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## My Report!

First year review report

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#### Abstract

Giving a short overview of the work in your project.[1]

#### 0.1 Natural Numbers

First, we define the type of natural numbers inductively:

Here,  $\mathbb{N}$  is the type of natural numbers, with two constructors:

- zero represents 0.
- suc represents the successor function (i.e., n+1).

#### 0.2 Addition

Next, we define addition recursively:

```
-+ : \mathbb{N} \to \mathbb{N} \to \mathbb{N}

\operatorname{zero} + n = n

\operatorname{suc} m + n = \operatorname{suc} (m + n)
```

This definition states:

- 0 + n = n (base case).
- (m + 1) + n = (m + n) + 1 (recursive case).

### 0.3 A Simple Proof

We now prove that 2 + 2 = 4. First, we define the numbers:

```
\begin{split} & \text{two}: \, \mathbb{N} \\ & \text{two} = \text{suc (suc zero)} \\ & \text{four}: \, \mathbb{N} \\ & \text{four} = \text{suc (suc (suc (suc zero)))} \end{split}
```

Now, the proof reduces by computation:

```
proof : two + two \equiv four
proof = refl
```

Since Agda's definitional equality handles reduction, refl suffices.

#### 0.4 Conclusion

This example shows how Agda and LaTeX can be combined for formal proofs in papers. The full output is rendered with syntax highlighting.

# Bibliography

[1] ABBOTT, M., ALTENKIRCH, T., AND GHANI, N. Containers: Constructing strictly positive types. *Theoretical Computer Science 342*, 1 (2005), 3–27. Applied Semantics: Selected Topics.