# Lab Week #6 CIS 314

November 5, 2017

## 1 Y86 Assembly Simulator

After loading the Y86 simulator (https://xsznix.github.io/js-y86/), you'll be greeted with the following code:

```
.pos 0
Init:
    irmovl Stack, %ebp
    irmovl Stack, %esp
.pos 0x100
Stack:
```

The above creates some variable "Stack", and makes it an alias of the memory location 0x100. In addition, it has two instructions to set the stack and base pointer to the hex value 0x100.

From here, we can start doing useful things by adding a "main" procedure:

This procedure isn't very interesting – it only moves the value '2' in to the register eax...

### 2 An Add Procedure

Let's make a more interesting procedure, one to add two numbers:

```
.pos 0
Init:
   irmovl Stack, %ebp
   irmovl Stack, %esp
   call Main
   halt
Add:
       //int add(int a, int b)
   pushl %ebp
   rrmovl %esp, %ebp //prologue
   mrmovl 8(%ebp), %eax //get first argument
   mrmovl 12(%ebp), %ecx //get second argument
   addl %ecx, %eax
                         //add the two arguments
   popl %ebp
              //epilogue
   ret
Main:
   pushl %ebp
   rrmovl %esp, %ebp //prologue
   irmovl $2, %eax
   pushl %eax
   irmovl $3, %eax
   pushl %eax
   call Add
   brk
              //at this point, the result is stored in %eax
   popl %ebp
              //epilogue
   ret
.pos 0x100
Stack:
```

Uh oh! At the end of the execution, esp and ebp weren't put back to where they started! We need to clean up the stack in Main after we have fiddled with it!

```
.pos 0
Init:
   irmovl Stack, %ebp
   irmovl Stack, %esp
   call Main
   halt
       //int add(int a, int b)
Add:
   pushl %ebp
   rrmovl %esp, %ebp //prologue
   mrmovl 8(%ebp), %eax //get first argument
   mrmovl 12(%ebp), %ecx //get second argument
   addl %ecx, %eax
                          //add the two arguments
   popl %ebp
   ret
              //epilogue
Main:
   pushl %ebp
   rrmovl %esp, %ebp //prologue
   irmovl $2, %eax
   pushl %eax
   irmovl $3, %eax
   pushl %eax
   call Add
   rrmovl %ebp, %esp //at this point, the result is stored in %eax
   popl %ebp
              //epilogue
   ret
.pos 0x100
Stack:
```

That's better.

#### 3 Mult

Now for an even harder procedure. Since y86 doesn't have a multiplication instruction, if we want to multiply two numbers, we'll need to repeatedly add. Let's have a first attempt:

```
Mult: //int mult(int a, int b)
   pushl %ebp
   rrmovl %esp, %ebp

mrmovl 8(%ebp), %eax
mrmovl 12(%ebp), %ecx
irmovl $1, %edx

loop:
   addl %eax, %eax
   subl %edx, %ecx
   jg loop

popl %ebp
   ret
```

For some reason, this code actually serves to left shift! A moment's thought reveals that the problem is in the 'add' instruction – we're always doubling eax! Let's try this instead:

```
Mult: //int mult(int a, int b)
   pushl %ebp
   rrmovl %esp, %ebp
   pushl %ebx
                       //since ebx is callee save, we need to back it up
   mrmovl 8(%ebp), %eax
   rrmovl %eax, %ebx
   mrmovl 12(%ebp), %ecx
   irmovl $1, %edx
   loop:
       addl %ebx, %eax
       subl %edx, %ecx
       jg loop
   popl %ebx
   popl %ebp
   ret
```

That still doesn't work! To figure out why, consider the case where we multiply something by 1 – we'll always get in to the loop, which implies we will always multiply by at least two! The solution is to decrement the counter variable before entering the loop!

```
Mult: //int mult(int a, int b)
   pushl %ebp
   rrmovl %esp, %ebp
   pushl %ebx
   mrmovl 8(%ebp), %eax
   rrmovl %eax, %ebx
   mrmovl 12(%ebp), %ecx
   irmovl $1, %edx
   subl %edx, %ecx
   je end
   loop:
       addl %ebx, %eax
       subl %edx, %ecx
       jg loop
   end:
   popl %ebx
   popl %ebp
```

There are some obvious problems with the procedure – it won't work for 0 or anything negative – but it serves to show some fun things in y86.

#### 4 Pointers in Y86

Up until now, we've only seen how to deal with values... what about dealing with pointers? We can declare an array as follows:

```
.pos 0x104 //make the array start at location 0x104
.align 4 //make each element in the array occupy 4 bytes
array:
.long 0x1 //element 0
.long 0x2 //element 1
.long 0x3 //element 2
```

Neat! What if we modify the add procedure to take pointers, and we return the result of adding the dereferenced values together:

```
Add: //int add(int* a, int* b)

pushl %ebp

rrmovl %esp, %ebp //prologue

mrmovl 8(%ebp), %eax //get first argument

mrmovl 12(%ebp), %ecx //get second argument

mrmovl (%eax), %eax //dereference first argument

mrmovl (%ecx), %ecx //dereference second argument

addl %ecx, %eax //add the two arguments

popl %ebp

ret //epilogue
```

That's all fine and good, but how do we send a pointer to 'Add'?

```
Main:

pushl %ebp

rrmovl %esp, %ebp //prologue

irmovl array, %eax

irmovl $4, %ecx

pushl %eax

addl %ecx, %eax

pushl %eax

call Add

rrmovl %ebp, %esp

popl %ebp

ret //epilogue
```

The above will add array[0] + array[1]. Fun!