CS 461 Artificial Intelligence

Hill-Climbing Search

With randomly generated 8-queens starting states, the steepest-ascent hill climbing:

- ▶ 14% of the time it solves the problem
- 86% of the time it get stuck at a local minimum
- However...
 - Takes only 4 steps on average when it succeeds
 - And 3 on average when it gets stuck
 - (for a state space with $8^8 \approx 17$ million states)

Hill-Climbing: Variants

Stochastic hill-climbing:

- Chooses randomly among potential successors
- Sometimes better than steepest ascent

First-choice hill-climbing:

- Generates successors randomly and picks first
- Good for many successors

Random restart hill-climbing:

- Restarts from <u>randomly generated initial state</u> when failed
- Roughly 7 iterations with 8-queens problem

Hill-Climbing Search

- The success of hill climbing depends very much
 - on the shape of the state-space landscape
- If there are few local maxima and plateaux,
 - random-restart hill climbing will find a good solution very quickly.

Idea:

- escape local maxima by allowing some "bad" moves but gradually decrease the size and frequency of the bad moves,
- In thermodynamics, the probability to go from a state with energy E_1 to a state of energy E_2 is given by:

$$p = e^{\frac{(E_2 - E_1)}{kT}} = e^{\frac{-(E_1 - E_2)}{kT}}$$

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where

- e is Euler's number
- T is a "temperature" controlling the probability of downward steps
- k is Boltzmann's constant
 - (relating energy and temperature; with appropriate choice of units it will be equal to 1).

$$p = e^{\frac{(E_2 - E_1)}{kT}} = e^{\frac{-(E_1 - E_2)}{kT}}$$

- ▶ The idea is that probability decreases exponentially with $E_2 E_1$ increasing,
- The probability gets lower as temperature decreases
- ▶ If the *schedule* lowers *T* slowly enough, the algorithm will find a global optimum with probability approaching 1.

```
function SIMULATED-ANNEALING (problem, schedule) returns a solution state
   inputs: problem, a problem
             schedule, a mapping from time to "temperature"
   local variables: current, a node
                       next, a node
                        T, a "temperature" controlling prob. of downward steps
   current \leftarrow \text{Make-Node}(\text{Initial-State}[problem]) Similar to hill climbing,
   for t \leftarrow 1 to \infty do
                                                              but a random move instead
        T \leftarrow schedule[t]
                                                              of best move
        if T = 0 then return current
        next \leftarrow a randomly selected successor of current
                                                              case of improvement, make
        \Delta E \leftarrow \text{Value}[next] - \text{Value}[current]
                                                              the move
        if \Delta E > 0 then current \leftarrow next
        else current \leftarrow \underline{next} only with probability e^{\Delta E/T}
                             Otherwise, choose the move with probability that
```

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decreases exponentially with the "badness" of the move.

Simulated Annealing...Example

 Consider there are <u>3 moves</u> available, with changes in the objective function of

$$\Delta E_1 = -0.1, \Delta E_2 = 0.5, \Delta E_3 = -5$$

- ightharpoonup Suppose T = 1
- Pick a move randomly:
 - if ΔE_2 is picked, move there.
 - if ΔE_1 or ΔE_3 are picked, probability of move $=e^{\frac{\Delta E}{T}}$
 - move 1: prob1 = $e^{-0.1}$ = 0.9,
 - i.e., 90% of the time we will accept this move
 - move 3: prob3 = e^{-5} = 0.0067
 - i.e., 5% of the time we will accept this move

T = "temperature" parameter

- If T is high => the probability of "locally bad" move is higher
- If T is low => the probability of "locally bad" move is lower
- ightharpoonup typically, T is decreased as the algorithm runs longer
 - i.e., there is a "temperature schedule"

Convergence:

- With <u>exponential schedule</u>, will provably converge to global optimum
 - If <u>T decreases slowly</u> enough, then simulated annealing search will find a global optimum with probability approaching 1
- Few more precise convergence rate
 - Recent work on rapidly mixing Markov chains.
 - Surprisingly deep foundations.

- method proposed in 1983 by IBM researchers for solving VLSI layout problems.
 - theoretically will always <u>find the global optimum</u> (the best solution)
- Useful for some problems, but can be very slow
 - \circ slowness comes about because T must be decreased very gradually to retain optimality
- In practice <u>how to decide the rate</u> at which to decrease T? (this is a practical problem with this method)

Idea:

- Keep track of <u>k</u> states rather than just one
- Start with k randomly generated states
- At each iteration, all the successors of all k states are generated
- If any one is a goal state, stop;
- else select the k best successors from the complete list and repeat.

```
function Beam-Search(problem, k) returns a solution state
start with k randomly generated states
loop
```

generate all successors of all k states
if any of them is a solution then return it
else select the k best successors

A local beam search with k states might seem to be nothing more than running k random restarts in parallel instead of in sequence.

Local beam search with k = 1

- We would randomly generate 1 start state
- At each step we would generate all the successors, and retain the 1 best state
- Equivalent to HILL-CLIMBING

Local beam search with $k = \infty$

- 1 initial state and no limit of the number of states retained
- We start at initial state and generate <u>ALL</u> successor states (no limit how many)
- If one of those is a goal, we stop
- Otherwise, we generate all successors of those states (2 steps from the initial state), and continue
- ▶ Equivalent to **BREADTH-FIRST SEARCH** except that each layer is generated all at once.

Reading Material

- Artificial Intelligence, A Modern Approach Stuart J. Russell and Peter Norvig
 - Chapter 4.