

HINT: $\sum_{i=1}^n \frac{1}{i} = \Theta(\log n)$

1. (3 points) Honor Code

I promise that I will complete this quiz independently and will not use any electronic products or paper-based materials during the quiz, nor will I communicate with other students during this quiz.

I will not violate the Honor Code during this quiz.

☐ True ☐ False

2. (5 points) True or False

Determine whether the following statements are true or false.

- (a) (1') Quick-sort is an in-place sorting algorithm, while merge-sort is not since it requires $\omega(n)$ extra space. ☐ True ☐ False
- (b) (1') Randomized Quick-sort is unstable, while Quick-sort with deterministic method (choosing pivots) is stable. ☐ True ☐ False
- (c) (1') There exists a comparison-based sort algorithm that needs $O(1)$ extra space and takes $\Theta(n \log n)$ time. ☐ True ☐ False
- (d) (1') For an array $\{a_n\}$ with distinct elements, for fixed i, j , if $\forall a_k \neq a_i, a_k \neq a_j, (a_k - a_i)(a_k - a_j) > 0$, then a_i and a_j will be compared in any case when using randomized quick-sort to make $\{a_n\}$ sorted. ☐ True ☐ False
- (e) (1') When we use divide and conquer to solve a problem, we should divide the problem into one or more subproblems with the exact same scale, then recursively do them and merge their answers at last. ☐ True ☐ False

3. (5 points) Randomized quick-sort

- (a) (1') If we use randomized quick-sort (i.e. randomly choosing pivots) to sort the array $[3, 4, 6, 2, 1, 5, 8, 0]$, the probability of 2 and 5 are compared is ____.
- (b) (1') Use the same method as above to sort an array with n distinct elements, the probability of i -th largest and j -th largest element ($i \neq j$) are compared is ____.
- (c) (3') Prove that the expectation times of comparisons in the randomized quick-sort is $\Theta(n \log n)$.
Hint: The total expectation times can be obtained from the sum of the expectation of each comparison.

4. (12 points) Solving Recursion

Solve the recursion relation with $T_i(1) = 1, T_i(0) = 0$:

(a) (3') $T(n) = T(n-1) + \Theta(n^c)$ ($c > 0$ is a constant).

(b) (4') $T(n) = T(\frac{n}{2}) + \Theta(\log n)$

(c) (5') $T(n) = \Theta(n) + \frac{1}{n} \sum_{i=0}^{n-1} (T(i) + T(n-i))$, you can write $\Theta(n)$ as cn for your convenience.