Computer Architecture I Mid-Term

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P	inyin Name	e:											
S	tudent ID:												
E	-Mail @s	shangh	aitech.e	edu.cn:									
	Question:	1	2	3	4	5	6	7	8	9	10	Total	
	Points:	1	12	18	8	5	12	8	8	18	10	100	
	Score:												

- This test contains 12 numbered pages, including the cover page, printed on both sides of the sheet.
- Unless told otherwise always assume a 32bit machine.
- The total estimated time is 90 minutes.
- You have 90 minutes to complete this exam, 10 min for submission. The exam is open book.
- Under no circumstances are you allowed to communicate to anybody except the CS110 professors or TAs.
- There may be partial credit for incomplete answers; write as much of the solution as you can. We will deduct points if your solution is far more complicated than necessary. When we provide a blank, please fit your answer within the space provided.
- Write your answers on blank A4 pages. Include the correct header on every page (see Question 1 below). At the end scan/ photo all pages into a pdf (there are plenty of apps for that). Submit in time we will NOT accept late submissions!
- Each question from Q2 to Q10 will need to be started on a new A4 sheet so you'll need at least 9 A4 sheets!

1. Submitted pages format

Correctly write the header for ALL the pages that you submit. That includes: Your name in Chinese, your ShanghaiTech Email address (without '@shanghaitech.edu.cn'), the string "CS110 Midterm", this Exam number (see on the top left of the following pages), the page number (your own numbering). Also follow the rule that every new major question (e.g. Question 2, Question 3, etc.) is written on a new A4 paper. Submit correctly on gradescope or (only if gradescope fails) via email to soerensch@shanghaitech.edu.cn AND wangchd@shanghaitech.edu.cn .

		& Number Representations All floating point numbers follow IEEE 754 standard his question.
1	(a)	Name the 6 Great Ideas in Computer Architecture as taught in the lectures.
3	(b)	Consider this 8-bit binary pattern 0b10101010, please write down the number in decimal form if we are using the following representations:
		Unsigned binary
		Two's complement binary
		One's complement binary
2	(c)	For a 9-bit value, sign-magnitude integer, what are the largest AND smallest value you can represent in decimal?
2	(d)	What is the smallest positive integer that cannot be represented by a 32-bit floating point number?
2	(e)	(Multiple Choice) A 32-bit floating point number consists of:
		A. 1 sign bit, 8 exponent bits and 23 significant bits.
		B. 1 sign bit, 7 exponent bits and 24 significant bits.
		C. 1 sign bit, 6 exponent bits and 25 significant bits.
		D. 1 sign bit, 5 exponent bits and 26 significant bits.
2	(f)	(Multiple Choice) Which of the following choices about floating point number is true?
		A. Given that double precision number has 11 bits in its exponent part, then the

C. The distance between floating point numbers remains constant as the absolute

D. Max representable integer in floating point format is larger than the max integer

bias should be -1024.

B. Floating number addition is associative.

in unsigned int type (both in 32 bit).

value of the numbers increase.

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3. C Basics

2

(a) Suppose we have following macro:

```
1 #define DIFSQUARE(a, b) a * a - b * b
```

Write the result of DIFSQUARE(3 + 1, 1 + 2). You may find you did not get the result you want. What result did you expect? Fix the definition of DIFSQUARE(a, b).

(b) Consider the following code.

```
1 #include <stdio.h>
2 #define INT_NAME_VAL(name)
3 int main() {
4    int a = 1, b = 2, c = 3;
5    INT_NAME_VAL(a);
6    INT_NAME_VAL(b);
7    INT_NAME_VAL(c);
8    return 0;
9 }
```

Complete the definition of INT_NAME_VAL(name) using printf function, format specifiers (%d, %u, %x, %s, %c...) and #. So that the output is

a = 1 b = 2c = 3

(c) What is the output of the following program.

```
1 #include <stdio.h>
2
3 int some_bit_op(int x, int y, int z) {
4    return ((x >> 1) | (y)) ^ ((y) & (z << 1));
5 }
6
7 int main() {
8    printf("%d\n", some_bit_op(10, 14, 5));
9    return 0;
10 }</pre>
```

|4|

(d) Consider the following program.

We want to concatenate two 32bit integer to become one new integer. For the new integer, we want its 1^{st} byte to be the 4^{th} byte of number \mathbf{a} , its 2^{nd} and 3^{rd} byte to be the 2^{nd} and 3^{rd} byte of number \mathbf{b} , and its 4^{th} byte to be the 1^{st} byte of number \mathbf{a} .

```
1 #include <stdio.h>
  #include <stdint.h>
 3
  int byte concatenation(int a, int b) {
 5
       int new_a, c;
6
       printf("%x\n", new_a); /* Should print dc0000ab */
 7
8
 9
       return c;
10 }
11
12 int main() {
       int concatenated num = byte concatenation(0xABBACDDC, 0x14233341);
13
       printf("%x\n", concatenated num); /* Should print dc2333ab */
14
15
       return 0;
16 }
```

1. (Multiple Choice) At line 6, we want new_a to get the 1st and 4th byte of a and exchange their position, choose the right expression that can correctly print the result.

```
A. (a >> 24) & (a << 24 & 0xFF) B. (a >> 24 & 0xFF) | (a << 24) C. (a >> 24) | (a << 24 & 0xFF) D. (a >> 24 & 0xFF) & (a << 24)
```

2. At line 8, Write an expression that can correctly print the result using only bit operators including &, |, ^, ~, >>, <<. (You should not get 0xdc2333ab directly.)

7

(e) Memory in C

```
1 #include <stdio.h>
2 #include <string.h>
3 #include <stddef.h>
 4 #include <stdlib.h>
 5 #include <stdint.h>
 6 typedef struct data {
 7
      char
                 str[8];
8
      uint64_t num;
9
      union {
10
         short
                    у;
11
         char
                    z[4];
12
         uint64 t w;
13
      } InnerData;
14 } Data;
15
16 int main(){
17
      char* info = malloc(sizeof(char) * 8);
      Data testdata;
18
      sprintf(info, "%s", "Size: ");
printf("%s%s%lu\n", "Data", info, sizeof(testdata));
19
20
21
      printf("%s%s%lu\n", "InherData", info, sizeof(testdata.InnerData));
22
      return 0;
23 }
```

1. Fill in the correct memory section based on what the given C expressions evaluate to.

testdata	
info	
&info	

2. What is the expected output of the following code?

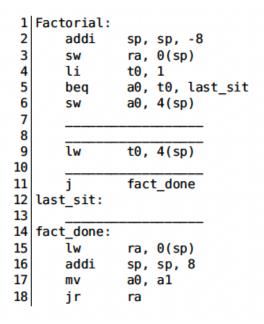
3. Does this simple program encounters memory problems? Point it out if there are such problem(s), or answer "No problem" if not.

4. RISC-V

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8 (

(a) Consider the following code snippet written in RISC-V. The function Factorial is to calculate the factorial of a given number. (i.e. $n! = n \cdot (n-1) \cdots 2 \cdot 1$)



Fill in the missing code below.

line 7: _____

line 8: _____

line 10: _____

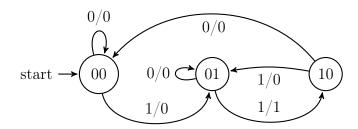
line 13: _____

5. **Logic**

Draw the **fully** simplified (with **fewest** number of primitive gates) circuit for the expression below. Write down the **steps** for simplification and draw the **final** circuit. You may use the following primitive gates: AND, OR, XOR and NOT.

$$\overline{(A+\overline{C}C)} + (\overline{A}+D)(A+BD)$$

6. **FSM**



(a) Copy and fill in the truth table for the FSM above.

state bit1	state bit0	input	next state bit1	next state bit0	output
0	1	0			
0	1	1			
0	0	0			
0	0	1			
1	0	0			
1	0	1			

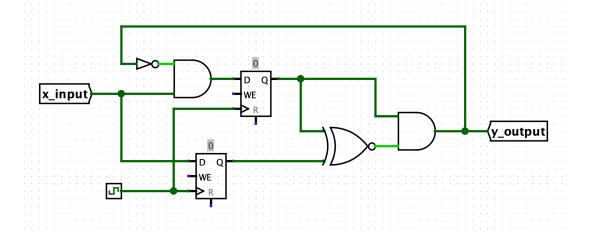
(b) What does the given FSM output with the input bit string '1111011010'?

(c) Draw a FSM that shares the same function with the above given FSM using only two states (Output the same as the given FSM for any input bit string).



7. **SDS**

We will be analyzing the following circuit, given the information below.



- For all registers, Clk-to-q delay time 4ns, Set up time 3ns.
- Gates Delay: AND gate 5ns, NOT gate 3ns, XOR gate 7ns, NXOR gate 10ns.
- x_input is directly attached to a source that changes its value (i.e. 1 to 0, or 0 to 1) 15ns after the rising edge of the clock.
- y output is directly attached to a register.
- All wires are ideal, i.e. they have zero delay.
- (a) What is the maximum possible hold time for registers to function properly in this circuit?
- 3 (b) What is the minimum possible clock period of this circuit?
- (c) Now we assume the clock frequency is 20MHz (which would be a proper frequency if you analyzed the circuit correctly) and the first rising edge of the clock is at 0ns. x_input is initialized to 0 at 0ns. Try to figure out at what time will y_output become 1?

8. Pipeline

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- 3
- (a) (Multiple Choice) Choose the correct statement(s) for a five stage pipelined processor. There may be more than one correct answer.
 - A. Pipeline rate is limited by the slowest pipeline stage.
 - B. Using pipeline is an effective method to reduce the latency of a single task.
 - C. When the two instructions add t0, t0, t1 and add t2, t0, t1 are executed consecutively, there is no hazard.
 - D. When instructions add t1, t0, t0, addi t2, t0, 5, and addi t4, t1, 5 are executed consecutively, the hazard can be solved by forwarding.
- (b) Consider the following instructions and complete the pipeline diagram for the first iteration. Register writes happen before register reads in the same clock cycle, branch comparison is done during the register read stage, and forwarding is implemented. Assume the branch is not taken in the first iteration.

```
1 loop:
2 lw t0, 0(a0)
3 beq t0, 0, end
4 addi t0, t0, 10
5 sw t0, 0, 0(a0)
6 end:
```

Note: we use F for fetch, D for decode, X for execute, M for memory, W for write back.

lw	F	D	X	Μ	W				
beq									
. dd:									
addi									
sw									

(c) Each stage in pipeline and the registers in pipeline has following timing:

F	D	X	M	W	Clock-to-Q	Hold time	Setup
175 ps	100ps	100ps	200ps	100ps	30 ps	20ps	30 ps

for the 5 stage pipeline, compute the minimum clock period, and show the calculation process.

9. Cache

5

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Notice: We assume a 32-bit machine by default.

(a) (True or False) Please fill your answer (T or F) in the table below.

Notice: You will get 1 point for a correct answer, -1 point for a wrong answer, and 0 point for not answering. The total point will not be less than 0.

- 1. Cache replacement policy is used for choosing which **cache line** should be evicted.
- 2. One way to decrease compulsory miss is enlarging the block size.
- 3. For a cache with fixed block size and length of the **tag** field, increasing or decreasing associativity will not influence the hit time.
- 4. Cache blocking can make transposition faster because it decreases compulsory miss.
- 5. For a 1-level, LRU cache, cache blocking will greatly speed up matrix transposition when the cache size is larger than the matrix size.

1	2	3	4	5

(b) Suppose the cache has the following settings:

Cache levels	1
Block size	16 bytes
Number of sets	4
Cache size	128 bytes
Block replacement policy	LRU

1.	Is this cache direct-mapped, set-associative, or fully-associative?	If it is asso-
	ciative, also write down its associativity.	

2.	Suppose	the	memory	addresses	are	32-bit	long,	what	is	the	length	of	the	tag
	field?													

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(c) Suppose the following code is running on a system with the above cache, where sizeof(int) == 4.

```
1 #define array size
2 #define repeat times 1
3 #define step size
4
5
  int main() {
6
       int array[array_size] = { };
7
8
       for (int r = 0; r < repeat times; <math>r++) {
9
           for (int i = 0; i < array size; i += step size) {
10
               array[i] = array[i] + 2333;
11
           }
12
       }
13
14
       return 0;
15 }
```

- 1. What is the total number of accesses to the cache?
- 2. What is the hit rate?
- 3. Which type(s) of miss occur(s)?
- 4. Suppose repeat_times goes to infinity (only for this question), what number will the hit rate converge to?
- 5. If repeat_times is changed to 2 (only for this question), try to swap two lines of the above code to maximize the hit rate without disturbing the results. Which two lines will you choose and what is the maximized hit rate?
- 6. For the modified code in the previous question, suppose again repeat_times goes to infinity (only for this question), what number will the hit rate converge to?

10. Threads and OpenMP

Consider the following code snippet written in C. The function abs_sum is to calculate the sum of absolute values in a given square matrix (i.e., $\sum_{i=1}^{n} \sum_{j=1}^{n} |mat[i][j]|$).

```
1 #include <omp.h>
2 #include <math.h>
3 #include <stdlib.h>
4 double abs sum(double** mat, int n){
       double ret = 0.0;
       #pragma omp parallel for
7
       for (int i = 0; i < n; i++){
8
           for (int j = 0; j < n; j++)
9
               ret += abs(mat[i][j]);
10
       }
11
       return ret;
12 }
```

(a) This code doesn't always give a correct result. Please give your explanation.

(b) Someone suggests to add a line #pragma omp atomic between line 8 and 9 to fix this issue. Does it work? Is it a efficient implementation? If so, give your explanation. If not, try to give a more efficient implementation.

(c) Assume we have fixed this bug now. If you modify line 9 to ret += abs(mat[j][i]), you will observe a significantly performance degradation. Try to explain it.