Robust Monopoly Pricing

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Introduction

- The optimal mechanisms under the theoretical literature have been widely used with the development of electronic marketplace and Internet trading.
- 2 How robust of such mechanism with respect to the model uncertainty?
- In the robust version, the seller only knows that the true distribution is in the neighborhood of a given model distribution.
- The optimal pricing policy of the seller in the presence of model uncertainty is an instance of decision-making with multiple priors.



Model

- The seller faces a single buyer with valuation $v \in [0,1]$ for an unit of good.
- The net utility of the buyer is v p.
- ullet The profit for the seller given price p and buyer's valuation v is

$$\pi(p,v) \triangleq p\mathbb{I}_{\{v \geqslant p\}}$$

• If the seller chooses a randomized policy $\Phi \in \Delta \mathbb{R}_+$, the buyer's utility given ν is

$$\pi(\Phi, v) \triangleq \int \pi(p, v) d\Phi(p)$$
.

ullet The seller's expected profit given price p and demand distribution F is

$$\pi(p,F) \triangleq \int \pi(p,v)dF(v).$$

• With randomizing pricing, the seller's expected profit is

$$\pi(\Phi, F) \triangleq \iint \pi(p, v) dF(v) d\Phi(p).$$



Model

The optimal random pricing policy given the demand distribution F solves

$$\Phi^*(F) \in \operatorname*{arg\,max}_{\Phi \in \Delta \mathbb{R}_+} (\Phi, F).$$

- A well-known result states that for every distribution F, there exists a deterministic price $p^*(F)$ that maximizes profits.
- The uncertainty is represented by a set of possible distributions. Given the model distribution F_0 , the ε neighborhood under the Prohorov metric, denoted by $\mathcal{P}_{\varepsilon}(F_0)$, is:

$$\mathcal{P}_{\varepsilon}\left(F_{0}\right)\triangleq\left\{ F\mid F(A)\leqslant F_{0}\left(A^{\varepsilon}\right)+arepsilon, \forall \text{ measurable } A\subseteq\left[0,1\right]
ight\} .$$

where the set A^{ε} denotes the closed ε neighborhood of any measurable set A:

$$A^{\varepsilon} \triangleq \Big\{ x \in [0,1] \mid \underbrace{\min_{y \in A} d(x,y)}_{d(x,A)} \leqslant \varepsilon \Big\}.$$

Model

• Example: if $A = [0, a] \subset [0, 1]$ and d(x, y) = |x - y|, so that

$$A^{\varepsilon} = \left\{ x \in [0,1] \mid d(x,A) \leqslant \varepsilon \right\} = [0,a+\varepsilon].$$

 The Prohorov metric is a metric for weak convergence of probability measures.

Definition (Topology Definition)

Let (M,d) be a metric space with its Borel sigma algebra $\mathcal{B}(M)$. Let $\mathcal{P}(M)$ denote the collection of all probability measures on the measurable space $(M,\mathcal{B}(M))$. For a subset $A\subseteq M$, define the ε -neighborhood of A by

$$A^{\varepsilon} := \{ p \in M \mid \exists q \in A, d(p,q) < \varepsilon \} = \bigcup_{p \in A} B_{\varepsilon}(p)$$

where $B_{\varepsilon}(p)$ is the open ball of radius ε centered at p. The Lévy-Prokhorov metric $\pi: \mathcal{P}(M)^2 \to [0, +\infty)$ is defined by setting the distance between two probability measures μ and ν to be $\pi(\mu, \nu) := \inf \left\{ \varepsilon > 0 \mid \mu(A) \leq \nu\left(A^{\varepsilon}\right) + \varepsilon \right\}$ and $\nu(A) \leq \mu\left(A^{\varepsilon}\right) + \varepsilon$ for all $A \in \mathcal{B}(M)$

Under maximin utility, the seller maximizes the minimum utility

$$\Phi_m \in \underset{\Phi \in \Delta \mathbb{R}_+}{\operatorname{arg \, max}} \min_{F \in \mathcal{P}_{\varepsilon}(F_0)} \pi(\Phi, F).$$

• Accordingly, we say that Φ_m attains maximin utility. We refer to F_m as a least favorable demand (for maximin utility) if

$$F_m \in \underset{F \in \mathcal{P}_{\varepsilon}(F_0)}{\operatorname{arg \, min}} \max_{\Phi \in \Delta \mathbb{R}_+} \pi(\Phi, F).$$

• Occasionally, it is useful to explicitly state the dependence of the optimal policies on the size of the neighborhood ϵ , by $\Phi_{m,\varepsilon}$ and $F_{m,\varepsilon}$.

• The regret of the monopolist at a given price p and valuation v is:

$$r(p, v) \triangleq v - p \mathbb{I}_{\{v \geqslant p\}} = v - \pi(p, v)$$

where v is the maximum profit the monopolist could make if she were to know the value v before setting her price and the profit she makes without this information.

- The regret is non-negative and can only vanish when p = v. Otherwise, if v > p, then r(p, v) = v p * 1 > 0, or v < p, then r(p, v) = v p * 0 > 0.
- The expected regret of a random pricing policy Φ given a demand distribution F is:

$$r(\Phi, F) \triangleq \iint r(p, v) d\Phi(p) dF(v) = \int v dF(v) - \int \pi(p, F) d\Phi(p).$$

• The pricing policy Φ_r attains minimax regret if it minimizes the maximum regret over all distributions F in the neighborhood of a model distribution F_0 :

$$\Phi_r \in \underset{\Phi \in \Delta\mathbb{R}_+}{\text{arg min}} \max_{F \in \mathcal{P}_{\varepsilon}(F_0)} r(\Phi, F).$$

 \bullet F_r is called a least favorable demand if

$$F_r \in \mathop{\arg\min}_{F \in \mathcal{P}_\varepsilon(F_0)} \mathop{\max}_{\Phi \in \Delta \mathbb{R}_+} r(\Phi, F) = \mathop{\arg\max}_{F \in \mathcal{P}_\varepsilon(F_0)} \left(\int v dF(v) - \mathop{\max}_{\Phi} \pi(\Phi, F) \right).$$

 In particular, the decision maker does not need the information to become available ex post to evaluate his expected regret.

Robust Policy

Definition (Robust Pricing Policy)

A family of pricing policies $\{\Phi_{\varepsilon}\}_{{\varepsilon}>0}$ is called robust if, for each $\gamma>0$, there is ${\varepsilon}>0$ such that $F\in \mathcal{P}_{\varepsilon}(F_0)\Rightarrow \pi\left(\Phi^*(F),F\right)-\pi\left(\Phi_{\varepsilon},F\right)<\gamma$.

- We can check $\Phi^*(F_0)$ is not a robust pricing policy.
- Consider a Dirac distribution which puts probability one on valuation v, i.e., $F = \delta_{\{v\}}$.
- The optimal monopoly price $\Phi^*(F) = v$. If the $F_0 = \delta_{\{v \epsilon\}}$, then the revenue is 0. The difference is $v \to 0$.

- The pricing rule that attains maximin utility is the equilibrium strategy in a game between the seller and adversarial nature.
- The seller chooses a probabilistic pricing policy, a distribution $\Phi \in \Delta \mathbb{R}_+$, and nature chooses a demand distribution $F \in \mathcal{P}_{\varepsilon}(F_0)$.
- In this game, the payoff of the seller is the expected profit while the payoff of nature is the negative of the expected profit.
- A Nash equilibrium of this zero-sum game is a solution (Φ_m, F_m) to the saddle point problem:

$$\underbrace{\pi\left(\Phi,F_{m}\right)\leqslant}_{\text{seller maximizes profit}}\pi\left(\Phi_{m},F_{m}\right)\underbrace{\leqslant\pi\left(\Phi_{m},F\right)}_{\text{nature minimizes profit}},$$

for all $\Phi \in \Delta \mathbb{R}_+, \forall F \in \mathcal{P}_{\varepsilon}(F_0)$.

• For a given price p, the expected profit of the seller is

$$\pi(p,F)=\int \pi(p,v)dF(v)=p(1-F(p)).$$

- The profit minimization demand should minimize (1 F(p)) for all p within the neighborhood $\mathcal{P}_{\varepsilon}(F_0)$.
- This defines the order: first-order-stochastic dominance relationship. The profit minimizing demand is the **smallest** elements of $\mathcal{P}_{\varepsilon}\left(F_{0}\right)$ in FOSD sense
- This results in the distribution $F_{m,\varepsilon}$ given by:

$$F_{m,\varepsilon}(v) \triangleq \min \{F_0(v+\varepsilon) + \varepsilon, 1\}$$

• Recall the definition of Prohorov metric, we have

$$F([0,a]) \le F_0([0,a+\epsilon]) + \epsilon$$



- Given that the profit minimizing demand $F_{m,\varepsilon}$ does not depend on the offered prices, the monopolist acts as if the demand is given by $F_{m,\varepsilon}$.
- In consequence, the seller maximizes profits at $F_{m,\varepsilon}$ by choosing a **deterministic** price $p_{m,\varepsilon}$ where $p_{m,\varepsilon} \triangleq p^*(F_{m,\varepsilon})$.

Proposition (1. Maximin Utility)

For every $\varepsilon > 0$, there exists a pair $(p_{m,\varepsilon}, F_{m,\varepsilon})$, such that $p_{m,\varepsilon} \in [0,1]$ attains maximin utility and $F_{m,\varepsilon}$ is a least favorable demand.

Proposition (2. Pricing under Maximin Utility)

The price $p_{m,\varepsilon}$ responds to an increase in uncertainty at $\varepsilon=0$ by:

$$\left. \frac{dp_{m,\varepsilon}}{d\varepsilon} \right|_{\varepsilon=0} = -1 + \frac{1 - f_0\left(p_0\right)}{\partial \pi^2\left(p_0, F_0\right)/\partial p^2} < -\frac{1}{2}.$$

Proposition (3. Robustness)

The family of pricing policies $\{p_{m,\varepsilon}\}_{\varepsilon>0}$ is a robust family of pricing policies.

- We consider the minimax regret problem of the seller, where a (probabilistic) pricing policy Φ_r and a least favorable demand F_r are the equilibrium policies of a zero-sum game.
- In this zero-sum game, the payoff of the seller is the negative of the regret while the payoff to nature is regret itself.
- A Nash equilibrium (Φ_r, F_r) is a solution to the saddle point problem:

$$\underbrace{r\left(\Phi_{r},F\right)\leqslant}_{\text{nature maximizes regret}}\left(\Phi_{r},F_{r}\right)\underbrace{\leqslant r\left(\Phi,F_{r}\right)}_{\text{seller minimizes regret}},\quad\forall\Phi\in\Delta\mathbb{R}_{+},\forall F\in\mathcal{P}_{\varepsilon}\left(F_{0}\right).$$

- The deterministic policy can not resolve the follow confliction:
 - ▶ If setting a low price, nature can cause regret with a distribution which puts substantial probability on high valuation buyers.
 - ▶ If setting a high price, nature can cause regret with a distribution which puts substantial probability at valuations just below the offered price.
- To get more intuitions, we can consider the seller's expected regret given a single price p:

$$r(p,F) = \int v dF(v) - \int \pi(p,v) dF(v) = \underbrace{\int v dF(v)}_{\text{mean valuation}} - \underbrace{p(1-F(p))}_{\text{upper cumulative prob}}$$

• For nature can maximize the regret by increasing mean valuation or lower upper cumulative prob (choose a smaller F in FOSD). However, two actions are conflict!

Proposition (4. Existence of Minimax Regret)

A solution (Φ_r, F_r) to the saddlepoint condition exists.

Corollary

Suppose that $G=(X_i,u_i)_{i=1}^N$ is a compact, Hausdorff game. Then G possesses a mixed strategy Nash equilibrium if its mixed extension, \bar{G} , is better-reply secure. Moreover, \bar{G} is better-reply secure if it is both reciprocally upper semicontinuous and payoff secure.

Definition

A game $G = (X_i, u_i)_{i=1}^N$ is payoff secure if for every $x \in X$ and every $\epsilon > 0$, each player i can secure a payoff of $u_i(x) - \epsilon$ at x.

• For the monopolist: payoff secure requires that for each (F_r, Φ_r) with $F_r \in \mathcal{P}_{\epsilon}(F_0)$ and for every $\delta > 0$, there exists $\gamma > 0$ and $\bar{\Phi}$ such that $F \in \mathcal{P}_{\gamma}(F_r)$ implies

$$r(\bar{\Phi}, F) \leq r(\Phi_r, F_r) + \delta.$$

• For the seller: we need to show that for each (Φ_r, F_r) , with $F_r \in \mathcal{P}_{\epsilon}(F_r)$ and every $\delta > 0$, there exists $\gamma > 0$, and $\bar{F} \in \mathcal{P}_{\epsilon}(F_0)$, such that $\Phi \in \mathcal{P}_{\gamma}(\Phi_r)$ implies

$$r(\Phi, \bar{F}) \geq r(\Phi_r, F_r) - \delta.$$

Proposition (5. Minimax regret)

1 Given $\delta > 0$, for every ε sufficiently small, there exist a, b and c with 0 < a < b < c < 1 and $p_0 - \delta < a < p_0 < c < p_0 + \delta$ such that a minimax regret probabilistic price Φ_r is given by:

$$\Phi_r(p) = \begin{cases} 0 & \text{if } 0 \leqslant p < a \\ \ln \frac{p}{a} & \text{if } a \leqslant p < b \\ 1 - \ln \frac{c}{p} & \text{if } b \leqslant p \leqslant c \\ 1 & \text{if } c < p \leqslant 1 \end{cases}.$$

- ② The boundary points a, b and c respond to an increase in uncertainty at $\varepsilon = 0$:

 - $\lim_{\varepsilon \to 0} c'(0) = \infty;$



Corollary (1. Comparative Statics with Minimax Regret)

• The expected price $\mathbb{E}\left[p_{r,\varepsilon}\right]$ responds to an increase in uncertainty at $\varepsilon=0$ by:

$$\left. \frac{d}{d\varepsilon} \mathbb{E} \left[p_{r,\varepsilon} \right] \right|_{\varepsilon=0} = \begin{cases} -1 - \frac{f_0(p_0) + 1}{\partial \pi^2(p_0, F_0) / \partial p^2} > -1 & \text{if } p_0 \leqslant \frac{1}{2} \\ -1 - \frac{f_0(p_0) - 1}{\partial \pi^2(p_0, F_0) / \partial p^2} < -\frac{1}{2} & \text{if } p_0 > \frac{1}{2} \end{cases}$$

2 If ε is sufficiently small, then for any $v \in (a, c) \setminus b$,

$$\frac{d}{d\varepsilon}\mathbb{E}\left[p_{r,\varepsilon}\mid p_{r,\varepsilon}\leqslant v\right]<0.$$

Proposition (5. Robustness)

If $\{\Phi_{r,\varepsilon}\}_{\varepsilon>0}$ attains minimax regret at F_0 for all sufficiently small ε , then $\{\Phi_{r,\varepsilon}\}_{\varepsilon>0}$ is a robust family of pricing policies.

Discussion

- Axioms and behavioral implications:
 - ► Two criteria: the maximin utility and minimax regret criteria
 - ► The similarity and difference between maximin utility and minimax regret.
- The choice of metric (and neighborhoods):
 - different metrics: such as the Levy metric or the bounded Lipschitz metric which also metrize the weak topology.
 - the comparative static results near $\epsilon=0$ are unaffected by the specific notion for the metric.
- Beyond small neighborhoods: how about large neighbor?
- Beyond monopoly pricing: general mechanism design problem.

Thanks!