

## Algorithmic Game Theory: HW 2

- 1.
2. (a) Consider the utility maximizing game below starting with the the initial outcome  $(A_1, B_1)$ , the the game cycles forever which best-response dynamics, avoiding the pure Nash of  $(A_3, B_2)$ .

		P1		
		$A_1$	$A_2$	$A_3$
P2	$B_1$	4, 1	1, 2	0, 0
	$B_2$	0, 0	0, 0	5, 5
	$B_3$	3, 3	3, 2	0, 0

- (b)
- 3.
4. Let  $f_\epsilon(x) = (1 - \epsilon)^x$  and  $g_\epsilon(x) = 1 + \epsilon x$ , then

$$\begin{aligned}
 &f_\epsilon(0) = 1 = g_\epsilon(0) \\
 &f_\epsilon(1) = 1 - \epsilon = g_\epsilon(1) \\
 &\left. \begin{aligned} f'_\epsilon(x) &= (1 - \epsilon)^x \ln(1 - \epsilon) \\ g'_\epsilon(x) &= \epsilon \end{aligned} \right\} f'_\epsilon(0) = \ln(1 - \epsilon) < 0 = g'_\epsilon(0)
 \end{aligned}$$

also  $f_\epsilon$  is a convex function as  $f''_\epsilon(x) = (1 - \epsilon)^x [\ln(1 - \epsilon)]^2 > 0$  for  $\epsilon \in (0, 1/2]$ . This this proves  $f_\epsilon(x) \leq g_\epsilon(x)$  since the initial gradient of  $f_\epsilon$  is smaller then  $g_\epsilon$