

# Discrete Mathematics

## Lecture 13

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# Summary of Lecture 12

**$r$ -Combination of Set  $A = \{a_1, a_2, \dots, a_n\}$**

- Without repetition: an  $r$ -subset of  $A$ ;
- With repetition: an  $r$ -multiset of the form  $\{x_1 \cdot a_1, \dots, x_n \cdot a_n\}$

**$r$ -Combination of Multiset  $A = \{n_1 \cdot a_1, n_2 \cdot a_2, \dots, n_k \cdot a_k\}$**

- an  $r$ -subset of  $A$

**Binomial Transform:**  $b_n = \sum_{k=s}^n \binom{n}{k} a_k$

**Inverse Binomial Transform:**  $b_n = \sum_{k=s}^n (-1)^{n-k} \binom{n}{k} a_k$

- $a_n = \sum_{k=s}^n \binom{n}{k} b_k \Rightarrow b_n = \sum_{k=s}^n (-1)^{n-k} \binom{n}{k} a_k \quad (n \geq s)$

**Distributing Objects into Boxes:**

- Labeled/unlabeled objects + labeled/unlabeled box



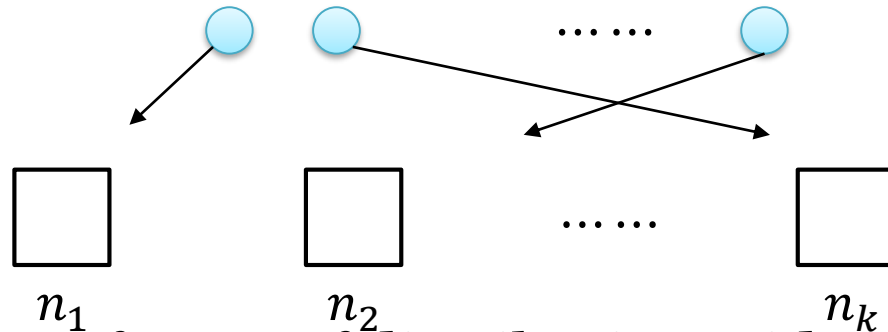
$$S_2(n, j)$$

**THEOREM:**  $S_2(n, j) = \frac{1}{j!} \sum_{i=0}^{j-1} (-1)^i \binom{j}{i} (j-i)^n$  when  $n \geq j \geq 1$ .

- $T(n, j)$ : the number of ways of distributing  $n$  labeled objects into  $j$  labeled boxes such that no box is empty
  - $T(n, j) = j! \cdot S_2(n, j)$ 
    - $T(n, j) = ?$
- $X$ : the set of ways of distributing  $n$  labeled objects into  $j$  labeled boxes.
  - By the product rule,  $|X| = j^n$
- $X_i \subseteq X$ : the set of ways where exactly  $i$  boxes are used,  $i = 1, 2, \dots, j$ 
  - $\{X_1, X_2, \dots, X_j\}$  is a partition of  $X$  and  $|X_i| = \binom{j}{i} T(n, i)$
  - $j^n = |X| = \sum_{i=1}^j |X_i| = \sum_{i=1}^j \binom{j}{i} T(n, i)$
  - $T(n, j) = \sum_{i=1}^j (-1)^{j-i} \binom{j}{i} i^n = \sum_{i=0}^{j-1} (-1)^i \binom{j}{i} (j-i)^n$  //inversion
- $S_2(n, j) = \frac{1}{j!} \cdot T(n, j) = \frac{1}{j!} \sum_{i=0}^{j-1} (-1)^i \binom{j}{i} (j-i)^n$

# Type 4

**Problem:** distributing  $n$  unlabeled objects into  $k$  unlabeled boxes



## Classifications

$$n_1 + n_2 + \cdots + n_k = n$$

$$n_1, n_2, \dots, n_k \in \mathbb{N}$$

$$n_1 \geq n_2 \geq \cdots \geq n_k$$

**EXAMPLE:** # of ways of distributing 4 identical books into 3 identical boxes.

- 4 0 0
- 3 1 0
- 2 2 0
- 2 1 1

**REMARK:** The schemes are determined by  $\{n_1, \dots, n_k\}$

# Partitions of Integers

整数的分解

**DEFINITION:**  $n = a_1 + a_2 + \cdots + a_j$  is called an  **$n$ -partition** with exactly  $j$  parts if  $a_1 \geq a_2 \geq \cdots \geq a_j$  are all positive integers.

- $p_j(n) = |\{(a_1, \dots, a_j): a_1 + \cdots + a_j = n, a_1 \geq a_2 \geq \cdots \geq a_j \geq 1 \text{ are integers}\}|$ 
  - $p_j(n)$ : # of ways of writing  $n$  as the sum of  $j$  positive integers.

**EXAMPLE:** The integer 4 has four different partitions:

- $4 = 4$
- $4 = 3 + 1$
- $4 = 2 + 2$
- $4 = 2 + 1 + 1$

**REMARK:** solution to the type 4 problem =  $\sum_{j=1}^k p_j(n)$

# Partitions of Integers

**THEOREM:** For  $n \in \mathbb{Z}^+, j \in [n]$ ,  $p_j(n + j) = \sum_{k=1}^j p_k(n)$

- $p_1(n) = 1, p_n(n) = 1$
- Let  $S_k = \{\text{partitions of } n \text{ into } k \text{ positive integers}\}, k \in [j]$
- Let  $S = \bigcup_{k=1}^j S_k$ .
  - $|S| = |S_1| + \dots + |S_j| = p_1(n) + \dots + p_j(n)$
- Let  $T = \{\text{partitions of } n + j \text{ into } j \text{ positive integers}\}$ 
  - $|T| = p_j(n + j)$
- $f: S \rightarrow T \quad (n_1, \dots, n_k) \mapsto (n_1 + 1, \dots, n_k + 1, \underbrace{1, \dots, 1}_{j-k})$ 
  - $f$  is bijective
  - $|T| = |S|$

**EXAMPLE:** determine  $p_3(6)$  and  $p_4(6)$  with the above theorem

- $p_3(6) = p_3(3 + 3) = p_1(3) + p_2(3) + p_3(3) = 1 + 1 + 1 = 3$
- $p_4(6) = p_4(2 + 4) = p_1(2) + p_2(2) + p_3(2) + p_4(2) = 1 + 1 + 0 + 0 = 2$

# Computing $p_j(n)$ Recursively

$$\begin{array}{cccccccccc}
& & & & & & & & & p_1(1) \\
& & & & & & & & & p_1(2) & p_2(2) \\
& & & & & & & & & p_1(3) & p_2(3) & p_3(3) \\
& & & & & & & & & p_1(4) & p_2(4) & p_3(4) & p_4(4) \\
& & & & & & & & & p_1(5) & p_2(5) & p_3(5) & p_4(5) & p_5(5) \\
& & & & & & & & & p_1(6) & p_2(6) & p_3(6) & p_4(6) & p_5(6) & p_6(6) \\
& & & & & & & & & p_1(7) & p_2(7) & p_3(7) & p_4(7) & p_5(7) & p_6(7) & p_7(7) \\
& & & & & & & & & p_1(8) & p_2(8) & p_3(8) & p_4(8) & p_5(8) & p_6(8) & p_7(8) & p_8(8) \\
& & & & & & & & & p_1(9) & p_2(9) & p_3(9) & p_4(9) & p_5(9) & p_6(9) & p_7(9) & p_8(9) & p_9(9)
\end{array}$$





# Principle of Inclusion–Exclusion

容斥原理

**Problem:**  $S$  is a finite set and  $A_1, A_2, \dots, A_n \subseteq S$ .

- $|\cup_{i=1}^n A_i| = ? \rightarrow \cap$
- $|\cap_{i=1}^n A_i| = ? \rightarrow \cup$

**EXAMPLE:** Let  $S$  be the set of permutations of  $[n]$ . Find  $|A|$  for

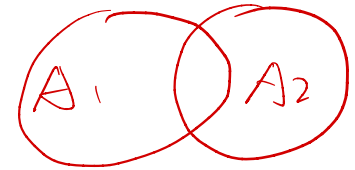
$$A = \{x_1 x_2 \dots x_n : x_1 x_2 \dots x_n \in S; x_i \neq i \text{ for all } i \in [n]\}.$$

- $A_i = \{x_1 x_2 \dots x_n : x_1 x_2 \dots x_n \in S; x_i = i\}, i = 1, 2, \dots, n$  错排问题
  - $A = S - \cup_{i=1}^n A_i$
  - $|S| = n!$
  - $|\cup_{i=1}^n A_i| = ?$

# Principle of IE (Two Sets)

**THEOREM:** Let  $S$  be a finite set. Let  $A_1, A_2$  be subsets of  $S$ . Then

- $|S - A_1| = |S| - |A_1|$ ;  $|A_1 - A_2| = |A_1| - |A_1 \cap A_2|$
- $S = A_1 \cup (S - A_1)$ ,  $A_1 \cap (S - A_1) = \emptyset$ ;
  - $\{A_1, S - A_1\}$  is a partition of  $S$ 
    - $|S| = |A_1| + |S - A_1|$ 
      - $|S - A_1| = |S| - |A_1|$
  - $A_1 - A_2 = A_1 - A_1 \cap A_2$ 
    - $|A_1 - A_2| = |A_1| - |A_1 \cap A_2|$
- $|A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2|$
- $A_1 \cup A_2 = (A_1 - A_2) \cup A_2$ ,  $(A_1 - A_2) \cap A_2 = \emptyset$ ;
  - $\{A_1 - A_2, A_2\}$  is a partition of  $A_1 \cup A_2$ 
    - $|A_1 \cup A_2| = |A_1 - A_2| + |A_2| = |A_1| - |A_1 \cap A_2| + |A_2|$
- $|A_1 \cap A_2| = |A_1| + |A_2| - |A_1 \cup A_2|$



# Principle of IE (Three Sets)

**THEOREM:** Let  $S$  be a finite set. Let  $A_1, A_2, A_3$  be subsets of  $S$ .

Then  $|\bigcup_{i=1}^3 A_i| = \sum_{t=1}^3 (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq 3} |A_{i_1} \cap \dots \cap A_{i_t}|$

- $|\bigcup_{i=1}^3 A_i| = |(A_1 \cup A_2) \cup A_3| = |A_1 \cup A_2| + |A_3| - |(A_1 \cup A_2) \cap A_3|$

- $|A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2|$

- $|(A_1 \cup A_2) \cap A_3| = |(A_1 \cap A_3) \cup (A_2 \cap A_3)|$

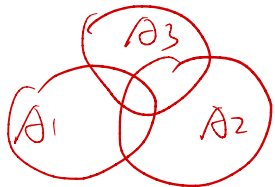
$$= |A_1 \cap A_3| + |A_2 \cap A_3| - |(A_1 \cap A_3) \cap (A_2 \cap A_3)|$$

$$= |A_1 \cap A_3| + |A_2 \cap A_3| - |A_1 \cap A_2 \cap A_3|$$

- $|\bigcup_{i=1}^3 A_i| = |A_1| + |A_2| - |A_1 \cap A_2| + |A_3|$

$$= |(A_1 \cup A_2) \cup A_3| - (|A_1 \cap A_3| + |A_2 \cap A_3| - |A_1 \cap A_2 \cap A_3|)$$

- $|\bigcap_{i=1}^3 A_i| = \sum_{t=1}^3 (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq 3} |A_{i_1} \cup \dots \cup A_{i_t}|$



# Principle of IE ( $n$ Sets)

**THEOREM:** Let  $S$  be a finite set. Let  $A_1, A_2, \dots, A_n$  be subsets of  $S$ .

Then  $|\bigcup_{i=1}^n A_i| = \sum_{t=1}^n (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq n} |A_{i_1} \cap \dots \cap A_{i_t}|$

- $n = 1: |A_1| = |A_1|$  t 组合
- $n = 2: |A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2|$
- $n = 3: |A_1 \cup A_2 \cup A_3| = |A_1| + |A_2| + |A_3| - |A_1 \cap A_2| - |A_1 \cap A_3| - |A_2 \cap A_3| + |A_1 \cap A_2 \cap A_3|$
- 数学归纳法 **Induction hypothesis:** the identity holds for  $n \leq k$  ( $k \geq 3$ )
- Need to show the identity for  $n = k + 1$
- $|A_1 \cup \dots \cup A_{k+1}| = |A_1 \cup \dots \cup A_k| + |A_{k+1}| - |(A_1 \cup \dots \cup A_k) \cap A_{k+1}|$   
 $= |\bigcup_{i=1}^k A_i| + |A_{k+1}| - |\bigcup_{i=1}^k (A_i \cap A_{k+1})|$

# Principle of IE ( $n$ Sets)

- $|\cup_{i=1}^k A_i| = \sum_{t=1}^k (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq k} |A_{i_1} \cap \dots \cap A_{i_t}|$
- $|\cup_{i=1}^k (A_i \cap A_{k+1})| = \sum_{t=1}^k (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq k} |(A_{i_1} \cap A_{k+1}) \cap \dots \cap (A_{i_t} \cap A_{k+1})|$
- $|\cup_{i=1}^{k+1} A_i| = \sum_{t=1}^k (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq k} |A_{i_1} \cap \dots \cap A_{i_t}| + |A_{k+1}| -$

$$\sum_{t=1}^k (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq k} |(A_{i_1} \cap A_{k+1}) \cap \dots \cap (A_{i_t} \cap A_{k+1})|$$

$$= \sum_{t=1}^{k+1} (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq k+1} |A_{i_1} \cap \dots \cap A_{i_t}|$$

**THEOREM:** Let  $S$  be a finite set. Let  $A_1, A_2, \dots, A_n$  be subsets of  $S$ .

$$\text{Then } |\cap_{i=1}^n A_i| = \sum_{t=1}^n (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq n} |A_{i_1} \cup \dots \cup A_{i_t}|$$

# Principle of Inclusion-Exclusion

全排列

**EXAMPLE:** Let  $S$  be the set of permutations of  $[n]$ . Find  $|A|$  for

$$A = \{x_1 x_2 \dots x_n : x_1 x_2 \dots x_n \in S; x_i \neq i \text{ for all } i \in [n]\}.$$

- $A_i = \{x_1 x_2 \dots x_n : x_1 x_2 \dots x_n \in S; x_i = i\}, i = 1, 2, \dots, n$  A<sub>i</sub> 第 i 个位置为 i

- $A = S - \bigcup_{i=1}^n A_i$

- $|S| = n!$

- $|\bigcup_{i=1}^n A_i| = ?$

- $|\bigcup_{i=1}^n A_i| = \sum_{t=1}^n (-1)^{t-1} \sum_{1 \leq i_1 < \dots < i_t \leq n} |A_{i_1} \cap \dots \cap A_{i_t}|$

- $|A_{i_1} \cap \dots \cap A_{i_t}| = (n - t)! \text{ for } t = 1, 2, \dots, n$

- $|A| = |S| - |\bigcup_{i=1}^n A_i|$

$$= n! - \left( \binom{n}{1} * (n - 1)! - \binom{n}{2} * (n - 2)! + \dots + (-1)^{n-1} * \binom{n}{n} * 1 \right)$$

$$= n! \left( 1 - \frac{1}{1!} + \frac{1}{2!} - \dots + (-1)^t \frac{1}{t!} + \dots + (-1)^n \frac{1}{n!} \right)$$





# Pigeonhole Principle

鸽巢原理

工作站

服务器

**EXAMPLE:** Connect 15 workstations  $W_1, \dots, W_{15}$  to 10 servers  $S_1, \dots, S_{10}$  such that any  $\geq 10$  workstations have access to all servers. How many cables are needed?

- **Solution 1:** Connecting every workstation directly to every server. 150
- **Solution 2:**  $S_i$  is connected to  $W_i$  for every  $i \in [10]$ ; and each of  $W_{11}, W_{12}, W_{13}, W_{14}, W_{15}$  is connected to all servers.
  - This solution requires 60 lines.
  - Is this solution optimal?

# Cover

**DEFINITION:** A **cover** of a finite set  $A$  is a family  $\{A_1, A_2, \dots, A_n\}$  of subsets of  $A$  such that  $\bigcup_{i=1}^n A_i = A$ .

**LEMMA:** Let  $\{A_1, A_2, \dots, A_n\}$  be a cover of a finite set  $A$ .

Then  $|A| \leq \sum_{i=1}^n |A_i|$ .

- $n = 1: |A| = |A_1|$
- $n = 2: |A| = |A_1 \cup A_2| = |A_1| + |A_2| - |A_1 \cap A_2| \leq |A_1| + |A_2|$
- Suppose true when  $n \leq k$  ( $k \geq 2$ ).
- When  $n = k + 1$ ,  
$$\begin{aligned} |A| &= \left| \bigcup_{i=1}^k A_i \cup A_{k+1} \right| \\ &\leq \left| \bigcup_{i=1}^k A_i \right| + |A_{k+1}| \\ &\leq \sum_{i=1}^k |A_i| + |A_{k+1}| \\ &= \sum_{i=1}^{k+1} |A_i| \end{aligned}$$

# Pigeonhole Principle

**THEOREM:** (simple form) Let  $A$  be a set with  $\geq n + 1$  elements.

Let  $\{A_1, A_2, \dots, A_n\}$  be a cover of  $A$ . Then  $\exists k \in [n], |A_k| \geq 2$ .

- Suppose that  $|A_i| \leq 1$  for every  $i \in [n]$ . Then  $n + 1 \leq |A| \leq \sum_{i=1}^n |A_i| \leq n$ .
  - If  $\geq n + 1$  objects are distributed into  $n$  boxes, then there is at least one box containing  $\geq 2$  objects.

**THEOREM:** (general form) Let  $A$  be a set with  $\geq N$  elements.

Let  $\{A_1, A_2, \dots, A_n\}$  be a cover of  $A$ . Then  $\exists k \in [n], |A_k| \geq \lceil N/n \rceil$ .

- If  $|A_i| < \lceil N/n \rceil$  for all  $i \in [n]$ , then  $N \leq |A| \leq \sum_{i=1}^n |A_i| < n \cdot N/n = N$ 
  - If we distribute  $\geq N$  objects into  $n$  boxes, then there is at least one box that contains  $\geq \lceil N/n \rceil$  objects.

# Pigeonhole Principle

**EXAMPLE:** Connect 15 workstations  $W_1, \dots, W_{15}$  to 10 servers  $S_1, \dots, S_{10}$  such that any  $\geq 10$  workstations have access to all servers. How many cables are needed?

- **Solution 2:**  $S_i$  is connected to  $W_i$  for every  $i \in [10]$ ; and each of  $W_{11}, W_{12}, W_{13}, W_{14}, W_{15}$  is connected to all servers. // 60 lines, optimal?
- Consider an **optimal scheme  $\Pi$** . 优化方案
  - Let  $A = \{(W_i, S_j) : i \in [15], j \in [10], W_i \text{ is not connected to } S_j\}$  in  $\Pi$
  - $A_t = \{(W_i, S_j) \in A : j = t\}$  for  $t = 1, 2, \dots, 10$ 
    - $\{A_1, A_2, \dots, A_{10}\}$  is a cover of  $A$
- If there are **< 60 lines in  $\Pi$** , then  $|A| > 150 - 60 = 90$ .
  - $\exists k \in [10]$  such that  $|A_k| \geq \lceil 91/10 \rceil = 10$ 
    - There are 10 workstations not connected to  $S_k$