Assignment #5
CIS 427/527

Group 2

February 22, 2016

2.2.1

Which of the following strings are formulas in predicate logic?

Solution

(a),(b),(f),(g) are formulas.

(c) isn't, as $f(m)$ is a term.
(d) isn't, as B is expecting two terms, yet $B(m,x)$ is a formula.
(e) isn't, as $B(m)$ doesn't have enough arguments.
(h) isn't, as $B(x)$ doesn't have enough arguments.

2.5.3

Provide proofs for the following sequents:

Solution

(a) $\forall x P(x) \vdash \forall y P(y)$
(b) $\forall x(P(x) \to Q(x)) \vdash (\forall x \neg Q(x)) \to (\forall x \neg P(x))$
(c) $\forall x(P(x) \to \neg Q(x)) \vdash \neg(\exists x(P(x) \land Q(x)))$

2.5.11

Prove the following sequents in predicate logic:

Solution

(a) $\forall x \forall y \forall z(S(x,y) \land S(y,z) \to S(x,z)), \forall x \neg S(x,x) \vdash \forall x \forall y(S(x,y) \to \neg S(y,x))$
(b) $\forall x(P(x) \lor Q(x)), \exists x \neg Q(x), \forall x(R(x) \to \neg P(x)) \vdash \exists x \neg R(x)$
(c) $\forall x(P(x) \to (Q(x) \lor R(x))), \neg \exists x(P(x) \land R(x)) \vdash \forall x P(x) \to Q(x)$
(e) $(\exists x P(x)) \to \forall y P(y) \vdash \exists x \forall y((P(x) \to P(y)) \land (P(y) \land P(x)))$
(f) $\exists x(P(x) \land Q(x)), \forall y(P(y) \to R(y)) \vdash \exists x R(x) \land Q(x)$

2.6.1

Consider the formula

$$\phi = \forall x\ \exists y\ Q(g(x,y), g(y,y), z)$$

Find two models M and M' with respective environments l and l' such that $M \models_l \phi$ but $M' \not\models_l \phi$

Solution

M
$A = \{a, b\}$
$Q^M = \{(a, a, b)\}$
$g^M(x, y) = a$ for all input
$z^M = b$

Then, $\forall x\ \forall y\ Q(g(x, y), g(y, y), z) = Q(a, a, b)$, which holds, so $M \models \phi$.

M'
$A = \{a, b\}$
$Q^{M'} = \{(a, a, c)\}$
$g^{M'}(x, y) = c$ for all input
$z^{M'} = a$

Then, $\forall x\ \forall y\ Q(g(x, y), g(y, y), z) = Q(c, c, a)$, but $(c, c, a) \notin Q^{M'}$, so $M' \not\models \phi$.

2.6.2

Consider the sentence

$$\phi = \forall x\ \exists y\ \exists z\ (P(x, y) \wedge P(z, y) \wedge (P(x, z) \rightarrow P(z, x)))$$

Which of the following models satisfies ϕ?

(a) $P^M = \{(m, n) | m < n\}$
(b) $P^{M'} = \{(m, 2 * m) | m \text{ natural number}\}$
(c) $P^{M''} = \{(m, n) | m < n + 1\}$

Solution

(a) This model does not satisfy ϕ, because we either need to force $P(x, z)$ to be false by requiring z to be smaller than x (in which case we can escape the natural numbers), or by having $x < z \wedge z < x$, which cannot happen.
(b) Yes, because the first two properties say $y = 2 * x$ and $y = 2 * z$, which means $x = z$ making $P(x, z)$ always false.
(c) Yes, let $y = z = x$, then all the properties hold.

2.6.3

Let P be a predicate with two arguments. Find a model M which satisfies $\forall x\ \neg P(x, x)$. Find also a model M' such that $M' \not\models \forall x\ \neg P(x, x)$.

Solution

- Let $P^M = \{(x, y) | x < y\}$, then we have $\forall x\ \neg P(x, x)$.
- Let $P^{M'} = \{(x, y) | x = y\}$, then we have $M' \not\models \forall x\ \neg P(x, x)$ as desired.

2.7.5

Show $\forall x(P(x) \bigvee Q(x)) \not\models \forall x P(x) \bigvee \forall x Q(x)$. Thus, find a model which satisfies $\forall x(P(x) \bigvee Q(x))$, but not $\forall x P(x) \bigvee \forall x Q(x)$

Solution

- Let $A =$ natural numbers, and let $P^M = \{x | x \leq 5\}$, $Q^M = \{x | x > 5\}$.

Then $\forall x. x \leq 5 \bigvee x > 5$ holds, but it is not true that $\forall x. x \leq 5 \bigvee \forall x. x > 5$.