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A COMPARISON OF THE SUCCESSIVE OVERRELAXATION METHOD AND SEMI-ITERATIVE METHODS USING CHEBYSHEV POLYNOMIALS*

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1. Introduction. It is the main purpose of this paper to compare the mean rates of convergence of two well-known schemes for solving self-adjoint partial difference equations of elliptic type: the Young-Frankel [6, 1] successive overrelaxation method, and the semi-iterative Chebyshev polynomial method as described by Lanczos [2, p. 42], Stiefel [3], and others. More generally, the analysis is applicable to any matrix equation of the form

$$\mathbf{A}\mathbf{u} = \mathbf{k},$$

provided the matrix $\mathbf{A} = || a_{i,j} ||$ is symmetric and positive definite, and, in the sense of Young [6, p. 93], satisfies property (A).

For semi-iterative methods, one considers iterates \mathbf{u}_i , where

(2)
$$\mathbf{u}_{i+1} = \mathbf{M}\mathbf{u}_i + \mathbf{k}$$
 $(i = 0, 1, 2, \cdots),$

M being a specific matrix. Then, one forms from the sequence of vectors \mathbf{u}_i a new sequence of vectors

(3)
$$\mathbf{t}_n = \sum_{j=0}^n \nu_j(n) \; \mathbf{u}_j \qquad (n = 0, 1, 2, \cdots),$$

the constants $\nu_j(n)$ being real numbers. Such a procedure is called a semiiterative method with respect to the matrix **M**.

While it is known [8, p. 293] that the successive overrelaxation method converges at least twice as fast as any semi-iterative method with respect to the Jacobi method, we shall give a different proof of this result, which generalizes to semi-iterative methods with respect to the Gauss-Seidel method. For the Gauss-Seidel method, the result is that the successive overrelaxation method converges at least as fast as any semi-iterative method with respect to the Gauss-Seidel method.

When solution by high-speed computing machines is involved, it should be noted that the successive overrelaxation method has the further advantage of requiring no auxiliary storage of extra iterates \mathbf{u}_m , whereas semi-iterative methods require that a few iterates \mathbf{u}_m be stored, along with suitable coefficients.

In a final section, an analogous discussion is given for semi-iterative

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methods applied to matrices whose (complex) eigenvalues are known to be confined to a *circle* in the complex plane. The result is that the best semi-iterative method with respect to the successive overrelaxation method, assuming only that all eigenvalues λ of the successive overrelaxation method satisfy $|\lambda| \leq \rho < 1$, is simply the basic method repeated n times, where n is the order of the semi-iterative method.

2. Description of methods. Equation (1) is rewritten in the form

$$\mathbf{u} = \mathbf{B}\mathbf{u} + \mathbf{f}$$

where the $N \times N$ matrix $\mathbf{B} = ||b_{i,j}||$ and the vector \mathbf{f} are given by

(5)
$$b_{i,j} = \begin{cases} -a_{i,j}/a_{i,i}, & i \neq j \\ 0, & i = j \end{cases}, \qquad f_i = k_i/a_{i,i},$$

for $i = 1, 2, \dots, N$. For the successive overrelaxation method, one forms the sequence of vectors defined by

(6)
$$u_i^{(m+1)} = \omega \left\{ \sum_{j=1}^{i-1} b_{i,j} u_j^{(m+1)} + \sum_{j=i+1}^{N} b_{i,j} u_j^{(m)} + f_i \right\} + (1 - \omega) u_i^{(m)}$$

where the "overrelaxation factor" ω is a parameter which is fixed throughout the course of iteration. The equation above may be written symbolically as

(7)
$$\mathbf{u}_{m+1} = \mathbf{L}_{\sigma,\omega}[\mathbf{u}_m] + \mathbf{g},$$

where g is a fixed vector, and $\mathbf{L}_{\sigma,\omega}$ denotes a linear operator.

Let \mathbf{u}_0 be a "trial solution" of (1), and let the sequences of vectors \mathbf{u}_j and \mathbf{t}_n be defined, respectively, by (2) and (3). If \mathbf{u}_0 is the unique solution of (1), then each \mathbf{t}_n is also a solution of (1) if and only if

(8)
$$\sum_{j=0}^{n} \nu_{j}(n) = 1 \qquad (n = 0, 1, 2, \cdots).$$

If \mathbf{v} denotes the unique solution of (1), and the *n*th error vector $\mathbf{t}_n - \mathbf{v}$ is denoted by $\boldsymbol{\varepsilon}_n^*$, then

(9)
$$\mathbf{\epsilon}_n^* = p_n(\mathbf{M})\mathbf{\epsilon}_0, \qquad p_n(x) \equiv \sum_{i=0}^n \nu_i(n)x^i,$$

and $\,\epsilon_0\,=\,u_0-v$ is the error vector associated with the trial solution $u_0\,.$ As

¹ Since **A** is symmetric and positive definite, (1) has a unique solution, and $a_{i,i} > 0$ for all i.

² For polynomials similarly normalized, see [2, p. 41], and [3, p. 63].

a consequence of (8), we shall henceforth assume $p_n(1) = 1$. For comparison, we note that in the successive overrelaxation method

(10)
$$\mathbf{\varepsilon}_n = L_{\sigma,\omega}^n[\mathbf{\varepsilon}_0],$$

which corresponds to the choice $p_n(x) = x^n$, and $\mathbf{M} = \mathbf{L}_{\sigma,\omega}$.

If the matrix **M** has real eigenvalues λ_k , and the interval $a \leq x \leq b$ is the smallest interval containing the λ_k 's, we define

(11)
$$\mu[p_n(\mathbf{M})] = \max_{a \le x \le b} \{ \mid p_n(x) \mid^{1/n} \},$$

$$R[p_n(\mathbf{M})] = -\log \mu[p_n(\mathbf{M})].$$

The quantities $\mu[p_n(\mathbf{M})]$ and $R[p_n(\mathbf{M})]$ are respectively the average spectral norm and the average rate of convergence³ at the n^{th} step of the semi-iterative method with respect to the matrix \mathbf{M} . For $p_n(x) = x^n$, we see that $\mu[p_n(\mathbf{M})]$ and $R(p_n(\mathbf{M})]$ are independent of n, and are, for matrices \mathbf{M} with real eigenvalues, respectively the usual definition of the spectral norm $\bar{\mu}$ and rate of convergence R of \mathbf{M} [6, p. 96].

3. Basis of Comparisons. In this section, we shall compare the rate of convergence of the successive overrelaxation method with that of the semi-iterative method with respect to the Jacobi method. We choose the matrix \mathbf{M} of (2) to be the matrix \mathbf{B} defined by (5). Under our initial assumptions, \mathbf{A} is symmetric and positive definite, and satisfies property (A). For this case, it is known [6] that all the eigenvalues λ of \mathbf{B} are real, and lie in the symmetric interval $-\bar{\mu} \leq \lambda \leq +\bar{\mu} < 1$, where $\bar{\mu}$ is the spectral norm of \mathbf{B} . Without loss of generality, we may assume that \mathbf{A} is consistently ordered [6, p. 93]. The best choice [6] of ω is given by

(12)
$$\omega_b = 1 + \left[\frac{\bar{\mu}}{1 + \left[1 - \bar{\mu}^2 \right]^{\frac{1}{2}}} \right]^2,$$

and

(13)
$$R[\mathbf{L}_{\sigma,\omega_b}] = -\log(\omega_b - 1).$$

To select the best semi-iterative method with respect to the matrix **B**, we choose the polynomial $p_n(x)$ such that

(14)
$$\max_{-\bar{\mu} \le x \le \bar{\mu}} |p_n(x)|$$

$$\| (p_n(\mathbf{M}))^k \mathbf{\epsilon}_0 \| \leq e^{-n} \| \mathbf{\epsilon}_0 \|,$$

where || x || refers to the Euclidean norm of the vector x.

³ Specifically, if the largest degree of the elementary divisors [5, Ch. III] of the matrix $p_n(\mathbf{M})$, for n fixed, is unity, then the reciprocal of $R[p_n(\mathbf{M})]$ is an estimate of the least positive integer k for which

is smallest. It is well-known that the solution of this problem is in Chebyshev polynomials, and we have explicitly

(15)
$$p_n(x) = \frac{T_n(x/\bar{\mu})}{T_n(1/\bar{\mu})},$$

where $T_n(x) = \cos[n \cos^{-1} x]$ is the Chebyshev polynomial of degree n. We shall call this particular method the Chebyshev semi-iterative method with respect to the Jacobi method. By definition, we have

$$\max_{-\bar{\mu} \leq x \leq \bar{\mu}} \left| p_n(x) \right| = \frac{\max\limits_{-1 \leq x \leq 1} \left| T_n(x) \right|}{\left| T_n(1/\bar{\mu}) \right|} = \frac{1}{T_n(1/\bar{\mu})},$$

since $T_n(\alpha) > 1$ for $\alpha > 1$. Forming the ratio of the average rates of convergence of the successive overrelaxation method and this Chebyshev semi-iterative method, we have

$$R[\mathbf{L}_{\sigma\omega_b}]/R[p_n(\mathbf{B})] = \left[\frac{2\log\left[1+(1-\overline{\mu}^2)^{\frac{1}{2}}\right]}{\overline{\mu}}\right] / \left(\frac{\log T_n(1/\overline{\mu})}{n}\right).$$

Since, for $\alpha > 1$, $\log T_n(\alpha)/n$ increases monotonically to $\cosh^{-1}(\alpha)$, and since $\cosh^{-1}(\alpha) = \log (\alpha + \sqrt{\alpha^2 - 1})$ for $\alpha > 1$, we have, replacing α by $1/\bar{\mu}$, $R[\mathbf{L}_{\sigma,\omega_b}]/R[p_n(\mathbf{B})] \geq 2$. This proves in a different manner the following theorem of Young [8, p. 293].

Theorem 1. The successive overrelaxation method with the optimum overrelaxation factor converges at least twice as fast as the Chebyshev semi-iterative method with respect to the Jacobi method, and therefore at least twice as fast as any semi-iterative method with respect to the Jacobi method. Furthermore, as the number of iterations tends to infinity, the successive overrelaxation method becomes exactly twice as fast as this Chebyshev semi-iterative method.

We finally remark that for values of ω near the optimum ω , ω_b , the successive overrelaxation method still has a faster rate of convergence than the Chebyshev semi-iterative method with respect to the Jacobi method, and therefore a faster rate of convergence than any semi-iterative method with respect to the Jacobi method.

4. Extension to polynomials⁴ in $\mathbf{L}_{\sigma,1}$. In this section, we merely replace the Jacobi matrix \mathbf{B} in the previous section by the Gauss-Seidel or Liebmann operator $\mathbf{L}_{\sigma,1}$ [1, 6]. Assuming \mathbf{A} to be consistently ordered, it is then known [6, p. 100] that the eigenvalues λ_k of $\mathbf{L}_{\sigma,1}$ satisfy $0 \leq \lambda_k \leq \overline{\mu}^2 < 1$, where $\overline{\mu}$ is the spectral norm of the matrix \mathbf{B} . As before, the polynomial $p_n(x)$ of degree n, normalized so that $p_n(1) = 1$, which has the property

⁴ See also [4, Chapter VI].

that $\max_{0 \le x \le \bar{\mu}^2} |p_n(x)|$ is smaller than all other such normalized polynomials, will be

(16)
$$p_n(x) = T_n \left(\frac{2x}{\bar{\mu}^2} - 1 \right) / T_n \left(\frac{2}{\bar{\mu}^2} - 1 \right).$$

The semi-iterative method based on these polynomials will be called the Chebyshev semi-iterative method with respect to the operator $L_{\sigma,1}$. The average rate of convergence of this iterative scheme at the *n*th step is

(17)
$$R[p_n(\mathbf{L}_{\sigma,1})] = \log T_n\left(\frac{2}{\overline{\mu}^2} - 1\right) / n.$$

If we form the ratio of $R[\mathbf{L}_{\sigma,\omega_b}]$ to $R[p_n(\mathbf{L}_{\sigma,1})]$, then using the monotone property of $\log T_n(\alpha)/n$ and the previously used identity for $\cosh^{-1}(\alpha)$ for $\alpha > 1$, we obtain

(18)
$$R[\mathbf{L}_{\sigma,\omega_{\delta}}]/R[p_{n}(L_{\sigma,1})] \\ \geq 2\log\left(\frac{1}{\overline{\mu}} + \left(\frac{1}{\overline{\mu}^{2}} - 1\right)^{\frac{1}{2}}\right) / \log\left(\frac{2}{\overline{\mu}^{2}} - 1 + \left(\left(\frac{2}{\overline{\mu}^{2}} - 1\right)^{2} - 1\right)^{\frac{1}{2}}\right).$$

But the right hand side of the inequality above reduces identically to unity for $\bar{\mu} < 1$. This proves

Theorem 2. The successive overrelaxation method with the optimum overrelaxation factor converges at least as fast as the Chebyshev semi-iterative method with respect to the operator $\mathbf{L}_{\sigma,1}$, and therefore at least as fast as any semi-iterative method with respect to the operator $\mathbf{L}_{\sigma,1}$. Furthermore, as the number of iterations tends to infinity, the successive overrelaxation method becomes exactly as fast as the Chebyshev semi-iterative method with respect to the operator $\mathbf{L}_{\sigma,1}$.

5. Extensions to polynomials in the operator $\mathbf{L}_{\sigma,\omega}$. If we have, as before, that the eigenvalues of \mathbf{B} are real and lie in $-\bar{\mu} \leq x \leq \bar{\mu}$, then we can formulate the problem of finding the best polynomial of degree n, normalized so that $p_n(1) = 1$, having the smallest absolute value on the interval $-\bar{\mu} \leq x \leq \bar{\mu}$, and we are naturally led to Chebyshev polynomials. With these polynomials, we then defined the Chebyshev semi-iterative method with respect to the matrix \mathbf{B} , which was, in some sense, the optimum semi-iterative method with respect to the matrix \mathbf{B} . The same is true if we consider, rather than the matrix \mathbf{B} , the linear operator $\mathbf{L}_{\sigma,1}$ whose eigenvalues λ_k are also real and satisfy $0 \leq \lambda_k \leq \bar{\mu}^2$, and optimize the selection of a sequence of normalized polynomials whose absolute value on the interval $0 \leq x \leq \bar{\mu}^2$ is smallest. The resulting semi-iterative method defined by this sequence of polynomials was called the Chebyshev semi-iterative method with respect to the operator $\mathbf{L}_{\sigma,1}$. As we pass to the case where

 $1 < \omega < 2$, the operator $\mathbf{L}_{\sigma,\omega}$ does not have all real eigenvalues [6, p. 101], and the selection of a sequence of normalized polynomials to define a semi-iterative method with respect to the operator $\mathbf{L}_{\sigma,\omega}$ is not immediate. As before, we have

(19)
$$\mathbf{\epsilon}_n^* = \sum_{j=0}^n \nu_j(n) \mathbf{L}_{\sigma,\omega}^j \mathbf{\epsilon}_0 = p_n(\mathbf{L}_{\sigma,\omega}) \mathbf{\epsilon}_0,$$

where $p_n(1) = 1$. We now assume that the eigenvalues λ_k of $\mathbf{L}_{\sigma,\omega}$ satisfy $|\lambda_k| \leq \rho < 1$. If $g_n(z)$ is any complex polynomial of degree n, let $M_{g_n}(r)$ denote the maximum modulus function of $g_n(z)$, i.e.,

$$M_{g_n}(r) = \max_{|z| \leq r} |g_n(z)|.$$

Completely analogous to the previous sections, let S_n be the set of all polynomials $g_n(z)$ of degree n for which $g_n(1) = 1$, and consider

$$\min_{g_n \in S} \ \{M_{g_n}(r)\}.$$

The following theorem, due to E. H. Zarantonello⁵, seems to be of interest by itself.

THEOREM 3. For all r such that $0 \le r \le 1$, $\min_{g_n \in S} \{M_{g_n}(r)\} = r^n$ for all positive integers n.

Proof. As is well known, we have

$$M_{g_n}(r) = \lim_{p \to \infty} \left\{ \int_0^{2\pi} |g_n(re^{i\theta})|^{2p} d\theta \right\}^{1/2p}.$$

Clearly, we have for any positive integer p,

$$\min_{g_n \in S_n} \left\{ \int_0^{2\pi} \mid g_n(re^{i\theta}) \mid^{2p} d\theta \right\}^{1/2p} \, \geq \, \min_{Q \in S_{np}} \left[\left(\int_0^{2\pi} \mid Q(re^{i\theta}) \mid^2 d\theta \right)^{\frac{1}{2}} \right]^{1/p}$$

since if $g_n \in S_n$, then

$$Q(z) = [g_n(z)]^p \in S_{np}.$$

 \mathbf{If}

$$Q(z) = \sum_{k=0}^{np} a_k z^k,$$

then

$$\left(\int_{0}^{2\pi}\mid Q(re^{i heta})\mid^{2}d heta
ight)^{rac{1}{2}}=\sqrt{2\pi}\left(\sum_{k=0}^{np}\mid a_{k}\mid^{2}r^{2k}
ight)^{rac{1}{2}}.$$

⁵ Personal communication.

But

$$\sqrt{2\pi} \left(\sum_{k=0}^{np} |a_k|^2 r^{2k} \right)^{\frac{1}{2}} \ge \sqrt{2\pi} r^{np} \left(\sum_{k=0}^{np} |a_k|^2 \right)^{\frac{1}{2}}$$

since $0 \le r \le 1$. Since $Q \in S_{np}$, then $\sum_{k=0}^{np} a_k = 1$. Therefore, using Schwarz' inequality, we have

$$1 = \left| \sum_{k=0}^{np} a_k \right| \le \sum_{k=0}^{np} |a_k| \le (np + 1)^{\frac{1}{2}} \left(\sum_{k=0}^{np} |a_k|^2 \right)^{\frac{1}{2}},$$

and

$$\left\{ \int_{0}^{2\pi} |Q(re^{i\theta})|^{2} d\theta \right\}^{\frac{1}{2}} \geq \sqrt{2\pi} r^{np} \left(\sum_{k=0}^{np} |a_{k}|^{2} \right)^{\frac{1}{2}} \geq \left(\frac{2}{np+1} \right)^{\frac{1}{2}} r^{np}.$$

Thus,

$$\min_{Q \in \mathcal{S}_{np}} \left[\left(\int_0^{2\pi} |Q(re^{i\theta})|^2 d\theta \right)^{\frac{1}{2}} \right]^{1/p} \ge r^n \left(\frac{2}{np+1} \right)^{1/2p}.$$

Letting $p \to \infty$, we have

$$\min_{g_n \in S_n} M_{g_n}(r) \ge r^n.$$

Since $z^n \in S_n$, and $M_{z^n}(r) = r^n$, then

$$\min_{g_n \in S_n} M_{g_n}(r) = r^n \quad \text{for all} \quad 0 \le r \le 1,$$

for all positive integers n.

In view of this theorem, we have

Theorem 4. The best semi-iterative method with respect to the successive overrelaxation method, $\omega > 1$, which can be obtained, assuming only that eigenvalues λ_k of the successive overrelaxation method satisfy $|\lambda| \leq \rho < 1$, is simply the basic method repeated n times, where n is the order of the semi-iterative method.

We remark that, for $\omega > 1$, the error vector associated with the best semi-iterative method with respect to $\mathbf{L}_{\sigma,\omega}$ satisfies

(20)
$$\mathbf{\epsilon}_n^* = \mathbf{L}_{\sigma,\omega}^n \mathbf{\epsilon}_0 = \mathbf{\epsilon}_n .$$

This particularly simple form of a semi-iterative method with respect to $\mathbf{L}_{\sigma,\omega}$ incidentally has been used repeatedly to solve multigroup diffusion problems in two or more (space) dimensions [9, 10]. Apparently, the choice for the iteration method is not related to the theorem above, but rather to the inherent simplicity of the iteration method of (20).

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