

1. GCD Characterization
2. Extended Euclidean Algebra

$$d = \gcd(a, b)$$

$$\textcircled{1} d \mid a$$

$$\textcircled{2} d \mid b$$

$\textcircled{3}$ If $c \mid a$ and $c \mid b$, then $c \leq d$.

Proposition (GCD Characterization): For $a, b \in \mathbb{Z}$ not both 0, if d is a positive divisor of a, b and $ax + by = d$ has an integer for (x, y) , then $d = \gcd(a, b)$.

ex. $a = 12$ $b = 15$

$12x + 15y = \square$ If this has int solution, is \square the gcd? Only if $\square \mid 12$ and $\square \mid 15$.

$12x + 15y = 12$ $(1, 0)$ is a solⁿ, but $12 \neq \gcd$

$3 \mid 12, 3 \mid 15$, $(-1, 1)$ is an int solution

$$\gcd(12, 15) = 3.$$

$12x + 15y = 1$ $1 \neq \gcd$, conc is false \Rightarrow hyp. is false.

$1 \nmid 12, 1 \nmid 15$, so no solⁿ.

Proof: By assumption, $d \mid a$ and $d \mid b$. Let $c \in \mathbb{Z}$ such that $c \mid a$ and $c \mid b$. Suppose (x_0, y_0) is an int solution to $ax + by = d$.

By div. of int. comb, $c/(ax_0+by_0)$, so c/d . using bounds by div, $|c| \leq |d| = d$, since $d > 0$, so $c \leq d$, and $d = \gcd(a, b)$

The integers (x, y) serves as a certificate that d is the gcd of a, b . To verify $d = \gcd(a, b)$:

$$d|a \checkmark \quad d|b \checkmark \quad ax+by=d \checkmark$$

Extended Euclidean Algorithm (EEA)

Find x, y such that $ax+by = \gcd(a, b)$

Example: $\gcd(744, 264)$

$$= \gcd(264, 216)$$

$$= \gcd(216, 48)$$

$$= \gcd(48, 24)$$

$$= 24$$

$$744 = 2 \cdot 264 + 216 \quad (1)$$

$$264 = 1 \cdot 216 + 48 \quad (2)$$

$$216 = 4 \cdot 48 + 24 \quad (3)$$

Back substitution...

$$24 = 216 - 4 \cdot 48 \quad (3)$$

$$= 216 - 4(264 - 1 \cdot 216)$$

$$= 5 \cdot 216 - 4 \cdot 264 \quad (2)$$

$$= 5(744 - 2 \cdot 264) - 4 \cdot 264$$

$$= 5 \cdot 744 - 14 \cdot 264 \quad (1)$$

Proposition (EEA): Let $a, b \in \mathbb{Z}$. If $d = \gcd(a, b)$, then $ax + by = d$ has an int solution.

(Only prove for $a, b \in \mathbb{N}$)

Let $E(a, b)$ be the # of steps in the EA when finding $\gcd(a, b)$.

Proof: Induction on $E(a, b)$.

If $a = b$, then $\gcd(a, b) = a$. So $a \cdot 1 + b \cdot 0 = a$.

Without loss of generality, assume $a > b$.

Base Case: If $E(a, b) = 1$, then $b|a$, and $\gcd(a, b) = b$.

So $a \cdot 0 + b \cdot 1 = b$.

Ind. Hyp.: Assume result holds when $E(a, b) = k$ for some $k \in \mathbb{N}$.

Ind. Step: Suppose $E(a, b) = k + 1$

In the first step of E-A, we find $a = qb + r$, and $\gcd(a, b) = \gcd(b, r)$. We need k steps to find $\gcd(b, r)$, so $E(b, r) = k$. By ind. hyp, there exists $x_0, y_0 \in \mathbb{Z}$ such that $bx_0 + ry_0 = \gcd(b, r) = \gcd(a, b) = d$.

Replace $r = a - qb$, $d = bx_0 + (a - qb)y_0 = ay_0 + b(x_0 - qy_0)$.
So $ax + by = d$ has an int solution. $(y_0, x_0 - qy_0)$