



TRƯỜNG ĐẠI HỌC BÁCH KHOA HÀ NỘI
HANOI UNIVERSITY OF SCIENCE AND TECHNOLOGY

Database

Lesson 2. Relational databases

Viet-Trung Tran
ls.hust.edu.vn/~trungtv/

Learning Map

Sequence	Title
1	Introduction to Databases
2	Relational Databases
3	Relational Algebra
4	Structured Query Language – Part 1
5	Structured Query Language – Part 2
6	Constraints and Triggers
7	Entity Relationship Model
8	Functional Dependency
9	Normalization
10	Storage - Indexing
11	Query Processing
12	Transaction Management – Part 1
13	Transaction Management – Part 2

Outline

1. Relational data model
2. Constraints

Learning objectives

- Upon completion of this lesson, students will be able to:
 - Recall some basic concepts of relational data model.
 - Show some constraints of relational data model.

Keywords

Data model	A set of concepts used to describe the structure of a database: data types, relationships, constraints, semantics...
Relation	Is thought of as a table of values, each row in the table represents a collection of related data values.
Key	An attribute or a set of attributes in the relation, which can identify a tuple uniquely.
Integrity constraints	Provide a way of ensuring that changes made to the database by authorized users do not result in a loss of data consistency.

1. Relational data model

1.1. Introduction

1.2. Database Basic concepts

1.1. Different data models

- Hierarchical database model
- Network model
- Object-oriented database model
- Relational model
- Entity-relationship model
- Document model
- ...

1.2. Relational data model

- Is very simple model, was first introduced by Ted Codd of IBM Research in 1970.
- Used by most of commercial database systems.
- Query with high-level languages.
- Efficient implementations.
- Based on mathematical theory, closed to file structure and data structure, there are three sets of terminology:

Relation	Table	File
Tuple	Row	Record
Attribute	Column	Field

1.3. Basic concepts

Relations

- are saved in the format of tables, which have rows and columns.

Relation instance/state

- actual contents at given point in time. The lowercase letters q, r, s denote relation states.

Database

- a set of named relations (or tables).

1.3. Basic concepts [2]

Tuple

- a single row of a table, which contains a single record for that relation. The letters t , u , v denote tuples.

Cardinality

- is the number of tuples in a relation.

Degree (arity)

- is the number of attributes in a relation.

1.4. Relational schema

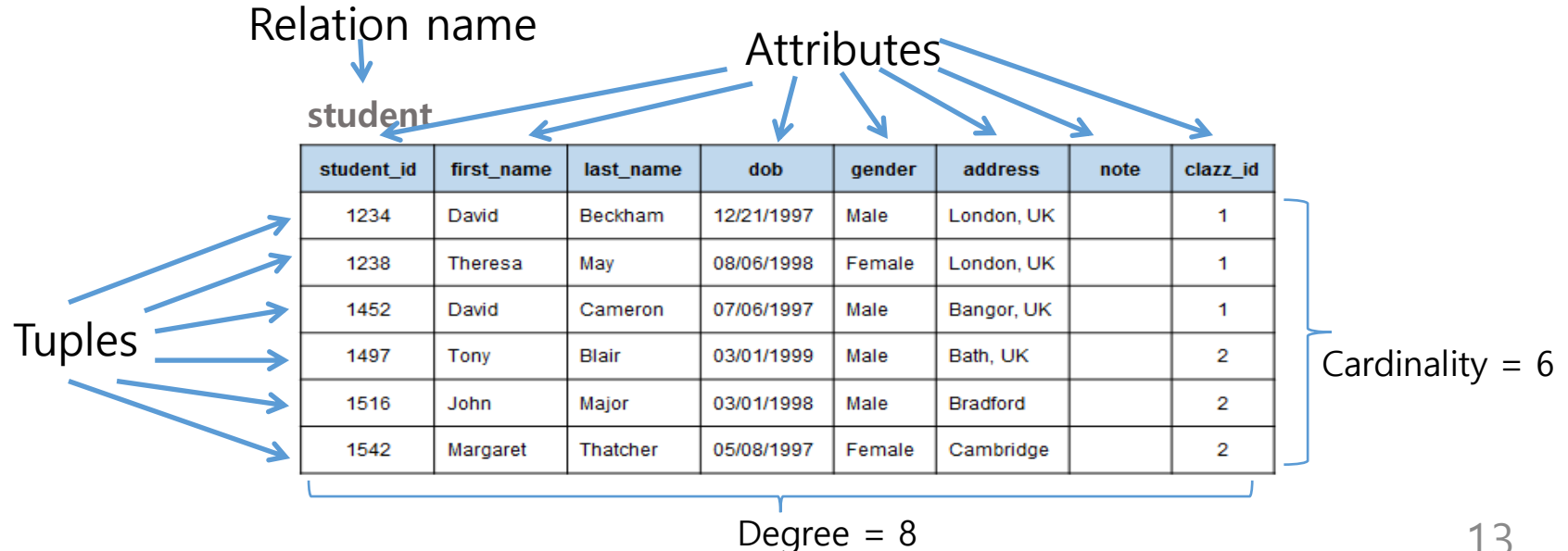
- Relational schema: structural description of relations in database.
- A relation schema R of degree n , denoted by $R(A_1, A_2, \dots, A_n)$, is made up of a relation name R and a list of attributes A_1, A_2, \dots, A_n
- Each attribute A_i has values belong to domain D_i of A_i , denoted by $\text{dom}(A_i)$
- An n -tuple t in a relation $r(R)$ is denoted by $t = \langle v_1, v_2, \dots, v_n \rangle$, where v_i is the value corresponding to attribute A_i . Both $t[A_i]$ and $t.A_i$ (and sometimes $t[i]$) refer to the value v_i in t for attribute A_i

1.4. Relational schema [2]

- Notice that the uppercase letters Q, R, S denote relation names.
- A relation (or relation state) r of the relation schema $R(A_1, A_2, \dots, A_n)$, also denoted by $r(R)$, is a set of n -tuples $r = \{t_1, t_2, \dots, t_m\}$. Each n -tuple t is an ordered list of n -values $t = \langle v_1, v_2, \dots, v_n \rangle$, where each value v_i , $1 \leq i \leq n$, is an element of $\text{dom}(A_i)$ or is a special NULL value.
- A relation (or relation state) $r(R)$ is a mathematical relation of degree n on the domains $\text{dom}(A_1), \text{dom}(A_2), \dots, \text{dom}(A_n)$, which is a subset of the Cartesian product of the domains that define R :
 - $$r(R) \subseteq (\text{dom}(A_1) \times \text{dom}(A_2) \times \dots \times \text{dom}(A_n))$$

1.5. An example

student(student_id, first_name, last_name, dob, gender, address, note, clazz_id)



2. Constraints

2.1. Introduction

2.2. Types of constraints

2.3. An example

2.1. Introduction

- Every relation has some conditions that must hold for it to be a valid relation.
- These conditions are called **Relational Integrity Constraints**.
- Provide a way of ensuring that changes made to the database by authorized users do not result in a loss of data consistency.

2.2. Types of constraints

- Key constraints
- Domain constraints
- Referential integrity constraints

2.2.1. Key constraints

- A key is an attribute or a set of attributes in the relation, which can identify a tuple uniquely.
- Key constraints enforce the following
 - in a relation with a key, no two tuples can have identical values for key attributes.
 - a key can not have NULL values.
 - Key constraints are also referred to as Entity Constraints.

2.2.1. Key constraints [2]

- Superkey: An attribute, or a set of attributes, that uniquely identifies a tuple within a relation.
- Candidate Key:
 - Superkey (K) such that no proper subset is a superkey within the relation
 - In each tuple of the relation, values of K uniquely identify that tuple (uniqueness)
 - No proper subset of K has the uniqueness property (irreducibility)
- Primary Key: Candidate key selected to identify tuples uniquely within a relation. Each key attribute of primary key has its name underlined.

2.2.1. Key constraints [3]

- Alternate Keys: Candidate keys that are not selected to be the primary key.
- Minimal key: a minimal set of attributes that can be used to identify a single tuple.
- Foreign Key:
 - Attribute, or set of attributes, within one relation that matches candidate key of some relation
 - Used to model relationships between relations
 - Each key attribute of foreign key has its name italic

2.2.2. Domain constraints

- Attributes have specific values in real-world scenario. Every attribute is bound to have a specific range of values.
- Within each tuple, the value of each attribute A must be an atomic value from the domain $\text{dom}(A)$.
- The data types associated with domains:
 - standard numeric data types for integers (short integer, integer, and long integer) and real numbers (float, double precision float).
 - Characters, Booleans, fixed-length strings, and variable-length strings, date, time, timestamp, and money, or other special data types.
 - a subrange of values from a data type .
 - an enumerated data type in which all possible values are explicitly listed.

2.2.2 Domain constraints [2]

- Null value
 - Represents value for an attribute that is currently unknown or not applicable for any tuple;
 - deals with incomplete or exceptional data;
 - represents the absence of a value and is not the same as zero or spaces

2.2.2. Domain constraints [3]

- Referential integrity constraints
 - Referential integrity constraints work on the concept of Foreign Keys. A foreign key is a key attribute of a relation that can be referred in other relation.
 - Referential integrity constraint states that if a relation refers to a key attribute of a different or same relation, then that **key element must exist**.

2.3. An example

- student(**student_id**, first_name, last_name, dob, gender, address, note, *clazz_id*)
- subject(**subject_id**, name, credit, percentage_final_exam)
- enrollment(*student_id*, *subject_id*, *semester*, midterm_score, final_score)

2.3. An example [2]

student

<u>student_id</u>	first_name	last_name	dob	gender	address	note	clazz_id
1234	David	Beckham	12/21/1997	Male	London, UK		1
1238	Theresa	May	08/06/1998	Female	London, UK		1
1452	David	Cameron	07/06/1997	Male	Bangor, UK		1
1497	Tony	Blair	03/01/1999	Male	Bath, UK		2
1516	John	Major	03/01/1998	Male	Bradford		2
1542	Margaret	Thatcher	05/08/1997	Female	Cambridge		2

subject

<u>subject_id</u>	name	credit	percentage_final_exam
IT3090	Databases	3	0.7
IT4843	Data integration	3	0.7
IT4868	Web mining	2	0.6
IT2000	Introduction to ICT	2	0.5
IT3020	Discrete Mathematics	2	0.7
IT3030	Computer Architectures	3	0.7

enrollment

<u>student_id</u>	<u>subject_id</u>	<u>semester</u>	midterm_score	final_score
1234	IT3090	20171	7	8
1238	IT3090	20171	9	8
1452	IT3090	20171	6	6
1234	IT2000	20162	5	8
1234	IT3020	20171	8	9
1452	IT3030	20171	7	9
1238	IT3020	20162	7	7

Primary key

Foreign key

Foreign key

Remark

- Relational data model
- Constraints
 - Key constraints
 - Domain constraints
 - Referential integrity constraints

Quiz 1.

Quiz Number	1	Quiz Type	OX	Example Select
Question	Relations: are saved in the format of			
Example	A. Tables B. Rows C. Columns D. Tuples			
Answer	A			
Feedback	In relational data model, relations are saved in the format of Tables. This format stores the relation among entities. A table has rows and columns, where rows represents records and columns represent the attributes.			

26

Quiz 2.

Quiz Number	2	Quiz Type	OX	Example Select
Question	Each key attribute has its name underlined			
Example	A. Super key B. Primary key C. Foreign key D. Candidate key			
Answer	B			
Feedback	Primary key is a chosen candidate key. Each relation has at most one primary key. By convention, all attributes of a primary key must be underlined			

Summary

- Relational data model
 - Relations, relation instance/state, relation schema
 - Database, tuple
 - Cardinality, degree
- Constraints
 - Key constraints
 - Domain constraints
 - Referential integrity constraints

Next lesson: Relational Algebra

- Raghu Ramakrishnan and Johannes Gehrke, Database Management Systems, 3rd edition, Mc Graw Hill, 2003.
- Elmasri and Navathe, Fundamentals of Database Systems, 6th edition, Addison-Wesley, 2011.