quad a_{s(p)}(\tilde{y}^p) = a_{s(q)}(y_{s(q)}^{q+1}; \tilde{y}^p) \geq a_{s(q)}(y_{s(q)}^{q+1}; \tilde{y}^q) \geq 0, \quad \text{if } K_{s(p)}^p \neq \emptyset.$$

We note that the relations (3.22) and (3.27) extend the inequality (3.16) to rank $p+1$. The relations (3.7)–(3.9), (3.22), and (3.27) imply the existence of a unique

y^{p+1} . By the definition of the a -supermapping F^a , it follows from the relations (3.7) and (3.19) that

$$(3.28) \quad y_{s(p)}^{p+1} \leq \tilde{y}_{s(p)}^p = y_{s(p)}^p.$$

Thus, it follows from the relations (3.7) and (3.28) that

$$(3.29) \quad y^{p+1} \leq y^p.$$

We note that (3.29) extends the inequality (3.17) to rank $p+1$. By the off-diagonal antitonicity of a and the definition of the a -supermapping F^a , it follows from the relations (3.29), (3.9), and (3.7) that

$$(3.30) \quad a_{s(p)}(y^{p+1}) \geq a_{s(p)}(y_{s(p)}^{p+1}; y^p) \geq a_{s(p)}(y_{s(p)}^{p+1}; \tilde{y}^p) \geq 0.$$

Moreover, by the off-diagonal antitonicity of a it follows from (3.7), (3.29), and (3.18) that

$$(3.31) \quad a_j(y^{p+1}) = a_j(y_j^p; y^{p+1}) \geq a_j(y^p) \geq 0, \quad \text{if } j \neq s(p).$$

The inequalities (3.30) and (3.31) imply

$$(3.32) \quad a(y^{p+1}) \geq 0.$$

The inequality (3.32) extends (3.18) to rank $p+1$. Thus, the induction is complete. By the surjectivity of a , there exists $\bar{y} \in E$ such that

$$(3.33) \quad a(\bar{y}) = 0.$$

Moreover, we have

$$(3.34) \quad a(y^p) \geq 0, \quad \text{for all } p.$$

By the inverse isotonicity of a it follows from the relations (3.33) and (3.34) that $\bar{y} \leq y^p$, for all p .

Hence, there exists y^* such that $\bar{y} \leq y^*$, $y^* \leq y^p$ for all p , and

$$(3.35) \quad \lim_{p \rightarrow \infty} y^p = y^*.$$

By the continuity of the mapping a it follows from the relations (3.34) and (3.35) that

$$a(y^*) \geq 0.$$

Thus, y^* is an a -supersolution. □

4. TWO CLASSES OF a -SUPERMAPPINGS AND ASSOCIATED CONVERGENCE RESULTS

4.1. a -supermappings of the first kind.

Definition 9. Let F^a and F^b be two a -supermappings associated with F with components F_i^a and F_i^b respectively, defined on the domain $\{y \in E \mid a_i(y) \geq 0\}$. Then F^a and F^b satisfy the relation

$$F^a \alpha F^b,$$

if for all $i \in \{1, \dots, \alpha\}$ and y element of the domain $\{y \in E \mid a_i(y) \geq 0\}$, we have

$$(4.1) \quad F_i^b(y) \in \langle F_i^a(y), y_i \rangle_i.$$

Definition 10. The mapping F^a with components F_i^a is an a -supermapping of the first kind associated with F if F^a is an a -supermapping associated with F , there exists an a -supermapping F^b associated with F such that $F^a \alpha F^b$, and

$$(4.2) \quad y^p \downarrow y^*, p \rightarrow \infty \text{ implies } F_i^b(y^p) \downarrow F_i^b(y^*), p \rightarrow \infty \text{ for all } i \in \{1, \dots, \alpha\},$$

where the notation $y^p \downarrow y^*, p \rightarrow \infty$ means that

$$y^* \leq \dots \leq y^{p+1} \leq y^p \leq \dots \leq y^0 \quad \text{and} \quad \lim_{p \rightarrow \infty} y^p = y^*.$$

Remark 5. The relation (4.2) can be interpreted as a property of continuity of F^b related to the partial ordering.

Proposition 7. Let the assumptions of Proposition 6 hold. Let F^a be an a -supermapping of the first kind associated with F and assume that the steering \mathcal{S} satisfies

$$(4.3) \quad \text{for all } i \in \{1, \dots, \alpha\}, \text{ the set } \{p \in N \mid s(p) = i\} \text{ is infinite and}$$

$$(4.4) \quad \text{for all } i \in \{1, \dots, \alpha\}, \quad \lim_{p \rightarrow \infty} \rho_i(p) = +\infty.$$

Then the sequence $\{y^p\}$ defined by (3.7)–(3.9) satisfies $y^p \downarrow \bar{y}$, where \bar{y} is the unique solution of problem (2.5).

Proof. It follows from (3.3), (3.4), and (3.7) that for all $p \geq 1$ such that $s(p) = i$ and $K_i^p \neq \emptyset$, we have

$$(4.5) \quad y_i^{\rho_i(p)} = y_i^p = \dots = y_i^{q+1} = F_i^a(\tilde{y}^q),$$

where

$$(4.6) \quad q = \max_{k \in K_{s(p)}^p} k.$$

We introduce now the sequence $\{z^m\}$ such that for all $i \in \{1, \dots, \alpha\}$,

$$(4.7) \quad z_i^{2m} = y_i^p, \quad m \geq 0,$$

$$(4.8) \quad z_i^{2m-1} = F_i^b(\tilde{y}^q), \quad m \geq 1,$$

where $m = \text{card } K_i^p$, $s(p) = i$, and $q = \max_{k \in K_i^p} k$. It follows from Proposition 6 and (4.7) that

$$(4.9) \quad z_i^{2m} \downarrow y_i^*, m \rightarrow \infty \quad \text{for all } i \in \{1, \dots, \alpha\},$$

where y^* is an a -supersolution. It follows from the relations (4.8), (4.1), (3.19), (4.5), and (4.7) that for all $i \in \{1, \dots, \alpha\}$ and $m \geq 1$, we have

$$(4.10) \quad z_i^{2m-1} = F_i^b(\tilde{y}^q) \in \langle F_i^a(\tilde{y}^q), y_i^q \rangle_i = \langle z_i^{2m}, z_i^{2m-2} \rangle_i.$$

Consider now the following notation. For all $i \in \{1, \dots, \alpha\}$ and $l \in N$,

$$(4.11) \quad z_i^l = \begin{cases} z_i^{2m} & \text{if } l = 2m, \\ z_i^{2m-1} & \text{if } l = 2m-1. \end{cases}$$

The relations (4.9)–(4.11) imply that

$$(4.12) \quad z_i^l \downarrow y_i^*, l \rightarrow \infty, \quad \text{for all } i \in \{1, \dots, \alpha\}.$$

We note that $\lim_{p \rightarrow \infty} \text{card } K_i^p = +\infty$ since

$$K_i^p = \{k \in N \mid s(k) = i \text{ and } 0 \leq k < p\} \subset \{k \in N \mid s(k) = i\},$$

and the set $\{k \in N \mid s(k) = i\}$ is infinite from (4.3). Hence $q = \max_{k \in K_i^p} k \rightarrow +\infty$, if $m = \text{card } K_i^p \rightarrow +\infty$. So, it follows from the relations (4.12) and (4.8) that for all $i \in \{1, \dots, \alpha\}$,

$$(4.13) \quad y_i^* = \lim_{m \rightarrow \infty} z_i^{2m-1} = \lim_{q \rightarrow \infty} F_i^b(\tilde{y}^q).$$

It follows from (3.9) and (4.4) that

$$(4.14) \quad \tilde{y}^q \downarrow y^*, \quad q \rightarrow \infty.$$

Thus, the relations (4.13), (4.14), and (4.2) imply

$$y_i^* = \lim_{q \rightarrow \infty} F_i^b(\tilde{y}^q) = F_i^b(y^*), \quad \text{for all } i \in \{1, \dots, \alpha\}.$$

It follows from Definition 4 that $F_i^b(y^*) = y_i^*$ implies $F_i(y^*) = y_i^*$, for all $i \in \{1, \dots, \alpha\}$. Hence, $y^* = \bar{y}$ which is the unique solution of problem (2.5). \square

4.2. a -supermappings of the second kind.

Definition 11. The mapping F^a with components F_i^a is an a -supermapping of the second kind associated with F if F^a is an a -supermapping associated with F and for all $i \in \{1, \dots, \alpha\}$ and y element of the domain of definition of $F_i^a: \{y \in E \mid a_i(y) \geq 0\}$, there exists $\delta > 0$ such that

$$(4.15) \quad \|y_i - F_i^a(y)\|_i \geq \delta \|y_i - F_i(y)\|_i,$$

where $\|\cdot\|_i$ is a norm defined in E_i .

Proposition 8. Let the assumptions of Proposition 7 hold and assume that F^a is an a -supermapping of the second kind associated with F . Then the sequence $\{y^p\}$ defined by (3.7)–(3.9) satisfies $y^p \downarrow \bar{y}$, where \bar{y} is the unique solution of problem (2.5).

Proof. It follows from Proposition 6 that $y^p \downarrow y^*$. Then for all $i \in \{1, \dots, \alpha\}$ and $\varepsilon > 0$, there exists $p(\varepsilon) \in N$, such that for all $p \geq p(\varepsilon)$ with $s(p) = i$, we have

$$\|y_i^p - y_i^*\|_i \leq \frac{\varepsilon}{2} \delta \quad \text{and} \quad \|y_i^{p+1} - y_i^*\|_i \leq \frac{\varepsilon}{2} \delta.$$

Then

$$(4.16) \quad \|y_i^p - y_i^{p+1}\|_i \leq \varepsilon \delta.$$

It follows from (4.16), (3.7), (3.3), (3.9), and (4.15) that for all $i \in \{1, \dots, \alpha\}$ and $\varepsilon > 0$, there exists $p(\varepsilon) \in N$, such that for all $p \geq p(\varepsilon)$ with $s(p) = i$, we have

$$(4.17) \quad \|\tilde{y}_i^p - F_i(\tilde{y}^p)\|_i \leq \varepsilon.$$

It follows from Proposition 6, (3.9), and (4.4) that

$$(4.18) \quad \tilde{y}^p \downarrow y^*, \quad p \rightarrow \infty.$$

Let $i = s(p)$. We now introduce the following notation: $u^p(i) = \{\dots, u_j^p(i), \dots\}$, where

$$u_j^p(i) = \begin{cases} \tilde{y}_j^p & \text{if } j \neq i, \\ F_i(\tilde{y}^p) & \text{if } j = i. \end{cases}$$

From (2.7) and (2.8) we find that

$$(4.19) \quad a_i(u^p(i)) = 0, \quad \text{for all } p \in N \text{ and } i \in \{1, \dots, \alpha\} \text{ such that } i = s(p).$$

It follows from (4.17) and (4.18) that for all $i \in \{1, \dots, \alpha\}$, we have

$$(4.20) \quad \lim_{p \rightarrow \infty} u^p(i) = y^*.$$

By the continuity of a it follows from (4.19) and (4.20) that

$$a_i(y^*) = 0, \quad \text{for all } i \in \{1, \dots, \alpha\}. \quad \square$$

5. APPLICATION TO A CLASS OF NONLINEAR SIMULTANEOUS EQUATIONS

We consider in this section an M -function obtained by a diagonal monotone perturbation of an M -matrix A

$$(5.1) \quad a(x) = Ax - b + \varphi(x),$$

where $b \in R^n$, $\varphi(x) = \text{diag}\{\dots, \varphi_i(x_i), \dots\}$, and the point-to-point functions φ_i are monotone increasing. Moreover, the functions φ_i are also continuous. Consider now the solution of the following nonlinear simultaneous equations

$$(5.2) \quad a(x) = 0,$$

by asynchronous iterations with order intervals derived from the Schwarz alternating subdomain method.

Remark 6. The results of this section can be extended to the case where $\varphi(x)$ is a diagonal monotone maximal and possibly multivalued operator (see [1]). Then we must solve the following nonlinear algebraic problem

$$(5.3) \quad b - Ax \in \varphi(x).$$

5.1. Connection with the discrete Schwarz alternating procedure. To simplify the presentation we consider in this subsection the case of two discrete subdomains that may overlap. We introduce m_1, m_2 and n which are three integers, n is the number of discretized points, m_1 denotes the last index of the first subdomain, and m_2 the first index of the second subdomain, then

$$1 < m_2 < m_1 < n.$$

We introduce also the sets

$$\mathcal{T}_1 = \{1, \dots, m_1\}, \quad \mathcal{T}_2 = \{m_2, \dots, n\}, \quad \mathcal{T}' = \mathcal{T}_1 \cap \mathcal{T}_2 = \{m_2, \dots, m_1\}.$$

Let $\delta m = m_1 - m_2 + 1$ and $m = n + \delta m$. We will consider in the sequel augmented vectors of R^m and use the following notations:

$$\begin{aligned} \tilde{\mathcal{T}}_1 &= \{1, \dots, m_1\}, & \tilde{\mathcal{T}}_2 &= \{m_1 + 1, \dots, m\}, \\ \tilde{\mathcal{T}}'_1 &= \{m_2, \dots, m_1\}, & \tilde{\mathcal{T}}'_2 &= \{m_1 + 1, \dots, m_1 + \delta m\}, \end{aligned}$$

where $\tilde{\mathcal{T}}'_1$ and $\tilde{\mathcal{T}}'_2$ correspond to the overlapping. We define the mappings $\tilde{p}^1, \tilde{p}^2: R^n \rightarrow R^m$ as follows. For all $u \in R^n$, the mappings $\tilde{p}^1(u), \tilde{p}^2(u)$, with components

$\tilde{p}_l^1(u)$, $\tilde{p}_l^2(u)$, respectively, are given by

$$\tilde{p}_l^1(u) = \begin{cases} u_l & \forall l \in \tilde{T}_1, \\ 0, & \forall l \in \tilde{T}_2', \\ u_{l-\delta m}, & \forall l \in \tilde{T}_2 \ominus \tilde{T}_2', \end{cases}$$

$$\tilde{p}_l^2(u) = \begin{cases} u_l, & \forall l \in \tilde{T}_1 \ominus \tilde{T}_1', \\ 0, & \forall l \in \tilde{T}_1', \\ u_{l-\delta m}, & \forall l \in \tilde{T}_2. \end{cases}$$

We define also the mappings $r^1, r^2: R^m \rightarrow R^n$ as follows: for all $\tilde{u} \in R^m$, the mappings $r^1(\tilde{u}), r^2(\tilde{u})$ with components $r_l^1(\tilde{u}), r_l^2(\tilde{u})$, respectively, are given by

$$r_l^1(\tilde{u}) = \begin{cases} \tilde{u}_l, & \forall l \in T_1, \\ \tilde{u}_{l+\delta m}, & \forall l \in T_2 \ominus T', \end{cases}$$

$$r_l^2(\tilde{u}) = \begin{cases} \tilde{u}_l, & \forall l \in T_1 \ominus T', \\ \tilde{u}_{l+\delta m}, & \forall l \in T_2. \end{cases}$$

Remark 7. For all $i \in \{1, 2\}$ and $u \in R^n$, we have $r^i(\tilde{p}^i(u)) = u$. The mappings $\tilde{p}^1, \tilde{p}^2(r^1, r^2)$ transform a vector of size n (m) into a vector of size m (n).

We consider now the following system of equations

$$(5.4) \quad \tilde{a}(\tilde{u}) = 0,$$

where the mapping \tilde{a} from R^m into R^m with components $\tilde{a}_k(\tilde{u})$, $k = 1, \dots, m$, is defined as follows:

$$(5.5) \quad \begin{cases} \tilde{a}_k(\tilde{u}) = a_k(r^1(\tilde{u})), & \forall k \in \tilde{T}_1, \\ \tilde{a}_k(\tilde{u}) = a_{k-\delta m}(r^2(\tilde{u})), & \forall k \in \tilde{T}_2. \end{cases}$$

Proposition 9 (see [17]). *Let $A \in \mathcal{L}(R^n; R^n)$ be an M -matrix. Then the matrix $\tilde{A} \in \mathcal{L}(R^m; R^m)$ which results from the above augmentation process is also an M -matrix.*

Remark 8. The nonlinear diagonal operator φ from R^n onto itself can be extended similarly.

Corollary 1. *Under the assumptions of this section, the mapping \tilde{a} defined by (5.5) is a surjective M -function.*

Remark 9. The cases with α subdomains ($\alpha > 2$) where no components belong to more than two subdomains can be reduced to the case where $\alpha = 2$ by considering a banded subdomain decomposition with red-black ordering.

5.2. Application to the solution of nonlinear partial differential equations. An illustration of the above theoretical study is given in this subsection. We consider two examples.

Example 1. Nonlinear diffusion problem. We shall primarily be concerned with the following problem.

Find u such that

$$(5.6) \quad \begin{cases} -\Delta u + a \frac{\partial u}{\partial x} + b \frac{\partial u}{\partial y} + e^{cu} = f & \text{everywhere in } \Omega, \\ u/\partial\Omega = 0, \end{cases}$$

where Ω is an open domain of R^2 (or R^3), $\partial\Omega$ is the boundary of Ω , a, b, c are real numbers, and c is positive. We consider the discretization of the problem (5.6) using five-point difference equations for the Laplace operator and only one-sided backward or forward difference equations, according to the sign of a and b , for the first derivatives of u . Then we have the system

$$(5.7) \quad \sum_{j=1}^{\alpha} A_{ij}x_j + \varphi_i(x_i) = b_i, \quad \text{for all } i \in \{1, \dots, \alpha\},$$

where $x_i, b_i \in E_i$, and

$$(5.8) \quad \varphi_i(x_i) \text{ is isotone, convex, and continuously differentiable.}$$

We note that the matrix A is block-partitioned. Moreover it follows from the above discretization that A is an M -matrix.

Remark 10. We can obtain a discretized system with the same properties using suitable P_1 finite element methods.

The problem (5.7) is similar to (2.5). Let $a_i(z_i; w) = A_{ii}z_i + \varphi_i(z_i) + \sum_{j \neq i} A_{ij}w_j - b_i$. We can define implicitly a fixed point mapping F from R^n into R^n satisfying (2.8) and associated with problem (5.7).

Let w, x^0 , and z^0 be three vectors of R^n such that

$$a_i(w) \geq 0, \quad z_i^0 = w_i, \quad \text{and} \quad x_i^0 = z_i = F_i(w), \quad \text{for all } i \in \{1, \dots, \alpha\}.$$

Under the above assumptions the mapping a_i is an M -function on the order interval $\langle x_i^0, z_i^0 \rangle_i$.

We introduce the block-diagonal matrix $C(w)$, derived from Newton's method, with diagonal blocks $C_i(w)$ given by

$$C_i(w) = A_{ii} + \varphi'_i(w).$$

The matrix $C_i(w)$ is an M -matrix since φ_i is convex and the matrix A_{ii} is an M -matrix. Consider now the mappings F^a with components F_i^a defined by

$$(5.9) \quad F_i^a(w) = z_i^k, \quad \text{for all } i \in \{1, \dots, \alpha\},$$

where k is an integer, $z_i^0 = w_i$, and z_i^k is the k th iteration of the following algorithm:

$$(5.10) \quad z^{k+1} = z^k - C^{-1}(z^{\rho'(k)}) \cdot a(z^k), \quad k = 0, 1, \dots,$$

where $0 \leq \rho'(k) \leq k$. We first consider the mapping F^b with components F_i^b given by $F_i^b(w) = z_i^1$, for all $i \in \{1, \dots, \alpha\}$.

Proposition 10. *The mapping F^b is an a -supermapping associated with F .*

Proof. It follows from the definition of z^0 that $a(z^0) \geq 0$. Moreover $C^{-1}(z^0) \geq 0$ since $C(z^0)$ is an M -matrix. Then

$$z^1 - z^0 = -C^{-1}(z^0)a(z^0) \leq 0.$$

It follows from the convexity of φ (see [31, p. 448]) that

$$(5.11) \quad a(z^0) - a(\bar{z}) \leq C(z^0)(z^0 - \bar{z}).$$

Thus if \bar{z} is the solution of the problem $a(z) = 0$, then

$$\bar{z} = \bar{z} - C^{-1}(z^0)a(\bar{z}) = z^1 - (z^0 - \bar{z}) + C^{-1}(z^0)(a(z^0) - a(\bar{z}))$$

and

$$z^1 - (z^0 - \bar{z}) + C^{-1}(z^0)(a(z^0) - a(\bar{z})) \leq z^1 - (I - C^{-1}(z^0) \cdot C(z^0))(z^0 - \bar{z}) = z^1.$$

Thus

$$\bar{z} \leq z^1.$$

Moreover it follows from the convexity of φ that

$$0 = a(z^0) + C(z^0)(z^1 - z^0) \leq a(z^1).$$

We can show analogously by induction (see [37]) that

$$(5.12) \quad \bar{z} \leq \dots \leq z^k \leq z^{k-1} \leq \dots \leq z^0 \quad \text{and} \quad a(z^k) \geq 0.$$

Thus the mappings F^a are a -supermappings associated with F . Moreover it follows from (5.10), (5.12), and Definition 9 that

$$F^a \alpha F^b.$$

We recall that $F_i^b(w) = z_i^1 = w - C^{-1}(w) \cdot a(w)$, where $C(w)$ is a block-diagonal matrix with blocks $C_i(w) = A_{ii} + \varphi'_i(w)$. It follows from (5.8) that the mapping $a(w)$ given by

$$a(w) = Aw + \varphi(w) - b,$$

is continuous. Moreover we have

$$(A_{ii} + \varphi'_i(w_i))^{-1} \leq A_{ii}^{-1},$$

since A is an M -matrix and φ' is positive (see [31, p. 55]). It follows that the spectral radius of $C^{-1}(w)$ is less than or equal to the spectral radius of the inverse of the block-diagonal matrix with blocks A_{ii} which is an M -matrix. Thus the linear application associated with the matrix $C^{-1}(w)$ is Lipschitz continuous and $C^{-1}(w)$ is uniformly Lipschitz continuous. By the continuity of $a(w)$, $C(w)$ and the fact that $C^{-1}(w)$ is uniformly Lipschitz continuous, we find that the mapping F^b is continuous. So, the mappings F^a defined by (5.9) are a -supermappings of the first kind associated with F .

Remark 11. Taking into account the norm convergence properties of Newton's method quoted in subsection 13.3.4 of [31], we can show that the above a -supermappings F^a and F^b are also a -supermappings of the second kind associated with F .

Example 2. The discretized and linearized Hamilton-Jacobi-Bellman problem. We consider the following problem

$$(5.13) \quad \begin{cases} \text{Find } u \text{ such that} \\ \max\{\mathbf{A}^1 u - f^1, \mathbf{A}^2 u - f^2\} = 0, \text{ everywhere in } \Omega, \\ u/\partial\Omega = 0, \end{cases}$$

where \mathbf{A}_1 and \mathbf{A}_2 are two elliptic operators of second order satisfying the Maximum Principle and f_1, f_2 are elements of $L^2(\Omega)$.

Under appropriate discretization of problem (5.13) by finite differences and assuming in particular that the incidence matrices B^1 and B^2 associated with the discretization matrices are equal, we obtain the following discretized problem

$$(5.14) \quad \begin{cases} \text{Find } x \text{ solution of} \\ \max(A^1 \cdot x - b^1, A^2 \cdot x - b^2) = 0, \end{cases}$$

where $b^1, b^2 \in R^n$ and A^1, A^2 are matrices of size $n \times n$ with entries a_{ij}^1, a_{ij}^2 , respectively, which satisfy

$$(5.15) \quad a_{ii}^r > 0, \quad a_{ij}^r \leq 0, \quad i = 1, \dots, n, \quad j = 1, \dots, n, \quad j \neq i, \quad r = 1, 2,$$

$$(5.16) \quad \sum_j a_{ij}^r \geq 0, \quad i = 1, \dots, n, \quad r = 1, 2,$$

$$(5.17) \quad \text{there exists at least one } i \text{ such that } \sum_j a_{ij}^1 > 0 \text{ and } \sum_j a_{ij}^2 > 0,$$

$$(5.18) \quad \text{the matrices } A^1 \text{ and } A^2 \text{ are irreducible.}$$

Remark 12. We recall that the incidence matrix B^1 (B^2) of size $n \times n$ is generated from the matrix A^1 (A^2) by replacing the entries a_{ij}^1 (a_{ij}^2) by 1 if $a_{ij}^1 \neq 0$ ($a_{ij}^2 \neq 0$).

We note that the matrices A^1 and A^2 are diagonally dominant. Under the above assumptions A^1 and A^2 are also M -matrices (see [38]). The problem (5.14) can be linearized as follows:

$$(5.19) \quad a(x) = C(x) \cdot x - b(x) = 0,$$

where $b(x) \in R^n$ and $C(x)$ is a matrix of size $n \times n$. If $(A_1 x - b_1)_i$ is greater than $(A_2 x - b_2)_i$, then the i th row of matrix $C(x)$ is equal to the i th row of matrix A_1 else it is equal to the i th row of matrix A_2 . The vector $b(x)$ is defined analogously. It follows from the above assumptions that the matrix $C(x)$ is an irreducible diagonally dominant matrix; moreover $C(x)$ is an M -matrix (see [38]). Thus a is an M -function.

The mapping $a_i(z_i; w)$ defined by

$$a_i(z_i; w) = \max \left(A_{ii}^1 z_i - \left(b_i^1 - \sum_{j \neq i} A_{ij}^1 w_j \right), A_{ii}^2 z_i - \left(b_i^2 - \sum_{j \neq i} A_{ij}^2 w_j \right) \right) = 0,$$

is a continuous surjective M -function since $C(x)$ is an M -matrix. We can define a fixed point mapping F from R^n into R^n satisfying (2.8) and associated with problem (5.19).

Let w, x^0, z^0 be three vectors of R^n such that

$$a_i(w) = \max \left(\sum_j A_{ij}^1 w_j - b_i^1, \sum_j A_{ij}^2 w_j - b_i^2 \right) \geq 0,$$

$z_i^0 = w_i$, and $x_i^0 = z_i = F_i(w)$, for all $i \in \{1, \dots, \alpha\}$. Let x, z be two vectors of R^n such that $x_i, z_i \in \langle x^0, z^0 \rangle_i$, for all $i \in \{1, \dots, \alpha\}$. Consider the following notation: for all $i \in \{1, \dots, \alpha\}$,

$$a_i(x) = \max_{t \in \{1, 2\}} (A^t x - b^t)_i = (A^r x - b^r)_i,$$

$$a_i(z) = \max_{t \in \{1, 2\}} (A^t z - b^t)_i = (A^s z - b^s)_i.$$

We have for all $i \in \{1, \dots, \alpha\}$, $(A^s x - b^s)_i \leq a_i(x)$. Then

$$(5.20) \quad (A^s x - b^s)_i - (A^r x - b^r)_i \leq 0.$$

Hence we have

$$(5.21) \quad (A^s z - b^s)_i - (A^r x - b^r)_i - (A^s(z - x))_i \leq 0.$$

Thus for all $i \in \{1, \dots, \alpha\}$,

$$(5.22) \quad a_i'(z) - a_i(x) \leq (C(z)(z - x))_i,$$

since the i th line of $C(z)$ is the i th line of $A^s(z)$. So the mapping a is order-convex. We now define the mapping F^a with components F_i^a as follows:

$$F_i^a(w) = z_i^k, \quad \text{for all } i \in \{1, \dots, \alpha\},$$

where k is an integer and the vector z^k is given by the following algorithm:

$$(5.23) \quad z^{k+1} = z^k - C^{-1}(z^{\rho(k)})a(z^k),$$

which is derived from the fixed point equation $z = z - C^{-1}(z)a(z)$. We analogously define the mapping F^b with components F_i^b as follows:

$$F_i^b(w) = z_i^1, \quad \text{for all } i \in \{1, \dots, \alpha\}.$$

By using an argument similar to the one developed in the proof of Proposition 10 it follows from (5.22) that the mappings F^a and F^b are two a -supermappings associated with F which satisfy

$$F^a \alpha F^b.$$

The continuity of the mapping F^b follows from the continuity of the mapping a since the convex hull of continuous functions of R^n is continuous. Thus the mapping F^a is an a -supermapping of the first kind associated with F .

Remark 13. The reader is referred to [23] for the solution of the Hamilton-Jacobi-Bellman problem via several sequential algorithms. The first two algorithms proposed in [23] when used for the solution of the subproblems presented above can lead to a -supermappings of the first kind; in particular the second algorithm corresponds exactly to the linearization method (5.19). Note also that the methods presented in [23] are not related to relaxation or subdomain methods.

Remark 14. The obstacle stationary problem associated to a second order elliptic operator satisfying the maximum principle can classically be written as a Hamilton-Jacobi-Bellman problem, moreover the algorithms presented in this paper can be used for the numerical solution of this problem. Nevertheless the above study cannot be applied to this case since assumption (5.18) is not satisfied. The obstacle problem is important, but we have restricted attention to the case treated in Example 2 in order to give a short presentation. The reader is referred to [37] for a detailed study of the obstacle stationary problem.

5.3. Numerical experiments. In this subsection we briefly present computational experience using a shared memory multiprocessor IBM 3090 with up to six vector processors.

We have considered the two problems presented in subsection 5.2 (i.e. the non-linear diffusion problem and the discretized Hamilton-Jacobi-Bellman equations). The Schwarz alternating procedure with overlapping has been used. The efficiency of parallel iterative algorithms is reported in Table 1 using the classical definition of efficiency: $e = \frac{t_1}{t_p} \cdot \frac{1}{p}$, where t_p denotes the computing time using p processors. Results are given for discretized domains with 25000 points. From Table 1 it can be seen that the efficiency of asynchronous iterations with order intervals is better than the efficiency of parallel synchronous iterations. Idle time due to synchronization and overheads of synchronization reduce the efficiency of parallel iterative

methods. Moreover the efficiency decreases as the number of processors increases. Details of the implementation and other computational results can be found in [37].

TABLE 1. Efficiency of parallel algorithms

	<i>nonlinear diffusion problem</i>		<i>Hamilton-Jacobi-Bellman problem</i>	
number of processors	synchronous algorithms	asynchronous algorithms	synchronous algorithms	asynchronous algorithms
3	0.87	0.99	0.86	0.90
6	0.65	0.81	0.85	0.90

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