

一

$$A \vee B, A \rightarrow C = A$$

前提 $\rightarrow P \vee (q \rightarrow r) . q \rightarrow (r \rightarrow s) . P . q \quad (3, 4)$

附加前提引入

前提引入

10 P 引入

② $\frac{P}{\rightarrow P \vee (q \rightarrow r)}$

③ $\frac{q \rightarrow (r \rightarrow s)}{q \rightarrow (r \rightarrow s)}$

④ $\frac{P}{P}$ 假设引入

⑤ $\frac{}{\rightarrow q \vee r}$

⑥ $r.$

⑦ $\frac{}{\rightarrow q \vee (\rightarrow r \vee s)}$

⑧ $\frac{}{\rightarrow r \vee s}$

⑨ $s.$

②④ 合取三步

①⑤ 折衷二步

(3). 选择二步

①⑦.

⑥⑧.

二. 附加前提 $P \rightarrow (q \rightarrow s) . q . P \vee r . ; r \quad 得出 s.$

①. $r \quad X$

附加前提 $A \vee (B \wedge C)$

② q

前提引入 $= (A \vee B) \wedge (A \vee C)$

③ $P \vee r . \quad X$

前提引入 $A \wedge (B \vee C)$

④ $P \rightarrow (q \rightarrow s)$

前提引入 $= (A \wedge B) \wedge (A \wedge C)$

⑤ $(P \vee r) \wedge r .$

①③合取:

$(P \wedge r) \vee (q \rightarrow r \wedge r)$

P \wedge r

⑥ P

⑤ 什向

⑦ r

③ 什向

⑧ q \rightarrow s

④ ⑥ 假言

⑨ $\neg q \vee s$

⑧ 蕊合乎角

⑩ $(\neg q \vee s) \wedge q$

② ⑨ 行取二段式

$(\neg q \wedge q) \vee (s \wedge q)$

)

0 $\vee (s \wedge q)$

$s \wedge q$

⑪ s.

⑩ 什向

$\vdash P \rightarrow \neg q \quad \neg r \vee q \quad r \wedge s \vdash P$

- ① $p \rightarrow q$ x (前提引入) .
 ② $\neg r \vee q$ x. (前提引入)
 ③ $r \wedge \neg s$ (前提引入) $(A \vee B) \wedge (\neg p) \Rightarrow A$.
 ④ p x (否定引入)
 ⑤ $\neg p \vee \neg q$ x ①蕴含
 ⑥ $\neg q$ ④⑤析取三段论.
 ⑦ $\neg r$. ②⑥析取三段论.
 ⑧ $\neg r \wedge \neg s = 0$. ③⑦合取.

故为成假命题

- III $p \rightarrow (q \rightarrow r)$, $r \rightarrow (q \rightarrow s)$ | p | $q \rightarrow s$.
- ① $p \rightarrow (q \rightarrow r)$ x 前提引入
 ② $r \rightarrow (q \rightarrow s)$ 前提引入
 ③ p x 附加前提引入
 ④ $q \rightarrow r$. ①③假言推理.
 ⑤ $\neg r \vee \neg q \vee s$ v ②蕴含等值
 ⑥ $\neg q \vee r$. ④蕴含
 ⑦ $\neg q \vee r \vee s$. v ⑥或
 $(\neg q \vee s) \wedge (\neg r \vee r)$ ⑦合
 $\neg q \vee s$
 ⑧ $q \rightarrow s$. ⑦蕴含

五

P: A 到过受害者房间

Q: 11点前离开.

S: A 犯罪

T: A 被人看见

前提 $(P \wedge \neg Q) \rightarrow S \quad | \quad P$. $\neg Q \rightarrow T, \neg T : ! \quad | \quad S$.

① $(P \wedge \neg Q) \rightarrow S$

\neg 前提引入

② P

前提引入

③ $\neg Q \rightarrow T$

前提引入

④ $\neg T$.

前提引入

⑤ $\neg Q$

③④ 抵推

⑥ $P \wedge \neg Q$.

②⑤ 合取

⑦ S .

①⑦ 假言证法.

六

P: 今天是星期六

Q: 去颐和园

R: 去圆明园

T: 颐和园人多

前提 $P \rightarrow (Q \vee R)$, $T \rightarrow \neg Q$, $(P \wedge T) : ! \quad | \quad R$.

① $P \rightarrow (Q \vee R)$

\neg 前提引入

② $T \rightarrow \neg Q$

前提引入

③ $P \wedge t \quad X$

前提引入

④ P

⑤ 化简

⑤ $t \quad X$

⑥ 化简

⑥ $\neg q$

⑦ ⑧ 假设引入

⑦ $q \vee r$

⑨ ⑩ 假定原理

⑧ $\neg q \wedge (q \vee r)$

⑪ ⑫ 合取

$(\neg q \wedge q) \vee (\neg q \wedge r)$

$\neg q \wedge r$

⑬ r

⑭ 假设

↑ P: A 进场钻石项链

$\neg P$: B 进场钻石项链

q: 发生在营业时间

r: A 提供了正确证

s: 货柜上气泡

前提 ① $P \vee \neg P$ ② $\neg P \rightarrow \neg q$ ③ $r \rightarrow \neg s$

④ $\neg r \rightarrow q$ ⑤ s

① $\neg P \rightarrow \neg q \quad X$

前提引入

② $r \rightarrow \neg s \quad X$

前提引入

③ $\neg r \rightarrow q \quad X$

前提引入

④ s X

前提引入

- ⑤ $P \vee q$ ① 蕴含.
 ⑥ $\neg r \vee \neg s$ X ② 蕴含
 ⑦ $r \vee q$. ③ 蕴含
 ⑧ $\neg r$ ④ ⑥ 析取三段论
 ⑨ q ⑤ ⑧ 析取三段论.
 ⑩ $\neg P$ ⑥ ⑦ 析取三段论
 故 $\neg P$ 矛盾

八

$P: A$ 努力工作.

$q: B$ 感到困难

$r: C$ 感到愉快

$s: D$ 感到愉快

前提 $P \rightarrow (q \vee r)$, $q \rightarrow \neg P$
 $s \rightarrow \neg r$, 得 $P \rightarrow \neg s$.

- ① $P \rightarrow (q \vee r)$ X 前提引入.
 ② $q \rightarrow \neg P$ X 前提引入
 ③ $s \rightarrow \neg r$ 前提引入
 ④ P X 结论前提引入
 ⑤ $\neg q \vee \neg P$ X ② 蕴含.
 ⑥ $\neg q$. ④ ⑤ 析取三段论.

⑦ $\rightarrow q \vee q \vee r$

① 諸多

⑧ $\neg s \vee \neg r$

③ 諸多

⑨ $\neg s.$

⑦⑧折取三段

九. x 是个体域

设. $R(x)$: x 可以阅读,

$L(x)$: x 可以说话

$I(x)$: x 是有智力的

$D(x)$: x 是白痴.

$\forall x (R(x) \rightarrow L(x))$

$\forall x (D(x) \rightarrow \neg L(x))$

$\exists x (D(x) \wedge I(x))$

求证 $\exists x (I(x) \wedge R(x))$.

① $\exists x (D(x) \wedge I(x))$ 前提引入

② $\forall x (R(x) \rightarrow L(x))$ 前提引入

③ $\forall x (D(x) \rightarrow \neg L(x))$ 前提引入

④ $D(c) \wedge I(c)$ ① $\exists-$

⑤ $R(c) \rightarrow L(c)$ x ② $\forall-$

- ⑥ $D(C) \rightarrow \neg L(C) \times$ ③ A -
 ⑦ $\rightarrow R(C) \vee L(C)$ ④ 道路
 ⑧ $\neg D(C) \vee \neg L(C) \times$ ⑤ 有车
 ⑨ $D(C) \times$ ⑥ 行人
 ⑩ $I(C)$ ⑦ 行人
 ⑪ $\neg L(C) \times$ ⑧ ⑨ 满足行驶
 ⑫ $\neg R(C)$ ⑨ ⑩ 不满足行驶
 ⑬ $I(C) \wedge \neg R(C)$ ⑩ ⑪ 行人
 ⑭ $\exists x(I(x) \wedge \neg R(x))$ ⑫ 存在

+
 $F(x)$ 喜欢步行
 $G(x)$ 喜欢坐汽车.
 $H(x)$ 骑自行车.

前提: $\forall x(F(x) \rightarrow \neg G(x))$.

$\forall x(G(x) \vee H(x))$

$\exists x(\neg H(x))$

结论: $\exists x \neg F(x)$

① $\forall x(F(x) \rightarrow \neg G(x))$ 前提引入

② $\forall x(G(x) \vee H(x))$ 前提引入

③ $\exists x \neg H(x)$ 前提引入

- ④ $\neg h(x)$ ③ $\exists -$
 ⑤ $F(x) \rightarrow \neg G(x) x$ ① $A -$
 ⑥ $G(x) \vee h(x).$ ② $A -$
 ⑦ $\neg F(x) \vee \neg G(x)$ ⑤ 范例等值
 ⑧ $G(x)$ ④ ⑥ 析取三段论
 ⑨ $\neg F(x)$ ⑦ ⑧ 折取三段论
 ⑩ $\exists x \neg F(x).$ ⑨ $\exists +.$

+ -

c) x 是数字

$F(x)$: x 是有理数

$G(x)$: x 是实数

$h(x)$: x 是整数

前提 $\forall x(F(x) \rightarrow G(x))$, $\exists x(F(x) \wedge h(x))$.

结论 $\exists x(G(x) \wedge h(x))$

- ① $\forall x(F(x) \rightarrow G(x))$ 前提引入
 ② $\exists x(F(x) \wedge h(x))$ 前提引入
 ③ $F(x) \wedge h(x)$ $\exists -$
 ④ $F(x) \rightarrow G(x) x.$ $A -$
 ⑤ $\neg F(x) \vee G(x)$ ④ 范例等值式
 ⑥ $F(x)$ ⑤ 代入]
 ⑦ $G(x)$ ⑤ ⑥ 析取三段论.

⑧

$h(c)$

③ 代入

⑨
⑩

$G(c) \wedge h(c)$

⑦ ⑧ 合取

$\exists x (G(x) \wedge h(x))$

⑨ $\exists +$

$\vdash \vdash$

$P(x): x$ 是 15 倍数

$Q(x): x$ 是 3 倍数

$R(x): x$ 是 5 倍数

前提 $\forall x (P(x) \rightarrow Q(x))$, $\forall x (P(x) \rightarrow R(x))$

$\exists x P(x)$.

结论 $\exists x (R(x) \wedge Q(x))$

① $\forall x (P(x) \rightarrow Q(x))$

前提引入

② $\forall x (P(x) \rightarrow R(x))$

前提引入

③ $\exists x P(x)$ x

前提引入

④ $P(c)$

③ $\exists -$

⑤ $P(c) \rightarrow Q(c)$

① $\forall -$

⑥ $P(c) \rightarrow R(c)$

② $\forall -$

⑦ $Q(c)$

④ 假设前提

⑧ $R(c)$

⑤ 假设前提

⑨ $R(c) \wedge Q(c)$

⑦⑧ 合取

⑩ $\exists x (R(x) \wedge Q(x))$

⑨ $\exists +$

十三

前提

$$\forall x (\neg F(x) \rightarrow G(x))$$

$$\exists x (G(x) \wedge H(x))$$

$$\text{结论} \cdot \exists x (H(x) \wedge F(x))$$

$$\textcircled{1} \quad \forall x (\neg F(x) \rightarrow \neg G(x))$$

前提引入

$$\textcircled{2} \quad \exists x (G(x) \wedge H(x))$$

前提引入

$$\textcircled{3} \quad G(c) \wedge H(c)$$

\textcircled{2} \exists -

$$\textcircled{4} \quad \neg F(c) \rightarrow \neg G(c)$$

\textcircled{1} \forall -

$$\textcircled{5} \quad F(c) \vee G(c)$$

\textcircled{4} 蕴含等值式

$$\textcircled{6} \quad G(c)$$

\textcircled{3} 化简

$$\textcircled{7} \quad H(c)$$

\textcircled{3} 化简

$$\textcircled{8} \quad F(c)$$

\textcircled{5} \textcircled{6} 行取三段论

$$\textcircled{9} \quad H(c) \wedge F(c)$$

\textcircled{5} \textcircled{6} 合取

$$\textcircled{10} \quad \exists x (H(x) \wedge F(x))$$

\textcircled{9} \exists +

十四

$$\forall x (F(x) \rightarrow (H(x) \wedge G(x)))$$

$$\exists x (F(x) \wedge R(x))$$

结论: $\exists x (F(x) \wedge R(x) \wedge G(x))$

$$\textcircled{1} \quad \forall x (F(x) \rightarrow (H(x) \wedge G(x)))$$

\textcircled{1} 前提引入

$$\textcircled{2} \quad \exists x (F(x) \wedge R(x))$$

前提引入

③ $F(c) \wedge R(c)$

② $\exists -$

④ $F(c) \rightarrow (R(c) \wedge G(c))$

① $\forall -$

⑤ $F(c)$

③ 化简

⑥ $R(c) \wedge G(c)$

④ ⑤ 互反律推导

⑦ $F(c) \wedge R(c) \wedge G(c)$

③ ⑥ 合取

⑧ $\exists x(F(x) \wedge R(x) \wedge G(x))$

⑦ $\exists +$

$\frac{+}{\exists}$

$A = a, b \in N \times N$

$b=b \Leftrightarrow \langle a, b \rangle R \langle a, b \rangle$. 故自反.

$\forall \langle a, b \rangle, \langle c, d \rangle \in N \times N$

$\langle a, b \rangle R \langle c, d \rangle \Leftrightarrow b=d \Rightarrow d=b \Leftrightarrow \langle c, d \rangle R \langle a, b \rangle$

故具有对称性

$\forall \langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle \in N \times N$

$\langle a, b \rangle R \langle c, d \rangle \Leftrightarrow b=d$

$\langle c, d \rangle R \langle e, f \rangle \Leftrightarrow d=f$

$d=d \quad b=f$

规律

$\langle a, b \rangle R \langle c, d \rangle \wedge \langle c, d \rangle R \langle e, f \rangle$

④ $b=d \wedge d=f$ -

⑤ $b=f \Leftrightarrow \langle a, b \rangle R \langle e, f \rangle$. 传递.

十五

$\forall \langle a, b \rangle \in N \times N$

$a = a \Leftrightarrow \langle a, b \rangle R \langle a, b \rangle$ 向反.

$\forall \langle a, b \rangle, \langle c, d \rangle \in N \times N$

$\langle a, b \rangle R \langle c, d \rangle \rightarrow a = c \Rightarrow (=a \Leftrightarrow \langle c, d \rangle R \langle a, b \rangle)$

$\forall \langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle.$

$\langle a, b \rangle R \langle c, d \rangle \wedge \langle c, d \rangle R \langle e, f \rangle$

(2) $a = c \wedge c = e.$

$\Rightarrow a = e \Leftrightarrow \langle a, b \rangle R \langle e, f \rangle.$ 伝達.

$\forall \forall \langle a, b \rangle \in A \times A \quad a - b = c - d$

$a - b = a - b \Leftrightarrow \langle a, b \rangle R \langle a, b \rangle$ 向反.

$\forall \langle a, b \rangle, \langle c, d \rangle \in A \times A$

$\langle a, b \rangle R \langle c, d \rangle \Leftrightarrow a - b = c - d \Rightarrow c - d = a - b$

$\Leftrightarrow \langle c, d \rangle R \langle a, b \rangle$ 互換

$\forall \langle a, b \rangle, \langle c, d \rangle, \langle e, f \rangle \in A \times A$

$\langle e, b \rangle R \langle c, d \rangle \wedge \langle c, d \rangle R \langle e, f \rangle$

$\Leftrightarrow a - b = c - d \wedge c - d \wedge e - f$

$\Rightarrow a - b = e - f$

$\Leftrightarrow \langle a, b \rangle R \langle e, f \rangle$

十八 $A \oplus B = (A \cup B) - (A \cap B)$.

$$A \oplus (A \oplus B) = A \oplus (A \oplus C).$$

$$\Rightarrow (A \oplus A) \oplus B = (A \oplus A) \oplus C.$$

$$\Rightarrow \emptyset \oplus B = \emptyset \oplus C$$

$$\Rightarrow B \oplus \emptyset = C \oplus \emptyset$$

$$\Rightarrow B = C.$$

~~证~~ $\forall (x, y) \in A \times (B \cap C)$

则 $x \in A, y \in B \cap C$

$\therefore x \in A, y \in B \text{ 且 } y \in C$

由 $x \in A, y \in B$:

$(x, y) \in (A \times B)$.

由 $x \in A, y \in C$

$(x, y) \in (A \times C)$.

$\therefore (x, y) \in (A \times B) \cap (A \times C)$.

$\therefore A \times (B \cap C) \subseteq (A \times B) \cap (A \times C)$

②. $\forall (x, y) \in ((A \times B) \cap (A \times C))$

$(x, y) \in A \times B$

$(x, y) \in A \times C$

即 $x \in A, y \in B \text{ 且 } y \in C$.

$\exists \forall x \in A \quad y \in B \cap C$

$(x, y) \subseteq A \times (B \cap C)$

故 $. (A \times B) \cap (A \times C) \subseteq A \times (B \cap C)$

$\therefore A \times (B \cap C) = (A \times B) \cap (A \times C)$

二+(1)

$$ab = e \cdot bc = e \cdot$$

設 $a, b, c \in A$. b 是 a 的逆元 $\wedge c$ 是 b 的逆元.

$$e = b * c = b * c (a * b) * c$$

$$= (b * a) * c (b * c)$$

$$= (b * a) * e .$$

$$= b * a .$$

(2) -

設 A 有兩個逆元 b, c .

$$b = b * c = b * c (a * c) = (b * a) * c = e * c = c .$$

二+(~)

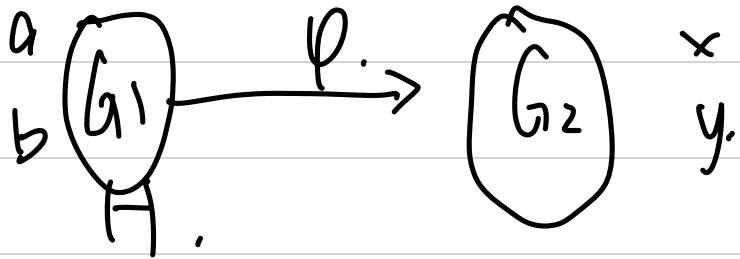
(1) H 是 G_1 子群 $\therefore H$ 非空. $\because \varphi(H)$ 非空.

$\forall a, b \in H$ 使 $f(a) = x \quad f(b) = y$.

$\because H$ 是 G_1 子群 $ab^{-1} \in H$.

$$xy^{-1} = f(a) \cdot f(b)^{-1} = f(ab^{-1}) \in f(H)$$

$\therefore f(H)$ 是 G_2 子群.



\therefore

H 是群 G 的子集, $a \in G$, 证 aHa^{-1} 是 G 的子集.

设 $aHa^{-1} = \{aha^{-1} ; h \in H\}$

$e \in G$. $e \in H$. H 是非空的.

$\forall a, a^{-1} \in H$. $aha^{-1} \in aHa^{-1}$

$$\begin{aligned} ah_1a^{-1} \cdot (ah_2a^{-1})^{-1} &= ah_1a^{-1} \cdot a \cdot h_2^{-1} \cdot a^{-1} \\ &= ah_1h_2^{-1}a^{-1} \in aHa^{-1}. \end{aligned}$$

\therefore

$a \in G$.

$H = \{x \in G \mid ax = xa\}$ 证明 H 是 G 的子集

$\forall x, y \in H$

$$ax = xa \quad \text{右乘 } a^{-1}$$

$$a \cdot x \cdot a^{-1} = x \quad H = \{x \in G \mid x = axa^{-1}\}.$$

$$a \cdot y \cdot a^{-1} = y$$

$$(xy^{-1})a = axa^{-1} \cdot a \cdot y^{-1}a \cdot a$$

$$= a x y^{-1} \quad .$$

$\therefore x y^{-1} \in H$. 可得.

二十一

如果其为素数阶群.

仅被 1 与 p 整除.

由拉格朗日定理

元素阶数只能为 1 或 p.

只有 1 是 1 阶元

其余全是生成元

二十五. 6 阶群必有 3 阶元

由拉格朗日定理 知.

G 中元素 1 阶 2 阶、3 阶、6 阶元

若 G 中含 6 阶元 则 a^2 为 3 阶元.

若不含 6 阶元 只含 1、2 阶元.

$$\forall a \in G, \quad a^2 = e.$$

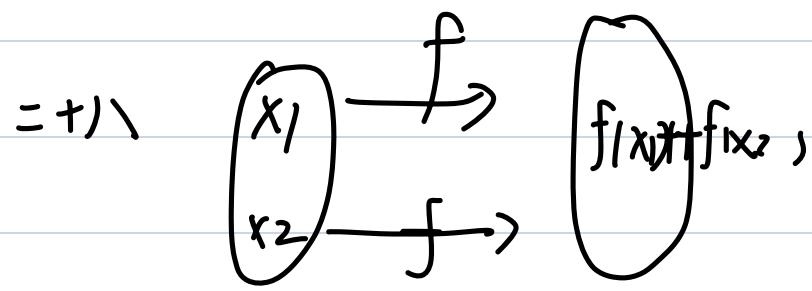
取 G 中两个不同的 2 阶元 a, b 令 $H = \{e, a, b, ab\}$.

H 是 G 子群 $|H| = 4 = |G|$ 矛盾.

故必含 3 阶元

二十二.

$\approx t$



设 $x_1 x_2$

$$f(x_1) + f(x_2) = f(x_1 x_2)$$

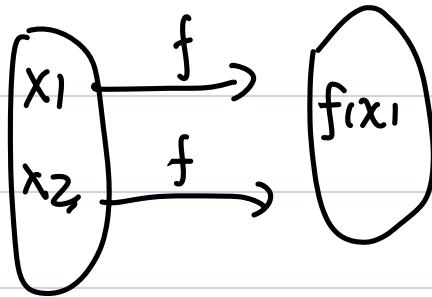
$$\ln x_1 + \ln x_2 = \ln x_1 + \ln x_2$$

同态 **单射** **满射**.

单射 $\forall x, y \in G \quad f(x) = f(y) \Rightarrow x = y$.

满射 $\forall y \in G \quad \exists x \in G \quad f(x) = y$.

=十九



$$f(x_1) \cdot f(x_2) = f(x_1 + x_2).$$
$$5^{x_1} \cdot 5^{x_2} = 5^{x_1+x_2} \quad \text{同态}$$

单射

$$\forall x, y \in G. \quad f(x) = f(y) \Rightarrow x = y.$$

满射

$$\forall y \in G \quad \exists x \in G \quad f(x) = y.$$

二十

设森林里每个树为 $S_1 \dots S_2 \dots S_k$.

阶数为 n_i ; ($i=1 \dots k$)

边数为 m_i ($i=1, 2 \dots k$).

$$S_1: n_1 = m_1 + 1$$

$$S_2: n_2 = m_2 + 1$$

$$S_k: n_k = m_k + 1$$

$$\sum_{i=1}^k n_i = \sum_{i=1}^k m_i + m.$$

$$n = m + k$$

$$m = n - k.$$

三十一.

假设两个奇度顶点不连通.

导致各自所对应的分支度数为奇数个

与握手定理矛盾

故必连通.

三十二

边数 m .

树叶 k .

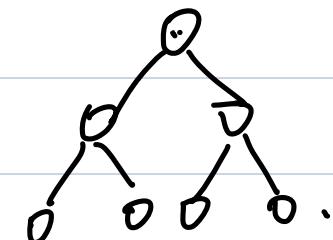
阶数 $m+1$.

由握手

出度为2

入度为1

度数为3.



$$2m = 2 + k + 3(m-k)$$

$f \in$

$$2m = 2 + k + 3m - 3k$$

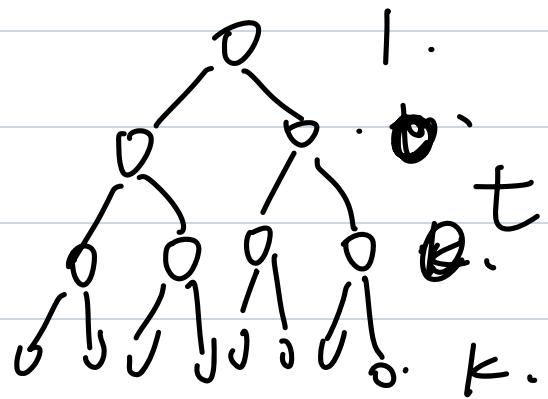
$$-m = 2 - 2k$$

$$m = 2(k-1).$$

三十三 共 n 个节点

树叶 t 个

3 度结点 k 个



推导:

$$\left\{ \begin{array}{l} n = t + k + 1 \quad t = n - k - 1 \\ 2(n-1) = 2 + 3t + k \end{array} \right.$$

$$2n - 2 = 2 + 3n - 3k - 3 + k$$

$$2n - 2 = 2 + 3n - 3k - 3 + k$$

$$-n = 4 - 3 - 2k$$

$$2k = n + 1$$

$$k = \frac{n+1}{2}$$

主析取项演消法原式极小项

解

成真赋值 · m.

$$1. (\underbrace{P \rightarrow Q} \wedge (Q \rightarrow R))$$

$$= (\neg P \vee Q) \wedge (\neg Q \vee R) \quad \text{主合取范式}$$

$$= (\neg P \vee Q \vee R) \wedge (\neg Q \vee \neg R) \wedge (P \vee \neg Q \vee R) \wedge (\neg P \vee \neg Q \vee \neg R)$$

要求主合取范式，要求其真值为1 求成假赋值 ·

100, 101, 010, 110.

即 M_4, M_5, M_2, M_6 .

主合取范式: $M_2 \wedge M_4 \wedge M_5 \wedge M_6$

主析取范式: $M_0 \vee M_1 \vee M_3 \vee M_7$.

$$2. \neg P \rightarrow Q \Leftrightarrow R$$

$$= (\neg P \vee Q) \Leftrightarrow R$$

$$= (\neg P \vee Q \rightarrow R) \wedge (R \rightarrow (\neg P \vee Q))$$

蕴含等值

等价等值

$$= (\neg(\neg P \vee Q) \vee R) \wedge (\neg R \vee (\neg P \vee Q))$$

$$= ((P \wedge \neg Q) \vee R) \wedge (\neg R \vee \neg P \vee Q)$$

$$= (P \vee R) \wedge (\neg Q \vee R) \wedge (\neg R \vee \neg P \vee Q) \rightarrow P \vee Q \vee \neg R$$

$$= (P \vee Q \vee R) \wedge \neg (P \vee \neg Q \vee \neg R) \wedge (P \vee \neg Q \vee R) \wedge (\neg P \vee Q \vee R) \wedge (\neg P \vee \neg Q \vee R)$$

$$= (P \vee Q \vee R) \wedge (\neg P \vee \neg Q \vee R) \wedge (\neg P \vee Q \vee \neg R)$$

主合取 成假. 000, 110, 101.

$M_0 \wedge M_6 \wedge S$

主合取式 $M_1 \vee M_2 \vee M_3 \vee M_4 \vee M_7$.

$$3. \rightarrow (p \rightarrow q) \wedge q \vee r$$

$$= \rightarrow (\rightarrow p \vee \neg q) \wedge q \vee r$$

$$= (p \wedge q) \wedge q \vee r$$

$$= (p \wedge q) \vee r.$$

$$= (p \vee r) \wedge (q \vee r)$$

$$= (p \vee q \vee r) \wedge (\neg p \vee q \vee r) \wedge (\neg q \vee r) \wedge (\neg r \vee r).$$

$$= (p \vee q \vee r) \wedge (\neg p \vee q \vee r) \wedge (\neg p \vee \neg q \vee r)$$

为主合取范式 为成假赋值 000, 010, 100.

$$M_0 \wedge M_2 \wedge M_4.$$

$$M_1 \vee M_3 \vee M_5 \vee M_6 \vee M_7.$$

4 A上 p

B上 q

C上 S.

前提 $p \rightarrow q, q \rightarrow r, \neg S \rightarrow (p \vee q)$.

$$= (p \rightarrow q) \wedge (q \rightarrow r) \wedge (\neg S \rightarrow (p \vee q))$$

$$= (\neg p \vee q) \wedge (\neg q \vee r) \wedge (\neg S \vee (p \vee q))$$

$$= (\neg p \vee q \vee S) \wedge (\neg q \vee r \vee S) \wedge (p \vee q \vee r \vee S) \wedge (\neg p \vee q \vee r \vee S) \wedge (p \wedge q \wedge r \wedge \neg S).$$

主合取 成假 100 110 011 111 000.

$$M_0 \wedge M_3 \wedge M_4 \wedge M_6 \wedge M_7.$$

成直主析取 $m_1 \vee m_2 \vee m_5$

BP 001, 010, 101.

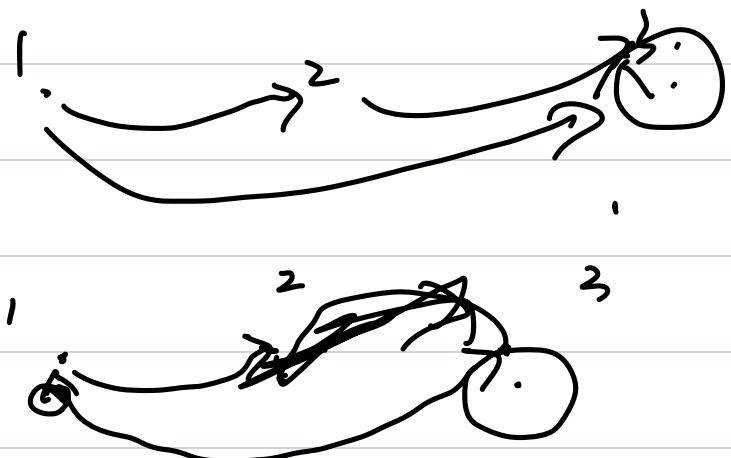
BP C上 AB不上

B上 AC不上

AC上 B不上

5.

$$M_R = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$



$r(R)$ 自反

$$\begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$

$s(R)$ 反称

$$\begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

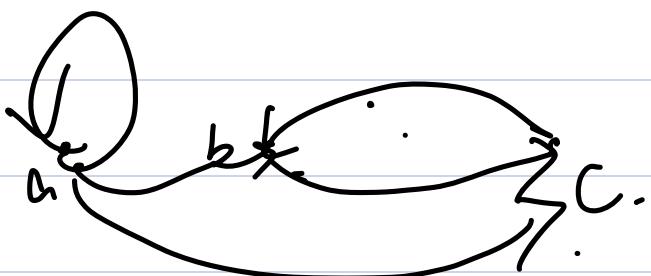
$t(R)$ 可传递

$$\begin{bmatrix} 0 & 1 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$

$$S(R) : \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

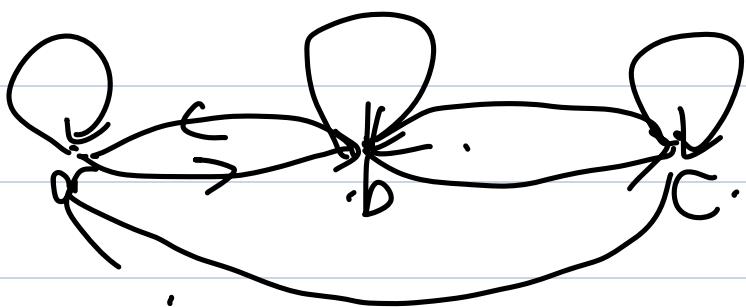
$$t(S(R)) = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}.$$

$$\begin{matrix} a & b & c \\ b & a & \begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \\ b & \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \\ c & \begin{bmatrix} 0 & 1 & 0 \end{bmatrix} \end{matrix}$$



$$r(R) \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix} \quad S(R) = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \quad t(R) \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

$$t(SR) = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$



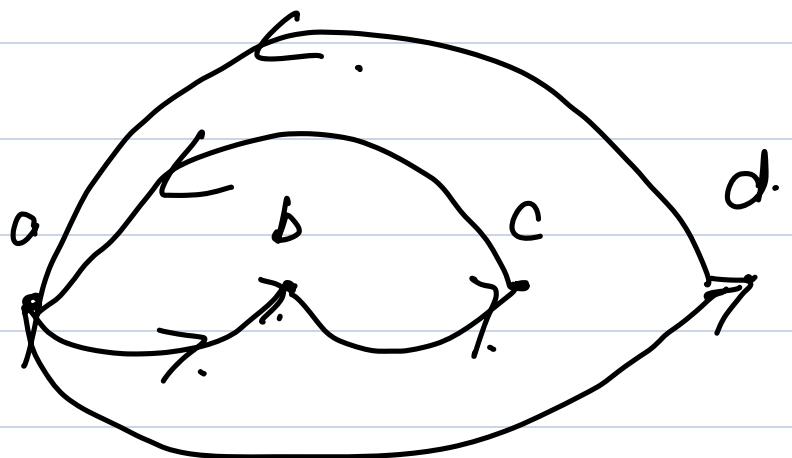
$$+ S \operatorname{r}(R) = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

7

$$\begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

8

$$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

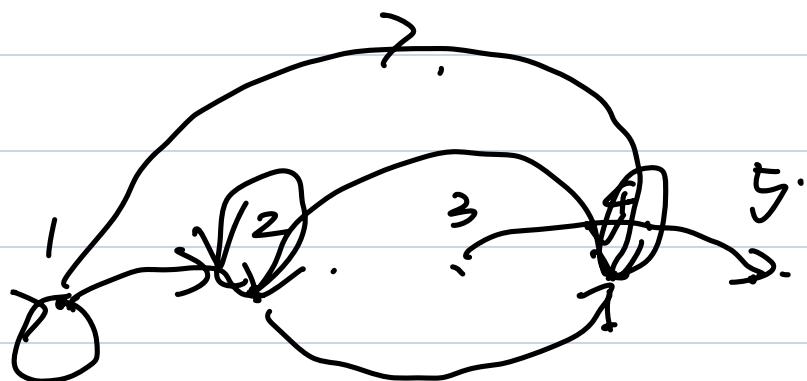


$t(R) = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$

1 | 1 | 1 | 1 |
 1 | 1 | 1 | 1 |

9.

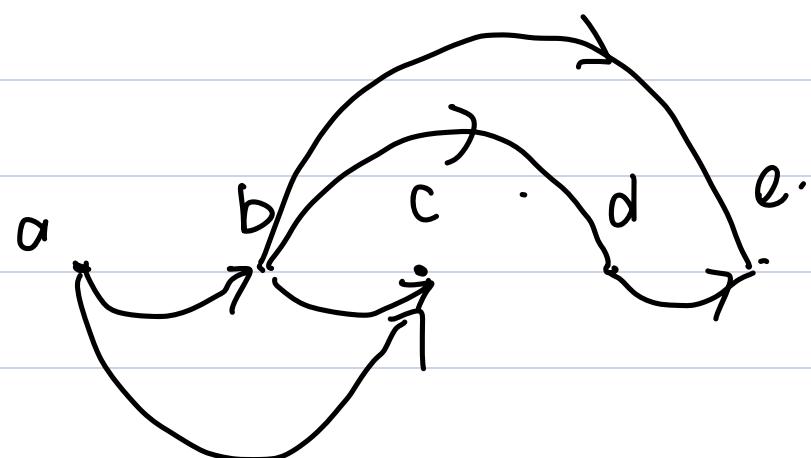
$$\begin{bmatrix} 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$



$t(CR) =$

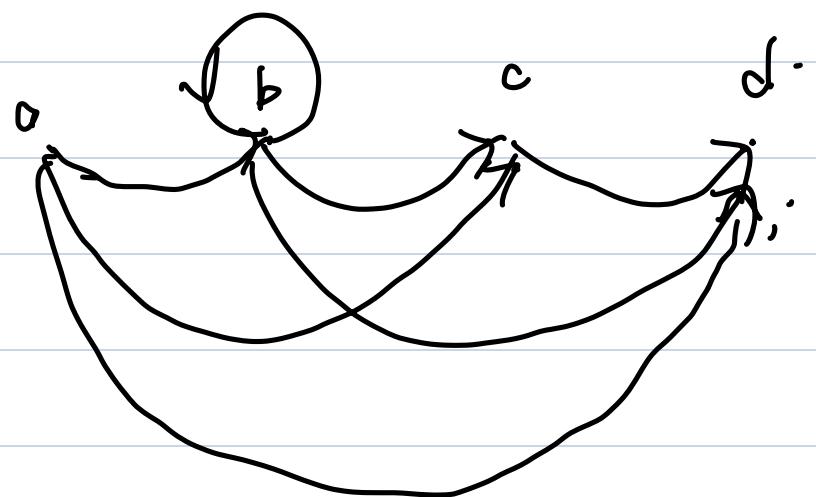
$$\begin{bmatrix} 1 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

	a	b	c	d	e
a	0	1	0	0	0
b	0	0	1	1	0
c	0	0	0	0	0
d	0	0	0	0	1
e	0	0	0	0	0



$$t(\alpha) = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\text{II.9} \quad \begin{bmatrix} a & b & c & d \\ 0 & 1 & 0 & 0 \\ b & 1 & 0 & 1 \\ c & 0 & 0 & 0 \\ d & 0 & 0 & 0 \end{bmatrix}$$



$$t(\beta) = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

[2. 证明自反、对称、~~可传递~~性质..]

$$R = \{ \langle x_1, y_1 \rangle, \langle x_2, y_2 \rangle \mid x_1 - y_1 = x_2 - y_2 \}.$$

对 $\forall \langle x, y \rangle \in R \quad \langle x, y \rangle \in R \quad \langle x, y \rangle \rightarrow x - y = x - y$ 自反.

对于 $\forall \langle x_1, y_1 \rangle, \langle x_2, y_2 \rangle \in R$

$$\langle x_1, y_1 \rangle R \langle x_2, y_2 \rangle \Leftrightarrow x_1 - y_1 = x_2 - y_2$$

$$\Rightarrow x_2 - y_2 = x_1 - y_1 \Leftrightarrow \langle x_2, y_2 \rangle R \langle x_1, y_1 \rangle \text{ 对称}.$$

对于 $\forall \langle x_1, y_1 \rangle, \langle x_2, y_2 \rangle, \langle x_3, y_3 \rangle \in X$:

$$\langle x_1, y_1 \rangle R \langle x_2, y_2 \rangle \wedge \langle x_2, y_2 \rangle R \langle x_3, y_3 \rangle$$

$$= x_1 - y_1 = x_2 - y_2 \wedge x_2 - y_2 = x_3 - y_3$$

$$\Rightarrow x_1 - y_1 = x_3 - y_3$$

$$\Leftrightarrow (x_1, y_1) R \langle x_3, y_3 \rangle \text{ 传递}$$

离散: $\left\{ \langle x_1 - y_1 = x_2 - y_2 = -2 \rangle, \langle x_1 - y_2 = x_2 - y_2 = -1 \rangle \right\}$.

$$+ \exists. \quad a - b = c - d.$$

$$2 - 5 = -3.$$

$$= \left\{ \langle 1, 4 \rangle, \langle 2, 5 \rangle, \langle 3, 6 \rangle, \langle 4, 7 \rangle, \langle 5, 8 \rangle, \langle 6, 9 \rangle \right\}.$$

④

$$\begin{bmatrix} 1 & .1 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

既不是自反也不是反自反、
不对称、
非传递。

$$\begin{bmatrix} 1 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 1 \end{bmatrix}$$

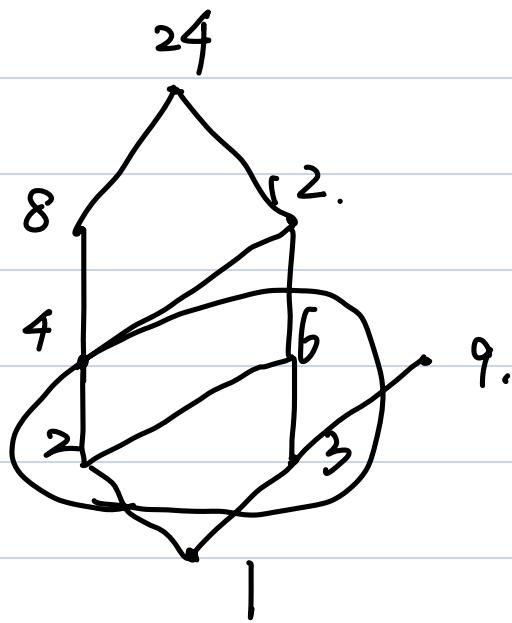
$$\begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

十五.



拔小 2, 3

拔大 6

~~最小~~ 大

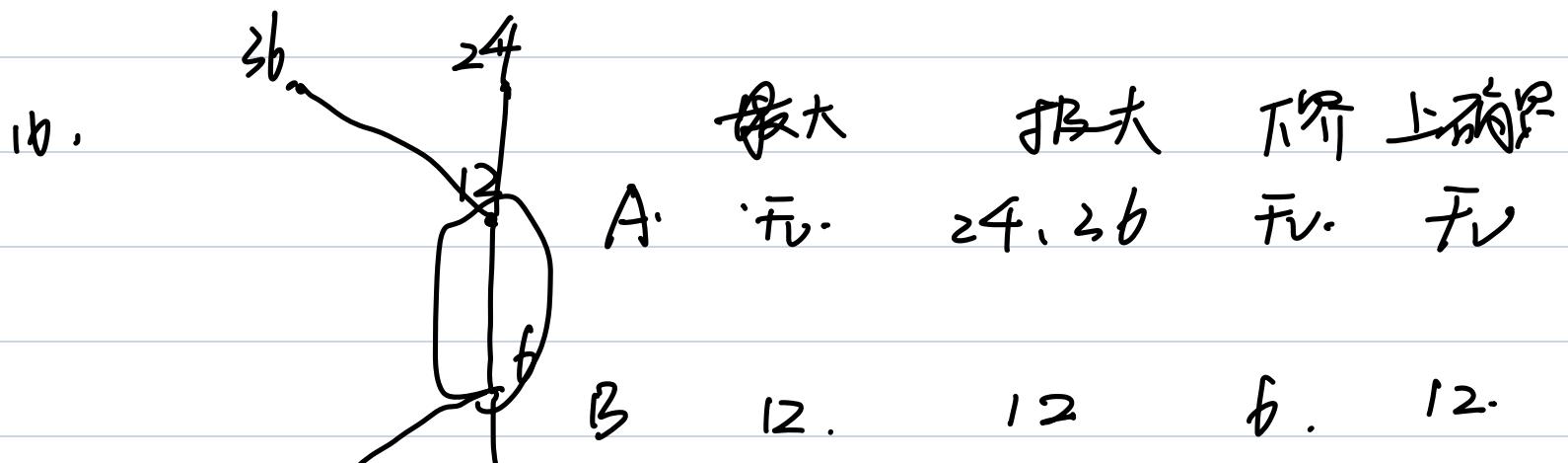
~~最大~~ 6

上界 6, 12, 24

下界: 1

~~最大上界~~: 6

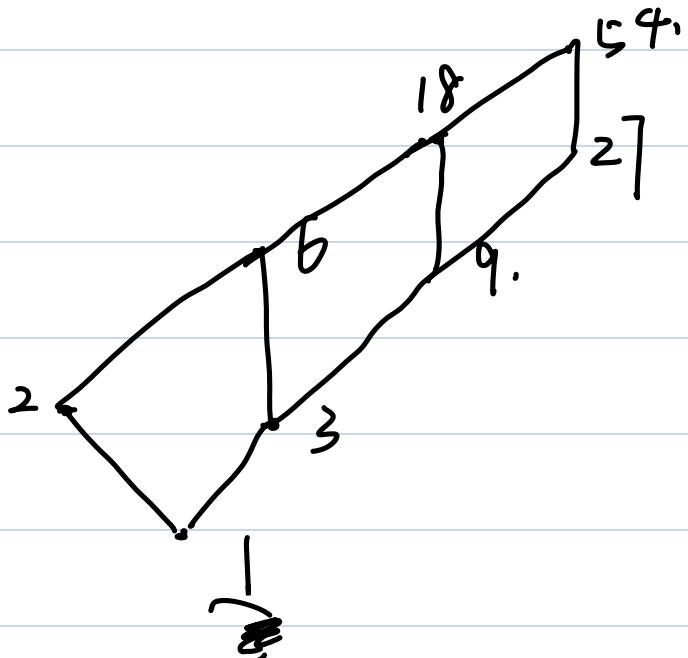
~~最小下界~~: 1.



2. 13 C f.. 1,2 F. 6,12.

17.

$$A = \{1, 2, 3, 6, 9, 18, 27, 54\}$$



最大 54

最小 1.

最大 54

TBW 1

$$18 + x, y \in \mathbb{Z} \quad x \circ y = x+y-1$$

得 $\forall x, y \in \mathbb{Z}$

$$\begin{aligned} (x \circ y) \circ z &= (x+y-1) \circ z \\ &= x+y-1+z-1 = x+y+z-2. \end{aligned}$$

$$x \circ (y \circ z) = x \circ (y+z-1)$$

$$= x+y+z-1-1 = x+y+z-2$$

$$(x \circ y) \circ z = x \circ (y \circ z) \text{ 故 } \text{成立}.$$

否单位元

$$x \circ y = x = x + y - 1$$

$1 = y$ 得 1 是此环的单位元.

可逆 x, x^{-1}, y, z 为 \mathbb{Z} 中的元素.

$$x \circ y = x + y - 1 = 1 \quad y = 2 - x \cdot = x^{-1} \cdot$$

$$x \circ y = x \circ x^{-1} = x \circ (2 - x) = x + 2 - x - (-1) \cdot$$

$$\text{即 } x^{-1} = 2 - x \cdot$$

且 $x, y \in \mathbb{Q}$. $x \circ y = x + y - c$. [可交换、可逆、含单位元]

$x, y, z \in \mathbb{Q}$.

$$(x \circ y) \circ z = (x + y - c) \circ z = x + y + z - c - c \cdot$$

$$x \circ (y \circ z) = x \circ (y + z - c) = x + y + z - c - c \cdot \text{ 互易性.}$$

$$x \circ y = x = x + y - c \quad y = c \cdot \text{ 单位元}$$

$x \quad x^{-1} \quad y \quad$ 可逆.

$$x \circ y = 1 = x + y - c \cdot$$

$$c + 1 - x \cdot = y \cdot$$

$$x \circ x^{-1} = x \circ (c + 1 - x) = c + 1 - x + x - c$$

$$= 1 \cdot$$

$\exists + \cdot (\mathbb{Z}_{24}, \oplus) = \{0 \rightarrow 23\} \cdot \langle 2n, \oplus \rangle$, $\mathbb{Z}_n = \{0 \dots n-1\}$ \oplus 为模 n 加法

1阶子群. $\langle \underline{\oplus} \rangle = \langle \underline{0} \rangle = \{0\}.$

2阶子群 $\langle \underline{12} \rangle = \{0, 12\}.$ $\frac{1}{12}x = v$

3阶子群 $\langle \underline{8} \rangle = \{0, 8, 16\}.$

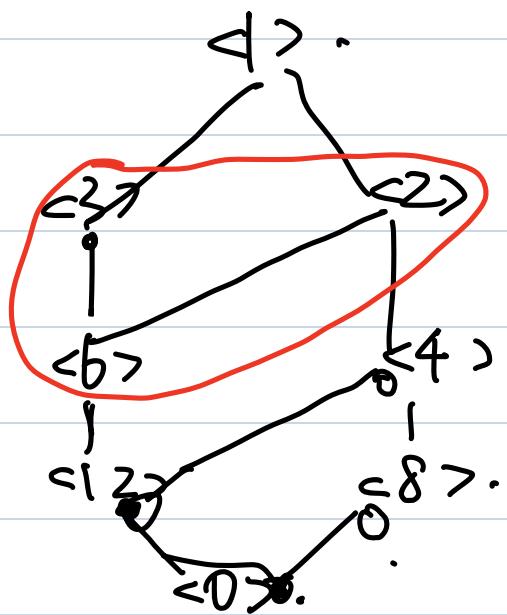
4阶子群 $\langle \underline{6} \rangle = \{0, 6, 12, 18\}.$

6阶子群 $\langle \underline{4} \rangle = \{0, 4, 8, 12, 16\}.$

8阶子群 $\langle \underline{3} \rangle = \{0, 3, 6, 9, 12, 15, 18, 21\}.$

12阶子群 $\langle \underline{2} \rangle = \{0, 2, 4, 6, 8, 10, 12, 14, 16, 18, 20\}.$

$2 \times P\Gamma_1$ $\langle \underline{1} \rangle = \mathbb{Z}_{24}.$



极大, <2>, <3>

极小, <6>

最大, <2>, <3>

极小, <6>

上界, <1>

下界, <6>, <12>, <0>.

上确界 <1>

下确界 <6>.

$$1\text{阶} = \langle 0 \rangle = \{0\}$$

$$2\text{阶} = \langle 24 \rangle = \{0, 24\}.$$

$$3\text{阶} = \langle 16 \rangle = \{0, 16\}.$$

$$4\text{阶} = \langle 12 \rangle = \{0, 12\}.$$

$$5\text{阶} = \langle 8 \rangle$$

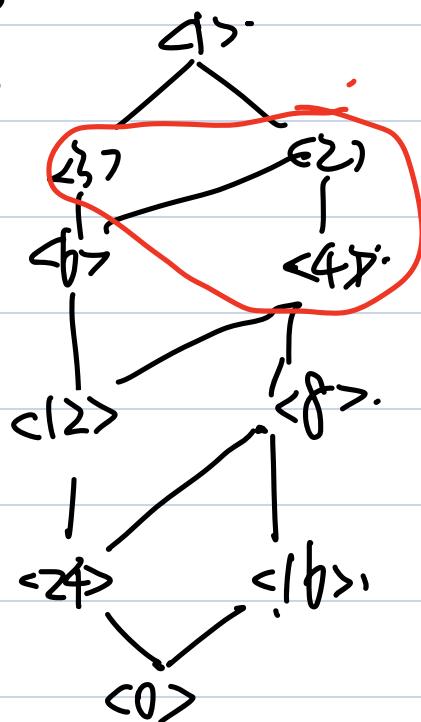
$$6\text{阶} = \langle 0 \rangle$$

$$12\text{阶} = \langle 4 \rangle$$

$$16\text{阶} = \langle 3 \rangle$$

$$24\text{阶} = \langle 2 \rangle$$

$$48\text{阶} = \langle 1 \rangle$$



最大: <2>, <3>

最小: <4>.

最大: <1>

最小: <4>.

上R: <1>

下R: <4>, <8>, <12>, <24>, <16>, <1>

上R₂R: <1>

下R₂R: <4>.

③

$$\begin{matrix} & v_1 & v_2 & v_3 & v_4 \\ v_1 & 1 & 1 & 0 & 0 \\ v_2 & 0 & 0 & 1 & 0 \\ v_3 & 1 & 0 & 0 & 1 \\ v_4 & 0 & 0 & 1 & 0 \end{matrix} \quad \text{消去法}$$

$\leq + \equiv$

1PTI

2PTI

$$\begin{bmatrix} 1 & 2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 1 & 2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 2 & 2 & 0 \\ 1 & 0 & 0 & 1 \\ 1 & 2 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

3PTI

$$\begin{bmatrix} 3 & 2 & 2 & 2 \end{bmatrix} \begin{bmatrix} 1 & 2 & 0 & 0 \end{bmatrix} \quad \begin{bmatrix} 5 & 6 & 5 & 2 \end{bmatrix} \quad \text{4PTI}$$

$$\left[\begin{array}{ccc|c} 1 & 2 & 1 & 0 \\ 2 & 2 & 2 & 1 \\ 1 & 2 & 1 & 0 \end{array} \right] \left[\begin{array}{cc} 0 & 0 \\ 1 & 0 \\ 0 & 1 \end{array} \right] = \left[\begin{array}{cccc} 2 & 2 & 2 & 1 \\ 4 & 4 & 3 & 2 \\ 2 & 2 & 2 & 1 \end{array} \right]$$

(1) 0 · 0 = 2 · 2 ·

(2) 1+3+7+11=22

二+1回

②

	v_1	v_2	v_3	v_4	v_5
v_1	0	0	0	0	1.
v_2	1	0	1	0	0.
v_3	0.	0.	0.	0	1
v_4	1	0	1	0.	0.
v_5	0	1	0	1.	0

简单。

1PTT

2PTT

$$\left[\begin{array}{ccccc} 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \end{array} \right] \left[\begin{array}{ccccc} 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \end{array} \right] = \left[\begin{array}{ccccc} 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 2 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 2 \\ 2 & 0 & 2 & 0 & 0 \end{array} \right] \left[\begin{array}{ccccc} 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \end{array} \right]$$

3PTT

4PTT

$$\left[\begin{array}{ccccc} 1 & 2 & 2 & 2 & 1 \\ 4 & 4 & 3 & 2 & . \\ 2 & 2 & 2 & 1 & . \end{array} \right] \left[\begin{array}{ccccc} 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{array} \right] = \left[\begin{array}{ccccc} 0 & 0 & 0 & 0 & 4 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{array} \right]$$

$$= \begin{bmatrix} 0 & 2 & 0 & 2 & 0 \\ 2 & 0 & -2 & 0 & 0 \\ 0 & 2 & 0 & 2 & 0 \\ 0 & 0 & 0 & 0 & 4 \end{bmatrix} \cdot \begin{bmatrix} 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \end{bmatrix} = \begin{bmatrix} 4 & 0 & 4 & 0 & 0 \\ 0 & 0 & 0 & 0 & 4 \\ 4 & 0 & -4 & 0 & 4 \\ 0 & 4 & 0 & 4 & 0 \end{bmatrix}$$

(2) 0, 2, 0, 0

~~(2)~~ 0 0 4 0

(3). 36

(4)-12

$$(5). \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \end{bmatrix}$$

二十五.

$$\begin{array}{c|ccccc} \textcircled{1} & v_1 & v_2 & v_3 & v_4 \\ \hline v_1 & 0 & 1 & 1 & 1 \\ v_2 & 0 & 0 & 1 & 0 \\ v_3 & 0 & 0 & 0 & 1 \\ v_4 & 0 & 1 & 0 & 0 \end{array} \quad \text{单向.}$$

$$\begin{bmatrix} 0 & 1 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 & 1 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \cdot \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

5条
0

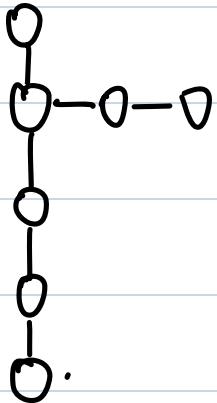
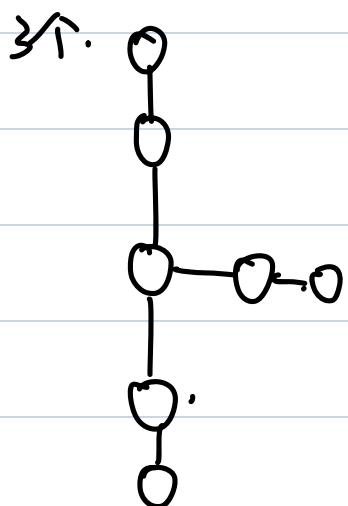
2b. 设有 n 条边. n+1 个点.

$$n \times 2 = (n+1-4) \times 1 + 2 \times 3 + 3$$

$$2n = n - 3 + 9$$

$$n = 6.$$

7个点.



27.

28 # $a: 2$

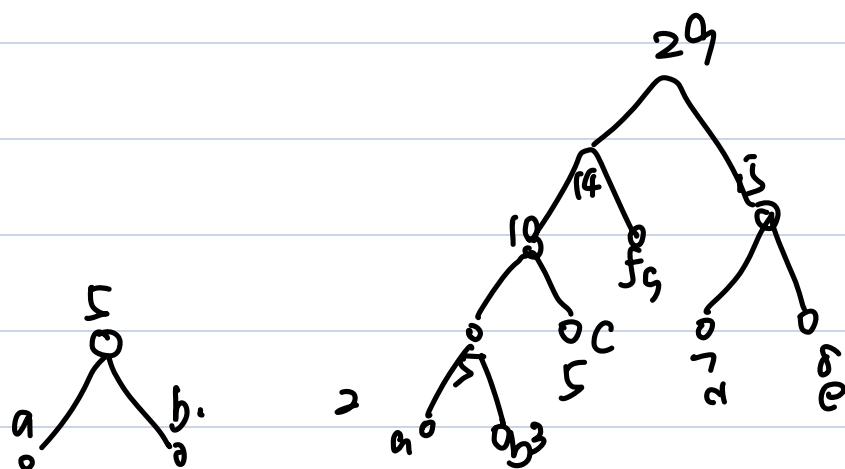
$b: 3$

$c: 5$

$d: 7$

$e: 8$

$f: 9$



$b[000],$

$$\begin{aligned} \text{WTI} &= 2 \times 4 + 3 \times 4 + 5 \times 3 + 5 \times 2 + 7 \times 2 + 8 \times 2 \\ &= 8 + 12 + 15 + 18 + 14 + 16 \\ &= 20 + 25 + 8 + 30 = 83. \end{aligned}$$

$a[0000] \ b[0001] \ c[001] \ f[0,1]$

$d[10] \ e[11].$

{0000, 0001, 001, 01, 10, 11}.

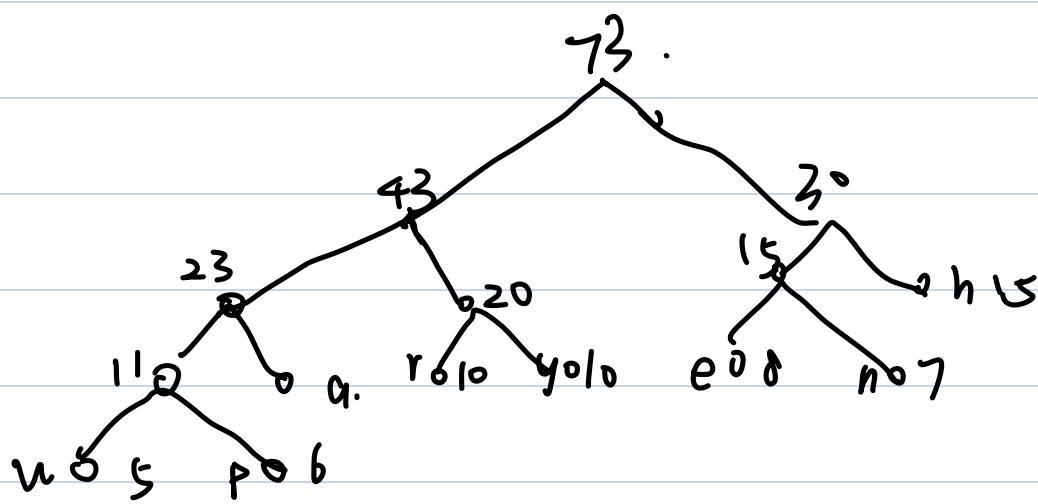
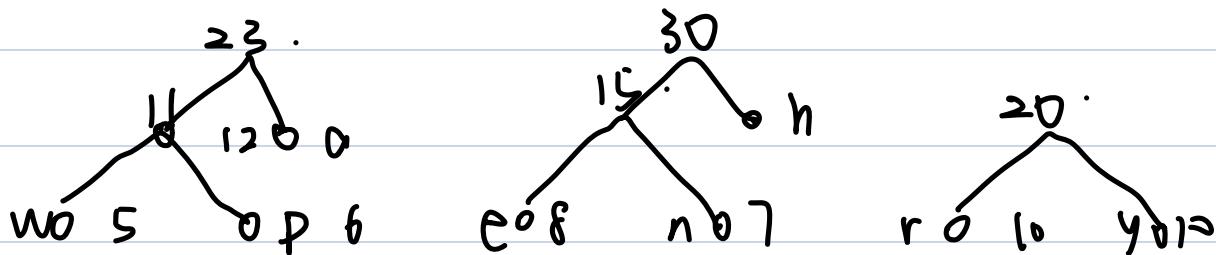
≈ 17

$a: 12$

$b: 15$

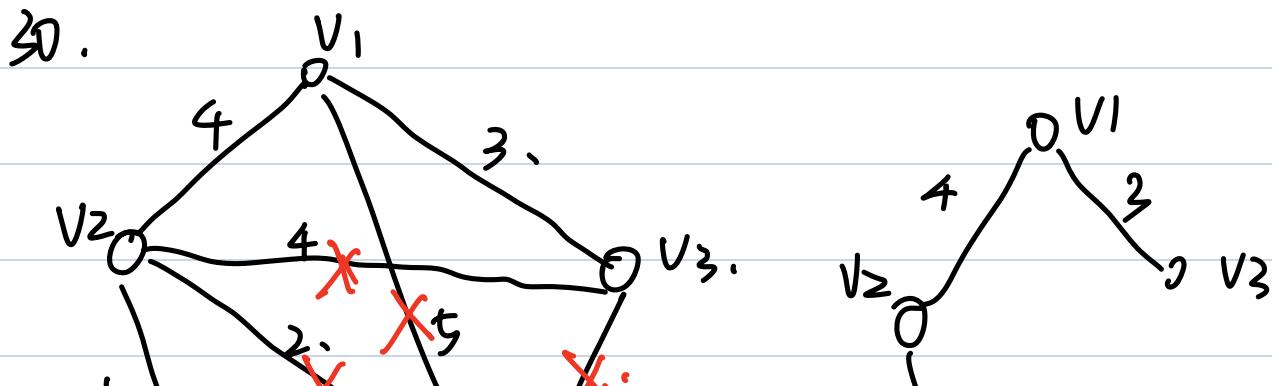
$\boxed{b} \cdot \boxed{5}$

~~e. 8~~ ~~n. 7~~ ~~r. 10~~ ~~y. 12~~.



$[0000, 0001, 0010, 0100, 011] \quad | \quad [100, 101, 11]$.

$[11, 001, 0001, 0001, 010, 101, 100, 0000, 011, 100, 001010]$



9.

