## cs5800 homework1

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- In each of the following situations, indicate whether f = O(g), or  $f = \omega(g)$ , or both (in which case  $f = \theta(g)$ ).
  - (a) There is C and C' make  $C*g(n) \le f(n) \le C'*g(n)$ , so there's f = O(g), f = Omega (g) and f = Theta(g)

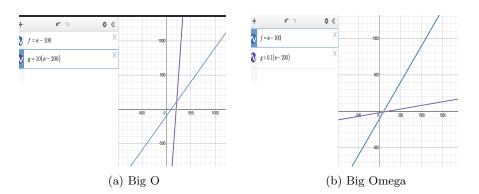
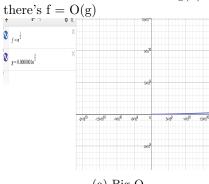


Figure 1: Image for a)

(d) There is C and C' make  $C * g(n) \le f(n) \le C' * g(n)$ , so there's f = O(g), f = O(g) and f = T(g)

(b) No matter how small C is C\*g(n) can always bigger than f(n) when  $n \ge n0$ , so there's f = O(g)



(a) Big O

Figure 2: Image for b)

(c) There is C and C' make  $C*g(n) \le f(n) \le C'*g(n)$ , so there's f = O(g), f = Omega (g) and f = Theta(g)

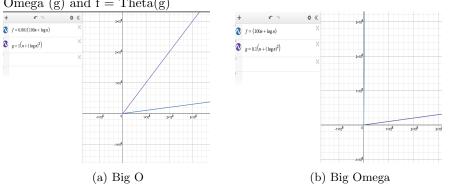


Figure 3: Image for c)

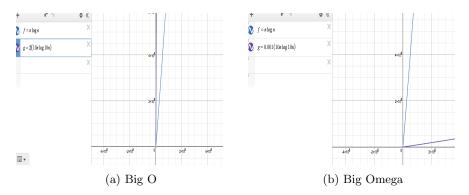


Figure 4: Image for d)

(e) There is C and C' make  $C*g(n) \le f(n) \le C'*g(n)$ , so there's f = O(g), f = Omega (g) and f = Theta(g)

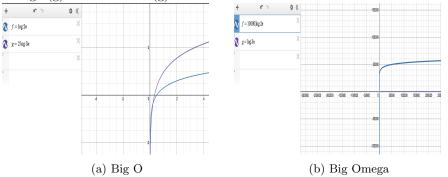


Figure 5: Image for e)

(f) There is C and C' make  $C*g(n) \le f(n) \le C'*g(n)$ , so there's f = O(g), f = Omega (g) and f = Theta(g)

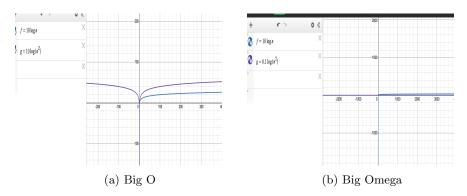


Figure 6: Image for f)

(g) There is C and C' make  $C*g(n) \le f(n) \le C'*g(n)$ , so there's f = O(g), f = Omega (g) and f = Theta(g)

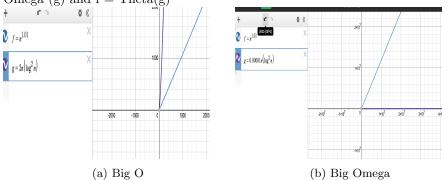


Figure 7: Image for g)

(h) No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\geq$ n0, so there's  $f=\Omega(g)$ 

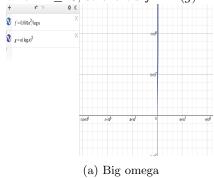
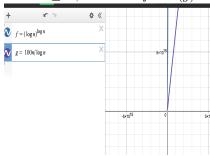


Figure 8: Image for h)

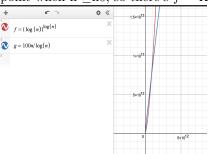
(i) No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\ge$  n0, so there's  $f=\Omega(g)$ 



(a) Big Omega

Figure 9: Image for i)

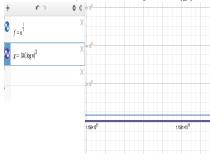
(j) this is j No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\geq$ n0, so there's  $f = \Omega(g)$ 



(a) Big Omega

Figure 10: Image for i)

(k) No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\ge$  n0, so there's  $f=\Omega(g)$ 



(a) Big Omega

Figure 11: Image for k)

(l) No matter how small C is C\*g(n) can always bigger than f(n) when n $\geq$ n0, so there's f = O(g)

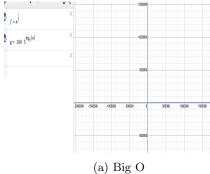


Figure 12: Image for l)

(m) No matter how small C is C\*g(n) can always bigger than f(n) when n≥n0, so there's f = O(g)

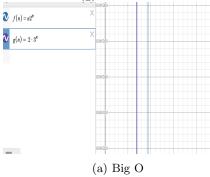


Figure 13: Image for m)

(n) There is C and C' make  $C * g(n) \le f(n) \le C' * g(n)$ , so there's f = O(g), f = O(g), and f = O(g).

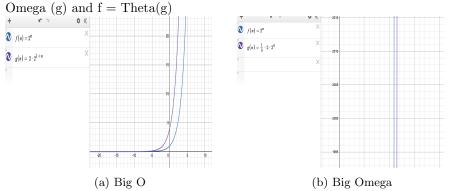


Figure 14: image for n

(o) No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\geq$ n0, so there's  $f=\Omega(g)$ 

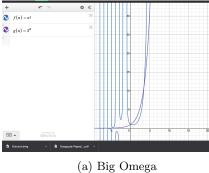


Figure 15: Image for i)

(p) No matter how small C is  $C^*g(n)$  can always bigger than f(n) when  $n \ge n0$ , so there's f = O(g)

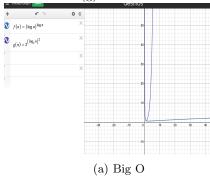


Figure 16: Image for i)

(q) No matter how big C is, C\*g(n) can always smaller than f(n) at some point when n  $\geq$ n0, so there's  $f=\Omega(g)$ 

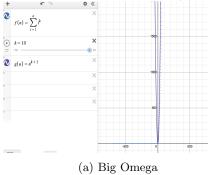


Figure 17: Image for i)