

# CENG424

## Logic for Computer Science

### Assignment 2

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#### Question 1

1.

1.	$p \Rightarrow q$	Premise
2.	$q \Rightarrow r$	Premise
3.	$p \Rightarrow \neg r$	Premise (by Deduction Theorem)
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4.	$(q \Rightarrow r) \Rightarrow (p \Rightarrow (q \Rightarrow r))$	Implication Introduction (II)
5.	$p \Rightarrow (q \Rightarrow r)$	Modus Ponens (4, 2)
6.	$(p \Rightarrow q) \Rightarrow (p \Rightarrow r)$	Implication Distribution (ID)
7.	$p \Rightarrow r$	Modus Ponens (6, 1)
8.	$(p \Rightarrow \neg r) \Rightarrow ((p \Rightarrow r) \Rightarrow \neg p)$	Contradiction Rule (CR)
9.	$(p \Rightarrow r) \Rightarrow \neg p$	Modus Ponens (8, 3)
10.	$\neg p$	Modus Ponens (9, 7)

It is shown that  $(p \Rightarrow \neg r) \Rightarrow \neg p$ .

2.

1.	$\neg s \vee \neg r \Rightarrow \neg t$	Premise
2.	$s \Rightarrow p \vee \neg r$	Premise
3.	$p \wedge s \Rightarrow q$	Premise
4.	$t$	Premise (by Deduction Theorem)
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5.	$\neg(\neg t) \Rightarrow \neg(\neg s \vee \neg r)$	Contraposition
6.	$t \Rightarrow s \wedge r$	Double Negation Rule
7.	$s \wedge r$	Modus Ponens (6,4)
8.	$s$	Conjunction elimination (7)
9.	$r$	Conjunction elimination (7)
10.	$p \vee \neg r$	Modus Ponens (2,7)
11.	$p$	Disjunction elimination (10)
12.	$p \wedge s$	Conjunction introduction (8,12)
13.	$q$	Modus Ponens (3,12)

It is shown that  $(\neg s \vee \neg r \Rightarrow \neg t) \wedge (s \Rightarrow p \vee \neg r) \wedge (p \wedge s \Rightarrow q) \Rightarrow (t \Rightarrow q)$ .

#### Question 2

$$(r \Rightarrow s) \wedge ((t \wedge q) \Rightarrow \neg s) \Rightarrow ((t \wedge q) \Rightarrow \neg r)$$

First, we express the given logical formula in conjunctive normal form (CNF) and negate the goal.

Premises:

$$r \Rightarrow s \equiv \neg r \vee s$$

$$(t \wedge q) \Rightarrow \neg s \equiv \neg t \vee \neg q \vee \neg s$$

$$(t \wedge q) \Rightarrow \neg r \equiv \neg t \vee \neg q \vee \neg r$$

Thus, the premises are:

1.  $\neg r \vee s$
2.  $\neg t \vee \neg q \vee \neg s$
3.  $\neg t \vee \neg q \vee \neg r$

We negate the goal  $(t \wedge q) \Rightarrow \neg r$ :

$$\neg((t \wedge q) \Rightarrow \neg r) \equiv (t \wedge q) \wedge r$$

So, the negated goal gives us the following clauses:

4.  $t$
5.  $q$
6.  $r$

The premises and the negated goal:

1.  $\neg r \vee s$
2.  $\neg t \vee \neg q \vee \neg s$
3.  $\neg t \vee \neg q \vee \neg r$
4.  $t$
5.  $q$
6.  $r$

We resolve the clauses step by step:

Resolving  $\neg r \vee s$  (1) and  $r$  (6) gives us:

7.  $s$

Resolving  $\neg t \vee \neg q \vee \neg s$  (2) and  $s$  (7) gives us:

8.  $\neg t \vee \neg q$

Resolving  $\neg t \vee \neg q \vee \neg r$  (3) and  $\neg t \vee \neg q$  (8) gives us:

9.  $\neg r$

Resolving  $r$  (6) and  $\neg r$  (9) gives us:

10.  $\square$  (Contradiction)

Since we reached a contradiction (an empty clause,  $\square$ ), the negation of the goal leads to a contradiction. Therefore, the original implication is valid.

## Question 3

1.

If Paul has holidays and it is snowing, he will go skiing. **(HOL  $\wedge$  SNO)  $\Rightarrow$  GSK**

He will go to France or Florida. **GFR  $\vee$  GFL**

There is no skiing in Florida.  **$\neg$ SFL**

There is skiing in France. **SFR**

Paul has holidays and he will go to Florida. **HOL  $\wedge$  GFL**

## 2.

Conjunctive Normal Form (CNF) of statements above:

$$\begin{aligned}(\mathbf{HOL} \wedge \mathbf{SNO}) \Rightarrow \mathbf{GSK} &= \neg (\mathbf{HOL} \wedge \mathbf{SNO}) \Rightarrow \mathbf{GSK} = \neg \mathbf{HOL} \vee \neg \mathbf{SNO} \vee \mathbf{GSK} \\(\neg \mathbf{HOL} \vee \neg \mathbf{SNO} \vee \mathbf{GSK}) \wedge (\mathbf{GFR} \vee \mathbf{GFL}) \wedge (\neg \mathbf{SFL}) \wedge (\mathbf{SFR}) \wedge (\mathbf{HOL}) \wedge (\mathbf{GFL})\end{aligned}$$

Statement we need to prove: It is not snowing.  $(\neg \mathbf{SNO})$

We assume  $\mathbf{SNO}$  (it is snowing) and will derive a contradiction.

*The resolution process:*

Resolving  $\mathbf{SNO}$  and  $(\neg \mathbf{HOL} \vee \neg \mathbf{SNO} \vee \mathbf{GSK})$ , end up with:  $(\neg \mathbf{HOL} \vee \mathbf{GSK})$

Resolve  $(\neg \mathbf{HOL} \vee \mathbf{GSK})$  with  $\mathbf{HOL}$ , end up with:  $\mathbf{GSK}$

Resolve  $\mathbf{GFL}$  with  $\mathbf{GSK}$ , end up with:  $\mathbf{GFL} \wedge \mathbf{GSK}$

Resolve  $\mathbf{GFL} \wedge \mathbf{GSK}$  with  $\neg \mathbf{SFL}$ , end up with a contradiction and an empty clause because Paul will go to Florida and he will go skiing, but there is no skiing in Florida.

Since we have reached a contradiction, we conclude that the assumption  $\mathbf{SNO}$  is false, and therefore  $\neg \mathbf{SNO}$ .