

Numerical Optimization, 2023 Fall

Homework 3

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Problem 1. Prove the dual of the dual of a linear programming (standard form) is itself. [25pts]

Consider a linear programming that is in standard form:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & A\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{aligned} \tag{1}$$

The Lagrangian of the above linear programming is:

$$L(\mathbf{x}, \boldsymbol{\lambda}) = \mathbf{c}^T \mathbf{x} + \boldsymbol{\lambda}^T (A\mathbf{x} - \mathbf{b}) \tag{2}$$

The dual problem is that:

$$\begin{aligned} \max_{\boldsymbol{\lambda}} \quad & \boldsymbol{\lambda}^T \mathbf{b} \\ \text{s.t.} \quad & A^T \boldsymbol{\lambda} \leq \mathbf{c} \\ & \boldsymbol{\lambda} \geq 0 \end{aligned} \tag{3}$$

The Lagrangian of the dual problem is:

$$L(\boldsymbol{\lambda}, \mathbf{x}) = \boldsymbol{\lambda}^T \mathbf{b} + \mathbf{x}^T (A^T \boldsymbol{\lambda} - \mathbf{c}) \tag{4}$$

The dual of the dual problem is that:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & A\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq 0 \end{aligned} \tag{5}$$

So above all, the dual of the dual of a linear programming (standard form) is itself.

Problem 2. Prove the dual objective increases after a pivot of the dual simplex method. [25pts]

Problem 3. Let $L(\mathbf{x}, \boldsymbol{\lambda})$ be the Lagrangian of a linear programming problem, and $(\mathbf{x}^*, \boldsymbol{\lambda}^*)$ be the optimal primal-dual solution. Prove that

$$L(\mathbf{x}, \boldsymbol{\lambda}^*) \geq L(\mathbf{x}^*, \boldsymbol{\lambda}^*) \geq L(\mathbf{x}^*, \boldsymbol{\lambda}),$$

for any primal feasible \mathbf{x} and dual feasible $\boldsymbol{\lambda}$. [25pts]

Problem 4. Construct a linear programming problem for which both the primal and the dual problem has no feasible solution. [25pts]

Construct a linear programming problem that is:

$$\begin{aligned}
 \min_{x_1, x_2} \quad & x_1 - 2x_2 \\
 \text{s.t.} \quad & x_1 - x_2 \leq 1 \\
 & x_1 - x_2 \geq 2 \\
 & x_1, x_2 \leq 0
 \end{aligned} \tag{6}$$

Since it is impossible to satisfy $x_1 - x_2 \leq 1$ and $x_1 - x_2 \geq 2$ at the same time, so the primal problem has no feasible solution.

And the dual problem is that:

$$\begin{aligned}
 \max_{\lambda_1, \lambda_2} \quad & \lambda_1 + 2\lambda_2 \\
 \text{s.t.} \quad & \lambda_1 + \lambda_2 \leq 1 \\
 & -\lambda_1 - \lambda_2 \leq -2 \\
 & \lambda_1 \leq 0, \lambda_2 \geq 0
 \end{aligned} \tag{7}$$

The second constrain $-\lambda_1 - \lambda_2 \leq -2$ can be written as $\lambda_1 + \lambda_2 \geq 2$.

Since it is impossible to satisfy $\lambda_1 + \lambda_2 \leq 1$ and $\lambda_1 + \lambda_2 \geq 2$ at the same time, so the dual problem has no feasible solution.

So above all, the above construction's primal and dual problem has no feasible solution.