

## L12: Images and Signals

Goal: To introduce the idea of image/signal compression.

### **Signal and Image Processing**

Digital images are stores as arrays of "picture elements", or pixels.

- Images are stored as an array of numbers

### **Graylevel ("Intensity") Images**

The stored numbers refer to the shade of gray.

With 256 shades of gray, a pixel can be represented by an 8-bit integer.  
a.k.a. "char" or "uint8"

### **Colour Images**

Consists of 3 intensity images: one for each of red, greed, and blue.  
(why only 3?)

Images (and videos) can take a lot of memory to store.

eg.

### **Data versus Information**

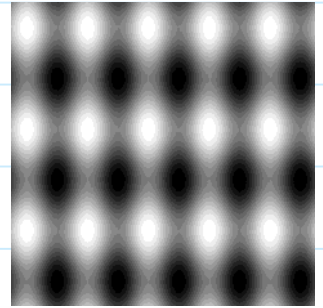
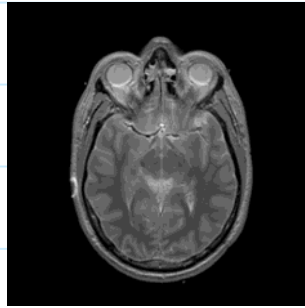
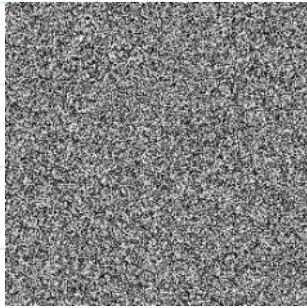
It would be nice to represent the image with less data. In reality, most images hold far less information than the raw memory requirements would suggest. The trick is to find a way to get rid of the unnecessary data.

Eg. Consider the space of all 256x256 pixel 8-bit images. How

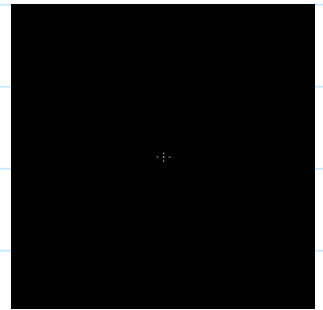
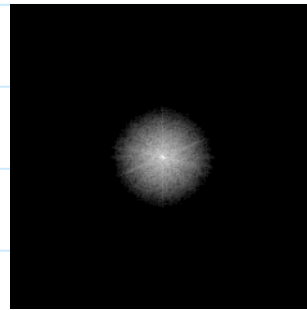
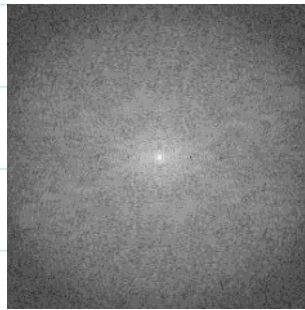
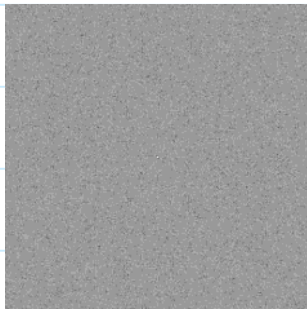
many are there?

Included in that set of images...

Image



Freq. Domain



### JPEG Compression

JPEG = "Joint Photographic Experts Group"

- Standard for storing digital image information
- file extensions .jpg or .jpeg

### MPEG Compression

MPEG = "Moving Picture Experts Group"

- Standard for storing digital videos.
- MP3 = "MPEG Audio Layer 3"

All these technologies use the Fourier transform to sort out what data to store, and what data to discard.

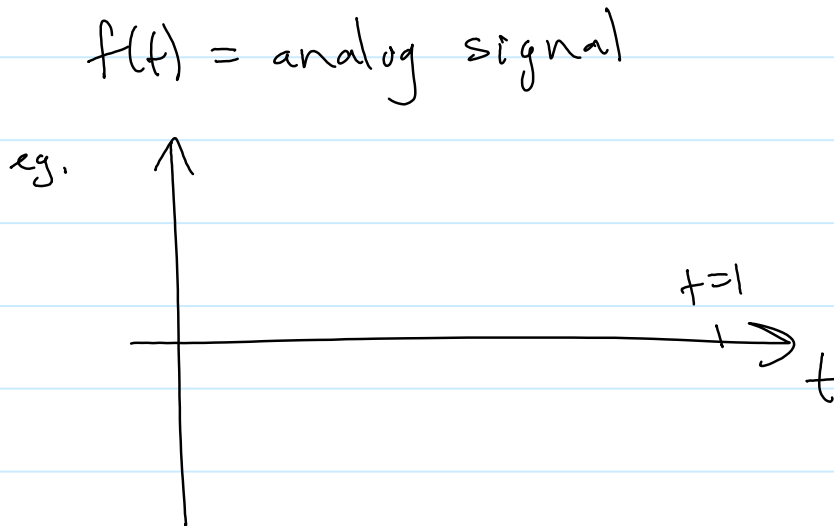
(audio compression demos)

### L13: Fourier Series

Goal: To introduce the use of trigonometric functions to approximate other functions.

To review complex numbers.

Consider  $f(t)$  a function of time,  $t$ .



Both sine and cosine (sin and cos) are periodic functions.

Let  $T = 1$  (second)

If  $f$  is periodic with period  $T$  seconds, then

The frequency is  $\frac{1}{T} = \omega$  Hz. (cycles per second)

eg. (1) the period of  $\sin t$  is  $2\pi$

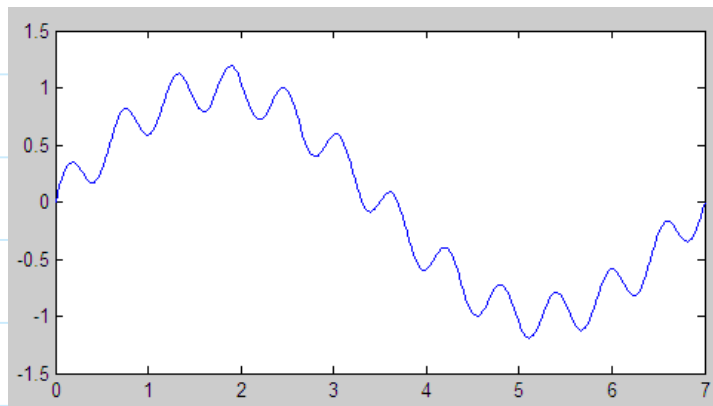
(2) the period of  $\sin\left(\frac{2\pi t}{T}\right)$  is  $T$

(3) the period of

$$\sin\left(\frac{2\pi tk}{T}\right) \text{ is } \frac{T}{k}$$

(4) the period of the sum of 2 sine functions

$$f(t) = \sin\left(\frac{2\pi t}{7}\right) + \frac{1}{5} \sin\left(\frac{2\pi t + 12}{7}\right) \text{ is}$$



Notes: 1)  $\sin\left(\frac{2\pi tk}{T}\right)$  and  $\cos\left(\frac{2\pi tk}{T}\right)$  are  $\frac{T}{k}$ -periodic.

They are also if  $k$  is an integer.

2) In fact, linear combinations of

can be used to represent "all"  $T$ -periodic functions.

Theorem: Suppose  $f$  is a "nice"  $N$ -periodic function.  
There exist coefficients  $a_k$  &  $b_k$  such that

$$f(x) =$$

This is known as a Fourier Series.

In practice, we approximate  $f$  with a truncated Fourier series,

$$f(x) =$$

Instead of treating the  $a$ 's and  $b$ 's separately, we can use the more sophisticated and compact complex notation,

$$f(x) =$$

Notice the sum is now from  $-m$  to  $m$ . Here's why. If  $f(x) \in \mathbb{R}$ , then we need to make sure all the imaginary parts cancel out.

$$f(x) = a_0 + \sum_{k=1}^m \left[ c_k \cos \frac{2\pi k x}{N} + i c_k \sin \frac{2\pi k x}{N} + c_{-k} \cos \frac{-2\pi k x}{N} + i c_{-k} \sin \frac{-2\pi k x}{N} \right]$$

=

Equating coefficients to  $\otimes$

$$\begin{cases} a_k = \\ b_k = \end{cases}$$

$$\textcircled{1} + i \textcircled{2} \Rightarrow a_k + i b_k =$$

$$\textcircled{1} + i\textcircled{2} = a_k + ib_k =$$

$$\textcircled{1} - i\textcircled{2} \Rightarrow a_k - ib_k =$$

Notice, then, that

The Fourier coefficients for a real-valued signal are

(Matlab demo - fourier\_series\_demo.m)

### Complex Numbers (remember those?)

Basic facts:

$$(1) \quad z = \quad \quad \quad a = \text{real part} \quad \quad b = \text{imaginary part}$$

$$(2) \quad \bar{z} = \quad \quad \quad \text{complex conjugate}$$

$$(3) \quad z_1 + z_2 =$$

=

$$(4) \quad z_1 z_2 =$$

=

=

$$(5) |z|^2 = z \bar{z} =$$

$|z|$  = modulus of  $z$

$$(6) z = re^{i\theta} = a + bi$$



### Euler's Identity

$$e^{i\theta} =$$

Conversely

$$\begin{cases} \cos \theta = \frac{e^{i\theta} + e^{-i\theta}}{2} \\ \sin \theta = \frac{e^{i\theta} - e^{-i\theta}}{2i} \end{cases}$$



This polar form is often very convenient. For example, let's multiply two complex numbers.

How about taking the complex conjugate?

⋮  
⋮  
⋮  
⋮  
⋮  
⋮  
⋮





## L14: Trigonometric Orthogonality

Goal: To derive the orthogonality behind what makes Fourier theory so useful.

### Complex Inner Product

Let  $a$  and  $b$  be complex vectors of length  $N$ .

$$\text{Then } \langle a, b \rangle = \sum_{k=1}^N a_k \bar{b}_k$$

$$\text{Thus } \langle a, b \rangle = \overline{\langle b, a \rangle}$$

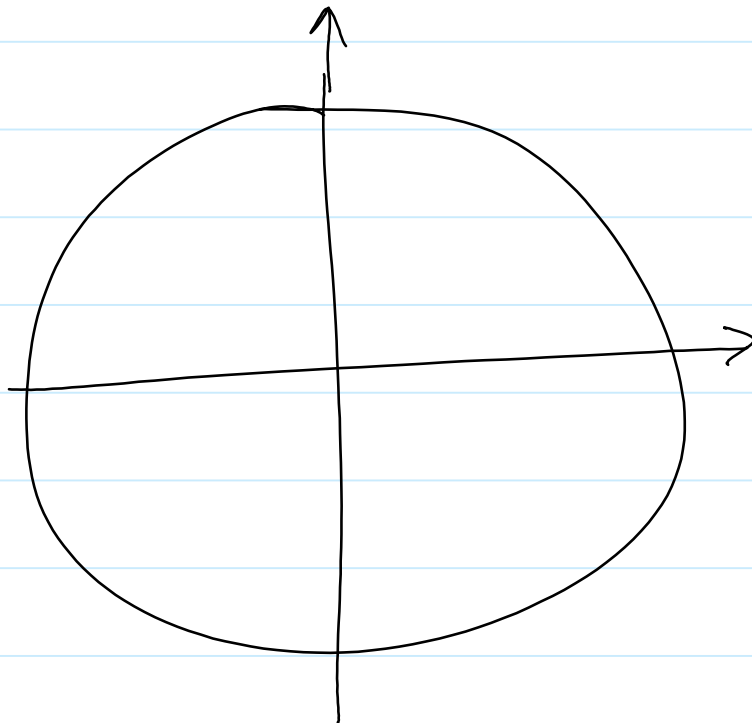
### The $N^{\text{th}}$ Root of Unity

Define the  $N^{\text{th}}$  root of unity,  $W_N$ , as

$$W_N =$$

Then,  $W_N^N$

eg.  $N=8$



Consider the vector defined by

$$W(k) = (1, W_N^k, W_N^{2k}, \dots, W_N^{(N-1)k})$$

Note:  $W_N^k = (e^{\frac{2\pi i}{N}})^k = e^{\frac{2\pi i k}{N}}$

$$W(k)_j = W_N^{jk} = e^{\frac{2\pi i j k}{N}}$$

Orthogonality of  $W(k)$  and  $W(l)$ .

$$\langle W(k), W(l) \rangle = \sum_{j=0}^{N-1} W(k)_j \overline{W(l)_j}$$

=

=

=

Formula:  $\sum_{j=0}^{N-1} r^j = \frac{1-r^N}{1-r}$   $r \neq 1$

In our case,  $r =$  (assume  $k \neq 1$ )

Then

$$\langle w(k), w(l) \rangle =$$

If  $k = l$

$$\langle w(k), w(l) \rangle =$$

$$\therefore \langle w(k), w(l) \rangle =$$

## L15: Discrete Fourier Transform

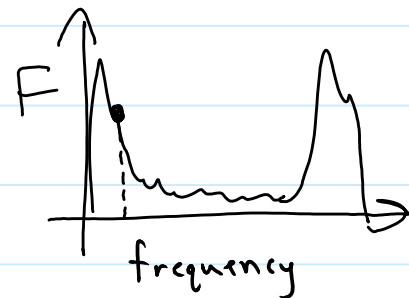
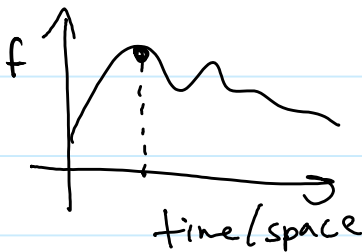
Goal: To derive the Discrete Fourier Transform (DFT) and its inverse.

### Discrete Fourier Transform

Given data samples  $\{f_n \in \mathbb{C} \mid n=0, \dots, N-1\}$ ,  $f_n$  can be written

$$f_n =$$

where  $W_N = e^{\frac{2\pi i}{N}}$  &  $F_k =$



(run fftgui.m)

### Matrix Form of the DFT

$$F_k = \sum_{n=0}^{N-1} f_n \bar{W}^{nk} \quad k=0, \dots, N-1$$

$$F_0 =$$

$$F_1 =$$

$$\vdots$$

$$F_{N-1} =$$

Using matrix-vector notation...

$$\begin{bmatrix} F_0 \\ F_1 \\ F_2 \\ f_3 \\ \vdots \\ F_{N-1} \end{bmatrix} = \begin{bmatrix} & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \end{bmatrix} \begin{bmatrix} f_0 \\ f_1 \\ f_2 \\ f_3 \\ \vdots \\ f_{N-1} \end{bmatrix}$$

Hence, we get  $F = M f$

Recall that

$$\bar{W}(k) = \begin{bmatrix} \bar{W}^{0k} \\ \bar{W}^{1k} \\ \bar{W}^{2k} \\ \vdots \end{bmatrix}$$

Recall that

$$\bar{w}(k) = \begin{bmatrix} \bar{w}^{0k} \\ \bar{w}^{1k} \\ \bar{w}^{2k} \\ \vdots \\ \bar{w}^{(N-1)k} \end{bmatrix}$$

so  $M = \begin{bmatrix} \bar{w}(0) & \bar{w}(1) & \dots & \bar{w}(N-1) \end{bmatrix}$

Notice that  $M = M^T =$

Consider  $M^T M =$

$$= \begin{bmatrix} \phantom{0} \\ \phantom{0} \\ \phantom{0} \\ \phantom{0} \\ \phantom{0} \end{bmatrix}$$

$$=$$

$$\overline{M}^T M =$$

Recall the DFT is  $F = M f$

This is the inverse DFT (IDFT). It takes Fourier coefficients, and converts them to the spatial or temporal domain. So, the IDFT can be written

$$\begin{bmatrix} f_0 \\ f_1 \\ \vdots \\ f_{N-1} \end{bmatrix} = \begin{bmatrix} \phantom{f_0} \\ \phantom{f_1} \\ \phantom{\vdots} \\ \phantom{f_{N-1}} \end{bmatrix} \begin{bmatrix} F_0 \\ F_1 \\ \vdots \\ F_{N-1} \end{bmatrix}$$

In summation notation,

$$f_n =$$



Given the Fourier coefficients,  $\{F_k \in \mathbb{C}, k=0, \dots, N-1\}$ , the IDFT gives the signal composed of the frequencies with amplitudes  $F_k$ .

(Matlab demo - dftMatrix\_demo.m)

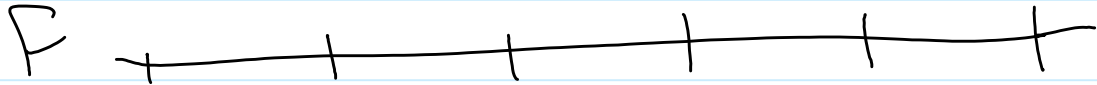
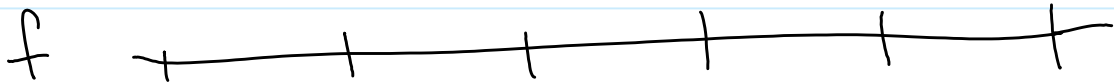


## L16: Properties of the Fourier Transform

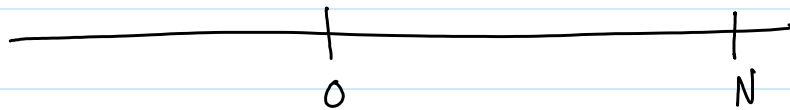
Goal: To see the amazing things that the Fourier Transform can do.

### Periodicity

By the nature of the DFT, both  $f$  and  $F$  are periodic. However, we only store and manipulate one period.



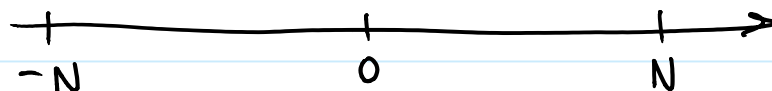
Proof: I will show that  $F_k = F_{N-k}$



$$F_{N-k} =$$

$=$

$=$



When looking at the Fourier coefficients, we can look at  $[0, N-1]$  or we can look at

$$N \text{ odd: } \left[ -\frac{N-1}{2}, \frac{N-1}{2} \right] \text{ eg. } [0 \ 1 \ 2 \ 3 \ 4] \xrightarrow{N=5} [-2 \ -1 \ 0 \ 1 \ 2]$$

$$N \text{ even: } \left[ -\frac{N}{2}, \frac{N}{2} - 1 \right] \text{ eg. } [0 \ 1 \ 2 \ 3 \ 4 \ 5] \rightarrow [-3 \ -2 \ -1 \ 0 \ 1 \ 2]$$

Matlab's commands "fftshift" and "ifftshift" do this transformation for you.

(Matlab demo - sound\_demos.m)

### Conjugate Symmetry

If your signal  $f$  is real-valued,

i.e.  $f_n = \overline{f_n}$ , then...

$$F_{N-k} = \sum_{n=0}^{N-1} f_n \overline{W_N^{n(N-k)}}$$

=

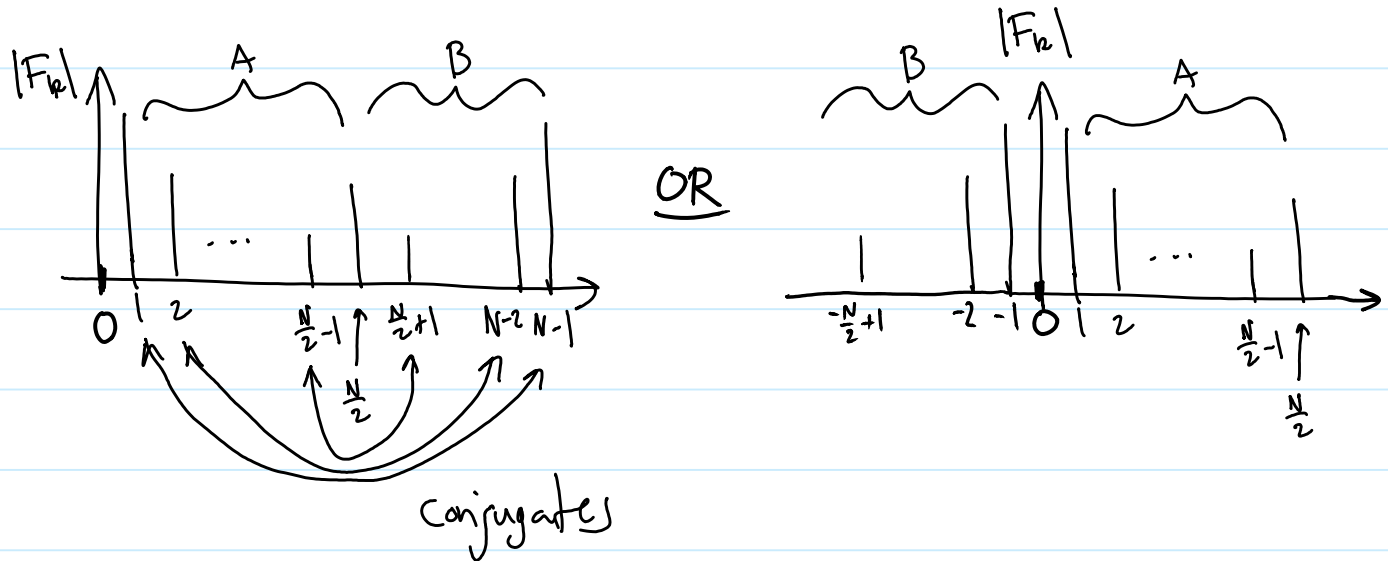
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$$\text{So, } F_{N-k} =$$

Consequently  $|F_{N-k}| = |F_k| = |F_k|$



## 2-Dimensional DFT

A graylevel image is represented by a 2-D array

$$f_{nj} \quad n=0, 1, \dots, N-1$$

$$j=0, 1, \dots, M-1$$

The 2-D DFT is defined as

$$\text{2D-DFT: } F_{kl} =$$

$$\text{2D-IDFT: } f_{nj} =$$

## How to Compute a 2-D DFT Using 1-D DFT

$$F_{kl} =$$

=

$H_{nl}$  = DFT of  $n^{\text{th}}$  row of  $f$ ,  $n=0,1,\dots,N-1$

$\Rightarrow N$  1D-DFT's

After that,

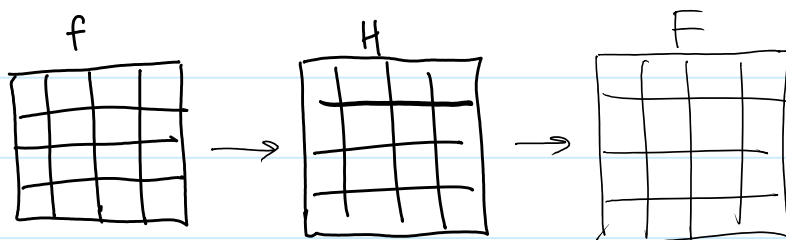
$$F_{kl} = \sum_{n=0}^{N-1} H_{nl} \overline{W_N^{nk}}$$

$l$  is fixed

$k=0,1,\dots,M-1$

This is the DFT of the  $l^{\text{th}}$  column of  $H$

$\Rightarrow M$  1D-DFT's



In Matlab, use  
fft2 & ifft2.

## L17: More Properties of the Fourier Transform

Goal: To see the amazing things that the Fourier Transform can do.

### Convolution

A convolution is a sum (integral) of the form

$$(f * g)_a =$$

Just for fun, let's take the DFT of  $(f * g)_a$ .

$$\text{DFT}(f * g)_k = \sum_{a=0}^{N-1} (f * g)_a \bar{W}^{ak}$$

=

=

Change of variables: Let  $b = a - n \Rightarrow a = b + n$

=

=

=

Thus, we can evaluate a convolution using the Fourier transform.

Since,  $\text{DFT}(f * g)_k =$

Example 1: One thing that convolution is used for is blurring.

$$f = \text{spikes}$$

$$g = \text{Gaussian function}$$
$$g(x) = \frac{1}{\sigma\sqrt{2\pi}} e^{-\frac{x^2}{2\sigma^2}}$$

$$f * g = \text{blurred spikes}$$

Example 2: It can also be used to compute the correlation coefficient.

$$CC(a; f, g) = \frac{\sum_{n=0}^{N-1} f_n g_{n-a}}{\|f\|^2 \|g\|^2} =$$

If  $CC = 0$

If  $CC = \pm 1$

## L18: Fast Fourier Transform

Goal: To learn how the FFT algorithm work to compute the DFT so efficiently.

### Fast Fourier Transform

The Fast Fourier Transform is one of the most important computational algorithms today. Here's how it works...

### Divide and Conquer

$$\begin{aligned} F_k &= \sum_{n=0}^{N-1} f_n \bar{W}_N^{nk} \\ &= \sum_{n=0}^{N/2-1} f_n \bar{W}_N^{nk} + \sum_{n=N/2}^{N-1} f_n \bar{W}_N^{nk} \end{aligned}$$

Change of index:  $m =$

$$\begin{aligned} (*) &= \sum_{m=0}^{N-1} \\ &= \sum_{n=0}^{N/2-1} \\ F_k &= \sum_{n=0}^{N/2-1} \end{aligned}$$

$$\bar{W}_N^{Nk/2} =$$

For even indices (of the form  $2l$ )

For even indices (of the form  $2k$ )

$$F_{2k} = \sum_{n=0}^{\frac{N}{2}-1}$$

$$k = 0, 1, \dots, \frac{N}{2} - 1$$

For odd indices (of the form  $2k+1$ )

$$F_{2k+1} = \sum_{n=0}^{\frac{N}{2}-1}$$

$$= \sum_{n=0}^{\frac{N}{2}-1}$$

$$k = 0, \dots, \frac{N}{2} - 1$$

Let  $g_n =$

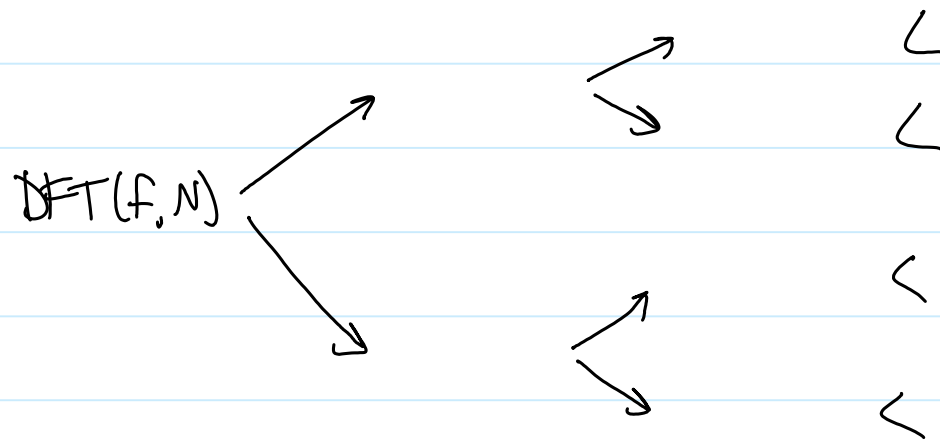
$$h_n =$$

Then

$$F_{2k} =$$

$$F_{2k+1} =$$





Eventually, you get to a DFT of a vector of length 1... which is

### Recursive FFT Algorithm

function  $\{F_n\} = \text{FFT}(\{f_n\}, N)$

if  $N=1$

else

.

.

.

.

.

.

end

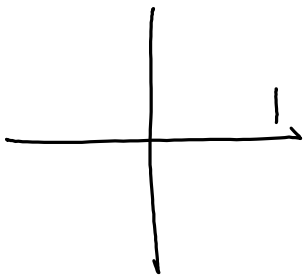
### Complexity

DFT by direct computation is

Using FFT algorithm:

•  
•  
•  
•  
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•  
•  
•

# FFT Butterfly Diagram



Decomposed down to  
1x1 vectors.

