

# Basic Relational Data Model **Keys & Constraints**

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#### **Objectives**

- Terminology
- Characteristics of Relations
- Relational Data Model Notations
- Key constraints
- Others Constraints



#### Terminology

- A relation is a table with columns and rows.
- Attribute is a named column of a relation.
- Domain is a set of allowable values for one or more attributes.
- Tuple is a row of a relation.
- Degree is a number of attributes in a relation.
- Cardinality is a number of tuples in a relation.
- Relationship is a link or dependency between relations.
- Relational Database is a collection of relations.

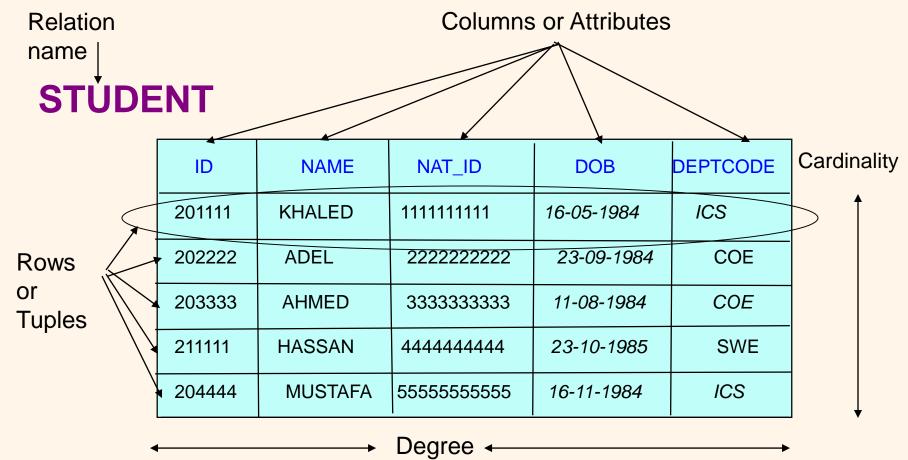


## Terminology

A relation is a table with columns and rows.

A table can be written as

STUDENT(ID, Name, Nat\_ID, DOB, DeptCode)





#### Characteristics of Relations

- Each relation in the same relational database schema has a distinct name
- Each value in a tuple is atomic
- Each attribute in a relation has a distinct name.
- Values of an attribute are all from the same domain.
- Each tuple is distinct.
- Order of attributes has no significance.
- Order of tuples has no significance, theoretically.



### Relational Data Model Notations

The letters Q, R, S denote the abstract relation names.

R(A<sub>1</sub>, A<sub>2</sub>, A<sub>3</sub>, ...., A<sub>n</sub>) denotes a relation schema R of degree n. Example: STUDENT(RollNo, Name, Phone, Address, Mobile, DOB) COURSE(<u>CourseID</u>, Title, CreditHours)

Both t[Ai] and t.Ai refers to the value vi in t for attribute Ai Example: in second tuple, both t[name] and t.name refers to "Adel"

Dom(DEPTCODE) = {'COE','ICS','SWE'}



#### Relation's Definition

```
Relation: A relation is a subset of Cartesian product of all domains
            R(r) \subset Dom(A_1) \times dom(A_2) \times ..... \times Dom(A_n)
            Where A1, A2, ...., An represent attributes of a relation r.
Evaluation:
            Given
                         A1 = \{1, 2\}, A2 = \{a\}, A3 = \{x, y\}
             Then relation
                         X = \{(1, a, y), (2, a, x), (2, a, y)\}; 3 tuples of a relation
            is a subset of
                         Y = Dom(A_1) \times dom(A_2) \times Dom(A_3)
= \{(1, a, x), (1, a, y), (2, a, x), (2, a, y)\}
Note: Degree of a relation is the number of attributes, in above case degree of relation is
3.
Alternate Definition of Relation
            For each tuple ti is a mapping from R to D
            Where D = Dom(A_1) \cup Dom(A_2) \cup ... \cup Dom(A_n)
            And t[Ai] \in Dom(Ai)
            Then a tuple can be considered as attribute and value.
For example: The following representation shows a tuple definition t = <(id, 21587), (Name, Muhammad,), (Hphone, Null), (Wphone, 866-1141)>
```



#### Data Models

- Object-Based Data Models
  - Entity-Relationship
  - Semantic
  - Functional
  - Object-Oriented
- Record-Based Data Models
  - Relational Data Model
  - Network Data Model
  - Hierarchical Data Model
- Physical Data-Models



#### Keys



A key is a combination of one or more columns that is used to identify rows in a relation

A composite key is a key that consists of two or

more columns



#### Super Key

A super key is a combination of columns that uniquely identifies any row within a relational database management system (RDBMS) table.

Roll#	FirstName	LastName	Address	City	NIC#	Deptno
12i-2565						CS
<del>(</del>	Ali				( )	CS EE
	All					CC

{Roll#}, {Roll#, NIC#}, {NIC#} {Roll#, FirstName, LastName, Address, City, NIC#, Deptno}



#### Example-1

Roll Number	First Name	Last Name
CSU0001	Shabbir	Bhimani
CSU0002	Jain	Rao
CSU0003	Nitin	Uday

#### Now we have the following as super keys

Roll Number	First Name	
Roll Number	First Name	Last Name



#### Example-2

Following table consists of four columns

EmployeeID, Name, Job, DeptID

Examples of superkeys in this table would be {employeeID, Name},

{employeeID, Name, job}, and

{employeeID, Name, job, departmentID}

In a real database we don't need values for all of those columns to identify a row

We only need, per our example, the set {EmployeeID}.

This is a minimal superkey

So, employeeID is a candidate key.

EmployeeID can also uniquely identify the tuples.

FirstName and FatherName (for NADRA)



#### Candidate Keys

A **candidate key** is a key that determines all of the other columns in a relation

Candidate key columns help in searching fewer duplicated or unique records.

Examples

In PRODUCT relation

Prod# is a candidate key

Prod\_Name is also a candidate key

In ORDER\_PROD

(OrderNumber, Prod#) is a candidate key

If BirthDate (or DrivingLicense#) is known then Name, address, telno etc. can be found from fewer searched records.



#### Identifying Candidate Keys

For Bank: AccountHolder (or Customer)

ACC#, Fname, Lname, DOB, CNIC#, Addr, City, TelNo, Mobile#, DriveLic#

For PTCL: Customer (single telno holders)

Consumer#, Fname, Lname, DOB, CNIC#, Addr, City, TelNo,
Mobile#

For NADRA: Citizen (CNIC#, Fname, Lname, FatherName, DOB, OldCNIC#, PAddr, PCity, TAddr, TCity, TelNo, Molile#)

TelNo, Mobile#:
Old value, Null,
fewer duplications,
changeable



# Identifying right record using Candidate Key – an Example

Let's analyze the following set of columns

{FirstName, LastName, DOB, TelNo, NatCardID}

First Name	Last Name	DOB	TelNo	NatCardID
Fiaz	Ali			
Fyaz	Ali			110116987231



#### Candidate Keys cont....

Example-1: Branch (branch-name, assets, branch-city) Candidate Key: {branch-name}

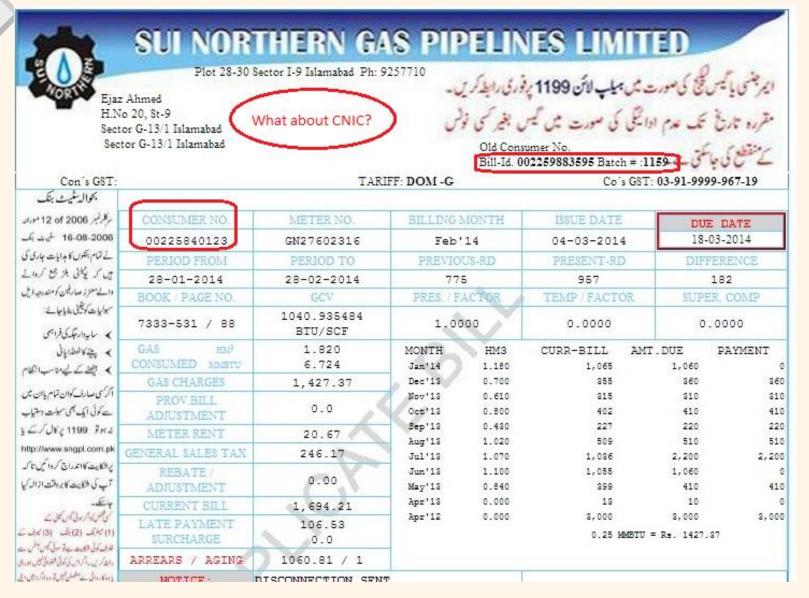
Example-2: Customer (cname, natid, address, city, telno) cname, address, city can be duplicated individually and cannot determine a record.

The following combinations distinguish customer records or tuples. {cname, telno} {natid} {natid, cname} As {natid} ⊆ {natid, cname}, then {natid, cname} is not candidate key and {natid} is a candidate key

Example-3: Employee(empno, name, birth\_date, address, city, telno, citizenship\_id) empno, telno, citizenship\_id are possible candidate keys???

#### **Exercise:**

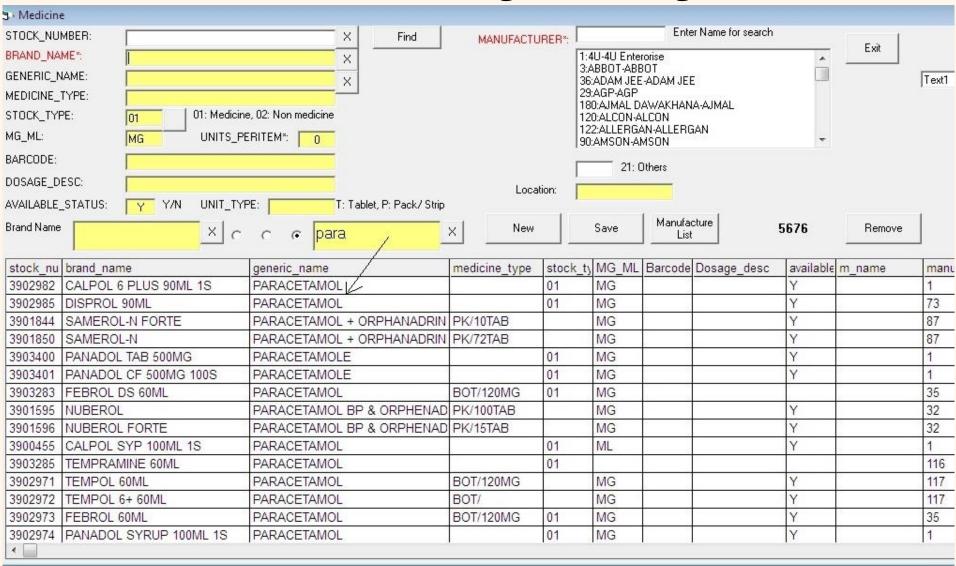
Course (cid, cname, deptno) Semester (sid, syear, startdate) ClassAllocation (sid, cid, sec#, building#, room#) Identify candidate keys in each of the above relations



Question: With one CNIC, there may be many consumer nos (houses/shops), explain?

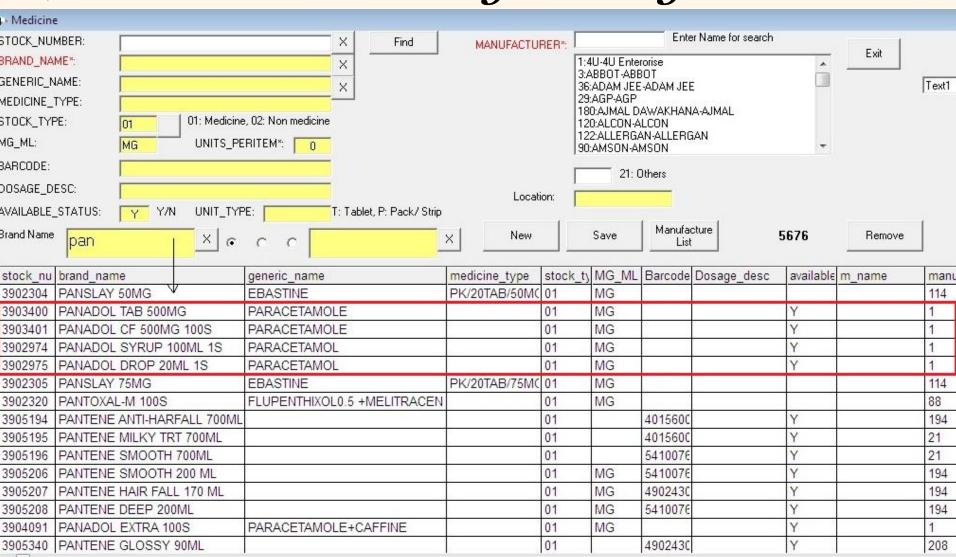


## Candidate Key during search





### Candidate Key during search ...



CS203 Database Systems



#### Dependents



For example: **Dependent** means family members (wife, son, daughter etc.), who depends on father. Father is an employee in some organization



#### Primary Keys

A **primary key** is a candidate key selected as the primary means of identifying rows in a relation:

There is one and only one primary key per relation

The primary key is NOT NULL, UNIQUE

The ideal primary key is short, numeric(alpha), indexed, fixed length and never changes

The primary key may be a composite key



#### Primary Keys cont....

A primary key is a minimal identifier that is used to identify tuples uniquely. This means that no subset of the primary key is sufficient to provide unique identification of tuples. NULL value is not allowed in primary key attribute.

#### Example:

STU\_ID

(Practically PK must have UNIQUE, NOT NULL, INDEXED on an attribute)

#### **STUDENT**

STU_ID	NAME	CID_ID	DOB	DEPTCODE
--------	------	--------	-----	----------

Although, candidate keys are STU\_ID, CIT\_ID. Organization ITSELF issues its own identifier PK. PK is a property of an organization. Driving license# is a property of Ministry of Motors/ transportation, it must not be used PK in university or other organizations. It is a candidate key for other organization for search purposes.

Note: It is not recommended to create PKs for columns like Quantity, Salary, Price and column that has datatype DATE



## Primary Keys cont...

STATE IDENTITY CARD

STUDENT

UNIVERSITY OF PANETRY

DISCONSTINA ASSESSMENT

DESCRIPTION ASSESSMENT

DESCRIPTION ASSESSMENT

DESCRIPTION

DESCRIPTION

TO BE STREETE

TO BE

**Example-1:** 

STUDENT(Stuld, FirstName, FamilyName, DOB, ...)

Example-2: Building (<u>B#</u>, BName, Location, Region), B# is a primary key. Although, BName is unique, not null but it is not short.

Example-3: Customer (cname, citizenid, address, city, telno)
This relation indicates the information about personal details. There is a chance that cname is duplicated, some may have citizenid and telno as null. This forces us to introduce new a attribute such as cust# that would be a primary key.

Customer (cust#, cname, citizenid, address, city, telno)

Example-4:

**BankBranch(Branch-Name, City, TotalAsset)** 

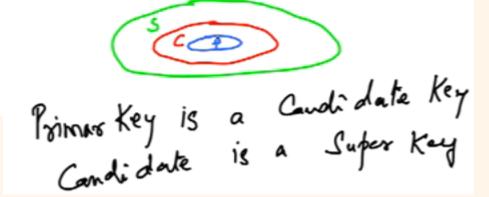
What will be a PK that you will suggest?
What is a candidate key?
Which attribute is unique
There may be many branches in one city then finalize this relation with possible constraints



#### Primary Keys cont...









## Primary Key Instances -Samples

**512**235665 610112654

Social Security Number

(By Birth Asian, Europian, HomeLand)

2003389

Student ID (With Registration year)

**DMM051** JUB004

Product ID (with Manufacturing Loc, branches)

312668852369 Bank Account# (followed by Branch Code)

100211 300411 SKU# (By Category) or a barcode

Serial#

Serial# starts with 1 and incremented by 1

005384

007004

Company has many branches with code 005, 007 etc., invoices issued from branches are recognized with branch codes







## Use of PK or Indexing

```
Table: PATIENT_VISIT
SQL> set timing on
                                                           Index: Yes, PK
SQL> select PATIENT_ID, max(PATIENT_VISITNO)
                                                           Test Query execution time: less than 1 second
    from PATIENT_VISIT
    where PATIENT_VISITNO = (select max(PATIENT_VISITNO)
                                                           Table: PATIENT_VISIT1
        from PATIENT_VISIT
                                                           Index: NO
       where DEPEND_SNO='0')
                                                           Test Query execution time:12 seconds
    group by PATIENT_ID
     faving count(*)>=1:
                                                           SQL> select count(*) Total_Records
                                                             2 from PATIENT_VISIT:
PATIENT_ID MAX(PATIENT_VISITNO)
                                                             Total_Records
9999993
                           1316
Elapsed: 00:00:00.12
                                                           Elapsed: 00:00:00.04
SQL> select PATIENT_ID, max(PATIENT_VISITNO)
    from PATIENT_VISIT1
    where PATIENT_VISITNO = (select max(PATIENT_VISITNO)
        from PATIENT_VISIT1
        where DEPEND_SNO='0')
   group by PATTENT_TD
     having count(*)>=1:
                                            Client
                                                                             Database.
PATIENT_ID MAX(PATIENT_VISITNO)
9999993
                           1316
Elapsed: 00:00:11.54
                                           PRIMARY KEY ( PATIENT_ID, DEPEND_SNO, PATIENT_VISITNO )
SQL>
```



#### Surrogate Keys

A **surrogate key** as an artificial column added to a relation to serve as a primary key:

DBMS supplied

Short, numeric and never changes – an ideal primary key!

Has artificial values that are meaningless to users

Normally hidden in forms and reports



#### Surrogate Keys cont....

RENTAL\_PROPERTY without surrogate key:

RENTAL\_PROPERTY (Street, City, State/Province, Zip/PostalCode, Country, Rental\_Rate)

RENTAL\_PROPERTY with surrogate key: RENTAL\_PROPERTY (<u>PropertyID</u>, Street, City, State/Province, Zip/PostalCode, Country, Rental\_Rate)

Other examples include Invoice#, ComplaintRef# ATMTransaction (Card#, Serial#, Amount, DrawDate)



#### Surrogate Keys cont....



#### Other examples include

Invoice#

**ComplaintRef**#

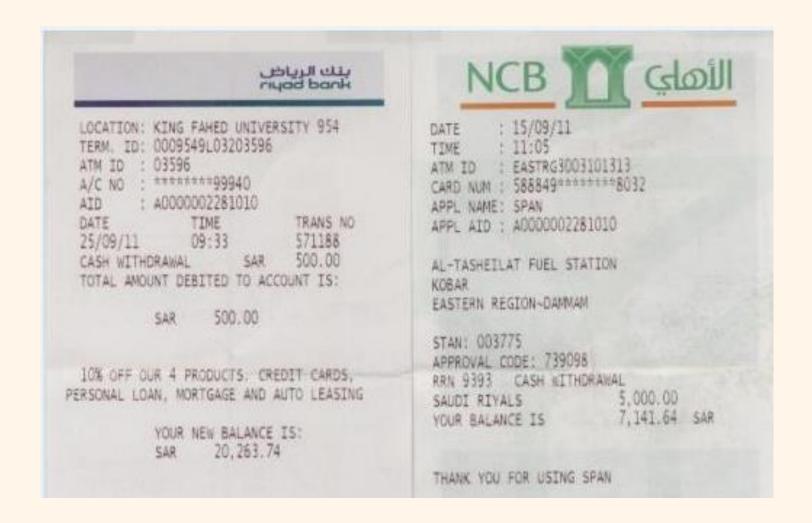


ATMTransaction (<u>Card#</u>, <u>Serial#</u>, Amount, TransDate, TransType)

InsurancePaid(<u>Policy#</u>, <u>PaidYear</u>, PaidDate, VehicleReg#, Amount)



#### ATM Transaction Sample





#### Surrogate Keys cont....

Surrogated Key

A tax payer with primary key NTN pays tax every year and tax is assessed in next or later year called ASSTT\_YEAR. An assessment or taxpayer status is identified every time with NTN and ASSTT\_YEAR.

TAXPAYER

NTN
FNAME
MNAME
MNAME
LNAME
ADDRESS
CITY
TELNO
......

DETAIL TAX
NTN
ASSTT YEAR
FILING\_DATE
AGG\_AMOUNT
PNY\_AMOUNT
.....

NTN and ASSTT\_YEAR are both primary keys in DETAIL\_TAX table, known as composite key of this table. NTN of DETAIL\_TAX table is a FK references to NTN of TAXPAYER table.

Exercise: What is different between PK and unique key?

Exercise: How to avoid that a record can be duplicated with all attributes other than PK attribute?

Exercise: Give three more examples from real life where composite keys can be applied

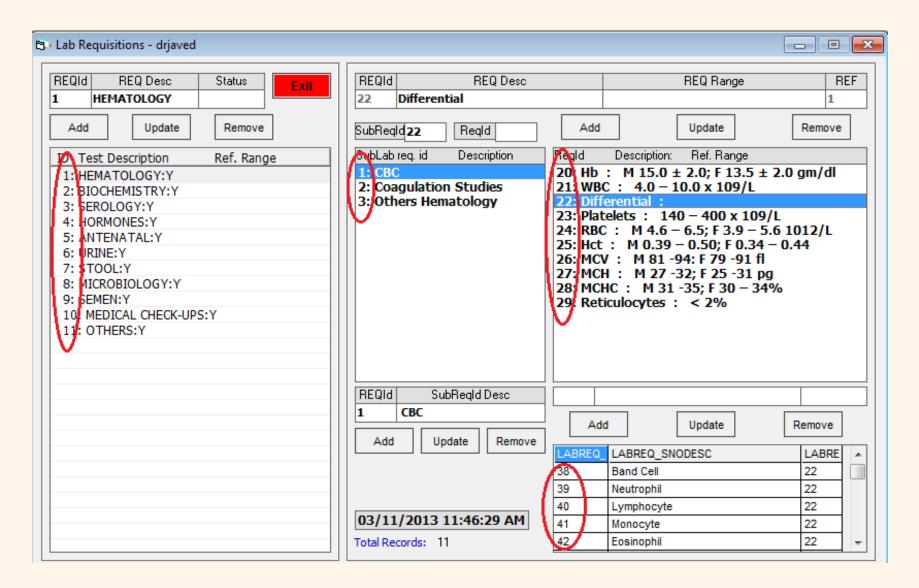


#### Surrogate Keys cont....

PatientID 558741 558741 558741	Visit#  1 2 3	Visit Date 
589210 589210 599211 599211	1 2 1 2	Better Solution using composite keys
PatientID 	Visit# 1 2 3 4 5 6 7	visit Date  



#### Example – End User View





#### Foreign Keys (FK)

A **foreign key** is an attribute that refers to a primary key of same or different relation to form a link (constraint) between the relations:

- A foreign key can be a single column or a composite key
- The term refers to the fact that key values are *foreign* to the relation in which they appear as foreign key values
- Ideally Data type, Length/ Size of FK and referring PK must be same



#### Foreign Keys

**NOTE:** The primary keys of the relations are <u>underlined</u> and any foreign keys are in *italics* in the relations below:

Name of FK column may be different from the name of referencing PK





**Example-1:** 

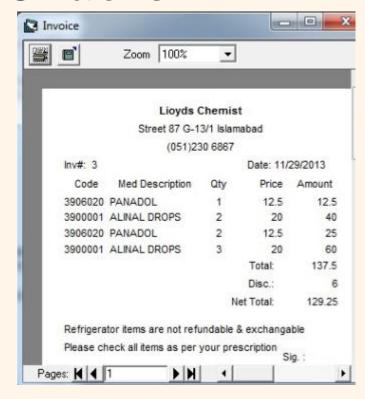
DEPARTMENT (<u>DeptID</u>, DepartmentName, BudgetCode, ManagerName)

**EMPLOYEE** (EmployeeNumber, EmployeeName, DeptID)



#### Customer-Orders





#### **Example-2:**

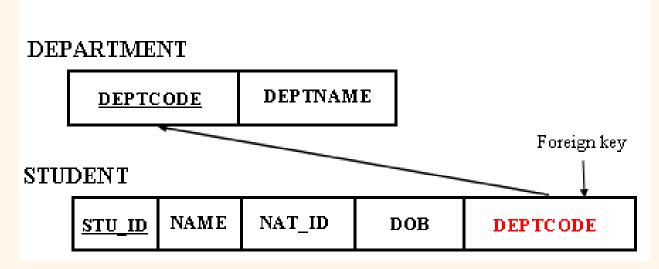
CUSTOMER(<u>CustId</u>, Name, Address, Telno, Email) ORDERS(<u>Invoice#</u>, InvType, InvDate, <u>CustomerId</u>)

Can we call Invoice# as a Surrogated Key?
Does Invoice# always contain serial numbers?



# Foreign Key with a Referential Integrity Constraint

An attribute or a set of attributes within one relation that matches the candidate key (or PK) of some (possibly the same) relation.

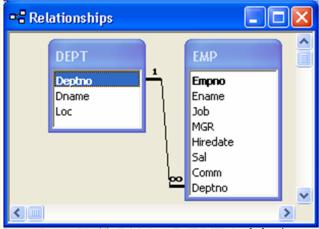


#### How to use FK in SQL:

select d.deptname, s.name, s.DOB
from Department d, Student s
where s.deptcode = e.deptcode;



# Foreign Key with a Referential Integrity Constraint



Deptno is a PK in DEPT table\*

Empno is a PK in EMP table

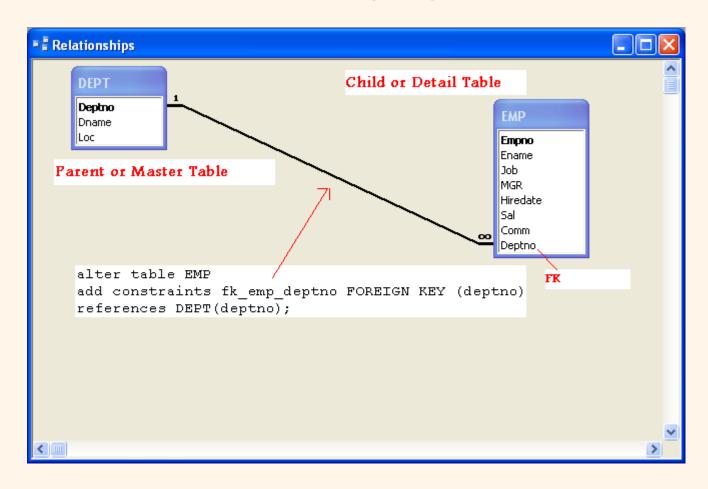
Deptno is a FK in EMP table\*

(it references to an attribute Deptno of DEPT table)

NOTE: An attribute within one relation that matches the key (Primary/same) in the same relation, is called recursive relation

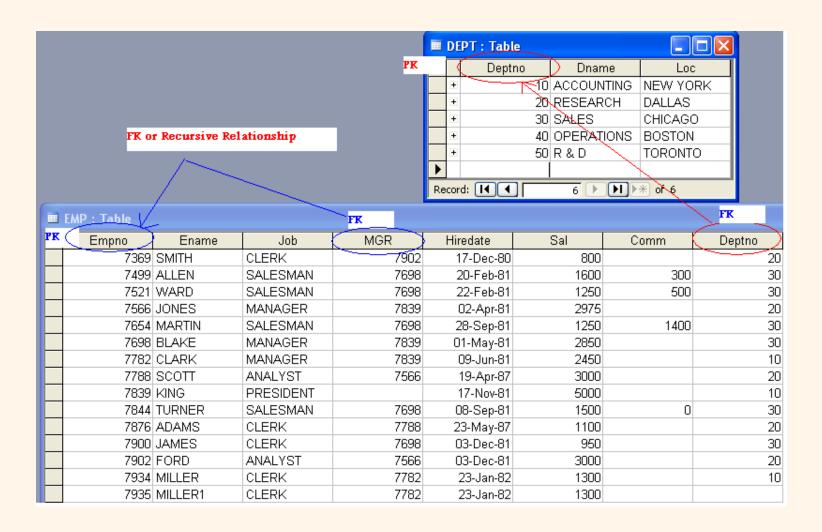


# Creating Foreign Key in MS Access & Oracle





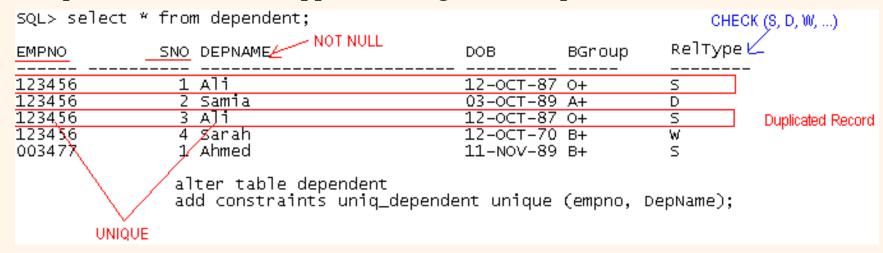
### FK and Recursive Relationship





#### Other Constraints Examples

Unique constraint can be applied on single or multiple columns



CHECK (Max\_Participants>=Enrolled\_Participants)

CHECK (Deptno>o)

CHECK (RelType IN ('S', 'D', 'W'))

Default 'Y'

**Database Triggers**