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# Contractive mappings with maximum norms: Comparison of constants of contraction and application to asynchronous iterations

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## *Abstract*

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In this paper, we give two extensions of Stein-Rosenberg's theorem. The first, which we name the general result, is an abstract nonlinear extension and can be described as follows: Given a first fixed point mapping on a Banach product space, we define a more implicit second fixed point mapping, possibly after a redecomposition of our product space, and we get that the new mapping has a constant of contraction lower or equal to the constant of contraction of the initial mapping. This result allows an efficient use of El Tarazi's theorem [5], about the convergence of asynchronous iterations. The second extension is close to the linear case and permits us to compare the constants of contraction using strict inequality. We give two applications of these results: the first, in a context near from the one studied by D.J. Evans and W. Deren, about a diagonal monotone perturbation of linear problems [6]. The second is a short example in a totally different framework about the formulation of asynchronous waveform relaxation for a system of ordinary differential equations with initial conditions. For another point of view concerning Stein-Rosenberg's theorem and asynchronous algorithms, see [7] and [10].

**Keywords.** Asynchronous iteration; diagonal monotone perturbation; Stein-Rosenberg's theorem; differential algebraic problems.

## 0. Introduction

In this paper, we give two extensions of Stein-Rosenberg's theorem. The first, which we name the general result, is an abstract nonlinear extension and can be described as follows: given a first fixed point mapping on a Banach product space, we define a more implicit second fixed point mapping, possibly after a redecomposition of our product space, and we get that the new mapping has a constant of contraction lower or equal to the constant of contraction of the initial mapping.

This result allows an efficient use of El Tarazi's theorem, about the convergence of asynchronous iterations. This use can be described as follows: If we are, for example, in the

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framework of the assumptions of El Tarazi's theorem for a by point decomposition of the finite dimensional problem, then our theorem implies that we are also in the same kind of framework for a more implicit, by block decomposition of the same fixed point problem.

Moreover, it is worthwhile to note that this abstract form of Stein-Rosenberg's theorem, as El Tarazi's theorem, deals with maximum norms which in application are often related by the way of the maximum principle, with order structures; this is the usual framework of Stein-Rosenberg's theorem.

The second extension is close to the linear case and permits us to compare the constants of contraction using strict inequality. We apply these points of view to two kinds of situations: First, in a context like the one studied by Evans and Deren [6] about a diagonal monotone perturbation of linear problems (in short: DMPL problems), associated to an  $M$ -matrix, which thanks to Evans and Deren's study, allows to include asynchronous Schwarz alternating methods. Besides the fact that our proofs of convergence are quite different from those of Evans and Deren, our statement includes here multivalued mappings (this is interesting for the case of constrained problems), and a somewhat different parameter: namely the spectral radius (or approximate spectral radius in reducible matrices) of the Jacobi associated matrix, instead of the norms used by these authors. Moreover, it appears that our general result has something to do with the classical Stein-Rosenberg's theorem: both allow an inequality:

- between constants of contraction for our theorem,
- between spectral radii for Stein-Rosenberg's theorem, the corresponding inequality being strict in this case, by the way of Perron-Frobenius theorem. The second extension leads us to take advantage of the matricial aspect of the DMPL problems, to improve the interval of convergence of the relaxation parameter and also the constant of contraction of the corresponding fixed point mapping.

Second, we give a short example in a totally different framework about the formulation of asynchronous waveform relaxation for a system of ordinary differential equations with initial conditions. In Bahi's thesis and a forthcoming paper (by M. Bahi, E. Griepentrog and J.C. Miellou), the same point of view is emphasized in order to study the waveform relaxation method for classes of differential algebraic problems.

This paper is organized in the following way: In Section 1 we give some backgrounds about asynchronous iterations. Section 2 is dedicated to the statement and the proof of the general result. In Section 3, we apply this result to the parallel treatment associated to a block composition of the nonlinear algebraic DMPL problems:

$$b - Ax \in M(x).$$

Section 4, we use the classical Stein-Rosenberg theorem in order to prove the second extension in the context of DMPL problems, which allows us to derive a best contraction constant and convergence condition on the relaxation parameter of our fixed point mapping. Section 5, we give the second application concerning asynchronous waveform relaxation for ordinary differential equations.

## Background: Asynchronous algorithms

Consider a Banach product space  $E = \prod_{i=1}^{\alpha} E_i$  and a fixed point mapping  $G$  defined on  $E$ . Introduce the following symbols:  $J = \{J(p)\}_{p \in \mathbb{N}}$  is a sequence of nonempty subsets of  $\{1, \dots, \alpha\}$ .  $S = \{(s_1(p), \dots, s_{\alpha}(p))\}_{p \in \mathbb{N}}$  is a sequence of  $\mathbb{N}^{\alpha}$  such that:

$\forall i \in \{1, \dots, \alpha\}$ , the subset  $\{p \in \mathbb{N}, i \in J(p)\}$  is infinite.

$\forall i \in \{1, \dots, \alpha\}, \forall p \in \mathbb{N}, s_i(p) \leq p$ .

$\forall i \in \{1, \dots, \alpha\}, \lim_{p \rightarrow \infty} s_i(p) = \infty$ .

The asynchronous (or synchronous) algorithm associated to  $G$  and noted  $(G, u^0, J, S)$  is defined as follows:

$$u_i^{p+1} = \begin{cases} G_i(\dots, u_k^{s_k(p)}, \dots) & \text{if } i \in J(p) \\ u_i^p & \text{if } i \notin J(p). \end{cases} \quad (1.1)$$

The following statement is useful in order to study the behaviour of asynchronous iterations.

**Theorem 1** [5]. *Let  $G$  be a mapping from  $D(G) \subset E$  in  $E$  and suppose that:*

- (a)  $D(G) \subset \prod_{i=1}^{\alpha} D_i(G)$
  - (b)  $\exists u^* \in D(G)$ , such that  $u^* = G(u^*)$
  - (c)  $\forall u \in D(G)$ ,  $\|G(u) - G(u^*)\| \leq \beta \|u - u^*\|$
- with  $0 < \beta < 1$ , then every asynchronous algorithm corresponds to  $G$  with a starting point  $u^0 \in D(G)$ , converges to the fixed point of  $G$ ,  $u^*$ .

**Proof.** See [5].  $\square$

## 2. An abstract nonlinear extension of Stein-Rosenberg's theorem

Let  $E_i$ ,  $i \in \{1, \dots, n\}$  be Banach spaces equipped with the norms:

$$\|\dots\|_i. \quad (2.1)$$

Consider  $E = \prod_{i=1}^n E_i$  equipped with the norm:

$$\|\dots\| = \max_{1 \leq i \leq n} \|\dots\|_i. \quad (2.2)$$

Consider a mapping:

$$T: E = \prod_{i=1}^n E_i \longrightarrow E = \prod_{i=1}^n E_i. \quad (2.3)$$

$$u \longrightarrow v = T(u).$$

Suppose that:  $\forall \alpha (\gamma_1, \dots, \gamma_n) \in \mathbb{R}_+^{*n}$  such that  $T$  is contractive in the norm:

$$\|\dots\| = \max_{1 \leq i \leq n} \frac{\|\dots\|_i}{\gamma_i}. \quad (2.4)$$

**Definition 1.** Let  $I_k (k \in \{1, \dots, \alpha\}, \alpha \leq n)$  be subsets of  $\mathbb{N}$  such that:

- (a)  $I_k \cap I_j = \emptyset, \forall k \neq j$
- (b)  $\bigcup_{k=1}^{\alpha} I_k = \{1, \dots, n\}.$

Define:

$$\mathbb{E}_k = \prod_{i \in I_k} E_i. \quad (2.5)$$

Thus, we can consider  $E$  as the product:

$$E = \prod_{k=1}^{\alpha} \mathbb{E}_k. \quad (2.6)$$

Now, let  $U^k \in \mathbb{E}_k$ , we have:

$$(U^k)_i = u_i, \quad \forall i \in I_k. \quad (2.7)$$

Let  $J_k$  be subsets of  $\{1, \dots, \alpha\}$ , and

$$\mathbb{I}_k = \bigcup_{l \in J_k} I_l.$$

For  $U \in E = \prod_{i=1}^{\alpha} \mathbb{E}_i$  and  $v \in E = \prod_{i=1}^n E_i$ , we define the vector  $\sigma_k(U, v)$  as follows:

$$\begin{aligned} \sigma_k(U, v) &= (w_1, \dots, w_n), \quad \text{where:} \\ \begin{cases} w_j = u_j & \text{if } j \in \mathbb{I}_k \\ w_j = v_j & \text{if } j \notin \mathbb{I}_k. \end{cases} \end{aligned} \quad (2.8)$$

Establishment of the mapping  $\mathbb{T}$ . We consider the following mapping:

$$\begin{aligned} \mathbb{T}: E = \prod_{i=1}^n E_i &\longrightarrow E = \prod_{i=1}^{\alpha} \mathbb{E}_i. \\ (v_1, \dots, v_n) &\longrightarrow (U^1, \dots, U^{\alpha}). \end{aligned}$$

such that:

$$T_i(\sigma_k(U, v)) = (U^k)_i. \quad (2.9)$$

**Remark 1.**

$$(\mathbb{T}_k(v))_i = T_i(\sigma_k(\mathbb{T}(v), v)), \quad (2.10)$$

$$(\mathbb{T}_k(u^*))_i = (u^*)_i \Leftrightarrow T_i(u^*) = u_i^*, \quad (2.11)$$

if  $T$  and  $\mathbb{T}$  are contractive then they have the same fixed point.

**Theorem 2** (General result). *If  $T$  is a contraction mapping in the norm:*

$$\| \dots \| = \max_{1 \leq i \leq n} \frac{\| \dots \|_i}{\gamma_i}, \quad \gamma_i > 0, \quad \forall i \in \{1, \dots, n\}, \quad (2.12)$$

then  $\mathbb{T}$  is also contractive in the norm:

$$\| \dots \| = \max_{1 \leq l \leq \alpha} \left( \max_{i \in \mathbb{I}_l} \frac{\| \dots \|_i}{\gamma_i} \right). \quad (2.13)$$

Moreover, if  $\nu$  and  $\nu'$  are the respective contraction constants of  $T$  and  $\mathbb{T}$  then:

$$\nu \geq \nu'. \quad (2.14)$$

**Proof.**

$$\|T(w) - T(v)\| \leq \nu \|w - v\|. \quad (2.15)$$

From (2.10) we have:

$$(\mathbb{T}_k(v))_i = T_i(\sigma_k(\mathbb{T}(v), v)), \quad (2.16(a))$$

$$(\mathbb{T}_k(w))_i = T_i(\sigma_k(\mathbb{T}(w), w)). \quad (2.16(b))$$

Then:

$$\frac{\|(\mathbb{T}_k(v))_i - (\mathbb{T}_k(w))_i\|_i}{\gamma_i} \leq \nu \max \left[ \max_{j \in \mathbb{I}_k} \frac{\|(\mathbb{T}_k(v))_j - (\mathbb{T}_k(w))_j\|_j}{\gamma_j}; \max_{1 \leq l \leq \alpha} \max_{j \in \mathbb{I}_l} \frac{\|v_j - w_j\|_j}{\gamma_j} \right].$$

Denote:

$$\rho_k = \max_{j \in \mathbb{I}_k} \frac{\|(\mathbb{T}_k(v))_j - (\mathbb{T}_k(w))_j\|_j}{\gamma_j} \quad (2.17)$$

then

$$\rho_k \leq \nu \left( \max \left( \rho_k; \max_{1 \leq l \leq \alpha} \max_{j \in \mathbb{I}_l} \frac{\|v_j - w_j\|_j}{\gamma_j} \right) \right).$$

Therefore, either

$$\rho_k \leq \nu \rho_k, \quad \text{so } \rho_k = 0, \quad (2.18)$$

or

$$\rho_k \leq \max_{1 \leq l \leq \alpha} \max_{j \in \mathbb{I}_l} \frac{\|v_j - w_j\|_j}{\gamma_j}. \quad (2.19)$$

Thus,

$$\max_{1 \leq k \leq \alpha} \max_{j \in \mathbb{I}_k} \frac{\|(\mathbb{T}_k(v))_j - (\mathbb{T}_k(w))_j\|_j}{\gamma_j} \leq \nu \max_{1 \leq l \leq \alpha} \max_{j \in \mathbb{I}_l} \frac{\|v_j - w_j\|_j}{\gamma_j},$$

that is

$$\|\mathbb{T}(v) - \mathbb{T}(w)\| \leq \nu \|v - w\|. \quad \square$$

**Remark 2.** Particular case. If  $\mathbb{I}_k = I_k (J_k = \{k\})$ , the mapping  $\mathbb{T}_k$  corresponds to the block Jacobi mapping which is the usual basic fixed point mapping which allows to generate block asynchronous iterations. We may use the previous theorem associated to this choice of  $\mathbb{I}_k$  in sections 3 and 5. In section 3.5 we present another kind of choice of  $\mathbb{I}_k$  in order to modelize mixed Gauss-Seidel or Jacobi asynchronous iterations.

### 3. The first application of the general result

In this section, we are interested in asynchronous parallel algorithms for solving nonlinear algebraic problems:

$$b - Ax \in M(x), \quad x \in D(M), \quad (3.1)$$

where:

$$A = a_{ij} \quad (3.2)$$

is an irreducible  $M$  matrix,

$$M(x) = (M_1(x_1), \dots, M_n(x_n))^T \quad (3.3)$$

is a diagonal, monotone, maximal and possibly multi-valued operator defined on  $D(M)$ . We

will subdivide problem (3.1) into  $\alpha$  subproblems which will be treated by  $\alpha$  processors. First we define a fixed point mapping whose fixed point coincides with the solution of (3.1).

### 1. The by point fixed point mapping $T$

We consider the following fixed point mapping:

$$\begin{aligned} T : D(M) \subset \mathbb{R}^n &\longrightarrow \mathbb{R}^n \\ x &\longrightarrow y \end{aligned}$$

such that:

$$\begin{cases} \lambda_i(y_i) + a_{ii}y_i = - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}x_j + b_i \\ \lambda_i(y_i) \in M_i(y_i). \end{cases} \quad (3.4)$$

**Remark 3.**

- )  $T$  is well defined, indeed:  $M_i$  is monotone maximal so  $y_i$  exists and is unique.
- ) If  $x^*$  is the fixed point of  $T$  then  $x^*$  is the unique solution of (3.1), indeed:

$$\begin{cases} \lambda_i(x_i^*) + a_{ii}x_i^* = - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}x_j^* + b_i \\ \lambda_i(x_i^*) \in M_i(x_i^*) \end{cases}$$

that

$$0 \in Ax^* + M(x^*) - b.$$

**Proposition 3.** *There exists a vector  $e > 0$  (all the components of  $e$  are  $> 0$ ), such that  $T$  is contractive with respect to the norm:*

$$\| \dots \| = \max_{1 \leq i \leq n} \frac{\| \dots \|_i}{e_i}.$$

**Proof.** We are here in a particular case of more general frameworks such as in El Tarazi [5], Miellou and Spiteri [8]; but we detail the proof for an easy reading: Let  $y = T(x)$  and  $\tilde{y} = T(\tilde{x})$ , then:

$$\lambda_i(y_i) + a_{ii}y_i = - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}x_j + b_i$$

and

$$\lambda_i(\tilde{y}_i) + a_{ii}\tilde{y}_i = - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}\tilde{x}_j + b_i$$

then

$$\lambda_i(y_i) - \lambda_i(\tilde{y}_i) + a_{ii}(y_i - \tilde{y}_i) = - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}(\tilde{x}_j - x_j). \quad (3.5)$$

Therefore:

$$(y_i - \tilde{y}_i)[\lambda_i(y_i) - \lambda_i(\tilde{y}_i) + a_{ii}(y_i - \tilde{y}_i)] = \left[ - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}(\tilde{x}_j - x_j) \right] (y_i - \tilde{y}_i).$$

From (3.3) we have:

$$(y_i - \tilde{y}_i)(\lambda_i(y_i) - \lambda_i(\tilde{y}_i)) \geq 0.$$

Then

$$(y_i - \tilde{y}_i)a_{ii}(y_i - \tilde{y}_i) \leq \left[ - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}(\tilde{x}_j - x_j) \right] (y_i - \tilde{y}_i) \quad (3.6)$$

so

$$(y_i - \tilde{y}_i) \leq - \frac{\sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}(x_j - \tilde{x}_j)}{a_{ii}}. \quad (3.7)$$

Let  $\|\dots\|_i$  be a norm on  $\mathbb{R}$ . From the Perron-Frobenius theorem, there exists  $e > 0$  such that:

$$J(e) \leq \rho(J) \cdot e,$$

where  $J$  is the Jacobi matrix of  $A$  and  $\rho(J)$  its spectral radius. Therefore, as in [9]:

$$\|y_i - \tilde{y}_i\|_i \leq - \frac{\sum_{\substack{j=1 \\ j \neq i}}^n a_{ij}e_j}{a_{ii}} \frac{\|x_j - \tilde{x}_j\|_j}{e_j},$$

so

$$\|y_i - \tilde{y}_i\|_i \leq \rho(J)e_i \frac{\|x_j - \tilde{x}_j\|_j}{e_j}. \quad (3.8)$$

Therefore

$$\frac{\|y_i - \tilde{y}_i\|_i}{e_i} \leq \rho(J) \frac{\|x_j - \tilde{x}_j\|_j}{e_j}, \quad (3.9)$$

so that

$$\|y - \tilde{y}\| \leq \rho(J) \|x - \tilde{x}\|.$$

Since  $A$  is an  $M$  matrix  $\rho(J) < 1$ ,  $T$  is contractive on  $\mathbb{R}^n$  endowed with the norm  $\|\dots\|$ .  $\square$

**Remark 4.** If  $A$  is reducible it is sufficient to replace  $\rho(J)$  by  $\rho(J) + \epsilon$ , where  $\epsilon$  is an arbitrary positive real number.

### 3.2. Decomposition of (3.1)

Let  $(n_1, \dots, n_\alpha) \in \mathbb{N}_+^\alpha$ , such that:

$$\sum_{i=1}^{\alpha} n_i = n, \quad (3.10)$$

decompose a vector  $x \in \mathbb{R}^n$  into  $\alpha$  blocks  $X_i$  of  $n_i$  components:

$$x = X = (X_1, \dots, X_\alpha)^T. \quad (3.11)$$

For an  $n \times n$  matrix  $A = a_{ij}$ , we denote:

$$A = A_{ij} \text{ the correspondent } \alpha \times \alpha \text{ matrix.} \quad (3.12)$$

we also denote:

$$M(x) = \mathbb{M}(X) = (\mathbb{M}_1(X_1), \dots, \mathbb{M}_\alpha(X_\alpha))^T \quad (3.13)$$

and:

$$b = B = (B_1, \dots, B_\alpha)^T. \quad (3.14)$$

(3.1) can be written after decomposition:

$$\begin{cases} 0 = \Lambda_i(X_i) + \sum_{j=1}^{\alpha} A_{ij}X_j - B_i \\ \Lambda_i(X_i) \in \mathbb{M}_i(X_i), \end{cases} \quad (3.15)$$

where

$$\lambda(x) = \Lambda(x) = (\Lambda_1(x_1), \dots, \Lambda_\alpha(x_\alpha))^T. \quad (3.16)$$

#### mark 5.

This decomposition is general, and it is worthwhile to note that the Schwartz alternating procedure can be viewed as a block asynchronous (or synchronous) relaxation algorithm applied to a problem of the kind of (3.1), but set on a space of greater dimension, see [6] for the explication of the correspondent extended matrices and problems, and for asynchronous forms of the Schwartz alternating method.

For a related point of view about redecomposition techniques, see Comte [4].

#### 4. The more implicit fixed point mapping $\mathbb{T}$

We consider the following fixed point mapping:

$$\begin{aligned} \mathbb{T}: \mathbb{R}^\alpha &\longrightarrow D(\mathbb{M}) \cap \mathbb{R}^\alpha \\ X &\longrightarrow Y \end{aligned}$$

such that

$$\begin{cases} \Lambda_i(Y_i) + A_{ii}Y_i = - \sum_{\substack{j=1 \\ j \neq i}}^{\alpha} A_{ij}X_j + B_i \\ \Lambda_i(Y_i) \in \mathbb{M}_i(Y_i). \end{cases} \quad (3.17)$$

**mark 6.** The same remark as in Remark 3 can be made.

**proposition 4.**  $\mathbb{T}$  is contractive, moreover its contraction constant is lower or equal to the contraction constant of  $T$ .

**proof.** We have only to use the particular case (Remark 2) of the general result applied to  $T$  defined in (3.4) and  $\mathbb{T}$  defined by (3.17).  $\square$



**Theorem 5.** Every asynchronous algorithm associated to  $\mathbb{T}$  with a starting point  $u^0 \in D(\mathbb{M})$ , converges to the unique solution of the problem (3.1).

**Proof.** Since  $\mathbb{T}$  is contractive in the norm:

$$\max_{1 \leq l \leq \alpha} \max_{j \in \mathbb{I}_l} \frac{\| \cdot \|_j}{\gamma_j} = \max_{1 \leq l \leq \alpha} \frac{\| \cdot \|_1}{\gamma_l}, \quad \text{where:}$$

$$\frac{\| \cdot \|_l}{\gamma_l} = \max_{j \in \mathbb{I}_l} \frac{\| \cdot \|_j}{\gamma_j},$$

it is sufficient to apply Theorem 1.  $\square$

### 3.4. The relaxed fixed point mappings

#### 3.4.1. The relaxed fixed point mapping $T_\omega$

Consider the mapping:

$$\begin{aligned} T_\omega: \mathbb{R}^n &\longrightarrow \mathbb{R}^n \cap D(M) \\ x &\longrightarrow y \end{aligned}$$

such that

$$y = (1 - \omega)x + \omega T(x) \quad (3.18)$$

where  $T$  is defined in (3.4).

Then:

**Theorem 6.** If  $\omega \in ]0, 2/1 + e(J)[$  where  $J$  is the Jacobi matrix of the matrix  $A$ , then every asynchronous algorithm correspondent to  $T_\omega$  converges to the unique solution of (3.1).

**Proof.** We will prove that  $T_\omega$  is contractive. Let  $y = T_\omega(x)$  and  $\tilde{y} = T_\omega(\tilde{x})$ , then:

$$\begin{aligned} \|y - \tilde{y}\| &= \|(1 - \omega)x + \omega T(x) - (1 - \omega)\tilde{x} - \omega T(\tilde{x})\| \\ &= \|(1 - \omega)(x - \tilde{x}) + \omega(T(x) - T(\tilde{x}))\| \\ &\leq |1 - \omega| \|x - \tilde{x}\| + |\omega| \rho(J) \|x - \tilde{x}\|. \end{aligned}$$

It is obvious that if  $0 < \omega < 2/1 + e(J)$ , then:

$$|1 - \omega| + |\omega| \rho(J) < 1,$$

so  $T_\omega$  is contractive. Now it is sufficient to apply Theorem 1 to prove Theorem 6.  $\square$

#### 3.4.2. The relaxed fixed point mapping $\mathbb{T}_\omega$

Consider the mapping  $\mathbb{T}_\omega$  defined as follows:

$$\begin{aligned} \mathbb{T}_\omega: \mathbb{R}^\alpha &\longrightarrow \mathbb{R}^\alpha \cap D(\mathbb{M}) \\ X &\longrightarrow Y \end{aligned}$$

such that

$$Y = (1 - \omega)X + \omega \mathbb{T}(X). \quad (3.19)$$

**Theorem 7.** If  $\omega \in ]0, 2/1 + e(J)[$  where  $J$  is the Jacobi matrix of the matrix  $A$ , then every asynchronous algorithm correspondent to  $\mathbb{T}_\omega$  converges to the unique solution of (3.1).

**oof.** It is sufficient to apply the general result applied to  $T_\omega$  defined by (3.18), and  $\mathbb{T}_\omega$  defined by (3.19).  $\square$

### 3. Mixed Gauss-Seidel asynchronous iterations

Let us assume a by point decomposition of each diagonal block  $A_{ii}$  of the form:

$$A_{ii} = D_i - L_i - U_i,$$

where with the usual notation  $D_i$  (resp  $L_i$ ;  $U_i$ ) is denoted the by point diagonal part of  $A_{ii}$  (resp triangular lower; triangular upper part of  $A_{ii}$ ). Let us define the following fixed point mapping: Given  $X = (X_1, \dots, X_\alpha)$  and  $Y = (Y_1, \dots, Y_\alpha)$ ,  $\mathbb{T}_{G,S}(X) = Y$  satisfies

$$A_i(Y_i) + (D_i - L_i)Y_i = U_i X_i - \sum_{j \neq i} A_{ij} X_j + B_i.$$

Let us introduce the subsets of indices:  $I'_i = \{l_i, \dots, l_{i+1} - 1\}$  defining the  $i$ th block  $A_{ii}$ ,  $i \in \{1, \dots, n\}$ ,  $\exists i$  such that  $l \in I'_i$ . Then we have:  $I_l = \{k, l_i \leq k \leq l\}$ , which by the way of (2.8) allows to define  $\sigma_l(Y, X)$ , and so the fixed point mapping  $\mathbb{T}_{G,S}$  satisfies the condition of Theorem 2, and so we have:

**Proposition 8.**  $\mathbb{T}_{G,S}$  is contractive with a constant of contraction lower or equal to the one of the point mapping  $T$  defined in (3.4).

**Remark 7.** In fact the kind of result given by the previous statement can also be obtained by considering the fact that all the asynchronous algorithms associated to the fixed point mapping  $\mathbb{T}_{G,S}$ , can be included in asynchronous methods associated to the by point fixed point mapping  $T$ .

### A second nonlinear extension of Stein-Roseinberg's theorem closer to the linear case

We are interested in the present section only in the case of strict inequalities between constants of contraction of fixed point mappings, because the case of non-strict inequalities is always solved by our previous Theorem 2. So we have to assume for the matrices involved in our problems and statement the irreducibility property. We have to use the following theorem (see [3]):

**Theorem 9.** Let  $A$  be an  $M$ -matrix. Let  $A = B_1 - C_1$  and  $A = B_2 - C_2$  be two regular splittings of  $A$  with:

$$C_2 < C_1 \quad (C_2 \leq C_1 \text{ and } C_2 \neq C_1) \quad (4.1)$$

and

$$B_1^{-1}C_1 \text{ irreducible.}$$

then

$$\rho(B_2^{-1}C_2) < \rho(B_1^{-1}C_1).$$

**oof.** See [3], p. 183, and [11] p. 57.  $\square$

Now let us consider that

$$A = B - C \quad (4.2)$$

is a regular decomposition of the  $M$ -matrix  $A$ , in which  $B$  is also an  $M$ -matrix, such that  $B^{-1}C$  is irreducible. Let us introduce the fixed point mapping  $T$  defined by

$$\begin{aligned} T(x) &= y \text{ such that} \\ M(y) + By &= Cx + b, \end{aligned} \quad (4.3)$$

where  $M$  is defined in (3.3). Let us introduce the norm

$$\|u\| = \max_l \frac{|u_l|}{e_l}, \quad (4.4)$$

where  $e$  is the eigenvector associated to the spectral radius of  $B^{-1}C$ , that is to say by Perron-Frobenius theorem:

$$B^{-1}Ce = \rho(B^{-1}C)e = \rho e. \quad (4.5)$$

Then we have:

**Proposition 10.**

$$\|T(v) - T(v')\| \leq \rho \|v - v'\|.$$

**Proof.** We have

$$\begin{aligned} \lambda(y) + By &= Cx + b \\ \lambda(y') + By' &= Cx' + b \end{aligned}$$

so

$$\lambda(y) - \lambda(y') + B(y - y') = C(x - x').$$

Then

$$(\lambda_l(y_l) - \lambda_l(y'_l))(y_l - y'_l) + \sum_k b_{lk}(y_k - y'_k)(y_l - y'_l) = \sum_k c_{lk}(x_k - x'_k)(y_l - y'_l).$$

Since  $b_k \leq 0$  for  $l \neq k$ , and  $b_{ll} > 0$ , we deduce (like in the proof of Proposition 3):

$$B\|y - y'\| \leq C\|x - x'\|.$$

Therefore,

$$\|y - y'\| \leq B^{-1}C\|x - x'\| \quad (B^{-1}C \geq 0).$$

Then with the use of The Perron-Frobenius theorem, we deduce:

$$\|y - y'\| \leq \rho(B^{-1}C)\|x - x'\|. \quad \square$$

Let us consider now two decompositions of  $A$ :

$$A = B_1 - C_1 \quad \text{and} \quad A = B_2 - C_2,$$

both satisfying the assumptions of Theorem 9, except the irreducibility which has to be satisfied only by  $B_1^{-1}C_1$ . Then if  $T_1, \|\cdot\|_1, \rho_1$  and  $T_2, \|\cdot\|_2, \rho_2$  are the fixed point mappings, the norms and the constants of contraction, corresponding by (4.3), (4.4), (4.5), to the respective decompositions:  $A = B_1 - C_1$ ; and  $A = B_2 - C_2$ , then:

**Corollary 11** (nonlinear variant of Stein-Rosenberg's theorem). *The constants of contraction of  $T_1$  and  $T_2$  can be compared and we have:*

$$\rho_2 < \rho_1.$$

**Proof.** Results both from Theorem 9 and Proposition 10.  $\square$

**position 12.** Under the assumptions of Theorem 9, suppose that  $A$  is irreducible in the block decomposition (3.15), then: if  $\omega \in ]0, 2/(1 + \rho(J'))[$  where  $J'$  is the Jacobi matrix correspondent to the block decomposition of the matrix  $A$ , then every asynchronous algorithm correspondent to  $\mathbb{T}_\omega$  converges to the unique solution of (4.1). Moreover:

$$\rho(J') < \rho(J),$$

where  $J$  is the by point Jacobi matrix associated to  $A$ .

**Proof.** It is sufficient to take:

$$B_1 = \text{diag}(\dots, a_{ii}, \dots); C_1 = B_1 - A$$

and

$$B_2 = \text{diag}(\dots, A_{ii}, \dots); C_2 = B_2 - A.$$

being irreducible,  $J$  is also irreducible, so we are in the framework of our previous irreducibility assumptions which ensure us of strict inequality.  $\square$

**Remark 8.** All the previous results can also be applied to the case in which we replace the assumption ' $A$  is an  $M$ -matrix' by ' $A$  is an  $H$ -matrix with positive diagonal elements', and in which we consider that:  $A = B - C$ , where  $B$  is obtained by cancelling extradiagonal elements of  $A$  (this is the case in each one of the examples associated both to block decompositions or to Gauss-Seidel blocks decompositions considered above).

## The second application of the general result

Consider an ordinary differential equation with initial condition:

$$\begin{cases} y'(t) = f(y(t), t) \\ y(t_0) = y_0, \end{cases} \quad (5.1)$$

where  $y(t) = (y_1(t), \dots, y_n(t))^T$ ;  $f = (f_1, \dots, f_n)^T$ ;  $y_0 = (y_{01}, \dots, y_{0n})^T$ . and  $t \in [t_0, t_0 + h]$  ( $h > 0$ ). Consider  $\|x\|_\infty = \max_i |x_i|$ , the maximum norm on  $\mathbb{R}^n$ . Suppose that  $f$  satisfies, for the sake of simplicity, a global Lipschitz condition:

$$\|f(y, t) - f(\tilde{y}, t)\|_\infty \leq L \|y - \tilde{y}\|_\infty,$$

where  $L$  does not depend on  $t$ . It is a classical result that for  $h$  sufficiently small the problem (5.1) has a unique solution, and that the following fixed point mapping:

$T(y) = x$  such that:

$$\begin{cases} \frac{dx_i(t)}{dt} = f_i(y_1(t), \dots, y_{i-1}(t), x_i(t), y_{i+1}(t), \dots, y_n(t), t) \\ x_i(t_0) = y_{0i} \end{cases} \quad (5.2)$$

is contractant and converges to the unique solution of (5.1). Now, we can write problem (5.1) in block decomposition:

$$\begin{cases} Y'(t) = F(Y(t), t) \\ Y(t_0) = Y_0, \end{cases} \quad (5.3)$$

where  $Y(t) = (Y_1(t), \dots, Y_\alpha(t))^T$ ;  $F = (F_1, \dots, F_\alpha)^T$ ;  $Y_0 = (Y_{01}, \dots, Y_{0\alpha})^T$ , and  $t \in [t_0, t_0 + h]$  ( $h > 0$ ).

$> 0$ ).  $Y_i$ ,  $F_i$  and  $Y_{0i}$  have  $n_i$  components, and  $\sum_{i=1}^a n_i = n$ . Let us consider the norm on  $C^n([t_0, t_0 + h])$ :

$$\|x\|_\infty = \max_{1 \leq i \leq n} \max_{t_0 \leq t \leq t_0 + h} |x_i(t)|.$$

Consider the following mapping:  $\mathbb{T}(Y) = X$  such that:

$$\begin{cases} \frac{dX_i(t)}{dt} = f_i(Y_1(t), \dots, Y_{i-1}(t), X_i(t), Y_{i+1}(t), \dots, Y_a(t), t) \\ X_i(t_0) = Y_{0i}, \end{cases} \quad (5.4)$$

then we have:

**Theorem 13.** *Every asynchronous algorithm associated to  $\mathbb{T}$  with a starting point  $Y_0$ , converges to the unique solution of the problem (5.1).*

**Proof.** It is a classical result that for  $h$  sufficiently small  $T$  is contractive with respect to the norm  $\|\dots\|_\infty$ . We deduce by the redecomposition theorem that  $\mathbb{T}$  is also contractive with respect to the same norm, and then it is sufficient to apply Theorem 1.  $\square$

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