

CS 4072 - Topics in CS Process Mining

Lecture # 09

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FAST - NUCES, CFD Campus

Dr. Rabia Maqsood

rabia.maqsood@nu.edu.pk

Today's Topics

- ▶ Alpha algorithm

Alpha Algorithm

Process discovery from event log

α -algorithm

- ▶ The α -algorithm *scans* the event log *for particular patterns*.
- ▶ Four *log-based ordering relations* are defined to capture relevant patterns in an event log.
 - ▶ Direct succession ($>$)
 - ▶ Causality (\rightarrow)
 - ▶ Parallelism ($||$)
 - ▶ Choice ($\#$)

Footprint of a log

- ▶ A matrix capturing relations between all the activities

Table 6.1 Footprint of L_1 :
 $a \#_{L_1} a$, $a \rightarrow_{L_1} b$, $a \rightarrow_{L_1} c$,
 etc.

$$L_1 = [<a, b, c, d>^3, <a, c, b, d>^2, <a, e, d>]$$

	a	b	c	d	e
a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

From footprints to process patterns

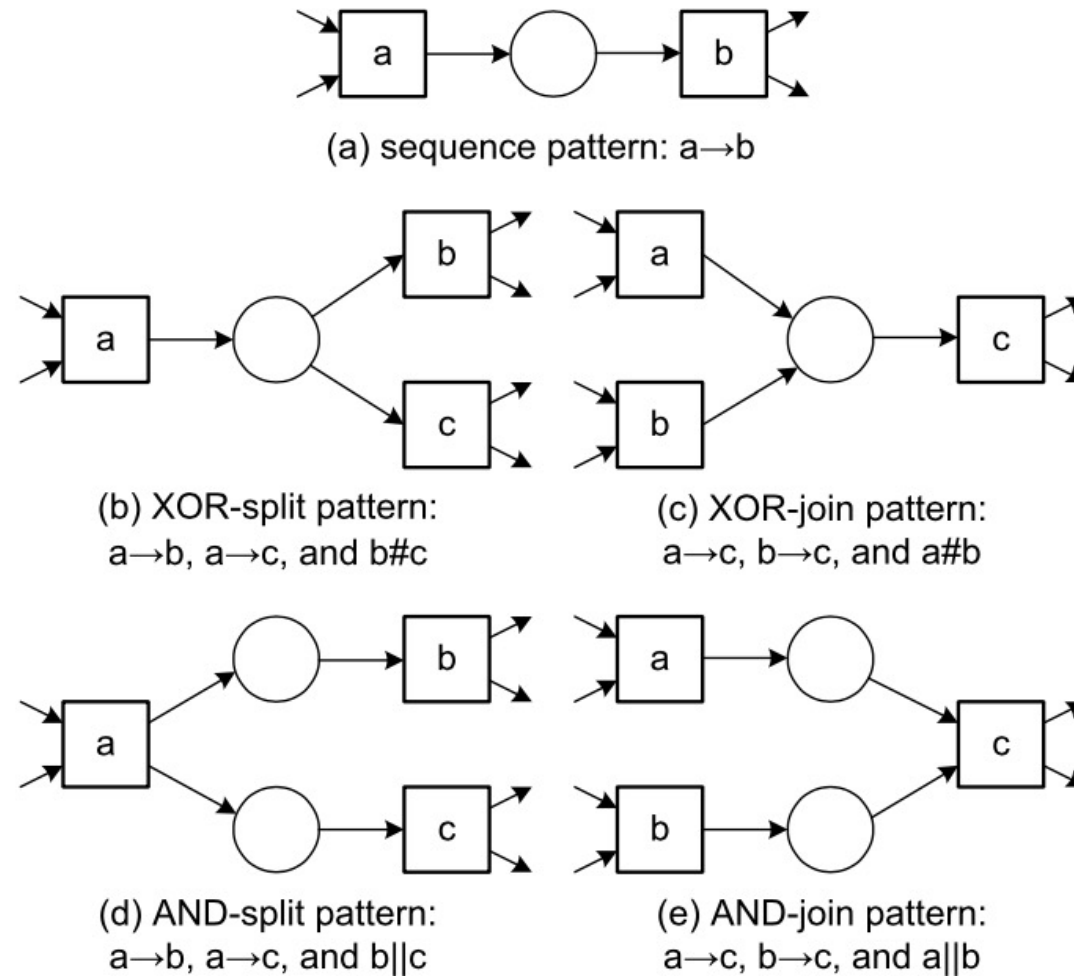


Fig. 6.4 Typical process patterns and the footprints they leave in the event log

Alpha Algorithm

► (α -algorithm): Let L be an event log over $T \subseteq \mathcal{A}$. $\alpha(L)$ is defined as follows:

1. $T_L = \{t \in T \mid \exists \sigma \in L \ t \in \sigma\}$,
2. $T_I = \{t \in T \mid \exists \sigma \in L \ t = \text{first}(\sigma)\}$,
3. $T_O = \{t \in T \mid \exists \sigma \in L \ t = \text{last}(\sigma)\}$,
4. $X_L = \{(A, B) \mid A \subseteq T_L \wedge A \neq \emptyset \wedge B \subseteq T_L \wedge B \neq \emptyset \wedge \forall a \in A \forall b \in B \ a \rightarrow_L b \wedge \forall a_1, a_2 \in A \ a_1 \#_L a_2 \wedge \forall b_1, b_2 \in B \ b_1 \#_L b_2\}$,
5. $Y_L = \{(A, B) \in X_L \mid \forall (A', B') \in X_L \ A \subseteq A' \wedge B \subseteq B' \implies (A, B) = (A', B')\}$,
6. $P_L = \{p_{(A, B)} \mid (A, B) \in Y_L\} \cup \{i_L, o_L\}$,
7. $F_L = \{(a, p_{(A, B)}) \mid (A, B) \in Y_L \wedge a \in A\} \cup \{(p_{(A, B)}, b) \mid (A, B) \in Y_L \wedge b \in B\} \cup \{(i_L, t) \mid t \in T_I\} \cup \{(t, o_L) \mid t \in T_O\}$, and
8. $\alpha(L) = (P_L, T_L, F_L)$.

Step 1: Each activity in L corresponds to a transition in $\alpha(L)$.

Step 2: The set of start activities includes the first elements of each trace.

Step 3: The set of end activities includes the last elements of each trace.

Example

Table 6.1 Footprint of L_1 :
 $a \#_{L_1} a, a \rightarrow_{L_1} b, a \rightarrow_{L_1} c$,
etc.

	$L_1 = [\langle a, b, c, d \rangle^3, \langle a, c, b, d \rangle^2, \langle a, e, d \rangle]$				
	a	b	c	d	e
a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

$T_L = \{a, b, c, d, e\}$
 $T_I = \{a\}$
 $T_O = \{d\}$

Alpha Algorithm

► (α -algorithm): Let L be an event log over $T \subseteq \mathcal{A}$. $\alpha(L)$ is defined as follows:

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4. $X_L = \{(A, B) \mid A \subseteq T_L \wedge A \neq \emptyset \wedge B \subseteq T_L \wedge B \neq \emptyset \wedge \forall a \in A \forall b \in B \ a \rightarrow_L b \wedge \forall a_1, a_2 \in A \ a_1 \#_L a_2 \wedge \forall b_1, b_2 \in B \ b_1 \#_L b_2\}$,
5. $Y_L = \{(A, B) \in X_L \mid \forall (A', B') \in X_L \ A \subseteq A' \wedge B \subseteq B' \implies (A, B) = (A', B')\}$,
6. $P_L = \{p_{(A,B)} \mid (A, B) \in Y_L\} \cup \{i_L, o_L\}$,
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8. $\alpha(L) = (P_L, T_L, F_L)$.

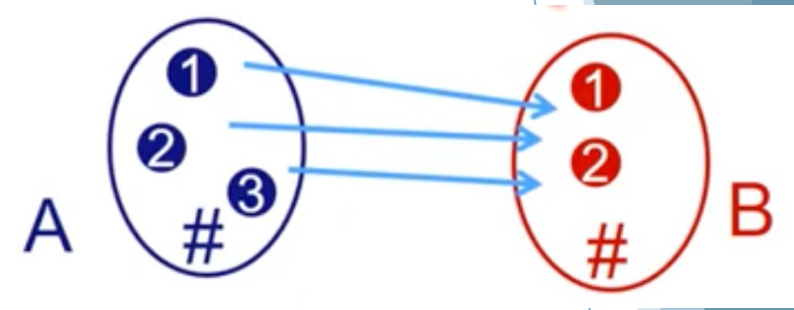
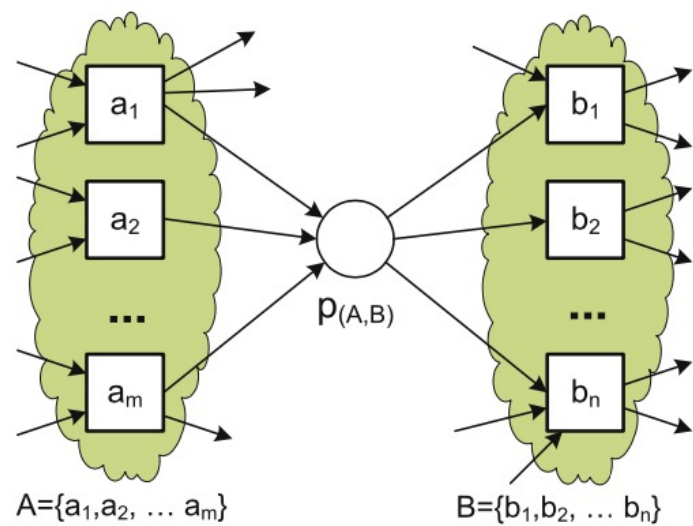
These are core steps of the α -algorithm.

The challenge is to determine the places of the WF-net and their connections.

We aim at constructing places named $p(A,B)$ such that A is the set of input transitions ($\bullet p(A,B) = A$) and B is the set of output transitions ($p(A,B) \bullet = B$) of $p(A,B)$.

Alpha Algorithm

Fig. 6.7 Place $p_{(A,B)}$ connects the transitions in set A to the transitions in set B



Check for the four quadrants

Table 6.4 How to identify $(A, B) \in X_L$? Rearrange the rows and columns corresponding to $A = \{a_1, a_2, \dots, a_m\}$ and $B = \{b_1, b_2, \dots, b_n\}$ and remove the other rows and columns from the footprint

	a_1	a_2	...	a_m	b_1	b_2	...	b_n
a_1	#	#	...	#	→	→	...	→
a_2	#	#	...	#	→	→	...	→
...
a_m	#	#	...	#	→	→	...	→
b_1	←	←	...	←	#	#	...	#
b_2	←	←	...	←	#	#	...	#
...
b_n	←	←	...	←	#	#	...	#

Example

Table 6.1 Footprint of L_1 :
 $a \#_{L_1} a, a \rightarrow_{L_1} b, a \rightarrow_{L_1} c$,
etc.

	$L_1 = [\langle a, b, c, d \rangle^3, \langle a, c, b, d \rangle^2, \langle a, e, d \rangle]$				
	a	b	c	d	e
a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

Identify all the pairs that meet the requirements of Step 4

Example

Table 6.1 Footprint of L_1 :
 $a \#_{L_1} a, a \rightarrow_{L_1} b, a \rightarrow_{L_1} c$,
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	$L_1 = [\langle a, b, c, d \rangle^3, \langle a, c, b, d \rangle^2, \langle a, e, d \rangle]$				
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a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

Following is the set of all pairs that meet the requirements of Step 4

$$X_{L_1} = \{(\{a\}, \{b\}), (\{a\}, \{c\}), (\{a\}, \{e\}), (\{a\}, \{b, e\}), (\{a\}, \{c, e\}), \\ (\{b\}, \{d\}), (\{c\}, \{d\}), (\{e\}, \{d\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

Inserting a place for any element in X_{L_1} would introduce too many places.

Therefore, only the “maximal pairs” (A,B) should be included.

Alpha Algorithm

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5. $Y_L = \{(A, B) \in X_L \mid \forall (A', B') \in X_L \ A \subseteq A' \wedge B \subseteq B' \implies (A, B) = (A', B')\}$,
6. $P_L = \{p_{(A, B)} \mid (A, B) \in Y_L\} \cup \{i_L, o_L\}$,
7. $F_L = \{(a, p_{(A, B)}) \mid (A, B) \in Y_L \wedge a \in A\} \cup \{(p_{(A, B)}, b) \mid (A, B) \in Y_L \wedge b \in B\} \cup \{(i_L, t) \mid t \in T_I\} \cup \{(t, o_L) \mid t \in T_O\}$, and
8. $\alpha(L) = (P_L, T_L, F_L)$.

Step 5 removes all non-maximal pairs.

Note that for any pair $(A, B) \in X_L$, non-empty set $A' \subseteq A$, and non-empty set $B' \subseteq B$, it is implied that $(A', B') \in X_L$.

Example

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c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
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e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

$$X_{L_1} = \{(\{a\}, \{b\}), (\{a\}, \{c\}), (\{a\}, \{e\}), (\{a\}, \{b, e\}), (\{a\}, \{c, e\}), \\ (\{b\}, \{d\}), (\{c\}, \{d\}), (\{e\}, \{d\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

X_{L_1} after the removal of non-maximal pairs:

$$Y_{L_1} = \{(\{a\}, \{b, e\}), (\{a\}, \{c, e\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

Alpha Algorithm

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8. $\alpha(L) = (P_L, T_L, F_L)$.

Step 6: Every element of $(A, B) \in Y_L$ corresponds to a place $p_{(A,B)}$ connecting transitions A to transitions B (in union with a special input and output place).

Example

Table 6.1 Footprint of L_1 :
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c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

$$X_{L_1} = \{(\{a\}, \{b\}), (\{a\}, \{c\}), (\{a\}, \{e\}), (\{a\}, \{b, e\}), (\{a\}, \{c, e\}),$$

$$(\{b\}, \{d\}), (\{c\}, \{d\}), (\{e\}, \{d\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

$$Y_{L_1} = \{(\{a\}, \{b, e\}), (\{a\}, \{c, e\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

Every element of $(A, B) \in Y_L$ corresponds to a place $p_{(A,B)}$ connecting transitions A to transitions B (in union with a special input and output place).

$$P_L = \{p_{(\{a\}, \{b,e\})}, p_{(\{a\}, \{c,e\})}, p_{(\{b,e\}, \{d\})}, p_{(\{c,e\}, \{d\})}, i_L, o_L\}$$

Alpha Algorithm

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4. $X_L = \{(A, B) \mid A \subseteq T_L \wedge A \neq \emptyset \wedge B \subseteq T_L \wedge B \neq \emptyset \wedge \forall a \in A \forall b \in B \ a \rightarrow_L b \wedge \forall a_1, a_2 \in A \ a_1 \#_L a_2 \wedge \forall b_1, b_2 \in B \ b_1 \#_L b_2\}$,
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7. $F_L = \{(a, p_{(A, B)}) \mid (A, B) \in Y_L \wedge a \in A\} \cup \{(p_{(A, B)}, b) \mid (A, B) \in Y_L \wedge b \in B\} \cup \{(i_L, t) \mid t \in T_I\} \cup \{(t, o_L) \mid t \in T_O\}$, and
8. $\alpha(L) = (P_L, T_L, F_L)$.

Step 7: The arcs of the WF-net are generated.

All places $p_{(A, B)}$ have A as input nodes and B as output nodes.

All start transitions in T_I have i_L as an input place and all end transitions T_O have o_L as output place.

Example

$$L_1 = [\langle a, b, c, d \rangle^3, \langle a, c, b, d \rangle^2, \langle a, e, d \rangle]$$

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	a	b	c	d	e
a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
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e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

$$\begin{aligned} T_L &= \{a, b, c, d, e\} \\ T_I &= \{a\} \\ T_O &= \{d\} \end{aligned}$$

Step 7: The arcs of the WF-net are generated.
 All places $p(A, B)$ have A as input nodes and B as output nodes.
 All start transitions in T_I have i_L as an input place and all end transitions T_O have o_L as output place.

$$Y_{L_1} = \{(\{a\}, \{b, e\}), (\{a\}, \{c, e\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

$$P_L = \{p(\{a\}, \{b, e\}), p(\{a\}, \{c, e\}), p(\{b, e\}, \{d\}), p(\{c, e\}, \{d\}), i_L, o_L\}$$

$$\begin{aligned} F_L = \{ & (a, p(\{a\}, \{b, e\})), (p(\{a\}, \{b, e\}), b), (p(\{a\}, \{b, e\}), e), \\ & (a, p(\{a\}, \{c, e\})), (p(\{a\}, \{c, e\}), c), (p(\{a\}, \{c, e\}), e), \\ & (b, p(\{b, e\}, \{d\})), (e, p(\{b, e\}, \{d\})), (p(\{b, e\}, \{d\}), d), \\ & (c, p(\{c, e\}, \{d\})), (e, p(\{c, e\}, \{d\})), (p(\{c, e\}, \{d\}), d), \\ & (i_L, a), \\ & (d, o_L) \} \end{aligned}$$

Alpha Algorithm

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8. $\alpha(L) = (P_L, T_L, F_L)$.

Step 8: resultant Petri-net that describes the behavior of input event log L .

Example

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c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

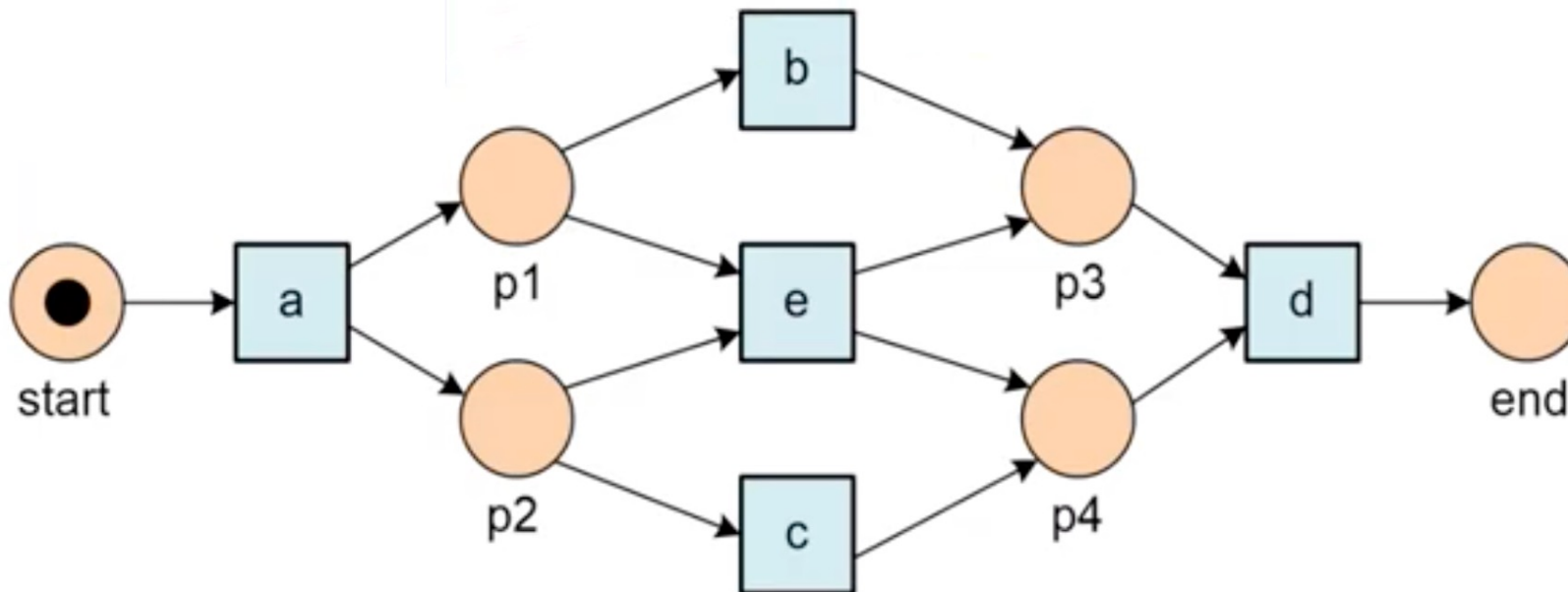
$$T_L = \{a, b, c, d, e\}$$

$$T_I = \{a\}$$

$$T_O = \{d\}$$

$$P_L = \{p(\{a\}, \{b, e\}), p(\{a\}, \{c, e\}), p(\{b, e\}, \{d\}), p(\{c, e\}, \{d\}), i_L, o_L\}$$

$$F_L = \{(a, p(\{a\}, \{b, e\})), (p(\{a\}, \{b, e\}), b), (p(\{a\}, \{b, e\}), e), \\ (a, p(\{a\}, \{c, e\})), (p(\{a\}, \{c, e\}), c), (p(\{a\}, \{c, e\}), e), \\ (b, p(\{b, e\}, \{d\})), (e, p(\{b, e\}, \{d\})), (p(\{b, e\}, \{d\}), d), \\ (c, p(\{c, e\}, \{d\})), (e, p(\{c, e\}, \{d\})), (p(\{c, e\}, \{d\}), d), \\ (i_L, a), \\ (d, o_L)\}$$

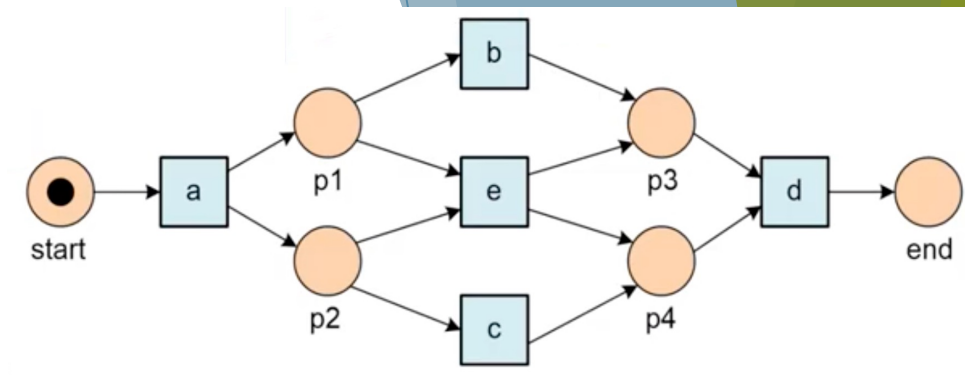


Example

Table 6.1 Footprint of L_1 :
 $a \#_{L_1} a, a \rightarrow_{L_1} b, a \rightarrow_{L_1} c$,
 etc.

	$L_1 = [\langle a, b, c, d \rangle^3, \langle a, c, b, d \rangle^2, \langle a, e, d \rangle]$				
	a	b	c	d	e
a	$\#_{L_1}$	\rightarrow_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}
b	\leftarrow_{L_1}	$\#_{L_1}$	\parallel_{L_1}	\rightarrow_{L_1}	$\#_{L_1}$
c	\leftarrow_{L_1}	\parallel_{L_1}	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$
d	$\#_{L_1}$	\leftarrow_{L_1}	\leftarrow_{L_1}	$\#_{L_1}$	\leftarrow_{L_1}
e	\leftarrow_{L_1}	$\#_{L_1}$	$\#_{L_1}$	\rightarrow_{L_1}	$\#_{L_1}$

1. $T_L = \{t \in T \mid \exists \sigma \in L \ t \in \sigma\}$,
2. $T_I = \{t \in T \mid \exists \sigma \in L \ t = \text{first}(\sigma)\}$,
3. $T_O = \{t \in T \mid \exists \sigma \in L \ t = \text{last}(\sigma)\}$,
4. $X_L = \{(A, B) \mid A \subseteq T_L \wedge A \neq \emptyset \wedge B \subseteq T_L \wedge B \neq \emptyset \wedge \forall a \in A \forall b \in B \ a \rightarrow_L b \wedge \forall a_1, a_2 \in A \ a_1 \#_L a_2 \wedge \forall b_1, b_2 \in B \ b_1 \#_L b_2\}$,
5. $Y_L = \{(A, B) \in X_L \mid \forall (A', B') \in X_L \ A \subseteq A' \wedge B \subseteq B' \implies (A, B) = (A', B')\}$,
6. $P_L = \{p_{(A, B)} \mid (A, B) \in Y_L\} \cup \{i_L, o_L\}$,
7. $F_L = \{(a, p_{(A, B)}) \mid (A, B) \in Y_L \wedge a \in A\} \cup \{(p_{(A, B)}, b) \mid (A, B) \in Y_L \wedge b \in B\} \cup \{(i_L, t) \mid t \in T_I\} \cup \{(t, o_L) \mid t \in T_O\}$, and
8. $\alpha(L) = (P_L, T_L, F_L)$.



$$T_L = \{a, b, c, d, e\}$$

$$T_I = \{a\}$$

$$T_O = \{d\}$$

$$X_{L_1} = \{(\{a\}, \{b\}), (\{a\}, \{c\}), (\{a\}, \{e\}), (\{a\}, \{b, e\}), (\{a\}, \{c, e\}), (\{b\}, \{d\}), (\{c\}, \{d\}), (\{e\}, \{d\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

$$Y_{L_1} = \{(\{a\}, \{b, e\}), (\{a\}, \{c, e\}), (\{b, e\}, \{d\}), (\{c, e\}, \{d\})\}$$

$$P_L = \{p_{\{a\}, \{b, e\}}, p_{\{a\}, \{c, e\}}, p_{\{b, e\}, \{d\}}, p_{\{c, e\}, \{d\}}, i_L, o_L\}$$

$$F_L = \{(a, p_{\{a\}, \{b, e\}}), (p_{\{a\}, \{b, e\}}, b), (p_{\{a\}, \{b, e\}}, e), (a, p_{\{a\}, \{c, e\}}), (p_{\{a\}, \{c, e\}}, c), (p_{\{a\}, \{c, e\}}, e), (b, p_{\{b, e\}, \{d\}}), (e, p_{\{b, e\}, \{d\}}), (p_{\{b, e\}, \{d\}}, d), (c, p_{\{c, e\}, \{d\}}), (e, p_{\{c, e\}, \{d\}}), (p_{\{c, e\}, \{d\}}, d), (i_L, a), (d, o_L)\}$$

Home Work

- Create the process model for the following event log using Alpha algorithm.

$L_5 = [\langle a, b, e, f \rangle^2, \langle a, b, e, c, d, b, f \rangle^3, \langle a, b, c, e, d, b, f \rangle^2, \langle a, b, c, d, e, b, f \rangle^4, \langle a, e, b, c, d, b, f \rangle^3]$

Table 6.5 Footprint of L_5

	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>
<i>a</i>	#	→	#	#	→	#
<i>b</i>	←	#	→	←		→
<i>c</i>	#	←	#	→		#
<i>d</i>	#	→	←	#		#
<i>e</i>	←				#	→
<i>f</i>	#	←	#	#	←	#

1. $T_L = \{t \in T \mid \exists \sigma \in L \ t \in \sigma\},$
2. $T_I = \{t \in T \mid \exists \sigma \in L \ t = \text{first}(\sigma)\},$
3. $T_O = \{t \in T \mid \exists \sigma \in L \ t = \text{last}(\sigma)\},$
4. $X_L = \{(A, B) \mid A \subseteq T_L \wedge A \neq \emptyset \wedge B \subseteq T_L \wedge B \neq \emptyset \wedge \forall a \in A \forall b \in B \ a \rightarrow_L b \wedge \forall a_1, a_2 \in A \ a_1 \#_L a_2 \wedge \forall b_1, b_2 \in B \ b_1 \#_L b_2\},$
5. $Y_L = \{(A, B) \in X_L \mid \forall (A', B') \in X_L \ A \subseteq A' \wedge B \subseteq B' \implies (A, B) = (A', B')\},$
6. $P_L = \{p_{(A, B)} \mid (A, B) \in Y_L\} \cup \{i_L, o_L\},$
7. $F_L = \{(a, p_{(A, B)}) \mid (A, B) \in Y_L \wedge a \in A\} \cup \{(p_{(A, B)}, b) \mid (A, B) \in Y_L \wedge b \in B\} \cup \{(i_L, t) \mid t \in T_I\} \cup \{(t, o_L) \mid t \in T_O\},$ and
8. $\alpha(L) = (P_L, T_L, F_L).$

Reading Material

- ▶ Chapter 6: Aalst