## CS188–Spring 2019 — Homework 1

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**Due:** Monday 2/4/2019 at 11:59pm (submit via Gradescope).

Leave self assessment boxes blank for this due date.

Self assessment due: Monday 2/11/2018 at 11:59pm (submit via Gradescope) For the self assessment, fill in the self assessment boxes in your original submission (you can download a PDF copy of your submission from Gradescope). For each subpart where your original answer was correct, write correct. Otherwise, write and explain the correct answer.

**Policy:** Can be solved in groups (acknowledge collaborators) but must be written up individually. **Submission:** Your submission should be a PDF that matches this template. Each page of the PDF should align with the corresponding page of the template (page 1 has name/collaborators, question 1 begins on page 2, etc.). **Do not reorder, split, combine, or add extra pages.** The intention is that you print out the template, write on the page in pen/pencil, and then scan or take pictures of the pages to make your submission. You may also fill out this template digitally (e.g. using a tablet.)

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Collaborators	None		

Q1(a)	DFS and BFS is wrong, the search is tie-breaking scheme, not alphabetically order				
Q1(b)(i)	The minimum distance of B to D is 1+3,rather than 5				
Q1(b)(ii)	correct				
Q1(b)(iii)	Calculate the $h_3(D)$ wrongly.				
Q2(a)	correct				
Q2(b)	correct				
Q2(c)	correct				
Q2(d)	correct				
Q2(e)(i)	correct				
Q2(e)(ii)	correct				

## Q1.Search

	Search Algorithm	A-B-D-G	A-C-D-G	A-B-C-D-F-G
(a)	Depth first search			
	Breadth first search	V		
	Uniform cost search			
	$A^*$ search with heuristic $h_1$			
	$A^*$ search with heuristic $h_2$			

- (b) (i) A heuristic h is admissible (optimistic) if: $0 \le h(n) \le x(n)$ , where x(n) is the true cost to a nearest goal. So  $h_3(B) \in [0, 13]$ , where 13 denotes 5+3+5.
  - (ii) Because h(A)-h(B) $\leq$ cost(A to B)=1 and h(B)-h(c) $\leq$ cost(B to C)=1, h(A)=10,h(C)=9, then  $h_3(B) \in [9, 10]$ .
  - (iii)  $h_3(A) + g(A) = 10 + 0 = 10, h_3(C) + g(C) = 9 + 4 = 13, h_3(B) + g(B) = H_3(B) + 5, h_3(D) + g(D) = 7 + 10 = 17$ , and if it is in order, then  $13_3(B) + 5 \le 17$ , then  $h_3(B) \in [8, 12]$ .

## Q2.n-pacmen search

- (a) M tuples  $((x_1, y_1), (x_2, y_2), ...(x_M, y_M))$  encoding the x and y coordinates of each parman.
- (b) The number of pacmen:n Number of squares where pacmen can go:M So the state space is  $M^n$ .
- (c)  $M^n$ .
- (d) Because there is a closed set, the same node can not be expanded twice. The bound is  $M^n$ .
- (e) (i) Not consistent and not admissible. Consider the situation that a square that there are three pacmen on its left, right, down. And the actual cost from this situation to goal is 1.(All of them move to the middle square) But  $h_1 = 1 + 1 + 1 = 3 > 1$ . So  $h_1$  is not consistent and not admissible.
  - (ii) Consistent and admissible. Because the every pacman moves by at most one unit vertically or horizontally at each time step. So  $1/2*2 \le 1$ . And  $h(A)-h(B) \le cost(A \text{ to } B)$ , satisfies consistent and admissible