



ĐẠI HỌC ĐÀ NẴNG

TRƯỜNG ĐẠI HỌC CÔNG NGHỆ THÔNG TIN VÀ TRUYỀN THÔNG VIỆT - HÀN
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Nhân bản – Phụng sự – Khai phóng

Linked Lists

Data Structures & Algorithms

- **Linked List Basics**
- **Singly linked lists**
- **Circularly linked list**
- **Doubly linked lists**
- **Applications**
 - Polynomial representation
 - Equivalence relations
 - Sparse matrices

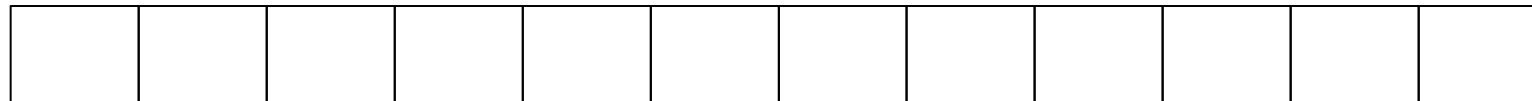
- **Linked List Basics**
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- Linked lists & arrays are similar - Both store collections of data.
- **Array**: features all follow from its strategy of allocating sequentially its elements.
- **Linked lists**: use an entirely different strategy: linked lists allocate memory for each element separately and only when necessary.
 - Linked lists are used to store a collection of information (like arrays)
 - A linked list is made of nodes that are pointing to each other
 - We only know the address of the first node (head)
 - Other nodes are reached by following the “next” pointers
 - The last node points to NULL

- **Linked List vs. Array**

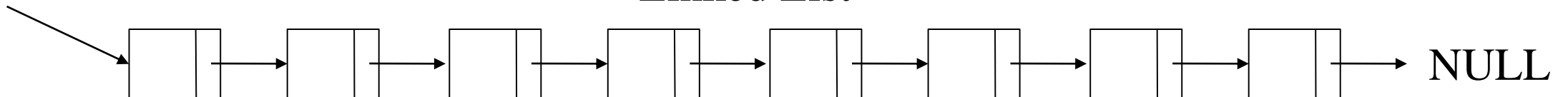
- Elements of array are contiguous
- In a linked list, nodes are not necessarily contiguous in memory (each node is allocated with a separate “new” call)

Array



head

Linked List



- **Linked List vs. Array**

Array:

- Advantages:

- Easy to use
- A good choice for a small list
- $O(1)$ access time

- Disadvantages:

- Fixed size
- Memory wasting
- Still space and time wasting in dynamic array

Linked List:

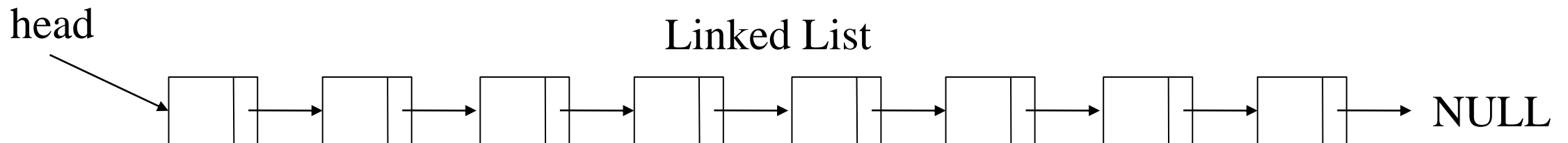
- Advantages:

- Arbitrary size
- No shift required

- Disadvantages:

- Necessity to allocate next
- $O(N)$ access time

- we use linked lists if...
 - The number of elements that will be stored cannot be predicted at compile time
 - Elements may be inserted in the middle or deleted from the middle
 - We are less likely to make random access into the data structure (because random access is expensive for linked lists)

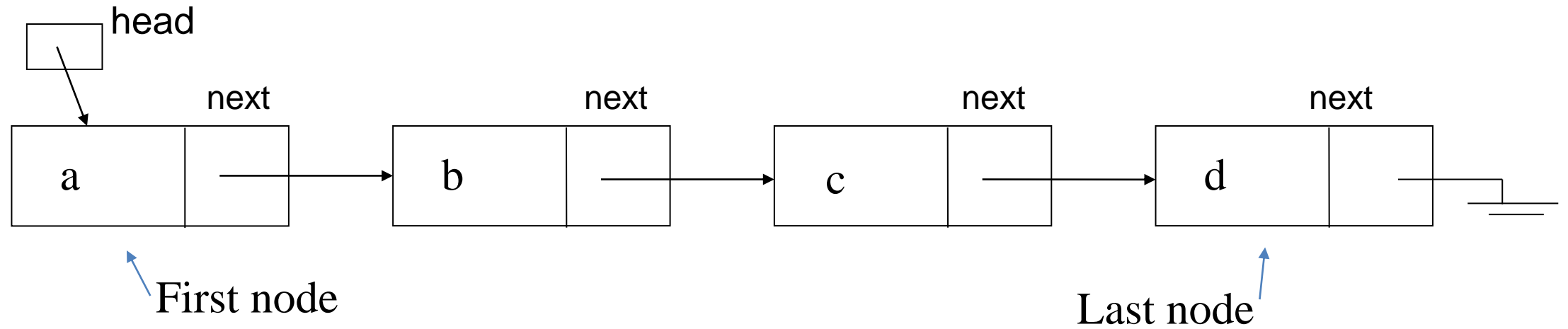


- **Operations on Linked List**

- **Generic methods:** size(), isEmpty()
- **Query methods:** isFirst(p), isLast(p)
- **Accessor methods:** first(), last()
before(p), after(p)
- **Update methods:** insertFirst(e), insertLast(e)
insertBefore(p,e), insertAfter(p,e)
removeAfter(p)
invert(p)
replaceElement(p,e)
swapElements(p,q),
...

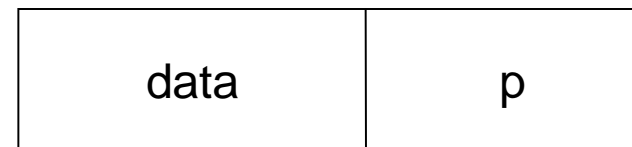
- Linked List Basics
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• Linked Lists



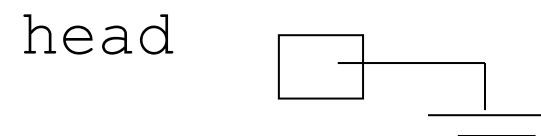
• Each node has (at least) 2 fields:

- Data
- Pointer to the next node



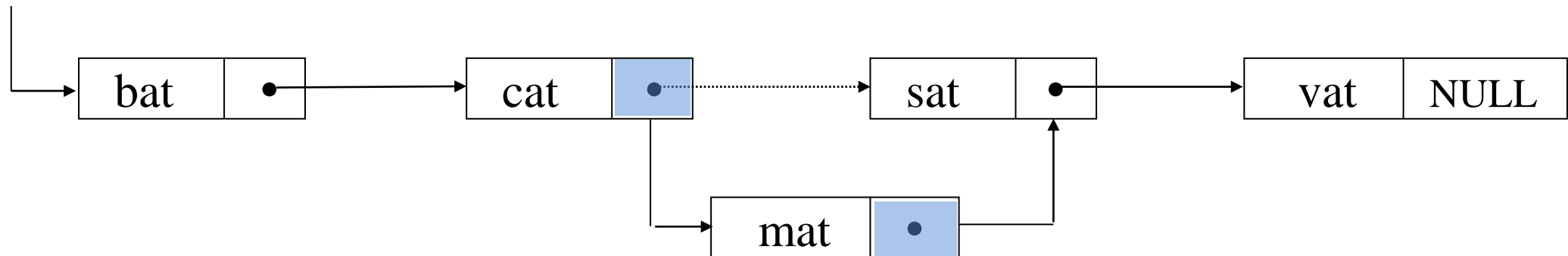
• Empty linked

`head = NULL;`



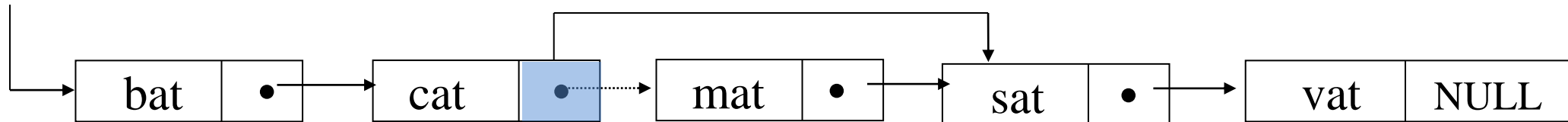
- Insert **mat** between **cat** & **sat**

1. Get a node that is currently unused; let its address be paddr.
2. Set the data field of this node to *mat*.
3. Set paddr's link field to point to the address found in the link field of the node containing *sat*.
4. Set the link field of the node containing *cat* to point to paddr.



- Delete **mat** from the list

- Find the element that immediately precedes **mat**, which is **cat**, and set its link field to point to **mat**'s link field



• Implementation

–Declaration

```
typedef struct node *pnode;
typedef struct node {
    char data [4];
    pnode next;
};
```

–Creation

```
pnode head =NULL;
```

–Testing

```
#define IS_EMPTY(ptr) (!(ptr))
#define IS_FULL(ptr) (!(ptr))
```

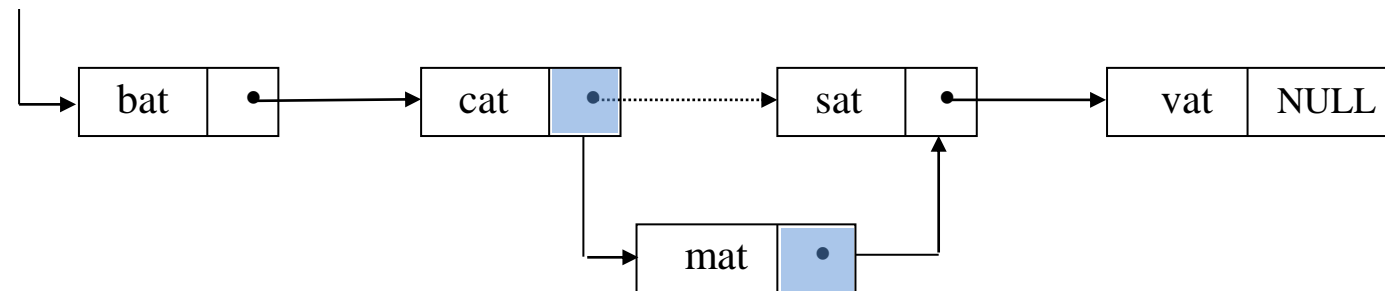
–Traverse a list

```
void traverseList(pnode head){
    printf("The list contains: ");
    for (pnode p = head ; p!=NULL ;
        p = p->next)
        printf("%s\n", p->data);
}
```

```
p=head;
while (p!=NULL){
    cout<<P->data;
    p=p->next;
}
```

- Implementation - **Insert after a specific position**

```
void insertAfter(pnode p, char* data){
    /* insert a new node with data into the list ptr after node */
    pnode temp;
    temp = (pnode) malloc(sizeof(node));
    if (IS_FULL(temp)){
        fprintf(stderr, "The memory is full\n");
        exit (1);
    }
    strcpy(temp->data, data);
    if (p) { //noempty list
        temp->next=p->next;
        p->next= temp;
    }else { //empty list
        temp->next= NULL;
        p =temp;
    }
}
```



- Implementation - **Delete a node after a specific position**

```
void removeAfter(pnode p){
```

```
    /* delete what follows after node p in the list */
```

```
    pnode tmp;
```

```
    if (p) {
```

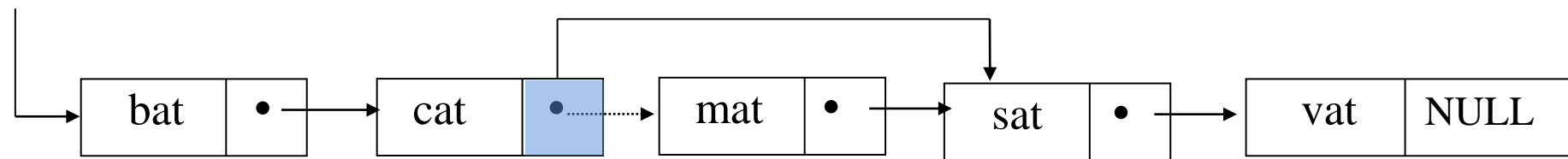
```
        tmp = p -> next;
```

```
        p->next = p->next->next;
```

```
        free(tmp);
```

```
    }
```

```
}
```



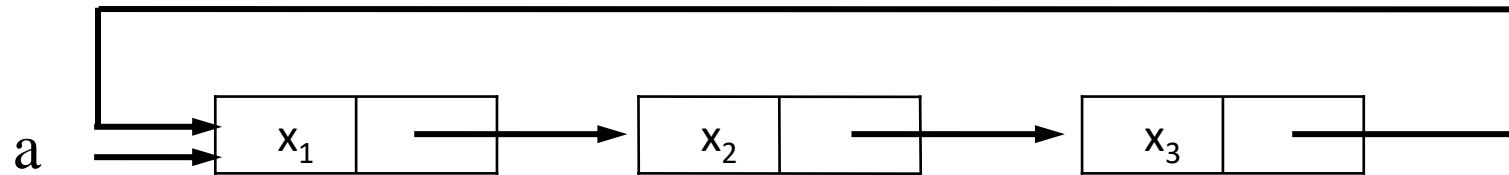
- Implementation - **Inverting a list**

```
pnode invertList(pnode lead){  
    /* invert the chain pointed to by lead */  
    pnode middle, trail;  
    middle = NULL;  
    while (lead) {  
        trail = middle;  
        middle = lead;  
        lead = lead->next;  
        middle->next = trail  
    }  
    return middle;  
}
```

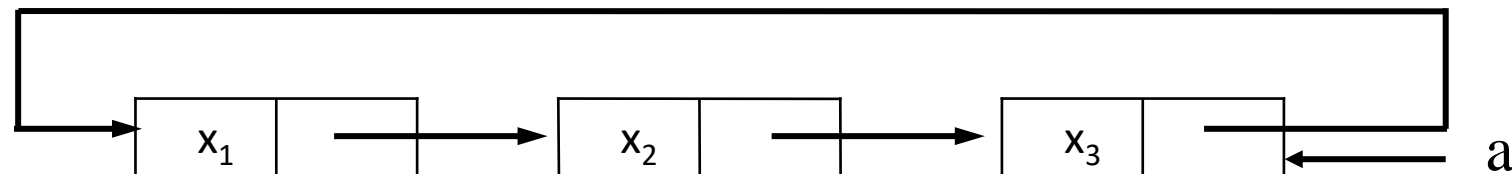

- Linked List Basics
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- **Circularly linked list**
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- Circularly linked list

- The link field of the last node points to the first node in the list



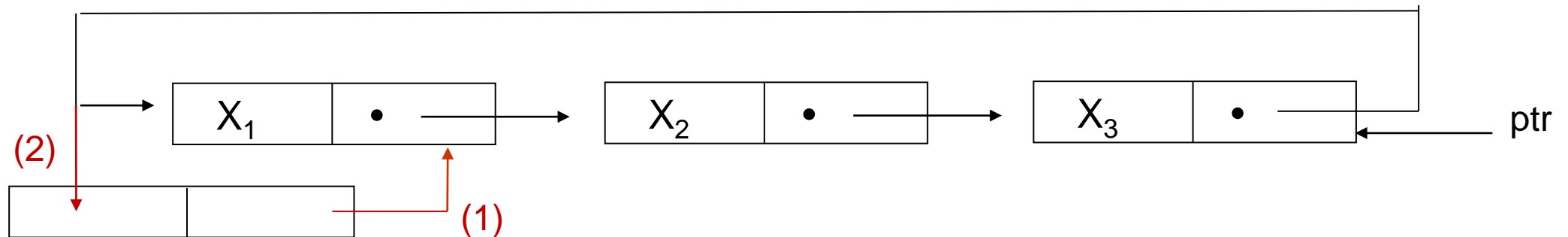
- It is more convenient when insert a new node if the name of the circular list points to the last node



• Insert a node

```
void insertFront (pnode* ptr, pnode node){
    /* insert a node in the list with head (*ptr)->next */
    if (IS_EMPTY(*ptr)){
        *ptr= node;
        node->next = node;    /* circular link */
    }
    else {
        node->next = (*ptr)->next;    (1)
        (*ptr)->next = node;          (2)
    }
}
```

```
typedef struct node *pnode;
typedef struct node {
    char data;
    pnode next;
};
```



- **List length**

```
int length(pnode ptr){
    pnode temp;
    int count = 0;
    if (ptr) {
        temp = ptr;
        do {
            count++;
            temp = temp->next;
        } while (temp!=ptr);
    }
    return count;
}
```

- **Print list**

```
void printList(pnode start, pnode ptr){
    if (start == ptr) return;
    if (ptr) printf("%c ", ptr->data);
    printList(start, ptr->next);
}
⇒ Use: printList(start, start->next);
```

- **Other operations**

- Create a node (with data)
- Delete a node (with data)
- Find a node

- Linked List Basics
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- **Singly linked lists**

- Some operations are expensive
 - insertLast, removeAfter
 - Why? ⇒ **need for traversing the list**

⇒ Solution: add **previous** link to elements

- **Doubly linked list has at least three fields:**

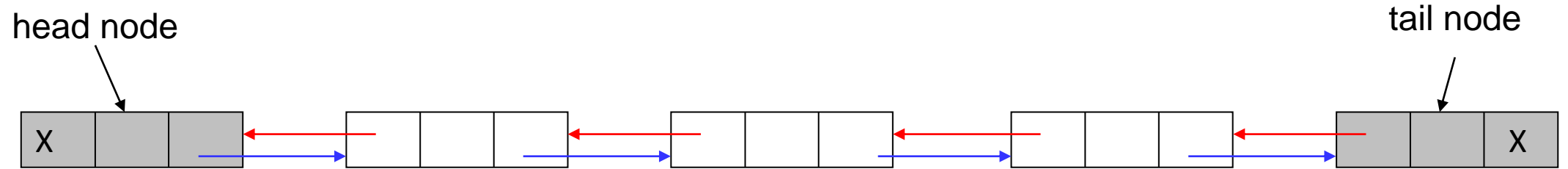
- a left link field (llink)
- a data field (item)
- a right link field (rlink)

- **Declarations**

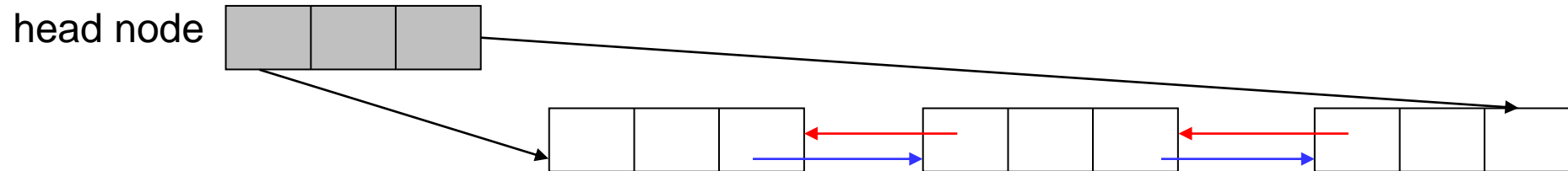
```
typedef struct node *node_pointer;  
typedef struct node {  
    node_pointer llink;  
    element item;  
    node_pointer rlink;  
}
```

- **Different uses**

- Doubly linked list with a head node and a tail node

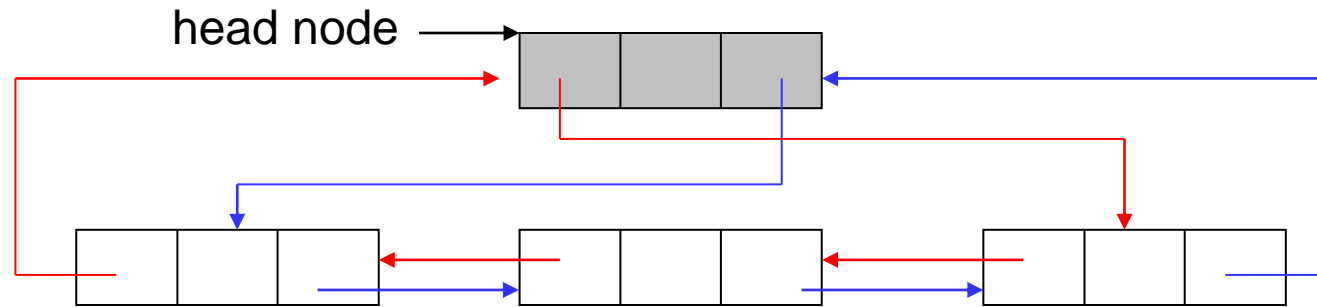


- Doubly linked list with a head node points to the first node in the list *and* to the last node in the list



- **Different uses**

- Doubly linked circular list with head node

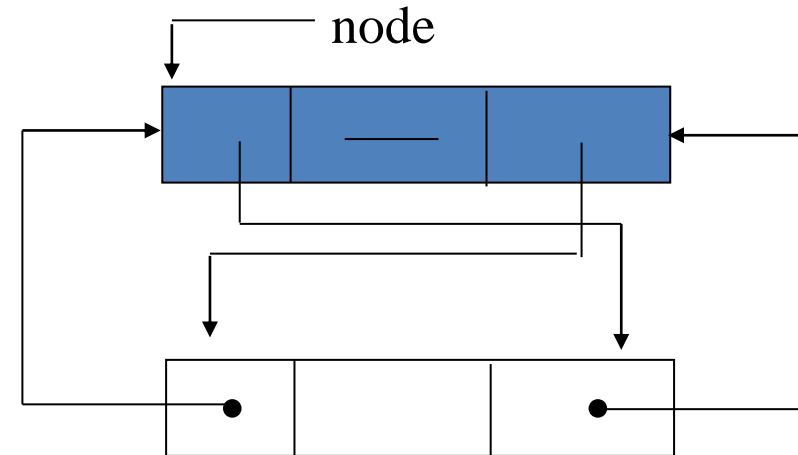
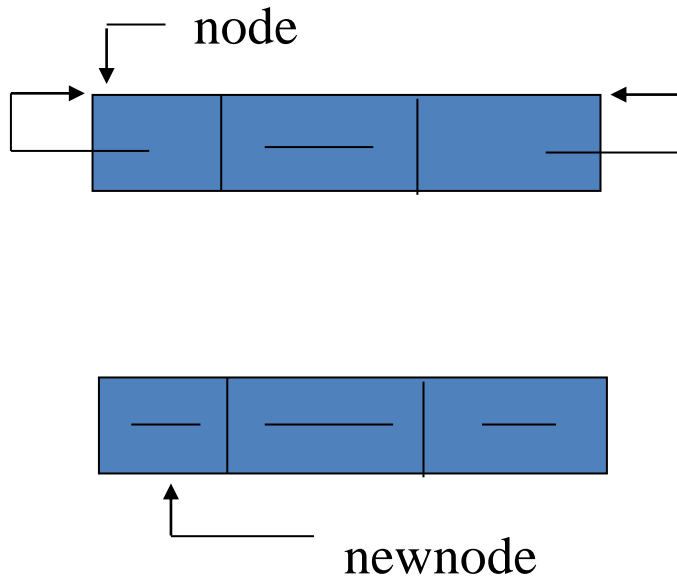


- Empty doubly linked circular list with head node



- Different uses

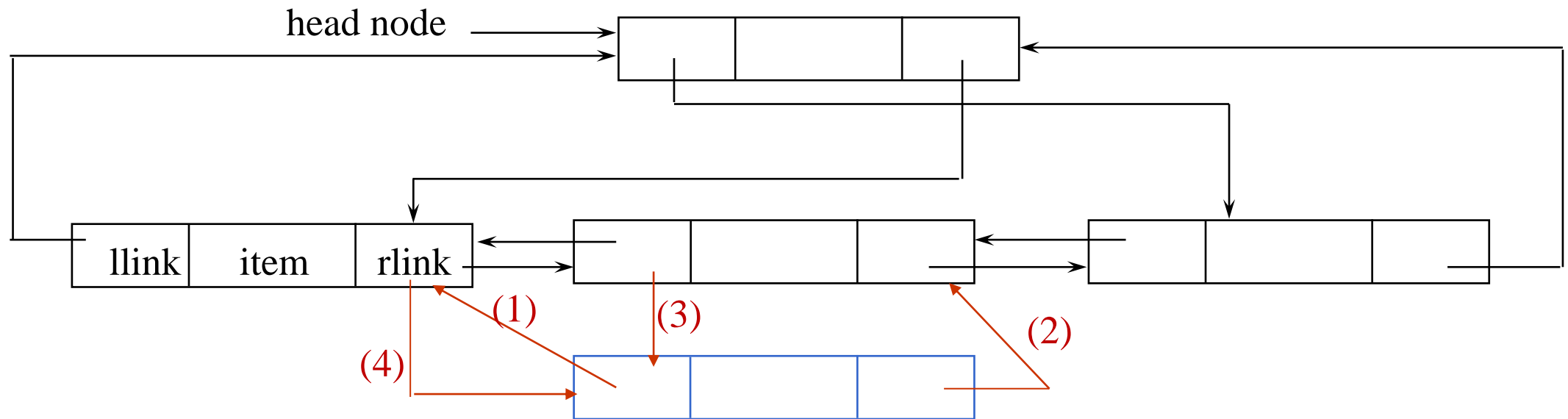
- Doubly linked circular list with head node
 - Insertion into an empty doubly linked circular list



- Different uses

- Doubly linked circular list with head node

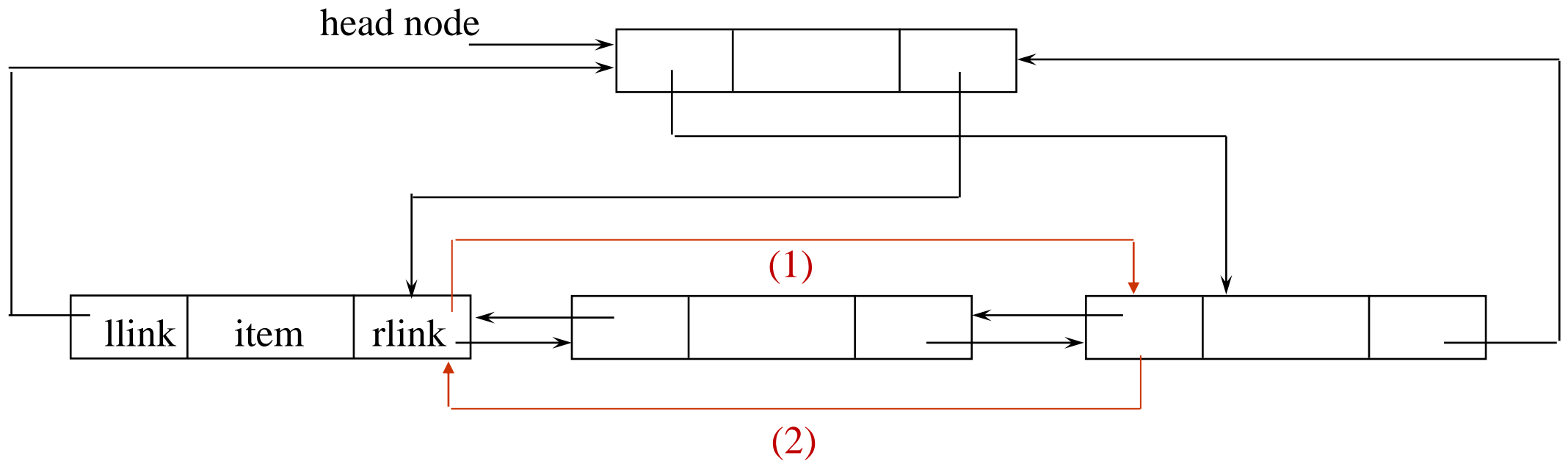
- Insertion into a doubly linked circular list



- Different uses

- Doubly linked circular list with head node

- Deletion from a doubly linked circular list



- Linked List Basics
- Singly linked lists
- Stacks and Queues
- Circularly linked list
- Doubly linked lists
- **Applications**
 - Polynomial representation
 - Equivalence relations
 - Sparse matrices

- **Polynomials representation**

- Representing polynomials as singly linked lists

$$A(x) = a_{m-1}x^{e_{m-1}} + a_{m-2}x^{e_{m-2}} + \dots + a_0x^{e_0}$$

a_i are nonzero coefficients, e_i are nonnegative integer exponents such that

$$e_{m-1} > e_{m-2} > \dots > e_1 > e_0 \geq 0$$

- Each term as a node containing coefficient , exponent, as well as a pointer to the next term

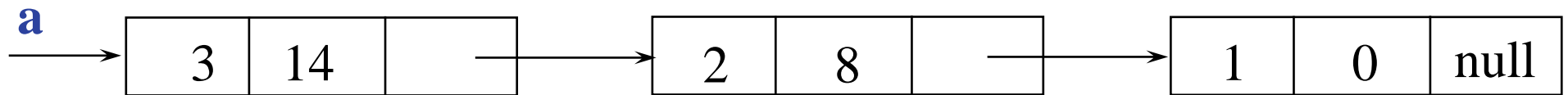
coef	expon	link
------	-------	------

- **Declarations**

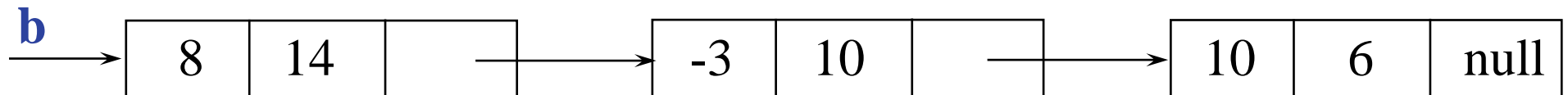
```
typedef struct poly_node  *poly_pointer;
typedef struct poly_node {
    int coef;
    int expon;
    poly_pointer link;
};
poly_pointer a , b, c;
```

- Polynomials representation - Example

$$a = 3x^{14} + 2x^8 + 1$$



$$b = 8x^{14} - 3x^{10} + 10x^6$$



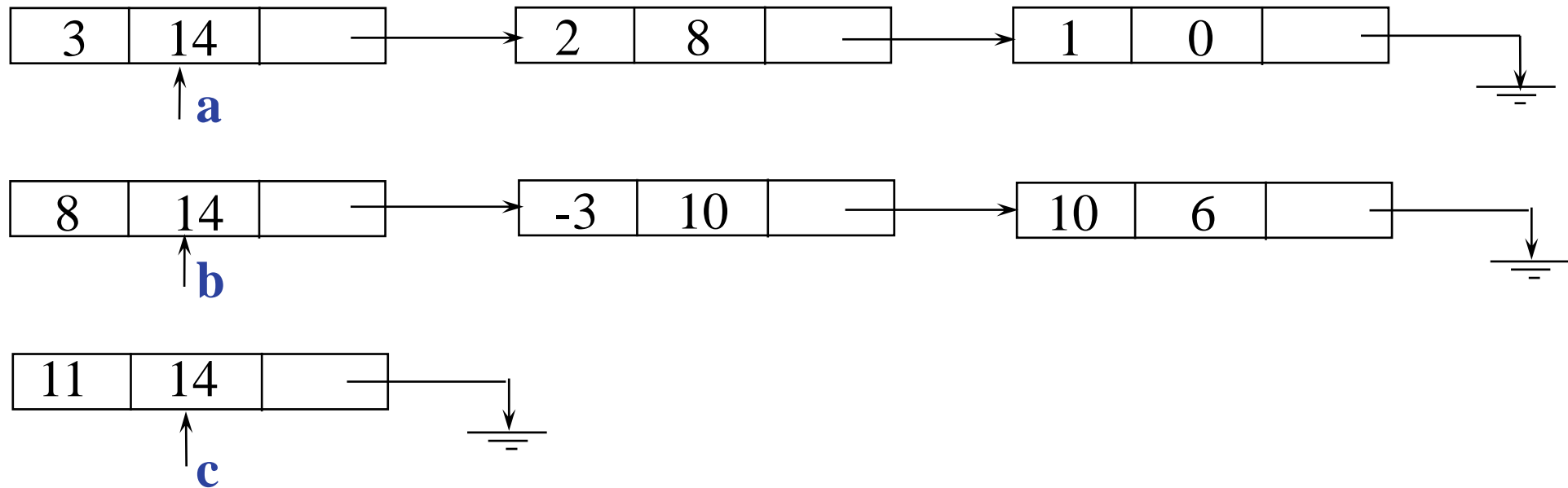
- Add two polynomials

To add 2 polynomials, we examine their terms starting at the nodes pointed to by **a** and **b**, there is 3 cases:

1. If the exponents of the two terms are equal, we add the two coefficients and create a new term **c** for the result
2. If the exponent of the current term in **a** is less than the exponent of the current term in **b**, then we create a duplicate term of **b**, attach this term to the result, called **c**, and advance the pointer to the next term in **b**.
3. Take a similar action on **a** if $a \rightarrow \text{expon} > b \rightarrow \text{expon}$

1. $a \rightarrow \text{expon} == b \rightarrow \text{expon}$

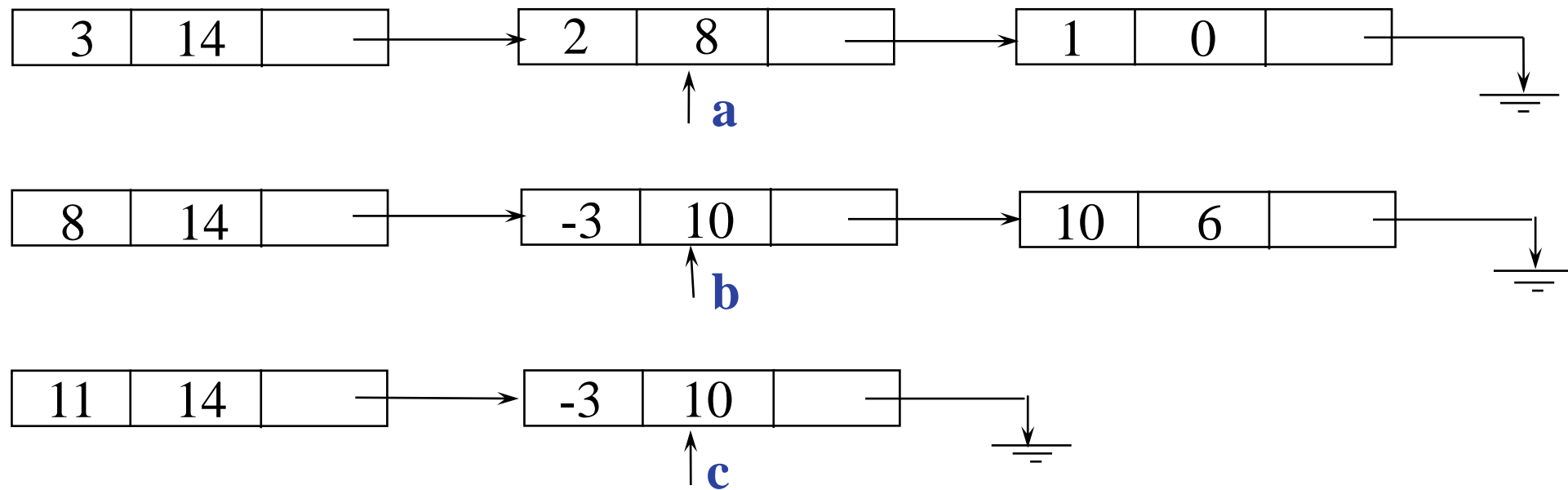
If the exponents of the two terms are equal, we add the two coefficients and create a new term c for the result



$$a \rightarrow \text{expon} == b \rightarrow \text{expon}$$

2. $a \rightarrow \text{expon} < b \rightarrow \text{expon}$

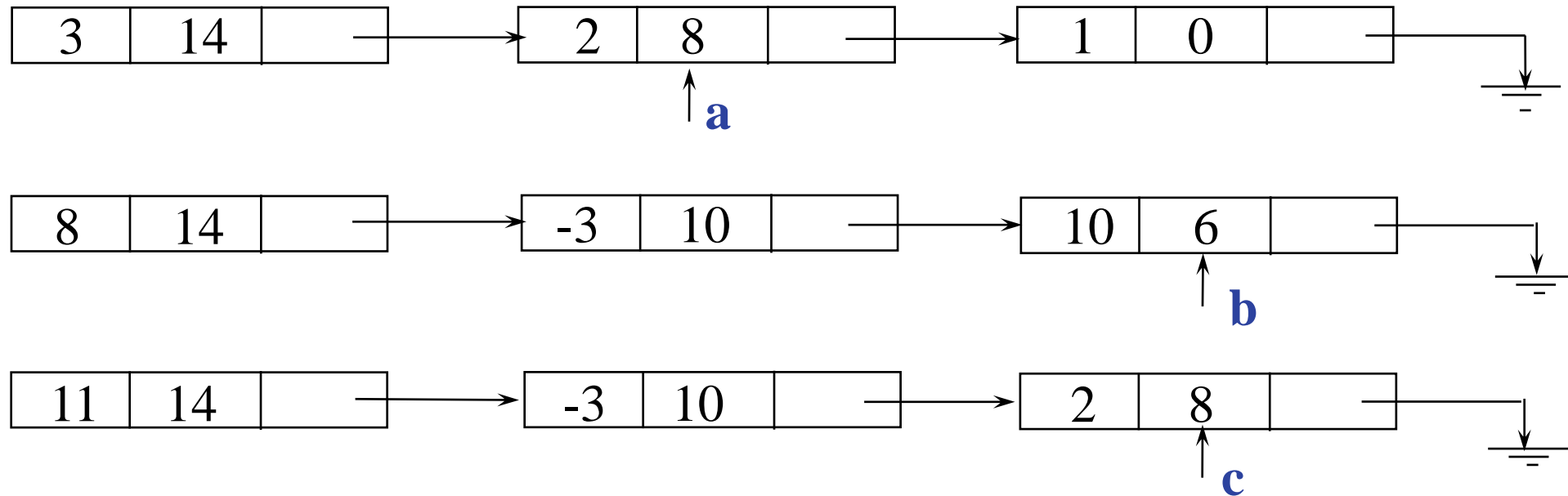
If the exponent of the current term in a is less than the exponent of the current term in b , then we create a duplicate term of b , attach this term to the result, called c , and advance the pointer to the next term in b .



$a \rightarrow \text{expon} < b \rightarrow \text{expon}$

3. $a \rightarrow \text{expon} > b \rightarrow \text{expon}$

Take a similar action on a if $a \rightarrow \text{expon} > b \rightarrow \text{expon}$



$a \rightarrow \text{expon} > b \rightarrow \text{expon}$

- Function Add two polynomials

```
poly_pointer addPoly(poly_pointer a, poly_pointer b) {  
    poly_pointer front, rear, temp;  
    int sum;  
    rear =(poly_pointer)malloc(sizeof(poly_node));  
    if (IS_FULL(rear)) {  
        fprintf(stderr, "The memory is full\n");  
        exit(1);  
    }  
    front = rear;  
    while (a && b) {  
        switch (COMPARE(a->expon, b->expon)) {
```

• Function Add two polynomials

```

case -1:          /* a->expon < b->expon */
    attach(b->coef, b->expon, &rear);
    b = b->link;
    break;

case 0:           /* a->expon == b->expon */
    sum = a->coef + b->coef;
    if (sum)       attach(sum, a->expon, &rear);
    a = a->link;   b = b->link;
    break;

case 1:           /* a->expon > b->expon */
    attach(a->coef, a->expon, &rear);
    a = a->link;

} // end switch
} //end while

```

- **Function Add two polynomials**

```
for (; a; a = a->link)
    attach(a->coef, a->expon, &rear);
for (; b; b = b->link)
    attach(b->coef, b->expon, &rear);
rear->link = NULL;
temp = front;
front = front->link;
free(temp);
return front;
} // end function
```

Delete extra initial node.

- **Attach a term**

```
void attach(float coefficient, int exponent, poly_pointer *ptr){  
    /* create a new node attaching to the node pointed to  
       by ptr. ptr is updated to point to this new node. */  
    poly_pointer temp;  
    temp = (poly_pointer) malloc(sizeof(poly_node));  
    if (IS_FULL(temp)) {  
        fprintf(stderr, "The memory is full\n");  
        exit(1);  
    }  
    temp->coef = coefficient;  
    temp->expon = exponent;  
    (*ptr)->link= temp;  
    *ptr = temp;  
}
```

- Analysis

- (1) coefficient additions

$$0 \leq \text{additions} \leq \min(m, n)$$

where m (n) denotes the number of terms in a (b).

- (2) exponent comparisons

extreme case

$$e_{m-1} > f_{m-1} > e_{m-2} > f_{m-2} > \dots > e_0 > f_0$$

$m+n-1$ comparisons

- (3) creation of new nodes

extreme case

$m + n$ new nodes

summary $O(m+n)$

- Erasing polynomials

```
void erase(poly_pointer *ptr){  
    /* erase the polynomial pointed to by ptr */  
    poly_pointer temp;  
    while (*ptr) {  
        temp = *ptr;  
        *ptr = (*ptr)->link;  
        free(temp);  
    }  
}
```


- **Equivalence relations**

- A relation over a set, S , is said to be an *equivalence relation* over S iff it is **symmetric**, **reflexive**, and **transitive** over S .
 - Reflexivity: $x=x$
 - Symmetry: if $x=y$, then $y=x$
 - Transitivity: if $x=y$ and $y=z$, then $x=z$

- **Example**

$0=4, 3=1, 6=10, 8=9, 7=4, 6=8, 3=5, 2=11, 11=0$

\Rightarrow three equivalent classes:

$\{0,2,4,7,11\}; \{1,3,5\}; \{6,8,9,10\}$

- Algorithm to find Equivalence Classes

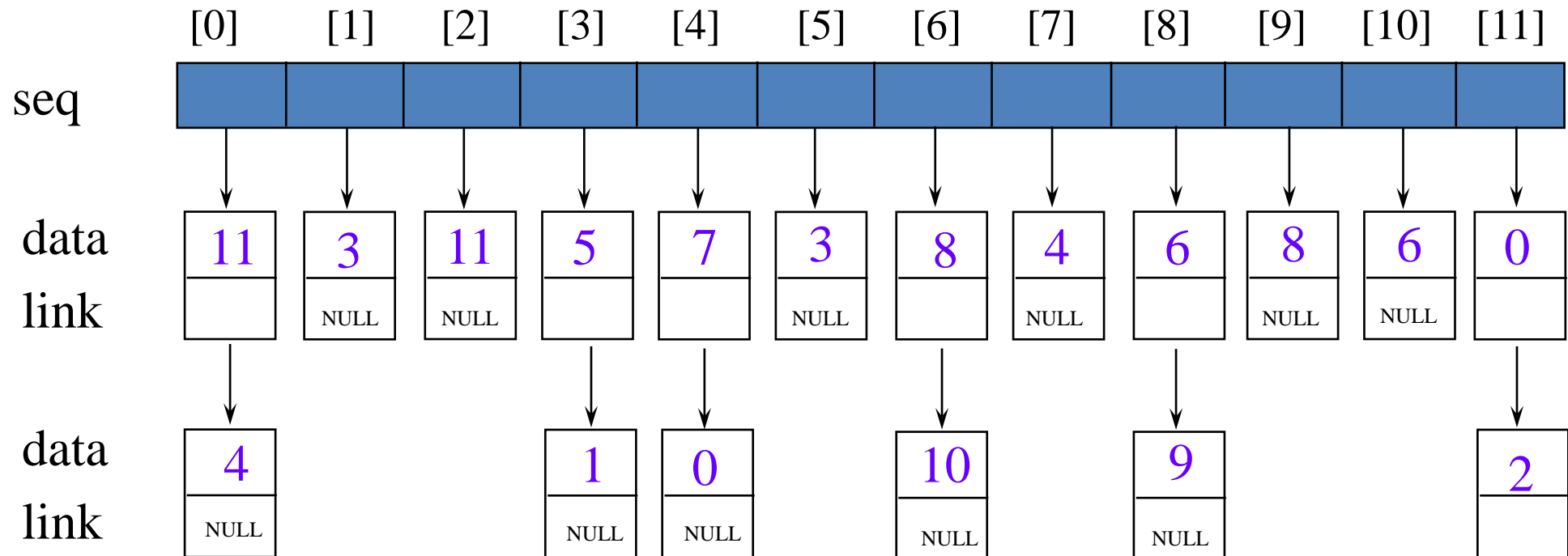
```
void equivalence() {  
    initialize data structures;  
    while (there are more pairs) {  
        read the next pair <i,j>;  
        process this pair;  
    }  
    initialize the output;  
    do {  
        output a new equivalence class;  
    } while (not done);  
}
```

- More detailed Algorithm to find Equivalence Classes

```
void equivalence() {  
    initialize seq to NULL and out to TRUE;  
    while (there are more pairs) {  
        read the next pair, <i,j>;  
        put j on the seq[i] list;  
        put i on the seq[j] list;  
    }  
    for (i=0; i<n; i++)  
        if (out[i]) {  
            out[i]= FALSE;  
            output this equivalence class;  
            compute indirect equivalence using transitivity by using stack;  
        }  
}
```

• Illutration

$0 \equiv 4$
 $3 \equiv 1$
 $6 \equiv 10$
 $8 \equiv 9$
 $7 \equiv 4$
 $6 \equiv 8$
 $3 \equiv 5$
 $2 \equiv 11$
 $11 \equiv 0$



Example: $0=4, 3=1, 6=10, 8=9, 7=4, 6=8, 3=5, 2=11, 11=0$
 \Rightarrow three equivalent classes: $\{0,2,4,7,11\}; \{1,3,5\}; \{6,8,9,10\}$

- **Program (1/4)**

```
#include <stdio.h>
#include <alloc.h>
#define MAX_SIZE 24
#define IS_FULL(ptr) (!(ptr))
#define FALSE 0
#define TRUE 1
typedef struct node *node_pointer ;
typedef struct node {
    int data;
    node_pointer link;
};
```

- Program (2/4)

```
void main(void) {  
    short int out[MAX_SIZE];  
    node_pointer seq[MAX_SIZE];  
    node_pointer x, y, top;  
    int i, j, n;  
    printf("Enter the size (<= %d) ", MAX_SIZE);  
    scanf("%d", &n);  
    for (i=0; i<n; i++) {  
        out[i]= TRUE;  seq[i]= NULL;  
    }  
    printf("Enter a pair of numbers (-1 -1 to quit): ");  
    scanf("%d%d", &i, &j);
```

- Program (3/4)

```
while (i>=0) {
    //Phase 1: input the equivalence pairs:
    x = (node_pointer) malloc(sizeof(node));
    if (IS_FULL(x))
        fprintf(stderr, "memory is full\n");
        exit(1);
    }
    x->data= j; x->link= seq[i]; seq[i]= x;    //Insert x to the top of lists seq[i]
    x = (node_pointer) malloc(sizeof(node));
    if (IS_FULL(x))
        fprintf(stderr, "memory is full\n");
        exit(1);
    }
    x->data= i; x->link= seq[j]; seq[j]= x;    //Insert x to the top of lists seq[j]
    printf("Enter a pair of numbers (-1 -1 to quit): ");
    scanf("%d%d", &i, &j);
}
```

•Program (4/4)

```

for (i=0; i<n; i++) {
    if (out[i]) {
        //Phase 2: output the equivalence classes
        printf("\nNew class: %5d", i);
        out[i]= FALSE;    //mark class as output
        x = seq[i];  top = NULL;    //initialize stack
        for (;;) { //find the entire class
            while (x) { //process a list
                j = x->data;
                if (out[j]) { // first time, visit this seq[j]
                    printf("%5d", j);
                    out[j] = FALSE;
                    y = x->link; x->link = top; //push the linked number to stack if it links to
                    top = x; x = y; // number another
                } else x = x->link;
            }
            if (!top) break; //stack empty
            x = seq[top->data]; top = top->link; //pop from stack to find the same class number
        } //for (;;)
    } // for (i=0; i<n; i++)
} //main()

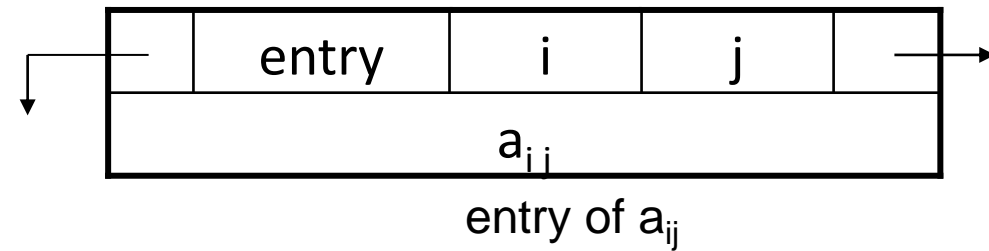
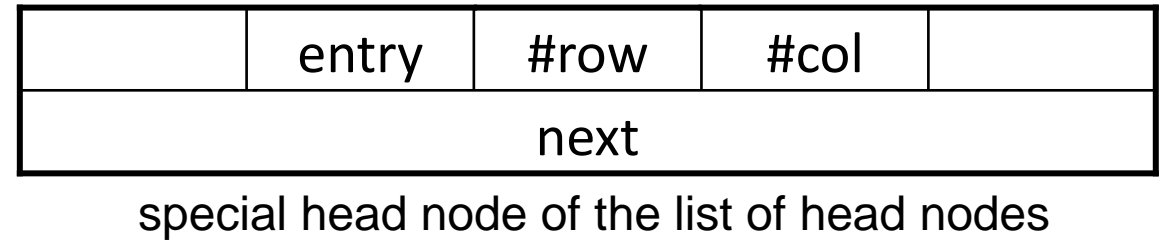
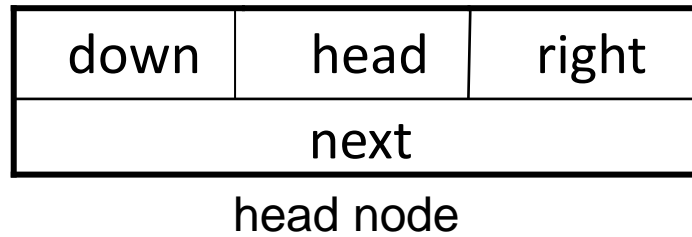
```


- Sparse matrices

- Each column of a sparse matrix is represented as a circularly linked list with a head node
- A similar representation for each row of a sparse matrix
- Each node has a tag field that is used to distinguish between head nodes and entry nodes
- Each **head node** has three fields: down, right, and next
 - down field: links into a column list
 - right field: links into a row list
 - next field: links the head nodes together
- The head node for row i is also the head node for column i , and the total number of head nodes is $\max \{\text{number of rows, number of columns}\}$

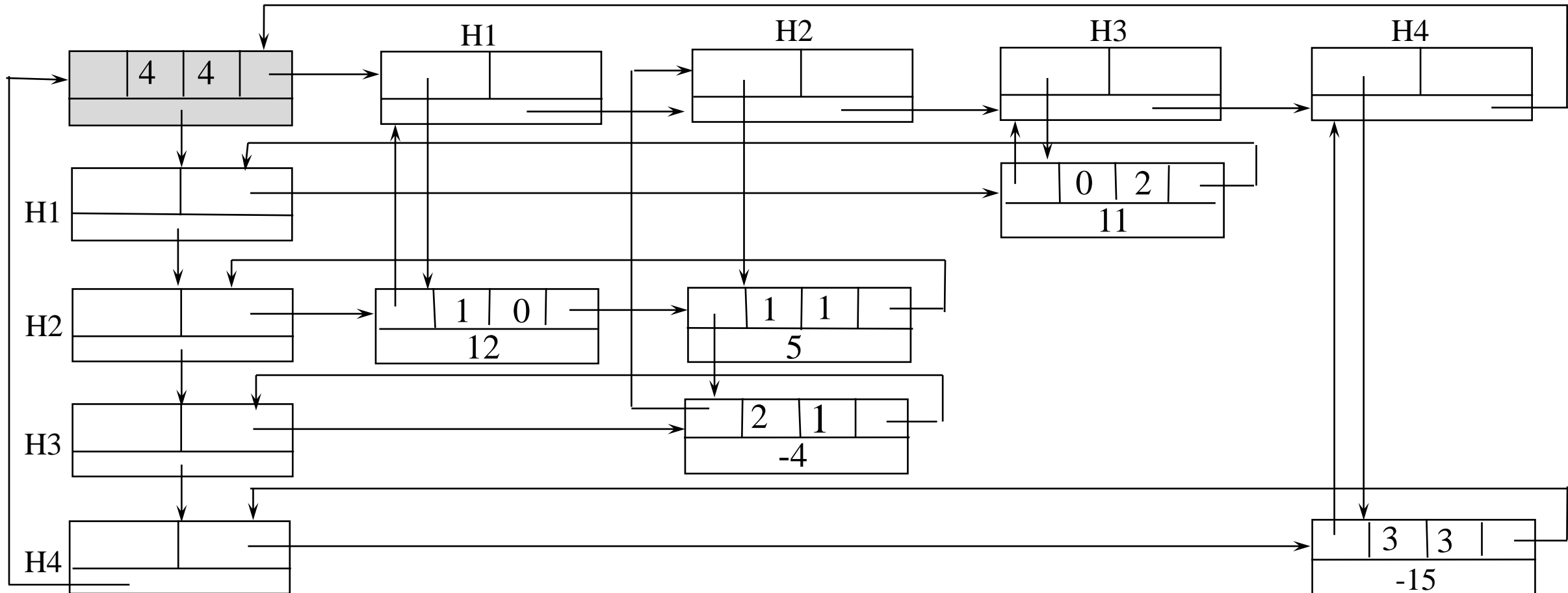
- Each entry node has 6 fields: tag, row, col, down, right, value.
 - down field: links to the next nonzero term in the same column
 - right field: links to the next nonzero term in the same row
 - tag field: entry
 - row field: row index
 - col field: column index
 - value field: nonzero value
- A $num_rows \times num_cols$ matrix with num_terms nonzero terms needs $\max\{num_rows, num_cols\} + num_terms + 1$ node
 - $\max\{num_rows, num_cols\}$: number of head nodes
 - num_terms : number of nonzero terms
 - 1 node: a special head node for the list of row and column head nodes contains the dimensions of the matrix
- Total storage will be less than $num_rows \times num_cols$ when num_terms is sufficiently small

• Sparse matrices



$$\# \text{ of head nodes} = \max \{ \# \text{ of rows}, \# \text{ of columns} \}$$

$$\begin{bmatrix} 0 & 0 & 11 & 0 \\ 12 & 5 & 0 & 0 \\ 0 & -4 & 0 & 0 \\ 0 & 0 & 0 & -15 \end{bmatrix}$$



• Program (1) - Declarations

```
#include <stdio.h>
#include <stdlib.h>
#define MAX_SIZE 25
typedef enum {head,entry} tagfield;
typedef struct MatrixNode *MatrixPointer;
struct EntryNode {
    int row;
    int col;
    int value;
};
```

```
struct MatrixNode {
    MatrixPointer down;
    MatrixPointer right;
    tagfield tag;
    union {
        MatrixPointer next;
        struct EntryNode entry;
    } u;
};
MatrixPointer HdNode[MAX_SIZE];
```

- **Program (2) - Read in a matrix and set up its linked representation**

```
MatrixPointer readM(void){
    /* read in a matrix and set up its linked representation.
    An auxilliary global array HdNode is used */
    int NumRows, NumCols, NumEntries, NumHeads, i;
    int row, col, value, CurrentRow;
    MatrixPointer temp,last,node;

    printf("Enter the number of rows, columns and entries: ");
    scanf("%d,%d,%d",&NumRows, &NumCols, &NumEntries);
    NumHeads = (NumCols > NumRows) ? NumCols : NumRows;
    /* set up head node for the list of head nodes */
    node = (MatrixPointer)malloc(sizeof(struct MatrixNode));
    node->tag = entry;
    node->u.entry.row = NumRows;
    node->u.entry.col = NumCols;
```

- **Program (3) - Read in a matrix and set up its linked representation**

```

if (!NumHeads)    node->right = node;        /* when list of head nodes is empty */
else {            /* initialize the head nodes */
    for (i = 0; i < NumHeads; i++) {
        temp = (MatrixPointer)malloc(sizeof(struct MatrixNode));
        HdNode[i] = temp;
        HdNode[i]->tag = head;
        HdNode[i]->right = temp;
        HdNode[i]->u.next = temp;
    }
    CurrentRow = 0;
    last = HdNode[0];
    for (i = 0; i < NumEntries; i++) {
        printf("Enter row, column and value: ");
        scanf("%d,%d,%d",&row,&col,&value);
        if (row > CurrentRow) {
            last->right = HdNode[CurrentRow];
            CurrentRow = row;
            last = HdNode[row];
        }
    }
}

```

- **Program (4) - Read in a matrix and set up its linked representation**

```

temp = (MatrixPointer)malloc(sizeof(struct MatrixNode));
temp->tag = entry;          temp->u.entry.value = value;
temp->u.entry.row = row;    temp->u.entry.col = col;
last->right = temp;        /* link into row list */
last = temp;
HdNode[col]->u.next->down = temp; /* link into column list */
HdNode[col]->u.next = temp;
} // for
last->right = HdNode[CurrentRow]; /* close last row */
for (i = 0; i < NumCols; i++)      /* close all column lists */
    HdNode[i]->u.next->down = HdNode[i];
for (i = 0; i < NumHeads-1; i++)    /* link all head nodes together */
    HdNode[i]->u.next = HdNode[i+1];
HdNode[NumHeads-1]->u.next = node;
node->right = HdNode[0];
} // if
return node;
}

```


- **Program (5) - Print out the matrix in each row**

```
void writeM(MatrixPointer node){ /* print out the matrix in row major form */
    int i;
    MatrixPointer temp;
    printf("\n\nNumRows = %d, NumCols = %d\n",
        node->u.entry.row, node->u.entry.col);
    printf(" The matrix by row, column, and value: \n\n");
    for (i = 0; i < node->u.entry.row; i++) /* print out the entries in each row */
        for (temp = HdNode[i]->right; temp != HdNode[i]; temp = temp->right)
            printf("%5d%5d%5d\n",temp->u.entry.row,
                temp->u.entry.col,  temp->u.entry.value);
}
```

• Program (6) - Erase the matrix

```

void merase(MatrixPointer *node) {
    heap */
    MatrixPointer x,y;
    int i, NumHeads;
    for (i = 0; i < (*node)->u.entry.row; i++) { /* free the entry pointers by row */
        y = HdNode[i]->right;
        while (y != HdNode[i]) {
            x = y;
            y = y->right;
            free(x);
        }
    }
    /* determine the number of head nodes and free these pointers */
    NumHeads = ((*node)->u.entry.row > (*node)->u.entry.col) ?
        (*node)->u.entry.row : (*node)->u.entry.col;

    for (i = 0; i < NumHeads; i++)
        free(HdNode[i]);
    *node = NULL;
}

```

- **Linked List Basics**
- **Singly linked lists**
- **Circularly linked list**
- **Doubly linked lists**
- **Applications**



Nhân bản – Phụng sự – Khai phóng



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