

CS 59300 – Algorithms for Data Science

Classical and Quantum approaches

Lecture 21 (11/25)

**Quantum Gibbs Sampling and Open Quantum
Systems**

https://ruizhezhang.com/course_fall_2025.html

Outline

- Motivation
- Description of open quantum system dynamics
- Quantum simulation algorithms
- Applications

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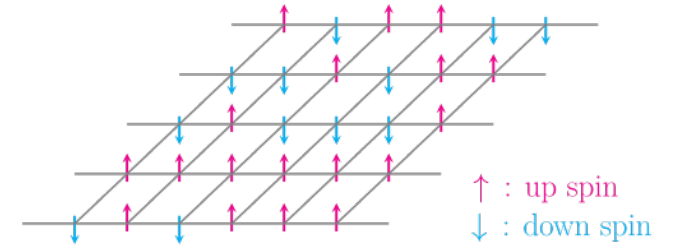
Classical Gibbs sampling: revisited

- Consider a classical spin system (e.g., the Ising model):

$$H_{\text{Ising}} = - \sum_{i \sim j} Z_i Z_j$$

- Eigenstates $\{0, 1\}^n$ and eigenvalues $\{E_x\}_{x \in \{0, 1\}^n}$
- The goal is to (approximately) sample from the Gibbs distribution:

$$\pi_\beta(x) = \frac{e^{-\beta E_x}}{Z_\beta}, \quad Z_\beta = \sum_x e^{-\beta E_x}$$



The Metropolis-Hastings algorithm

Set of **jump operators** $\{A^a\}$ such that A^a is selected with probability $p(a)$

- **Example:** flip a randomly chosen spin

Let x be the current state

- Pick A^a with probability $p(a)$
- $y \leftarrow A^a$ applied to x
- **Accept** move with probability $\gamma_\beta(x, y) := \min \left\{ 1, \exp \left(-\beta (E_y - E_x) \right) \right\}$
- If **reject**, stay at x

Fact. Under certain conditions, the Gibbs distribution π_β is the unique fixed point.

- **Detailed balance (DBC):** $e^{-\beta E_x} \gamma_\beta(x, y) = e^{-\beta E_y} \gamma_\beta(y, x)$

Quantum Gibbs sampling

- Let H be a quantum Hamiltonian with eigenstates $\{|\psi_j\rangle\}$ and eigenvalues $\{E_j\}$
- The goal is to (approximately) sample from the “quantum Gibbs distribution”:

$$\pi_\beta(|\psi_j\rangle) = \frac{e^{-\beta E_j}}{Z_\beta}$$

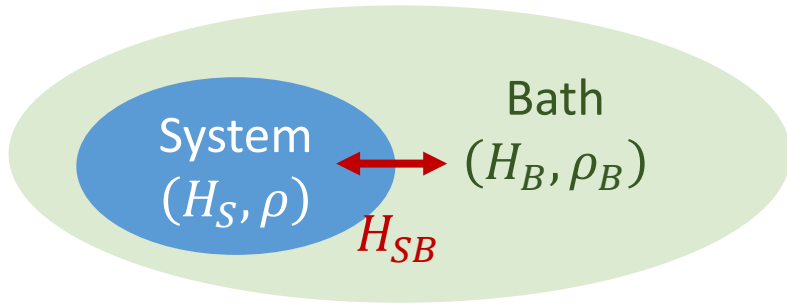
- Mixed state:

$$\rho_\beta = \sum \frac{e^{-\beta E_j}}{Z_\beta} |\psi_j\rangle\langle\psi_j| = \frac{\exp(-\beta H)}{\text{tr}[\exp(-\beta H)]} \quad \text{Gibbs state}$$

Challenges:

1. Given $|\psi_j\rangle$ cannot calculate the energy E_j exactly
2. Rejection requires backing up after a quantum measurement

Open quantum system



System dimension: $d \ll$ Bath dimension: D

- Total Hamiltonian $H = H_S + H_B + H_{SB}$
- The whole system evolution follows the [Liouville-von Neumann equation](#):

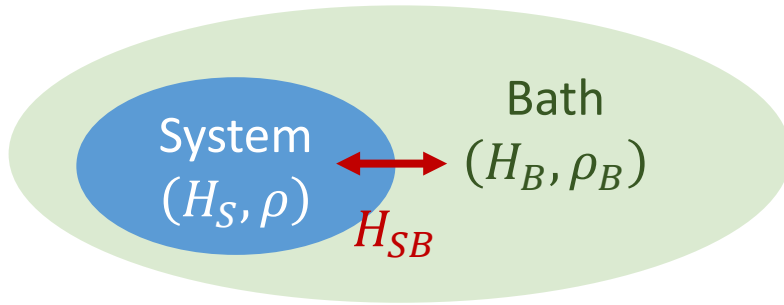
$$\frac{d\rho_{SB}}{dt} = -\mathbf{i}[H, \rho_{SB}]$$

- We only care the system part, i.e., $\rho = \text{tr}_B[\rho_{SB}]$

Is it possible to describe the dynamics of ρ **without** simulating the bath?

i.e. a kind of **dimension reduction**

Open quantum system



System dimension: $d \ll$ Bath dimension: D

Thermalization

- If you leave a quantum system in contact with a heat bath at temperature $T = 1/\beta$, then the state $\rho(t) = \text{tr}_B[\rho_{SB}(t)]$ **forgets** its initial condition and converges to the Gibbs state:

$$\rho_\beta = \frac{\exp(-\beta H_S)}{\text{tr}[\exp(-\beta H_S)]}$$

- A **physical** approach to quantum Gibbs sampling:

Put the quantum spin system into a refrigerator and wait for thermalization 😊

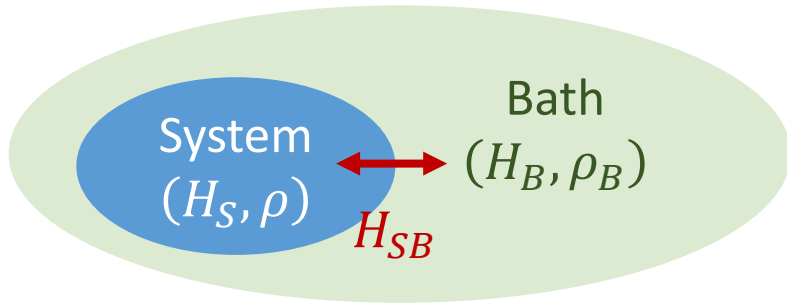
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Open quantum system



Is it possible to describe the dynamics of ρ without simulating the bath?

Gorini-Kossakowski-Sudarshan-Lindblad master equation:

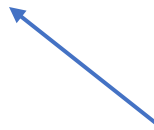
$$\frac{d\rho}{dt} = -\mathbf{i}[H, \rho] + \sum_{j=1}^m \left(L_j \rho L_j^\dagger - \frac{1}{2} \{L_j^\dagger L_j, \rho\} \right) =: \mathcal{L}\rho$$

Lindbladian /
Lindblad generator

- $H \in \mathbb{C}^{d \times d}$ is the system Hamiltonian (but is not equal to H_S)
- $L_j \in \mathbb{C}^{d \times d}$ are called the jump operators

Mathematical derivation of the master equation

We want to find a linear map \mathcal{L} such that $V_t = e^{\mathcal{L}t}$ is **completely-positive and trace-preserving (CPTP)**

- V is **CP** if and only if $V(\rho) = \sum_i L_i \rho L_i^\dagger$
 - V is **TP** if and only if $\sum_i L_i L_i^\dagger = I$
- Quantum Markov semigroup
- 

First consider the CP condition:

- $\mathcal{L} = \lim_{\Delta t \rightarrow 0} \frac{V_{\Delta t} - I}{\Delta t}$
- Expand \mathcal{L} under a generic basis K_1, K_2, \dots, K_{d^2} such that $\text{tr}[K_i^\dagger K_j] = \delta_{ij}$
 - Set $K_1 = \frac{1}{\sqrt{d}} I$ so that $\text{tr}[K_i] = 0 \quad \forall i \geq 2$
- $V_{\Delta t}(\rho) = \sum_i L_i \rho L_i^\dagger = \sum_{i,j} c_{i,j}(\Delta t) K_i \rho K_j^\dagger$

Mathematical derivation of the master equation

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Quantum Markov semi-group

First consider the CP condition:

$$\begin{aligned}\mathcal{L}\rho &= \sum_{i,j=2}^{d^2} g_{ij} K_i \rho K_j^\dagger + \underbrace{\sum_{i=2}^{d^2} (g_{i1} K_i \rho + g_{1i} \rho K_i^\dagger)}_{\rho G + G^\dagger \rho} + g_{11} \rho \\ &= \sum_{i,j=2}^{d^2} g_{ij} K_i \rho K_j^\dagger + \rho F + F \rho - \mathbf{i}[H, \rho]\end{aligned}$$

$$G = F + \mathbf{i}H$$

Hermitian anti-Hermitian

Mathematical derivation of the master equation

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Quantum Markov semi-group

Then consider the TP condition:

$$\mathcal{L}\rho = \sum_{i,j=2}^{d^2} g_{ij} K_i \rho K_j^\dagger + F\rho + \rho F - \mathbf{i}[H, \rho]$$

- $\text{tr} \left[\frac{d\rho}{dt} \right] = \text{tr}[\mathcal{L}\rho] = 0 \quad \Rightarrow \quad \text{tr} \left[\left(\sum_{i,j=2}^{d^2} g_{ij} K_j^\dagger K_i + 2F \right) \rho \right] = 0 \quad \forall \rho$
- $F = -\frac{1}{2} \sum_{i,j=2}^{d^2} g_{ij} K_j^\dagger K_i$

Mathematical derivation of the master equation

We want to find a linear map \mathcal{L} such that $V_t = e^{\mathcal{L}t}$ is **completely-positive and trace-preserving (CPTP)**

- V is **CP** if and only if $V(\rho) = \sum_i L_i \rho L_i^\dagger$
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Quantum Markov Semigroup

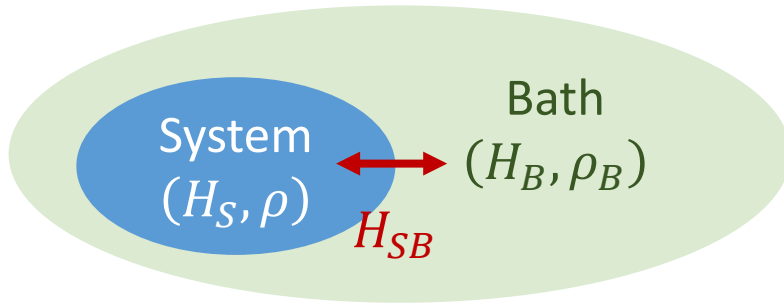


By diagonalization, we get that

$$\mathcal{L}\rho = -\mathbf{i}[H, \rho] + \sum_{j=1}^m \left(L_j \rho L_j^\dagger - \frac{1}{2} \{L_j^\dagger L_j, \rho\} \right)$$

- This derivation gives the most general form of Quantum Markov Semigroup (QMS)
- But H and L_j are **decoupled** from the physical system

Derivation of the master equation



System dimension: $d \ll$ Bath dimension: D

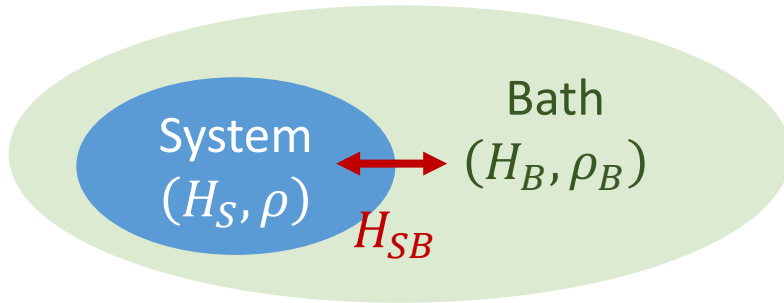
It remains open to derive the master equation for a given open quantum system from first principles

$$H(t) = H_S \otimes I_B + I_S \otimes H_B + g A_\alpha \otimes B_\alpha$$

Assumptions:

- **Born approximation:** $\rho_{SB}(t) = \rho(t) \otimes \rho_B$ and ρ_B is stationary
 - **Markov approximation:** memoryless
 - **Secular approximation (rotating wave approximation)**
- Weak coupling assumption** \leftarrow g is very small

Derivation of the master equation



System dimension: $d \ll$ Bath dimension: D

It remains open to derive the master equation for a given open quantum system from first principles

$$H(t) = H_S \otimes I_B + I_S \otimes H_B + g A_\alpha \otimes B_\alpha$$

Davies generator

$$\mathcal{L}\rho = -i[H_S + H_{LS}, \rho] + g^2 \sum_{\omega} \sum_a \left(L_{a,\omega}^\dagger \rho L_{a,\omega} - \frac{1}{2} \{L_{a,\omega}^\dagger L_{a,\omega}, \rho\} \right)$$

Lamb shift

Bohr frequency $\{\lambda_i - \lambda_j : \lambda_i, \lambda_j \in \text{spec}(H_S)\}$

Outline

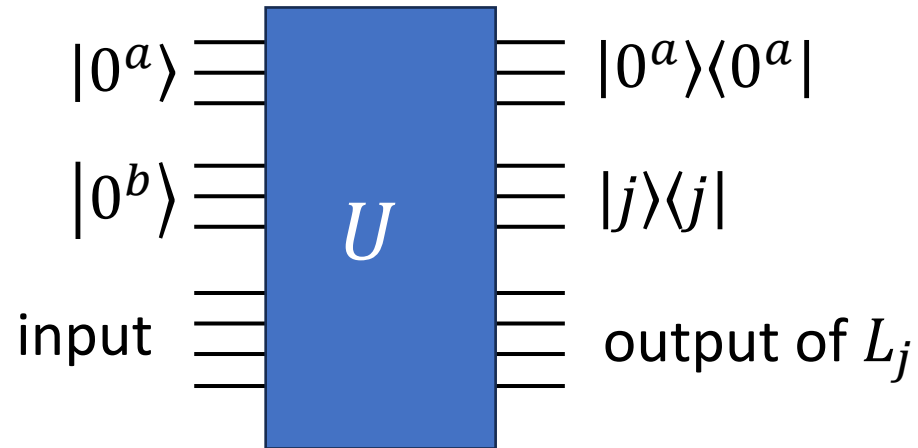
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Input model

Block-encoding for the jump operators $\{L_j\}$



$$\sum_j |j\rangle\langle j| \otimes L_j$$

The goal is to simulate $e^{\mathcal{L}t}$ on an initial state $\rho(0)$

$$\mathcal{L}\rho = -\mathbf{i}[H, \rho] + \sum_{j=1}^m \left(L_j \rho L_j^\dagger - \frac{1}{2} \{L_j^\dagger L_j, \rho\} \right)$$

Quantum simulation algorithms

1. “Quantizing” the classical numerical solvers for ODE/SDE
2. “Quantizing” the Metropolis-Hastings algorithm

A toy example

$$\mathcal{L}\rho = L\rho L^\dagger - \frac{1}{2}\{L^\dagger L, \rho\}$$

- Consider the dilated Hamiltonian: $\tilde{H} = \begin{bmatrix} & L^\dagger \\ L & \end{bmatrix}$

- We construct a quantum channel:

$$U_D = e^{-i\tilde{H}\sqrt{\Delta t}}, \quad \Phi(\rho) := \text{tr}_a[U_D(|0\rangle\langle 0|_a \otimes \rho)U_D^\dagger]$$

Hamiltonian
simulation

- Taylor expanding U_D to the second order:

$$U_D = I - i\sqrt{\Delta t} \begin{bmatrix} & L^\dagger \\ L & \end{bmatrix} - \frac{\Delta t}{2} \begin{bmatrix} L^\dagger L & \\ & LL^\dagger \end{bmatrix} + O(\Delta t^{1.5})$$

- We can check that

$$\|e^{\mathcal{L}\Delta t}(\rho) - \Phi(\rho)\|_1 \sim (\Delta t)^2$$

- For $T = r\Delta t$, we have $\|e^{\mathcal{L}T}(\rho) - \Phi^r(\rho)\|_1 \sim T\Delta t$

First-order accuracy

High-order Lindblad simulation

$$\begin{aligned}
 \mathcal{L}\rho &= -\mathbf{i}[H, \rho] + \sum_{j=1}^m \left(L_j \rho L_j^\dagger - \frac{1}{2} \{L_j^\dagger L_j, \rho\} \right) \\
 &= \underbrace{J\rho + \rho J^\dagger}_{\mathcal{L}_D \rho} + \underbrace{\sum_{j=1}^m L_j \rho L_j^\dagger}_{\mathcal{L}_J \rho} \\
 &\quad \text{“Drifting”} \qquad \qquad \text{“Jump”}
 \end{aligned}$$

$$J = -\mathbf{i}H_{\text{eff}}, \qquad H_{\text{eff}} := H + \frac{1}{2\mathbf{i}} \sum_{j=1}^m L_j^\dagger L_j$$

High-order Lindblad simulation

By Duhamel's principle,

$$\rho(t) = e^{\mathcal{L}t}\rho(0) = e^{\mathcal{L}_D t}\rho(0) + \int_0^t e^{\mathcal{L}_D(t-s)} \mathcal{L}_J \rho(s) ds$$

Apply Duhamel's principle again to $\rho(s)$:

$$\begin{aligned} \rho(t) = & e^{\mathcal{L}_D t}\rho(0) + \int_0^t e^{\mathcal{L}_D(t-s)} \mathcal{L}_J e^{\mathcal{L}_D s} \rho(0) ds \\ & + \int_{0 \leq s_1 \leq s_2 \leq t} e^{\mathcal{L}_D(t-s_2)} \mathcal{L}_J e^{\mathcal{L}_D(s_2-s_1)} \mathcal{L}_J \rho(s_1) ds_2 ds_1 \end{aligned}$$

truncation error

We can iterate this expansion for K -th order

High-order Lindblad simulation

$$\rho(t) = e^{\mathcal{L}_D t} \rho(0) + \sum_{k=1}^K \int_{0 \leq s_1 \leq \dots \leq s_k \leq t} e^{\mathcal{L}_D(t-s_k)} \mathcal{L}_J e^{\mathcal{L}_D(s_k-s_{k-1})} \mathcal{L}_J \dots e^{\mathcal{L}_D(s_2-s_1)} \mathcal{L}_J e^{\mathcal{L}_D s_1} \rho(0) ds_1 \dots ds_k + \text{truncation error}$$

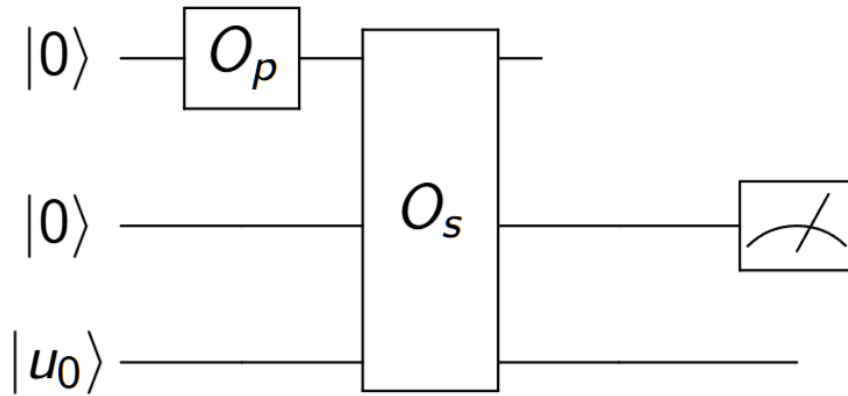
- The integrals can be discretized using **Gaussian quadrature**
- Need to block-encode $e^{\mathcal{L}_D t}$ and \mathcal{L}_J
- For the drifting part,

$$e^{\mathcal{L}_D t} \rho = e^{Jt} \rho e^{J^\dagger t} \approx \left(\sum_{\ell=0}^{K'} \frac{J^\ell t^\ell}{\ell!} \right) \rho \left(\sum_{\ell=0}^{K'} \frac{J^\ell t^\ell}{\ell!} \right)^\dagger \quad \left. \vphantom{\sum_{\ell=0}^{K'} \frac{J^\ell t^\ell}{\ell!}} \right\} \text{Kraus form}$$

- For the jump part, $\mathcal{L}_J = \sum_{j=1}^m L_j \rho L_j^\dagger$

Block-encoding Kraus operators

- Let A_1, A_2, \dots, A_m be Kraus operators and let U_i be the corresponding $(\alpha_i, n', 0)$ -block-encodings



$$O_p|0\rangle = \frac{1}{\|\alpha\|_2} \sum_{j=1}^m \alpha_j |j\rangle$$

$$O_s = \sum_{j=1}^m |j\rangle\langle j| \otimes U_j$$

- $|0\rangle|u_0\rangle \mapsto \frac{1}{\|\alpha\|_2} \sum_{j=1}^m |j\rangle A_j |u_0\rangle$
- Density matrix: $\rho_0 \mapsto \tilde{\rho} := \sum_{j,j'} |j\rangle\langle j'| \otimes A_j \rho_0 A_j^\dagger$
- Trace out the first ancilla register: $\text{tr}_a[\tilde{\rho}] = \sum_j A_j \rho_0 A_j^\dagger$

High-order Lindblad simulation

$$\begin{aligned}\rho(t) = & e^{\mathcal{L}_D t} \rho(0) \\ & + \sum_{k=1}^K \int_{0 \leq s_1 \leq \dots \leq s_k \leq t} e^{\mathcal{L}_D(t-s_k)} \mathcal{L}_J e^{\mathcal{L}_D(s_k-s_{k-1})} \mathcal{L}_J \dots e^{\mathcal{L}_D(s_2-s_1)} \mathcal{L}_J e^{\mathcal{L}_D s_1} \rho(0) ds_1 \dots ds_k \\ & + \text{truncation error}\end{aligned}$$

- The integrals can be discretized using **Gaussian quadrature**
- Use truncated Taylor series to approximate $e^{\mathcal{L}_D t}$
- Block-encode the Kraus forms and use a big LCU to combine all the terms
- Total complexity: $\mathcal{O}\left(t \|\mathcal{L}\|_{\text{be}} \log\left(\frac{t \|\mathcal{L}\|_{\text{be}}}{\epsilon}\right)\right)$ where $\|\mathcal{L}\|_{\text{be}} := \|H\| + \frac{1}{2} \sum_{j=1}^m \|L_j\|^2$

Li, Xiantao, and Chunhao Wang. "Simulating Markovian open quantum systems using higher-order series expansion."

Lindblad simulation via Hamiltonian simulation

$$\frac{d\rho}{dt} = J\rho + \rho J^\dagger + \sum_{j=1}^m L_j \rho L_j^\dagger, \quad J = -\mathbf{i}H - \frac{1}{2} \sum_{j=1}^m L_j^\dagger L_j$$

Unraveling

- Thinking about $\rho(t) = \mathbb{E}[|\psi_t\rangle\langle\psi_t|]$:

$$\frac{d\mathbb{E}[|\psi_t\rangle\langle\psi_t|]}{dt} = J\mathbb{E}[|\psi_t\rangle\langle\psi_t|] + \mathbb{E}[|\psi_t\rangle\langle\psi_t|]J^\dagger + \sum_{j=1}^m L_j \mathbb{E}[|\psi_t\rangle\langle\psi_t|] L_j^\dagger$$

- Reverse-engineer an **SDE** for $|\psi_t\rangle$ using Itô's lemma:

$$d|\psi_t\rangle = J|\psi_t\rangle dt + \sum_{j=1}^m L_j |\psi_t\rangle dW_t^j$$

independent 1d Brownian motion

Lindblad simulation via Hamiltonian simulation

$$d|\psi_t\rangle = J|\psi_t\rangle dt + \sum_{j=1}^m L_j |\psi_t\rangle dW_t^j$$

Euler-Maruyama scheme:

$$|\psi_{n+1}\rangle = |\psi_n\rangle + J|\psi_n\rangle\Delta t + \sum_{j=1}^m L_j |\psi_n\rangle \sqrt{\Delta t} W^j \stackrel{\sim \mathcal{N}(0,1)}{=} L_{1,\Delta t}(|\psi_n\rangle)$$

$$\begin{aligned} \mathbb{E}[|\psi_{n+1}\rangle\langle\psi_{n+1}|] &= \mathbb{E}[L_{1,\Delta t}(|\psi_n\rangle)L_{1,\Delta t}(|\psi_n\rangle)^\dagger] \\ &= \underbrace{(I + J\Delta t)}_{F_0} \mathbb{E}[|\psi_n\rangle\langle\psi_n|] (I + J\Delta t)^\dagger + \sum_{j=1}^m \underbrace{(\sqrt{\Delta t} L_j)}_{F_j} \mathbb{E}[|\psi_n\rangle\langle\psi_n|] (\sqrt{\Delta t} L_j)^\dagger \end{aligned}$$

Kraus form:

$$\rho_{n+1} = \sum_{j=0}^m F_j \rho_n F_j^\dagger$$

Lindblad simulation via Hamiltonian simulation

$$\rho_{n+1} = \sum_{j=0}^m F_j \rho_n F_j^\dagger =: \mathcal{K}[\rho_n]$$

- Since Euler-Maruyama is a first-order scheme,

$$\rho_{n+1} = \rho_n + \mathcal{L}\rho_n \Delta t + \mathcal{O}(\Delta t^2) = e^{\mathcal{L}\Delta t} \rho_n + \mathcal{O}(\Delta t^2)$$

Our goal is to find a dilated Hamiltonian \tilde{H} such that

$$\mathcal{K}[\rho] \approx \text{tr}_a \left[e^{-i\tilde{H}\sqrt{\Delta t}} (|0\rangle\langle 0| \otimes \rho) e^{i\tilde{H}\sqrt{\Delta t}} \right]$$

- **Stinespring representation** of a Kraus form:

Taylor expand and match the terms

$$\mathcal{K}[\rho] = \text{tr}_a \left[\begin{bmatrix} F_0 & * & * & * \\ F_1 & * & * & * \\ \vdots & * & * & * \\ F_m & * & * & * \end{bmatrix} (|0\rangle\langle 0| \otimes \rho) \begin{bmatrix} F_0 & * & * & * \\ F_1 & * & * & * \\ \vdots & * & * & * \\ F_m & * & * & * \end{bmatrix}^\dagger \right]$$

Lindblad simulation via Hamiltonian simulation

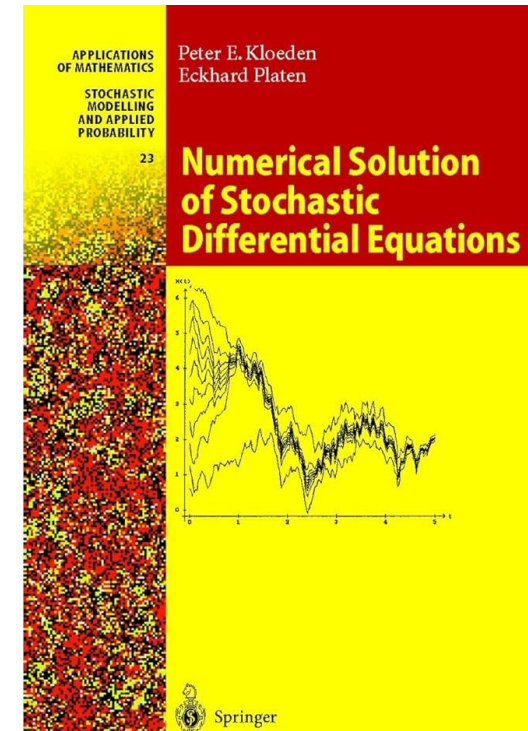
$$\tilde{H} = \begin{bmatrix} \sqrt{\Delta t} H & L_1^\dagger & \cdots & L_m^\dagger \\ L_1 & 0 & 0 & 0 \\ \vdots & 0 & 0 & 0 \\ L_m & 0 & 0 & 0 \end{bmatrix}, \quad \mathcal{K}[\rho] = \text{tr}_a \left[e^{-i\tilde{H}\sqrt{\Delta t}} (|0\rangle\langle 0| \otimes \rho) e^{i\tilde{H}\sqrt{\Delta t}} \right] + \mathcal{O}(\Delta t^2)$$

- This is a generalization of the toy example

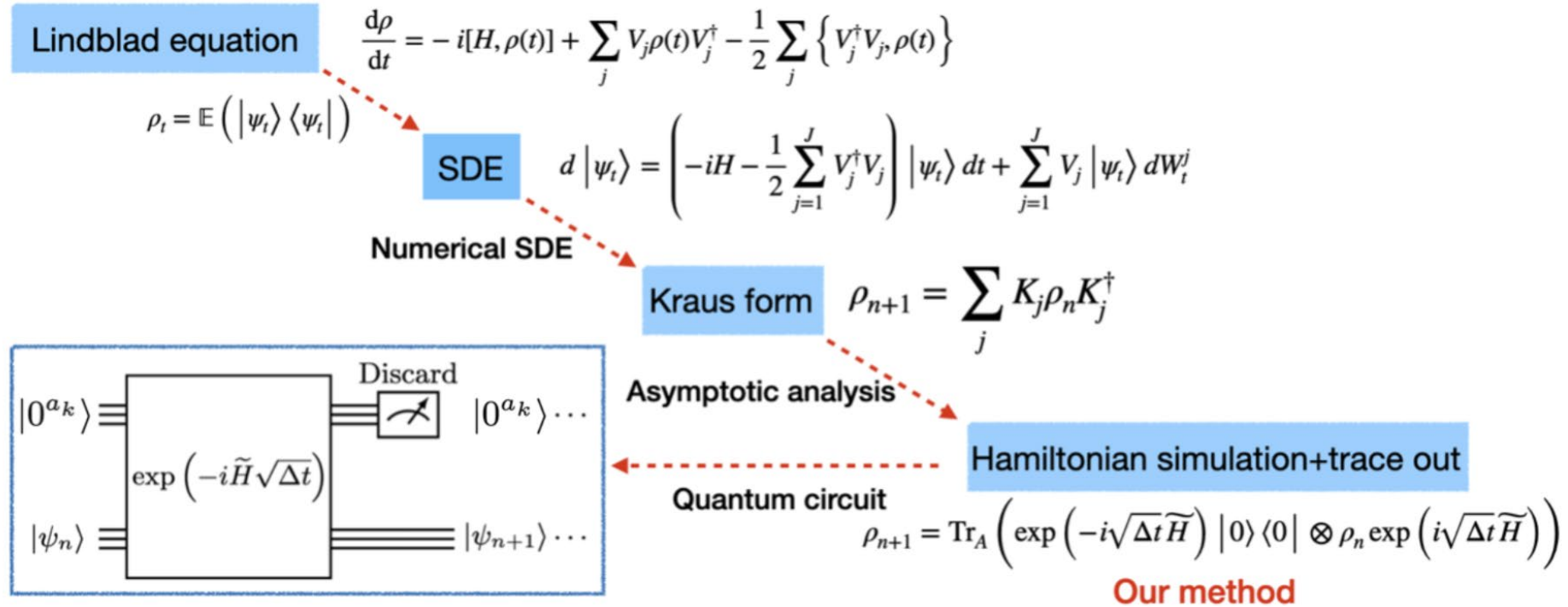
Three-step framework for k -th order Lindblad simulation

- **Step 1:** find a weak scheme of order k for the SDE (e.g., Itô-Taylor-expansion)
- **Step 2:** formulate the Kraus form
- **Step 3:** construct the dilated Hamiltonian

$$\tilde{H} = \begin{bmatrix} H_0 & H_1^\dagger & \cdots & H_S^\dagger \\ H_1 & 0 & 0 & 0 \\ \vdots & 0 & 0 & 0 \\ H_S & 0 & 0 & 0 \end{bmatrix}$$



Lindblad simulation via Hamiltonian simulation



For a k -th order scheme, the simulation error is $\|\rho_T - \rho_N\|_1 = \mathcal{O}(T\|\mathcal{L}\|_{\text{be}}\Delta t^k)$

Ding, Zhiyan, Xiantao Li, and Lin Lin. "Simulating open quantum systems using Hamiltonian simulations."

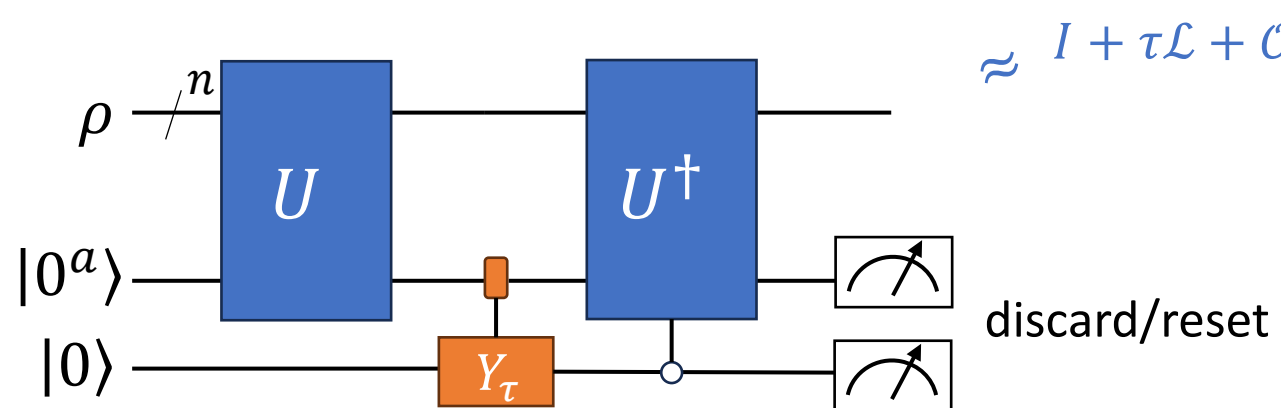
Quantum Metropolis algorithm

Let x be the current state

- Pick A^a with probability $p(a)$
- $y \leftarrow A^a$ applied to x
- **Accept** move with probability $\gamma_\beta(x, y) := \min \{1, \exp(-\beta(E_y - E_x))\}$
- If **reject**, stay at x

How to revert in quantum?

Weak measurement

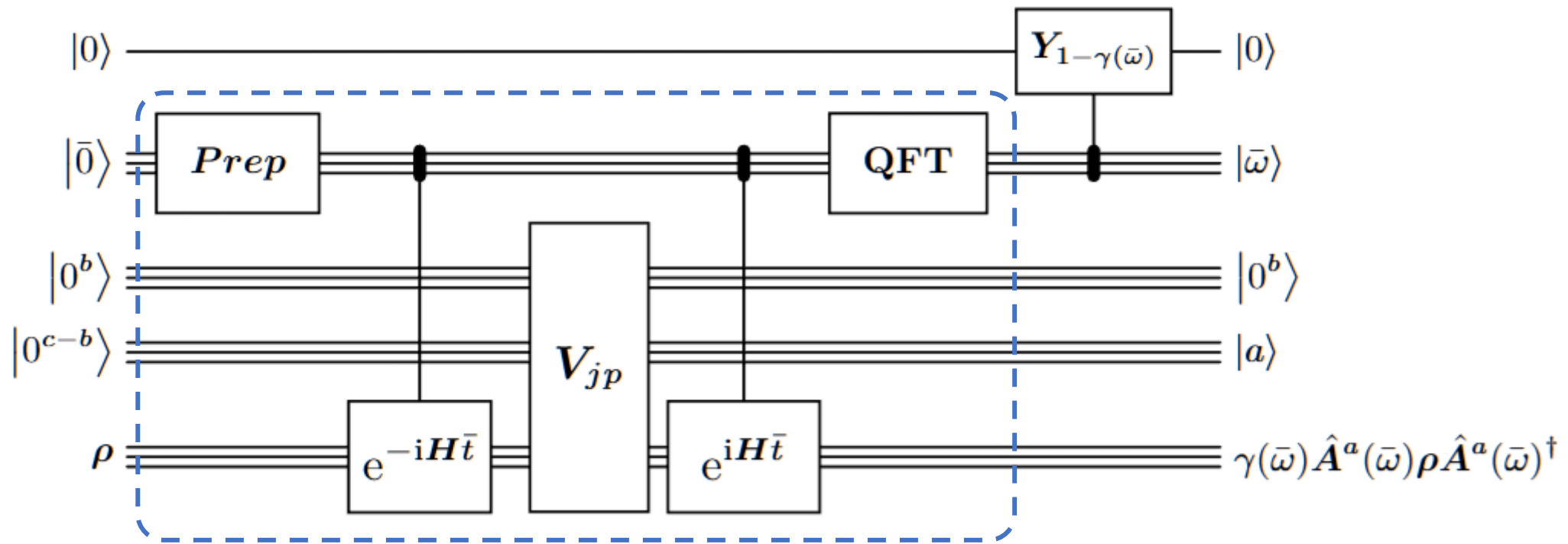


$$Y_\tau = |0^a\rangle\langle 0^a| \otimes \begin{bmatrix} \sqrt{1-\tau} & -\sqrt{\tau} \\ \sqrt{\tau} & \sqrt{1-\tau} \end{bmatrix} + (I - |0^a\rangle\langle 0^a|) \otimes I$$

Block-encoding the Lindbladian

$$\sum_{a, \bar{\omega}} \sqrt{\gamma(\bar{\omega})} |\bar{\omega}\rangle \otimes |a\rangle \otimes \hat{A}^a(\bar{\omega})$$

Calculate the “accept” probability based on the energy difference



a QPE-like circuit

Quantum Metropolis algorithm

Weak measurement implementation costs:

- Suppose the Lindblad simulation time is T
- #steps for weak measurement implementation is T/τ
- Total approximation error $= \frac{T}{\tau} \cdot \tau^2 = T\tau$
- We need to take $\tau \sim \frac{1}{T}$ for small error
- The **total #steps** $\sim T^2$
- It can be improved to $\tilde{\mathcal{O}}(T)$ using a complicated “**compression**” technique (Chen-Kastoryano-Brandão-Gilyén ’23)
- Simplified quantum Gibbs sampling algorithm using QPE (Irani-Jiang ’24)