

# Machine Learning Homework: Week 3 & 4

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## 1

### Statement

Prove that

$$\begin{aligned} \frac{1}{2} \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \mathbb{E}_z \phi(z) - \frac{1}{n} \sum_{i=1}^n \phi(z_i) \right| \geq 2\varepsilon \right) &\leq \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) \right| \geq \varepsilon \right) \\ &\leq 2 \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \mathbb{E}_z \phi(z) - \frac{1}{n} \sum_{i=1}^n \phi(z_i) \right| \geq \frac{\varepsilon}{2} \right) \end{aligned} \quad (1)$$

### Solution

$$\begin{aligned} &\mathbb{P} \left( \sup_{\phi \in \Phi} \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) \right| \geq \varepsilon \right) \\ &= \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) \right| > \varepsilon - \delta \right) \\ &\leq \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| > \frac{\varepsilon - \delta}{2} \text{ or } \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| > \frac{\varepsilon - \delta}{2} \right) \\ &\leq \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| > \frac{\varepsilon - \delta}{2} \right) + \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| > \frac{\varepsilon - \delta}{2} \right) \\ &= \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| \geq \frac{\varepsilon}{2} \right) + \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| \geq \frac{\varepsilon}{2} \right) \\ &= 2 \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \mathbb{E}_z \phi(z) - \frac{1}{n} \sum_{i=1}^n \phi(z_i) \right| \geq \frac{\varepsilon}{2} \right) \end{aligned} \quad (2)$$

which provides the proof for the inequality on the right side.

$$\begin{aligned}
& \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) \right| \geq \varepsilon \right) \\
&= \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) \right| > \varepsilon - \delta \right) \\
&\geq \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| > 2(\varepsilon - \delta) \text{ and } \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| < \varepsilon - \delta \right) \\
&\geq \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| > 2(\varepsilon - \delta) \right) \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| < \varepsilon - \delta \right)
\end{aligned} \tag{3}$$

By Chernoff bound we obtain that  $\mathbb{P} \left( \left| \frac{1}{n} \sum_{i=n+1}^{2n} \phi(z_i) - \mathbb{E}_z \phi(z) \right| < \varepsilon \right) \geq \frac{1}{2}$  when  $n \geq \frac{\ln 2}{\varepsilon^2}$ , which gives the proof for the inequality on the left side

$$\begin{aligned}
& \dots \geq \frac{1}{2} \mathbb{P} \left( \forall \delta > 0, \exists \phi \in \Phi \text{ s.t. } \left| \frac{1}{n} \sum_{i=1}^n \phi(z_i) - \mathbb{E}_z \phi(z) \right| > 2(\varepsilon - \delta) \right) \\
&= \frac{1}{2} \mathbb{P} \left( \sup_{\phi \in \Phi} \left| \mathbb{E}_z \phi(z) - \frac{1}{n} \sum_{i=1}^n \phi(z_i) \right| \geq 2\varepsilon \right)
\end{aligned} \tag{4}$$

## 2

### Statement

Prove that

$$\sum_{k=0}^d \binom{n}{k} < \left( \frac{en}{d} \right)^d$$

for any  $n \in \mathbb{N}^+$  and  $1 \leq d \leq n$ .

### Solution

By Taylor expansion we have

$$e^d = \sum_{k=0}^{\infty} \frac{d^k}{k!} > \sum_{k=0}^d \frac{d^k}{k!}$$

where immediately follows the proof

$$\begin{aligned}
\left( \frac{en}{d} \right)^d &> n^d \sum_{k=0}^d \frac{d^{k-d}}{k!} \\
&\geq n^d \sum_{k=0}^d \frac{n^{k-d}}{k!} \\
&= \sum_{k=0}^d \frac{n^k}{k!} \\
&\geq \sum_{k=0}^d \frac{n^{\underline{k}}}{k!} \\
&= \sum_{k=0}^d \binom{n}{k}
\end{aligned} \tag{5}$$

### 3

#### Statement

Prove that the following two programming problems are equivalent. For simplicity, we assume that the first problem has non-trivial solution, i.e. there exists  $t > 0$  satisfying  $y_i(\mathbf{w}^T \mathbf{x}_i + b) \geq t$  ( $\forall i \in [n]$ ).

$$\begin{array}{ll} \max_{\mathbf{w}, b, t} & t \\ \text{s.t.} & y_i(\mathbf{w}^T \mathbf{x}_i + b) \geq t \quad (\forall i \in [n]) \\ & \|\mathbf{w}\| = 1 \end{array} \qquad \begin{array}{ll} \min_{\mathbf{w}, b, t} & \frac{1}{2} \|\mathbf{w}\| \\ \text{s.t.} & y_i(\mathbf{w}^T \mathbf{x}_i + b) \geq 1 \quad (\forall i \in [n]) \end{array}$$

#### Solution

Let the solution to the two problems above be  $t_0$  and  $\mathbf{w}_0$  respectively. By the assumption, we have  $t_0 > 0$ .

By dividing the inequality in the first problem by  $t_0$  we obtain  $y_i \left( \frac{\mathbf{w}_0^T}{t_0} \mathbf{x}_i + \frac{b}{t_0} \right) \geq 1$ , so  $\mathbf{w}' = \frac{\mathbf{w}_0}{t_0}$  is a candidate for the second problem, and thus  $\|\mathbf{w}_0\| \leq \|\mathbf{w}'\| = \frac{1}{t_0}$ .

By dividing the inequality in the second problem by  $\|\mathbf{w}_0\|$  we obtain  $y_i \left( \frac{\mathbf{w}_0}{\|\mathbf{w}_0\|} \mathbf{x}_i + \frac{b}{\|\mathbf{w}_0\|} \right) \geq \frac{1}{\|\mathbf{w}_0\|}$ , thus  $t \geq \frac{1}{\|\mathbf{w}_0\|}$  since  $\left\| \frac{\mathbf{w}_0}{\|\mathbf{w}_0\|} \right\| = 1$ .

So we know that  $\|\mathbf{w}_0\|t_0 = 1$ , which shows the equivalence relationship between these two programming problems.